

1	2	$\sum (6+2)$

Blatt 8

(Abgabe am 14. December 2015)

Theoretical Assignment - *Suffix tree construction and application***a - *Suffix tree construction*****i**

If one uses the naive approach to construct the suffix tree of “CTAGTAGCAG”, the result would look like Figure 1.

ii**b - *Main data table of WOTD*****c - *Suffix tree application*****Theoretical Assignment - *Runtime and space complexity of suffix trees*****a**

Assume a text T of length n with n times the letter “a”. If one build a suffix tree for T using WOTD, each node will have just one c-group with all remaining suffixes in it. So evaluating the root node, one has to compute the longest common prefix of n suffixes, of $n - 1$ suffixes for the second node and so on. In each c-group the shortest suffix is of length 1, so for each node \bar{u} there are just $|R_a\{\bar{u}\}|$ numbers of comparisons.

Since T is of length n this will lead to an overall runtime of $\sum_{i=1}^n i = \frac{1}{2}n(n+1)$, which is in $O(n^2)$. \square

b

Suffix links are defined as an auxiliary edge that point from branching node \overline{bw} to the branching node \overline{w} , if it exists and to the root otherwise. If we now defined a suffix link, such that it points into the other direction (i.e. from branching node \overline{w} to branching node \overline{bw}) gives us an efficient method to find maximal unique matches. If we now find a branching node that indicates a unique match (for definition see script) and there is a suffix link that point to another branching node, we know that the current branching node is not a maximal unique match (because it is not left

maximal). So we just have to follow the suffix links until we reach a branching node without an outgoing suffix link and we have found a maximal unique match.

c

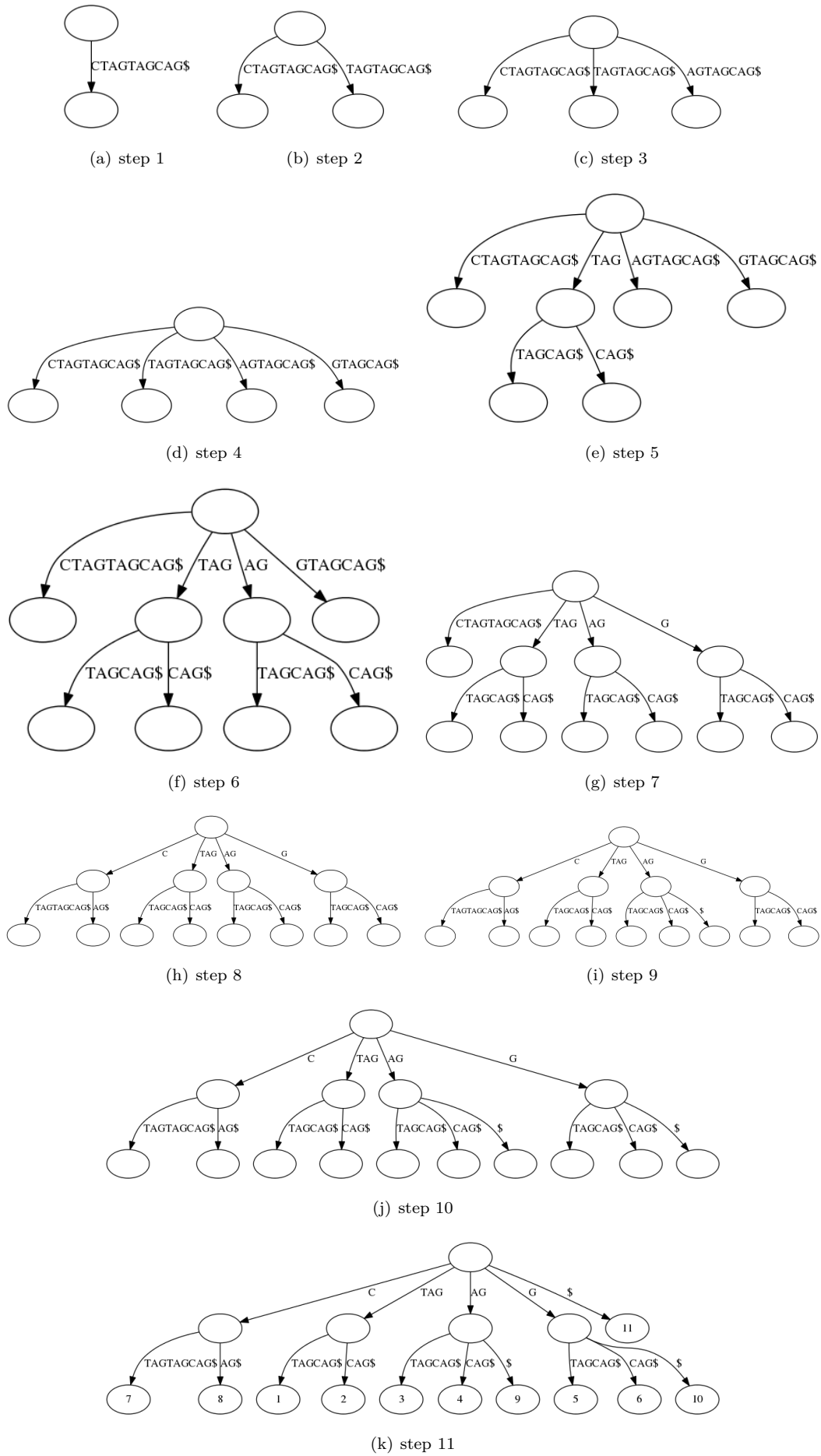


Figure 1: All steps of the naive implementation of suffix tree construction for the string “CTAGTAGCAG”.