

1	2	$\sum (6+2)$

**Blatt 8**

(Abgabe am 14. December 2015)

**Theoretical Assignment - *Suffix tree construction and application*****a - *Suffix tree construction*****i**

If one uses the naive approach to construct the suffix tree of “CTAGTAGCAG”, the result would look like Figure 1.

**ii****b - *Main data table of WOTD*****c - *Suffix tree application*****Theoretical Assignment - *Runtime and space complexity of suffix trees*****a**

Assume a text  $T$  of length  $n$  with  $n$  times the letter “a”. If one build a suffix tree for  $T$  using WOTD, each node will have just one c-group with all remaining suffixes in it. So evaluating the root node, one has to compute the longest common prefix of  $n$  suffixes, of  $n - 1$  suffixes for the second node and so on. In each c-group the shortest suffix is of length 1, so for each node  $\bar{u}$  there are just  $|R_a \bar{u}|$  numbers of comparisons.

Since  $T$  is of length  $n$  this will lead to an overall runtime of  $\sum_{i=1}^n i = \frac{1}{2}n(n+1)$ , which is in  $O(n^2)$ .  $\square$

**b****c**

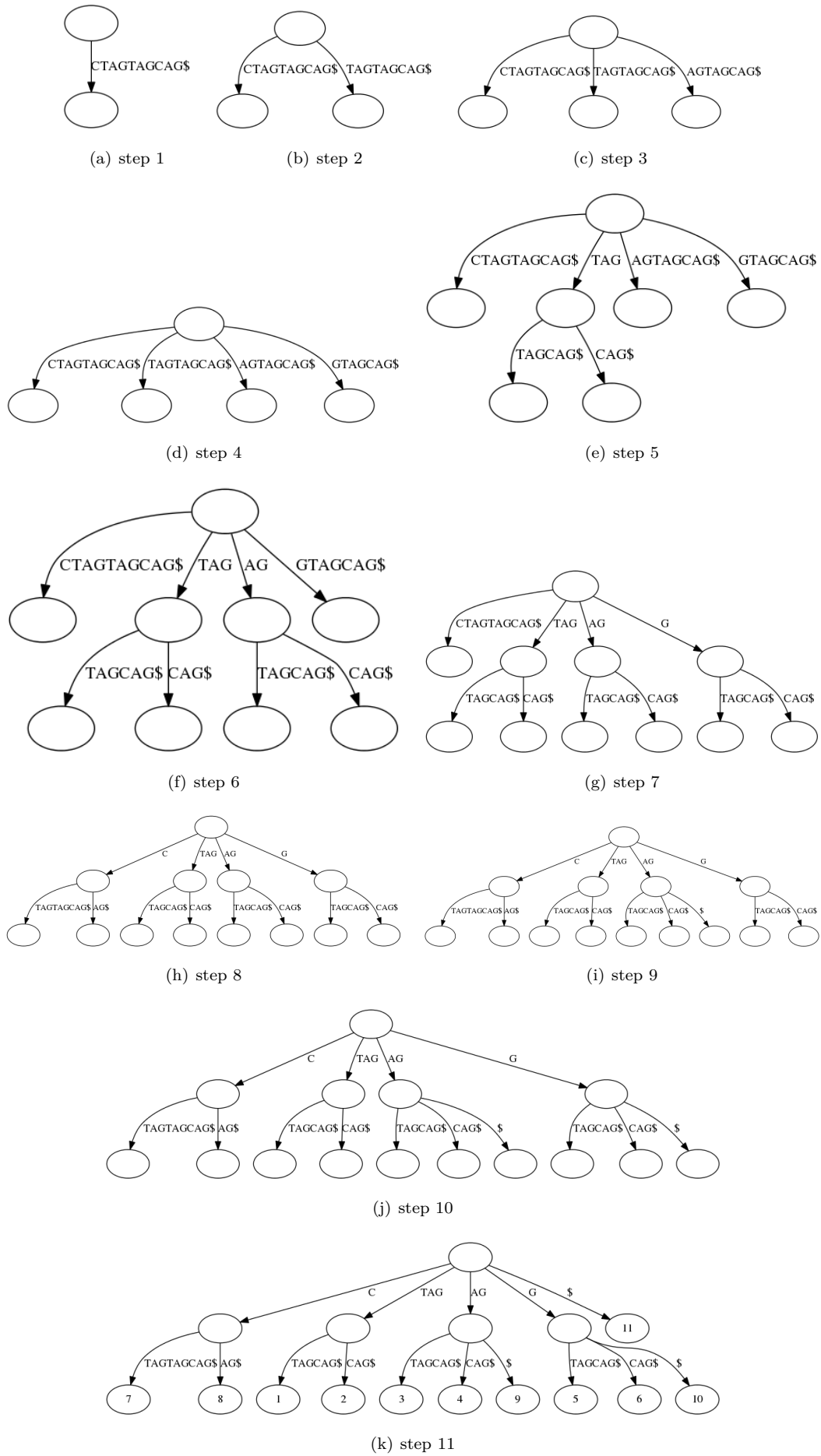


Figure 1: All steps of the naive implementation of suffix tree construction for the string “CTAGTAGCAG”.