# Computer Networks Congestion Control Fall 24-25, CS 3204, Section A

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#### Outline

- Congestion Control
- > TCP RENO

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## Congestion in networks



#### **□** Congestion

Congestion is a situation in a network in which the load on the network, the number of packets sent to the network, is greater than the capacity of the network, the number of packets a network can handle.

#### **□** Congestion control

Congestion control refers to the mechanisms and techniques to control the congestion and keep the load below the capacity [1].

## Congestion in networks...



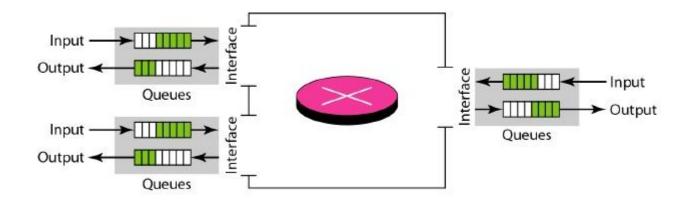
- Congestion in a network or internetwork occurs because routers and switches have queues-buffers that hold the packets before and after processing.
- A router, for example, has an input queue and an output queue for each interface.
- When a packet arrives at the incoming interface, it undergoes three steps before departing:
  - 1. The packet is put at the end of the input queue while waiting to be checked.
  - 2. The processing module of the router removes the packet from the input queue once it reaches the front of the queue and uses its routing table and the destination address to find the route.
  - 3. The packet is put in the appropriate output queue and waits its turn to be sent.

## Congestion in networks...



#### **\*** Two possible scenarios in a router:

- Packet arrival rate > packet processing rate,
   the input queues become longer and longer.
- 2. Packet departure rate < Packet processing rate, the output queues become longer and longer





Terminology

#### ☐ Congestion Window

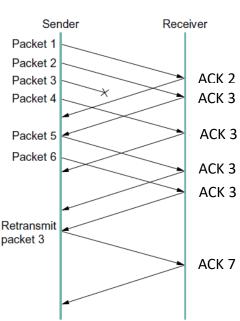
- The number of bytes the sender may have in the network at any time [2]
- $\triangleright$  LastByteSent LastByteAcked  $\leq$  cwnd

#### ☐ Sending Rate

- The congestion window size divided by the round-trip time (RTT<sup>1</sup>) of the connection [2].
- \* By adjusting the value of cwnd, the sender can therefore adjust the rate at which it sends data into its connection [3].

#### **Terminology**

- ☐ Loss Event
  - **Timeout** 
    - Sender does not receive any acknowledgement within a predefined interval
  - 3 duplicate Acknowledgement (3 duplicate ACK)
    - Reception of the same ACK four times.
    - Indicates that the channel is not congested that much as the receiver is still receiving segments.
    - Indicates that the segment 3 is lost as it receives the dupli ACK 3
    - Receiver is receiving segment out of order.



TCP Reno



- ☐ Steps of congestion window control
  - Slow Start
  - Congestion Avoidance

TCP Reno



#### **♦**Slow Start

Follows a greedy approach.

Starts sending data of size equals maximum sized segment (MSS).

If it receives an ACK for the previous transmission by next RTT  $\rightarrow$  it sends 2 MSS.

After next RTT, if it receives two ACKs → it sends 2+2=4 MSS in next round.

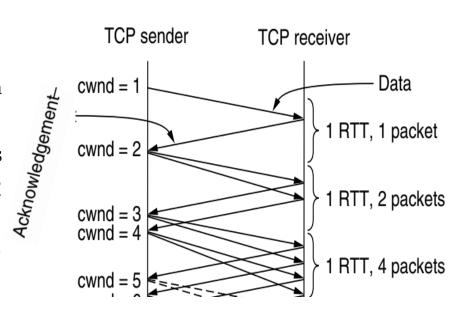


Fig. 1 Slow start from an initial congestion window of one segment [2].

TCP Reno



#### **❖**Slow Start (cont.)

That is, in each round, N+M MSS in a round if it sends N MSS in the last transmission round and receives M ACK before the current round.

If it receives ACK for all the N MSS (previously sent), it sends N+N MSS this round.

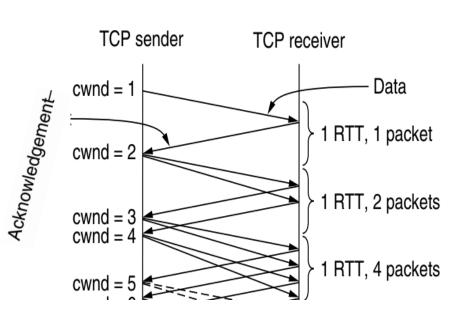


Fig. 1 Slow start from an initial congestion window of one segment [2].

TCP Reno



#### **Congestion Avoidance**

How long does cwnd continue to be increased exponentially?

 $if \ cwnd \ge threshold$ 

Increase cwnd by 1 per transmission round That is, increase cwnd linearly 1

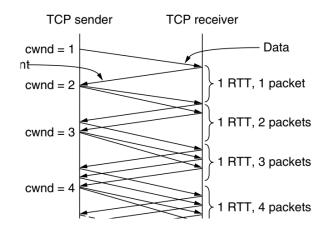


Fig. 2 Linear increase of cwnd

TCP Reno



## How long does cwnd continue to be increased linearly?

- Until one of two incidents happen
  - Timeout
  - 3 duplicate ACK
- Timeout

$$Threshold = \frac{cwnd}{2}$$

cwnd = 1 MSS

Start slow start

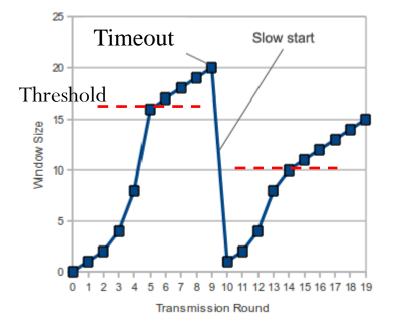


Fig. 3 Timeout



TCP Reno

#### O 3 duplicate ACK

$$Threshold = \frac{cwnd}{2}$$

cwnd = Threshold

Start congestion avoidance

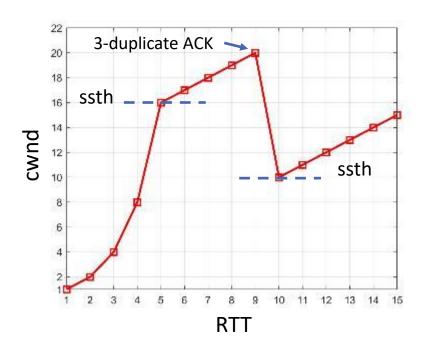


Fig. 4 3 duplicate ACK

TCP Reno



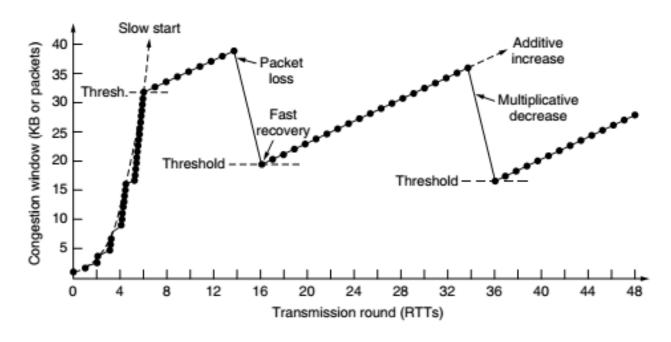
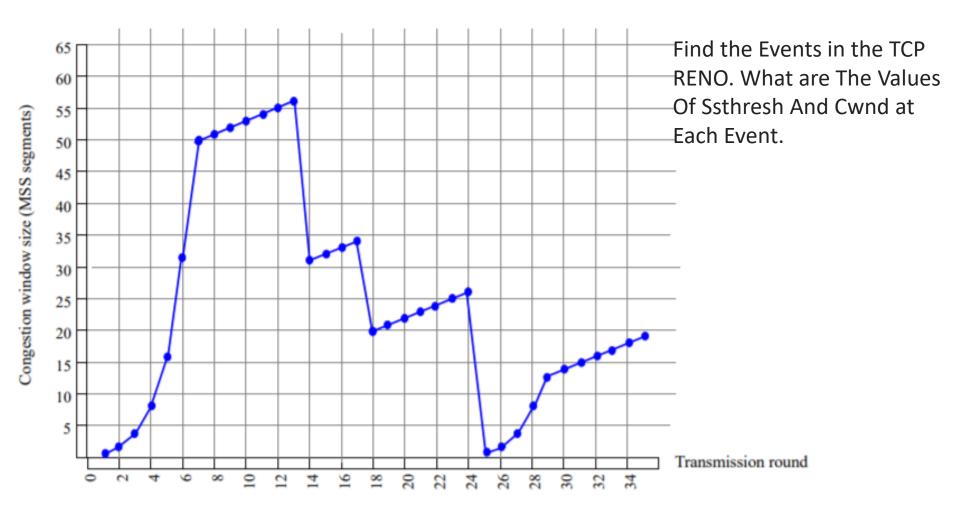


Fig. 5 TCP reno example

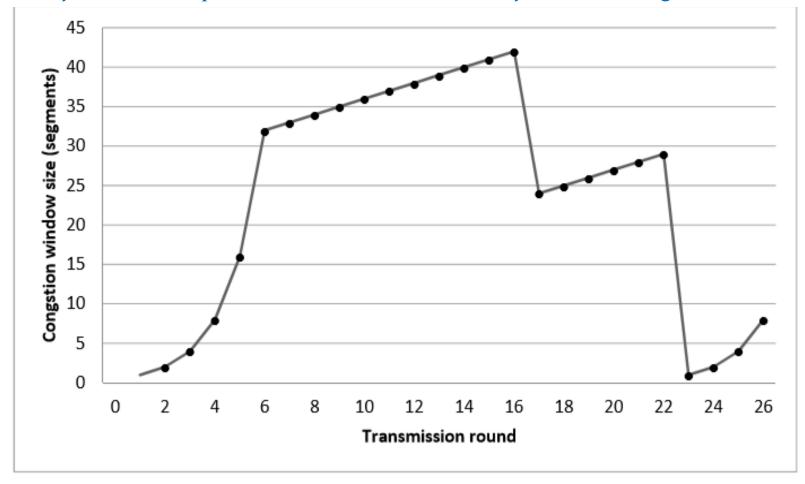
(3 points each) Consider the figure below. Assuming TCP Reno is the protocol experiencing the behavior shown above, answer the following questions. In all cases, you should provide a short discussion justifying your answer.



## Task 2



Identify Timeout/3 duplicate ACK incidence. Also identify slow start, congestion avoidance period.



### Flow control

The TCP stores the data that needs to be sent in the send buffer and the data to be received in the receive buffer. Flow control makes sure that no more packets are sent by the sender once the receiver's buffer is full as the messages will be dropped and the receiver won't be able to handle them. In order to control the amount of data sent by the TCP, the receiver will create a buffer which is also known as Receive Window.



#### Flow control



- Send window
  - O Sequence of byte numbers of which some are sent but yet to be acknowledged while others are waiting to be sent
  - Maintained by sender
- Receive window
  - Sequence of byte numbers which are expected to be received
  - Maintained by receiver

#### References



- [1] B. A. Forouzan, *Data Communication and Networking*, 4<sup>th</sup> ed., Tata McGraw Hill Companies, Inc., New Delhi, 2010, pp. 385.
- [2] A. S. Tanenbaum and D. J. Wetherall, *Computer Networks*, 5<sup>th</sup> ed., Pearson Education India, 2013, pp. 571-572.
- [3] J. F., Kurose and K. W. Ross, *Computer Networking: A Top-Down Approach*, 7<sup>th</sup> ed., Pearson Education, Inc., USA, 2017, pp. 274-277.
- [4] D. Medhi and K. Ramassamy, Network Routing Algorithms, Protocols and Archtectures, 2<sup>nd</sup> ed., Elsevier Inc., 2018, pp. 607-608.
- [5] D. Runemalm, D. M. Sarwar and M. Shalbaf, "Decreasing the Hybrid-ARQ bandwidth overhead through the Multiple Packet NAK (MPN) protocol", https://www.researchgate.net/publication/267554658, [Accessed: April. 22, 2020].
- [6] B. A. Forouzan, *TCP/IP Protocol Suite*, 4<sup>th</sup> ed., McGraw Hill Companies, Inc., USA, 2010, pp. 457-462.

#### **Recommended Books**



- **1. Data Communications and Networking**, *B. A. Forouzan*, McGraw-Hill, Inc., Fourth Edition, 2007, USA.
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#### References

> This is the Provided Materials.

> Online Website Research.

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