

- * $\text{JNZ} \rightarrow$ controls the loop
 - * Labels end with colon (:)
 - * CMP is just like SUB
 - * Signed Conditional Jumps \rightarrow JG or JNLE ($ZF=0, SF=OF$)
 (ZF, OF, SF) JGE or JNL ($SF=OF$)
 JL or JNGE ($SF \neq OF$)
 JLE or JNG ($ZF=1$ or $SF \neq OF$)
 - * Unsigned \rightarrow JA or JNBE ($ZF=0, CF=0$)
 (ZF, OF) JAF or JNB ($CF=0$)
 JB or JNAE ($CF=1$)
 JBE or JNA ($ZF=1$ or $CF=1$)
 - * Sign bit of a byte in a character is always zero
 - * Jump is primitive.



Single-Flag Jumps

JE or JZ	Jump if Equal Jump if equal to Zero	ZF = 1
JNE or JNZ	Jump if Not Equal Jump if Not Zero	ZF = 0
JC	Jump if Carry	CF = 1 CF = 0
JNC	Jump if no Carry	CF=0
JO	Jump if Overflow	CF=1 or ZF = 1
JNO	Jump if No Overflow	OF=1
JS	Jump if Sign Negative	SF = 1
JNS	Jump if Non-Negative Sign	SF =0
JP/JPE	Jump if Parity Even	PF=1
JNP/JPO	Jump if parity Odd	PF=0

□ What is Loop?

⇒ A loop is a sequence of instructions that is repeated.

There are 3 types of loop -

1. For loop
2. While loop
3. Repeat loop

□ What is for-loop?

⇒ In For loop, the loop statements are repeated a known numbers of times.

```
FOR loop-count times DO  
    Statements  
END-FOR
```

* Counter → Register CX
Decremented

* For loop execute → At least once

If CX = 0 , the Loop instruction causes CX to be decremented to FFFF h

→ To prevent this, JCXZ (Jump if CX is zero) used. before the loop.

JCXZ destination_label.

□ What is WHILE loop?

⇒ WHILE Loop depends on a condition. The loop is executes as long as the condition is true.

- * A while loop checks the terminating condition at the top of the loop.
- ** WHILE_label is used because WHILE is a reserved word.

WHILE condition DO

 statements

END WHILE

□ What is Repeat loop?

⇒ In a Repeat ... UNTIL loop, the statements are executed, and then the condition is checked.

- * If true → the loop terminates
- If false → Control branches to the top of the loop

REPEAT

 statements

UNTIL

 Condition

□ Difference between WHILE and REPEAT loop?

⇒ WHILE loop

1. Check condition before each iteration
2. If the condition is initially false, may not execute.
3. Code size is a little shorter. longer .
4. Has two jumps : Conditional JMP at the top
 — a JMP at the bottom,

REPEAT loop

1. Check condition after each iteration
2. Executes at least once.
3. Code size is a little longer. shorter
4. Has only one conditional jump .

□ Logic , shift and rotate → to do binary and hexadecimal (MCQ)

Lecture 7(a)

□ Logic instruction: AND, OR, XOR, NOT

→ Converting a lowercase letters to uppercase

→ Determining an 'Even' or 'ODD' numbers.

** Shifting is faster than Multiplication and Division (M&Q) **

*** Memory - memory operations are not allowed.

** Source bit pattern known as mask.

Effect on flags → SF, ZF, PF reflect the result

AF is undefined

CF, OF = 0

** Instead of SUB AL, 30H we can use AND AL, 0FH

→ Clear the high bits of AL

□ Converting Uppercase - Lowercase

→ SET / OR operation.

A → 41 → 0100 0001

a → 61 → 0100 0001

Set → 0010 0000 [Replace with 1, and remaining will be 0]

= 20 h = Mask

□ Converting lowercase - Uppercase

→ Clears / AND operations

a → 61 → 0110 0001

A → 41 → 0100 0001

Clears → 1101 1111 [Replace with 0 and remaining will be 1]
 = 0DF h
 ↳ Alphabet এর আগে '0' রয়েছে (MCQ)*

□ Clearing a Registers (MCQ)*

MOV AX, 0
 SUB AX, AX
 XOR AX, AX } Same

CMP CX, 0 ; To check the sign contents.

- * Test instructions performed AND operation.
- * Test instructions set the status flags.

OF, CF = 0

AF = Undefined

SF, ZF, PF reflect the result.

** For a shift → the bits shifted out are lost

** For intel processors → Allow the use of an 8-bit constant.

□ SHL

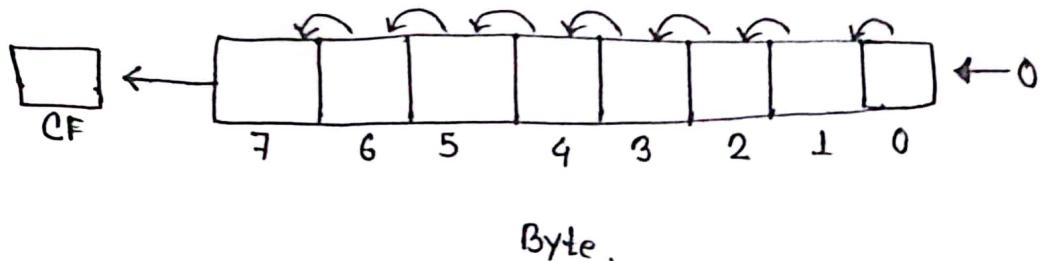
→ a 0 is shifted in to the right most position and msb is shifted into CF.

SHL destination, CL

Example: Suppose DH contains 8Ah and CL contains 3.

SHL DH, CL

⇒ DH → $\begin{matrix} 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{matrix}$ CF=1
 $\begin{matrix} 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \end{matrix}$ CF=0
 $\begin{matrix} 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \end{matrix}$ CF=0
 $\begin{matrix} 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \end{matrix}$ CF=0
= 50H



Multiplication by left shift

AL contains $5 = 0000\ 0101$

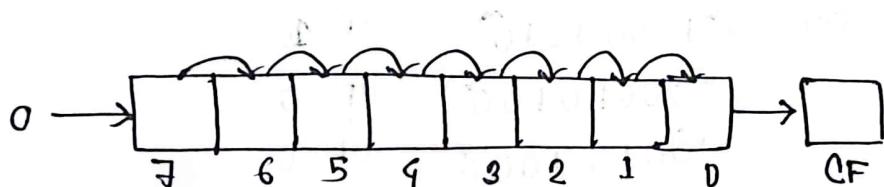
A left shift gives $0000\ 1010 = 10$; Doubling its value

Another " " " $0001\ 0100 = 20$; Doubled again.

** Overflow flag are not reliable for this.

SHR (Right shift)

- A 0 is shifted into the msb and the rightmost bit is shifted to the CF.



Lecture - 7(a) : Logic Instructions

Date : / /

1. AND Operation

$$\begin{array}{r} 10101010 \\ 11100000 \\ \hline 00000100 \end{array}$$

2. OR operation

$$\begin{array}{r} 01010101 \\ 00001111 \\ \hline 01011111 \end{array}$$

3. XOR operation [1]

$$\begin{array}{r} 10101010 \\ 00001111 \\ \hline 10110101 \end{array}$$

4. NOT Operation (1's complements)

$$\begin{array}{r} 10101010 \\ \rightarrow 01010101 \end{array}$$

- AND destination, source
- OR destination, source
- XOR destination, source

* The bits in the mask are selected to modify the corresponding destination bits when the instruction executes.

AND → Clear

- 0 mask bit clears
- 1 mask bit preserves

OR → Set

- 1 mask bit sets
- 0 mask bit preserves

XOR → Complement / Change

- 1 mask bit complements
- 0 mask bit preserves

Example: Clear the sign bit AL while leaving the other bit unchanged.

⇒

$$\begin{array}{r} 10111011 \\ 01111111 \rightarrow \text{mask} = 7Fh \\ \hline 00111011 \end{array}$$

mask value = 0

Thus, AND AL, 7Fh

Example: Set the most significant bit and least significant bits of AL while preserving the other bits.

⇒

$$\begin{array}{r} 10111011 \\ 10000001 \rightarrow \text{mask} = 81h \\ \hline 10111011 \end{array}$$

mask value = 1

Thus, OR AL, 81h

Example: Change the sign bit of DX

$$\begin{array}{r} 10111011 \\ 10000000 \rightarrow 800h \\ \hline 10111011 \end{array}$$

mask value = 1

Thus, XOR DX, 800H

Lecture : 7(b) → Rotate Instruction.

- Suppose DH contains 8A h and CL contains 2. What are values of DH and CF after the instruction SHR DH, CL is executed.

⇒ The binary value of DH = 1000 1010
 $\begin{array}{r} \textcircled{1} 000 \quad \textcircled{0} 101 \\ \textcircled{0} 100 \quad \textcircled{1} 010 \\ \textcircled{0} 010 \quad \textcircled{0} 010 \end{array}$
 $= 22 \text{ H}$

CF = 01

- DH = 8AH , CL = 3 ; SHL DH, CL

DH → 1000 1010
 $\begin{array}{r} \textcircled{1} 000 \quad \textcircled{1} 010 \\ \textcircled{0} 001 \quad \textcircled{0} 100 \\ \textcircled{0} 010 \quad \textcircled{1} 000 \\ \textcircled{0} 101 \quad \textcircled{0} 000 \end{array}$
 $CF = 000$
 $= 50 \text{ H}$

- SAR: Shift Arithmetic operates like SHR with 1 difference.
 → The msb retains its original value.

SAR destination, 1
 SAR destination, CL

Division by SHR

→ For even no. → Divided by 2

→ For odd no. → Rounds to the nearest integers.

exⁿ: BL = 00000101 = 5

SHR BL , BL = 0000 0010 = 2

* For left shift → Multiply by 2

Difference between Signed and Unsigned Division →

⇒ For unsigned interpretation → SAR should be used.

For signed interpretation → SAR must be used

↳ It preserves the sign.

Problem: Use right shifts to divide the unsigned numbers 65143 by 4, put the quotient in AX.

⇒ MOV AX, 65143

MOV CL, 2 ; two Right shifts [2+2]

SHR AX, 2 ; Dividend is unsigned.

Example: If AL contains -15, give the decimal value of AL after SAR AL,1 performed.

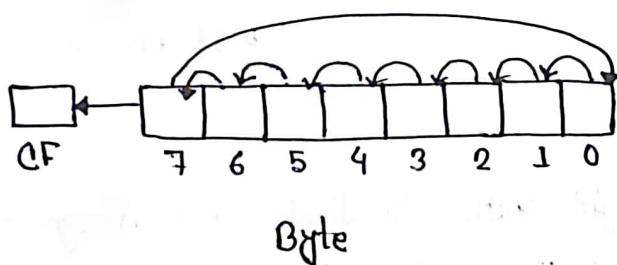
⇒ SAR AL,1 ; CL = 1

Old AL = -15

New AL = $-15/2 = -7.5 \approx -8$

□ ROL (Rotate Left)

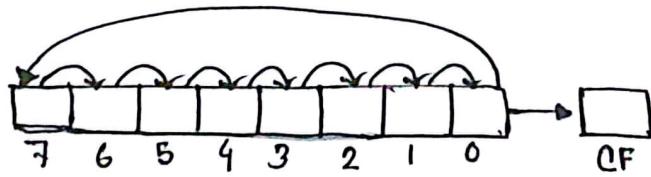
- Shifts bits to the left
- The msb shifts into the rightmost bit.
- The CF also gets the bit shifted out of the MSB.



ROL destination, CL

□ ROR (Rotate Right)

- Shifts bits to the right
- The lsb/rightmost bit is shifted into the msb and also into the CF



ROR destination, CL

- * ROL and ROR used to inspect the bits in a byte without changing the contents. (Mea)*

Example: Use ROL to count the numbers of 1 bits in BX, without changing BX. Put the answers in AX.

⇒ XOR AX, AX
MOV CX, 16

TOP:

ROL BX, 1
JNC NEXT

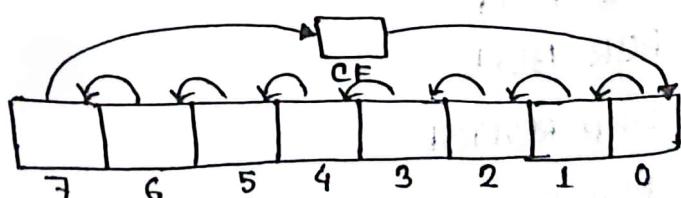
INC AX

NEXT:

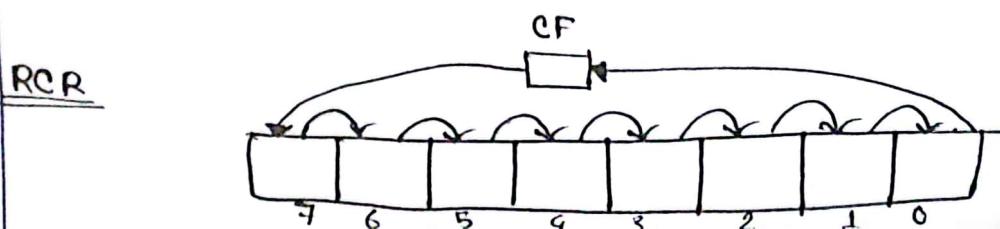
LOOP TOP

D RCL (Rotate Carrying Left)

- Shifts the bits to the left
- The msb is shifted into the CF and the previous value of CF is shifted into the rightmost bit.



RCL destination, CL



Example:

Suppose DH contains 87h, CF=1 and CL contains 3. What are the values of DH and CF after the instruction !

RCR DH, CL is executed ?

\Rightarrow

DH	CF	SF = 1
1000 1010	1	ZF = 0
11000 101	0	OF = 1
01100010	1	
1011 0001 = B1 h	0	

Example:

If AL contains 11011100, we want to make it 00111011

\Rightarrow

MOV CX, 8

REVERSE :

SHL AL, 1

RCR BL, 1

Loop REVERSE

MOV AL, BL

Lecture -09

□ MUL → Multiply for unsigned multiplication

IMUL → For signed multiplication

* Byte form → If two bytes are multiplied → 16 bits (word)

Word form → " " Words " " → 32 bits (doubleword)

** The source cannot be a constant (MCQ)*

* Doubleword product : DX : AX → least significant 16 bits
→ Most significant 16 bits.

Effect on Status Flags

→ SF, AF, ZF, PF → Undefined

→ CF/OF → MUL : CF/OF=0 if the upper half of the result is 0
CF/OF=1 ; Otherwise.

→ IMUL : CF/OF=0 ; If the bits of the upper half
are the same as the sign bit of
the lower half.

CF/OF=1 ; Otherwise .

Example

Ax contains 1 and Bx contains FFFF h.

Instructions	Hex Product	Dx	Ax	CF/OF
MUL BX	0000FFFF	0000	FFFF	0
IMUL BX	FFFFFFFF	FFFF	FFFF	0

→ Ax = 1, Bx = -1
2's complement → so product is -1

Example

Ax contains FFFF h and Bx contain FFFF h

Instruction	Decimal	Hex	Dx	Ax	CF/OF
MUL BX	65535×65535	FFFED001	FFFE	0001	1
IMUL BX	1	00000001	0000	0001	0

→ Ax and Bx both contain -1, so the product is 1

Example

Ax contains 0FFF h

MUL AX	16769025	00FFE001	00FF	E001	1
IMUL BX	16769025	00FF E001	00FF	E001	1

As, CF/OF = 1, this means the product is too big to fit in Ax.

Signed versus Unsigned Multiplication

Example of Multiplication



Example 9.4 Suppose AX contains 0100h and CX contains FFFFh:

<i>Instruction</i>	<i>Decimal product</i>	<i>Hex product</i>	<i>DX</i>	<i>AX</i>	<i>CF/OF</i>
MUL CX	16776960	00FFFF00	00FF	FF00	1
IMUL CX	-256	FFFFF00	FFFF	FF00	0

For MUL, the product FFFF00 is obtained by attaching two zeros to the source value FFFFh. Because the product is too big to fit in AX, CF/OF = 1.

For IMUL, AX contains 256 and CX contains -1, so the product is -256, which may be expressed as FF00h in 16 bits. DX has the sign extension of AX, so CF/OF = 0.

Signed versus Unsigned Multiplication

Example of Multiplication



Example 9.5 Suppose AL contains 80h and BL contains FFh:

<i>Instruction</i>	<i>Decimal product</i>	<i>Hex product</i>	<i>AH</i>	<i>AL</i>	<i>CF/OF</i>
MUL BL	128	7F80	7F	80	1
IMUL BL	128	0080	00	80	1

For byte multiplication, the 16-bit product is contained in AX.

For MUL, the product is 7F80. Because the high eight bits are not 0, CF/OF = 1.

For IMUL, we have a curious situation. 80h = -128, FFh = -1, so the product is 128 = 0080h. AH does not have the sign extension of AL, so CF/OF = 1. This reflects the fact that AL does not contain the correct answer in a signed sense, because the signed decimal interpretation of 80h is -128.

Signed versus Unsigned Multiplication

Simple Applications of MUL and IMUL



Example 9.5 Suppose AL contains 80h and BL contains FFh:

<i>Instruction</i>	<i>Decimal product</i>	<i>Hex product</i>	<i>AH</i>	<i>AL</i>	<i>CF/OF</i>
MUL BL	128	7F80	7F	80	1
IMUL BL	128	0080	00	80	1

For byte multiplication, the 16-bit product is contained in AX.

For MUL, the product is 7F80. Because the high eight bits are not 0, CF/OF = 1.

For IMUL, we have a curious situation. 80h = -128, FFh = -1, so the product is 128 = 0080h. AH does not have the sign extension of AL, so CF/OF = 1. This reflects the fact that AL does not contain the correct answer in a signed sense, because the signed decimal interpretation of 80h is -128.

MULTIPLICATION

Simple Applications (1) of MUL and IMUL



Example 9.6 Translate the high-level language assignment statement $A = 5 \times A - 12 \times B$ into assembly code. Let A and B be word variables, and suppose there is no overflow. Use IMUL for multiplication.

Solution:

MOV AX, 5	; AX = 5
IMUL A	; AX = 5 × A
MOV A, AX	; A = 5 × A
MOV AX, 12	; AX = 12
IMUL B	; AX = 12 × B
SUB A, AX	; A = 5 × A - 12 × B

Division

DIV (divide) → Unsigned division

IDIV (Integer-divide) → Signed division

DIV divisor

IDIV divisor

→ 8 bit registers [not be a constant]

* 16-bit dividend → AX

* 8-bit quotient → AL

* 8-bit remainders → AH

DX: AX

↓
Remainders → Quotient.

** All status flags are undefined (MCQ)*

Divide overflow occurs —

1. When the quotient doesn't fit in the specified destination (AL or AX)
2. If the divisor is much smaller than dividend.

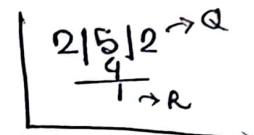
** Divisor and Dividend positive \Rightarrow DIV, IDIV result will be same.

Example

Suppose DX contains 0000h, AX contains 0005h and BX contains 0002h.

$$\Rightarrow \text{Dividend: } DX : AX \rightarrow 0000\ 0005 = 5$$

$$\text{Divisor: } BX \rightarrow 0002 = 2$$



Instruction	Decimal Quotient	Decimal Remainders	AX	DX
DIV BX	2	1	0002	0001
IDIV BX	2	1	0002	0001

Example

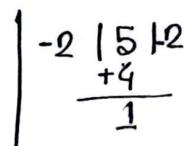
$$DX : AX \rightarrow 0000\ 0005 = 5$$

$$BX: FFFEh \rightarrow \begin{array}{r} 1111\ 1111\ 1111\ 1110 \\ 0000\ 0000\ 0000\ 0001 \\ +1 \\ \hline 0000\ 0000\ 0000\ 0010 \end{array} \quad (\text{1's complement})$$

$$= -2$$

	AX	DX
DIV BX	0	5
IDIV BX	-2	1

$$\left. \begin{array}{l} FFFF = -1 \\ FFFE = -2 \\ FFFD = -3 \\ FFFC = -4 \\ FFB = -5 \\ FFA = -6 \end{array} \right\}$$



Example 9.10 Suppose DX contains FFFFh, AX contains FFFBh, and BX contains 0002.

<i>Instruction</i>	<i>Decimal quotient</i>	<i>Decimal remainder</i>	<i>AX</i>	<i>DX</i>
IDIV BX	-2	-1	FFFFE	FFFF
DIV BX	DIVIDE OVERFLOW			

For IDIV, $DX:AX = FFFFFFFFh : FFFBh = -5$, $BX = 2$. -5 divided by 2 gives a quotient of -2 = FFFEh and a remainder of -1 = FFFFh.

For DIV, the dividend $DX:AX = FFFFFFFFh = 4294967291$ and the divisor = 2. The actual quotient is $2147483646 = 7FFFFFFEh$. This is too big to fit in AX, so the computer prints DIVIDE OVERFLOW and the program terminates. This shows what can happen if the divisor is a lot smaller than the dividend.



Signed versus Unsigned Division

Division Example

Example 9.11 Suppose AX contains 00FBh and BL contains FFh.

<i>Instruction</i>	<i>Decimal quotient</i>	<i>Decimal remainder</i>	<i>AX</i>	<i>AL</i>
DIV BL	0	251	FB	00
IDIV BL	DIVIDE OVERFLOW			

For byte division, the dividend is in AX; the quotient is in AL and the remainder in AH.

For DIV, the dividend is 00FBh = 251 and the divisor is FFh = 256. Dividing 251 by 256 yields a quotient of 0 and a remainder of 251 = FBh.

For IDIV, the dividend is 00FBh = 251 and the divisor is FFh = -1. Dividing 251 by -1 yields a quotient of -251, which is too big to fit in AL, so the message DIVIDE OVERFLOW is printed.

□ Find the max value among 3 numbers

• data

a db ?

b db ?

c db ?

d db "Enter 3 numbers: \$"

e db "Is the greatest \$"

• Code

Main Proc

```
mov ax, @data
mov ds, ax
```

```
Mov ah, 9
lea dx, d ; string
int 21 h
```

```
Mov ah, 1
int 21 h ; 1st input
```

mov a, al

```
int 21 h
mov b, al ; 2nd input
```

```
int 21 h
mov c, al ; 3rd input
```

mov ah, 2 ; Display

mov dl, 10

int 21 h

mov dl, 13

int 21 h

Cmp bl, al ; Cmp A
Jg greater

Greatest :

```
Cmp bl, cl ; Cmp A
Jg B-greater
Jl C-greater
```

Lesser :

```
Cmp bl, cl ; Cmp A
Mov cl, cl ; Cmp A
Cmp cl, al
```

```
Jg C-greater
Jl A-greater
```

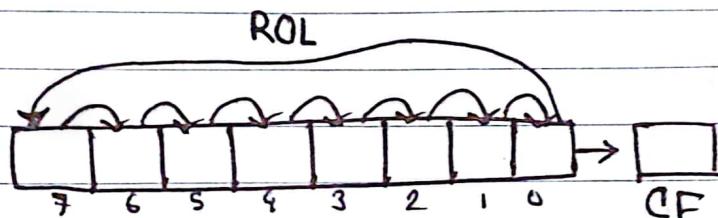
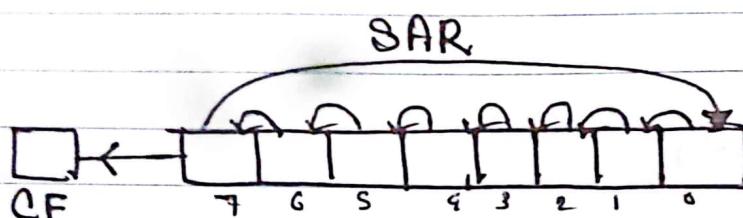
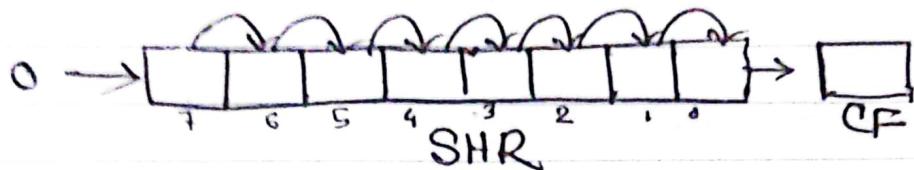
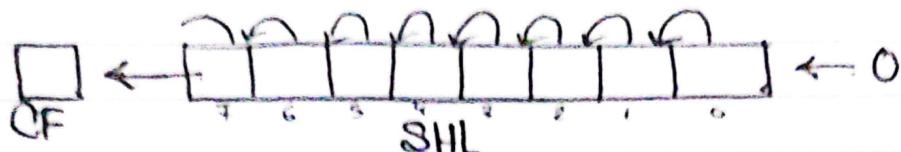
A-greater :

```
Mov ah, 2
Mov dl, a
Int 21 h
Mov ah, 9
Lea dx, e
Int 21 h
Jmp exit .
```

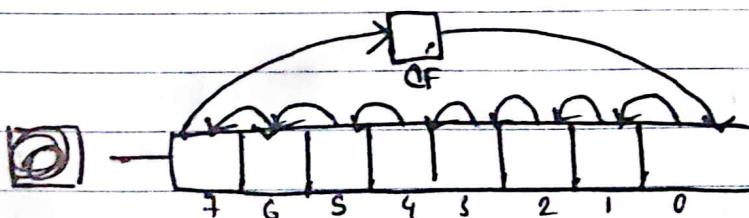
B-greater :

```
Mov ah, 2
Mov dl, b
Int 21 h
Mov ah, 9
Lea dx, e
Int 21 h
Jmp exit
```

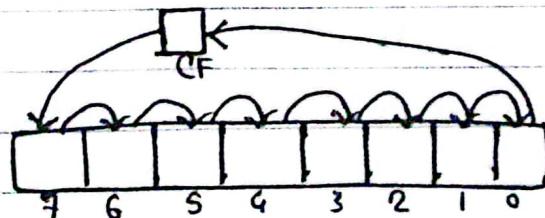
*SHL + 2X
SHR 2%*



ROR



RCL



RCR

Even-ODD

- Model small
- Stack 100 h
- data
 - A dw ?
 - B DB "Enter a numbers: \$"
 - C DB "The numbers is even \$"
 - D DB " " . . . ODD \$"
- code

Main PROC

```
move ax, @data.
mov ds, ax
```

```
MOV ah, 9
lea dx, B
int 21 h
```

```
mov ah, 1
int 21 h
sub al, 30 h
```

```
mov ah, 00h
mov A, ax
ADD ax, 1
JZ Even - lea dx, odd
JMP print
```

Print:

```
mov ah, 9h
int 21 h
```

```
mov ah, 9ch
int 21 h
main endp.
```

Even :

lea dx, Even

Even/Odd - Counting characters

- Model SMALL
- Stack 100h
- Data
- Code

Main Proc

```
mov ax, @data
mov ds, ax
mov cx, 0
```

Read-Char:

```
Mov Ah, 1
int 21h
cmp al, 0Dh ; with carriage return.
JE done.

inc cx
jmp read-char
```

done:

```
Mov dx, cx ; Move the count to dx
Add dx, 40'
Mov ah, 2
Int 21h
```

```
Mov ah, 4ch
Int 21h
main endp.
```

end.