ECU178 Computer Science: 210CT - Programming, Algorithms and Data Structures Portfolio

Due on Monday, December 15th, 2014

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Item 1: Week 3 - Linear Search and Duplicate Finder

1. Pseudocode for linear search

This Simple Algorithm demonstrates how to perform a linear search.

Input: This algorithm takes a populated array A and a value to search for v, as parameters. **Output:** The Algorithm is a boolean type and returns either True or False respective of whether the v was found in the list or not.

Algorithm 1 LinearSearch

```
1: procedure BOOL LINEARSEARCH(v, A[])
2: for each element i in A do
3: if A[i] = v then
4: return true
5: end if
6: end for
7: return false
8: end procedure
```

2. Pseudocode for finding duplicates in a list

This algorithm demonstrates how to examine if a list has duplicate entries using a linear search.

Input: This algorithm takes a populated array A as a parameter.

Output: This Algorithm is a boolean type and returns true or false respective of whether a duplicate value is found or not.

Algorithm 2 Examining for duplicates

```
    procedure BOOL EXFORDUPES(A[])
    for each element i in A[] do
    for each element j in A[] do
    if A[i] = A[j] then
    return true
    end if
    end for
    end procedure
```

Item 2: Week 4 - Time complexities and Big-O notation

1. Describe the runtime bounds of the linear search algorithm

Algorithm 3 LinearSearch 1: **procedure** BOOL LINEARSEARCH(item, list[]) for each element i in list do 3: (n) if list[i] = list then t (n) 4: return true (n) end if 6: end for 8: **return** false (1)9: end procedure

Collecting the line-by-line runtime data from the algorithms gives: n + n + n + 1 which is equivalent to: 3n + 1.

Therefore the time complexity of the algorithm is O(n).

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2. Describe the runtime bounds of the duplicate finder algorithm

```
Algorithm 4 Examining for duplicates
1: procedure BOOL EXFORDUPES(list[])
      for each element i in list[] do
2:
                                                (n)
          for each element j in list[] do
                                                 (n*n)
3:
             if list[i] = list[j] then
                                                 (n*n)
4:
            return true
                                                 (n*n)
5:
             end if
6:
          end for
7:
      end for
9: return false
                                                (1)
10: end procedure
```

Collecting the line-by-line runtime data from the algorithms gives: n + (n * n) + (n * n) + (n * n) + (n * n) + 1 which is equivalent to: $3n^2 + n + 1$.

Therefore the time complexity of the algorithm is $O(n^2)$

Additional work: Critical values of relative runtimes

Write a function that determines the critical value at which the relative runtime of two linear algorithms swap.

For this algorithm, I am assuming that k1 > k2 (Expression 1 > Expression 2, when n = 0). The Algorithm is very simple; While the value of Expression 1 is greater than Expression 2, increase the value of n.

When the Runtime of the algorithms swap, the while-loop exit condition is fulfilled and the current value of n is returned.

Algorithm 5 Relative runtime comparison algorithm

```
1: procedure CRITVAL(m1, k1, m2, k2)

2: while (m1 * n + k1) > (m2 * n + k2) do

3: n + +

4: end while

5: return n

6:
```

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7: end procedure

Item 3: Week 6 - Harmonic Series

1. Harmonic Series (Pseudocode)

Use pseudocode to specify a recursive algorithm to compute the nth value of the harmonic series, for some integer n.

The Harmonic series is as follows: 1 + 1/2 + 1/3 + 1/4 + 1/5 + ...1/n

Input: This algorithm takes two parameters t and n which are the total sum of the algorithm and the number of repetitions, respectively.

Output: This algorithm outputs the value t which is the total sum of the harmonic series.

This procedure uses a while-loop to control the number recursive iterations.

While the number of iterations left is above 0, add the next value to t, decrease the number of iterations by 1 and recursively call the procedure with the new values of t and n.

When the number of iterations left is no longer above 0, the final value of t is returned and the procedure ends.

Algorithm 6 Computing nth value of harmonic series

```
1: procedure HARM(float t, float n)
2: while n > 0 do
3: t \leftarrow t + (1/n)
4: n - -
5: HARM(t, n)
6: end while
7: return t
8: end procedure
```

2. Harmonic Series (JAVA Implementation)

The Harmonic Series computation algorithm implemented in Java

Listing 1: harms java class file

```
public class harms{
       public static void main(String[] args){
           System.out.println(f(0,3));
       public static float f(float t, float n) {
           t always has a value of 0 on the initial method call.
10
           n is the nth term, which decreases by 1 each recursive call
12
           When n = 0, stop recursive calling and return the value t.
13
           */
           while (n>0)
15
16
               t+= (1/n);
               f(t,--n);
18
           }
19
           return t;
20
21
22
24
```

Evidence of the Harmonic Series computation java implementation. The nth value passed to the method was 3.

```
Terminal × +

riglerr@university-work:~/workspace/210CT_Programming/Portfolio/Item_3/6_Code (master) $ java harms
1.8333334

riglerr@university-work:~/workspace/210CT_Programming/Portfolio/Item_3/6_Code (master) $
```

Item 4: Week 7 - Heapworksheet or RPN Calculator

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Item 5: Week 8 - Linked List Delete function

Below are the Linked List and Node classes that I have created as well as a Test class. This is a fully functional doubly Linked List capable of Inserting, Prepending, Deleting and Displaying items in the List. Please see the Evidence Screenshot below:

Listing 2: List Class Implementation: Java

```
package com.company;
2
  public class List {
        //Nodes to hold the head and tail of the list respectively
      Node head;
5
      Node tail;
      //----List Constructor to init values to null-----
      //-----
      public List() {
10
          head = null;
11
         tail = null;
12
13
14
      //----Insert an Element at an offset from head node-----
      //-----
16
17
      public void Insert(Node n, Node x) {
18
          //If previous node exists
19
          if (n != null) {
             //sWAP NODE POINTERS
21
             x.next = n.next;
22
             n.next = x;
             x.prev = n;
24
25
             // If new node isn't end of list
26
              if (x.next != null)
27
                 x.next.prev = x;
28
          }
29
          //Prepend to existing list
30
          if (head != null & n == null) {
31
             n = head;
32
             n.prev = x;
33
             x.next = n;
34
             x.prev = null;
35
             head = x;
36
          }
37
```

```
39
40
          // If new list (First item to be placed)
41
         else if (n == null & head == null)
42
          {
43
       //make new node equal to both head and tail
44
              head = tail =x;
45
              x.prev =x.next = null; // make head pointers null
46
          } //If all else exp
          else if (tail == n)
48
             tail =x;
49
      }
50
      //----
51
      //----DELETE A NODE FROM THE LIST------
      //----
53
      public void Delete(Node n) {
54
          //create temporary node
55
          Node i = head;
56
57
          while (i != null) { // Loop until null (tail.next)
58
59
              if (i.value == n.value) {
60
61
                  if (i == head) { //If node to remove is head node
62
                     head = i.next;
63
                     i.prev = null;
64
                     break;
65
                  }
66
67
            //if node to remove is tail node
68
                  else if (i == tail) {
69
                     tail = i.prev;
70
                     i.prev.next = null;
71
72
                  else { // If node to remove is not Head/Tail
73
                  i.prev.next = i.next;
74
                  i.next.prev = i.prev;
75
                  break; }
76
77
              i=i.next; // Increment
          }
79
80
```

```
81
82
83
     //----DISPLAY ALL NODES IN LIST-----
85
     //----
86
     public void display() {
         Node i = head;
88
89
         while(i != null) {
            System.out.print(i.value + ", ");
91
            i = i.next;
92
         }
93
94
95
96
97
```

Listing 3: Node Class Implementation: Java

```
package com.company;
2
   /**
    * Created by rob on 20/12/14.
5
  public class Node {
       int value;
       Node prev;
       Node next;
10
       public Node(int val){
12
13
                //Node Constructor
            value = val;
15
            next =null;
16
            prev =null;
17
18
19
20
21
```

Listing 4: Test class

```
package com.company;
2
 public class Main {
3
    public static void main(String[] args) {
5
       List obj1 = new List();
       obj1.Insert(null, new Node(42));
       obj1.Insert(obj1.head, new Node(32));
10
       obj1.Insert(obj1.head.next, new Node(102));
11
        //----
13
        //----Test Inserting Tail-----
14
        //----
       System.out.println("Insert 3 items:");
16
       obj1.display();
17
       System.out.println();
18
19
        // -----
20
        //----Test Insert Between-----
21
        //----
22
        //Insert 12 between the head node and the head.next
23
       System.out.println("Insert 12 after head node: ");
24
       obj1.Insert(obj1.head, new Node(12));
25
       obj1.display();
       System.out.println();
27
28
        //----
29
       //-----Test Prepend-----
30
       //----
31
        //Insert 13 at start of pre-existing list
32
       System.out.println("Prepend 13: ");
33
          obj1.Insert(null, new Node(13));
34
          obj1.display();
35
         System.out.println();
36
37
38
39
41
42
```

```
43
        //----Test Delete Head------
44
        //----
45
        //Delete the head node '13' from the list
        System.out.println("Delete Head: ");
47
        obj1.Delete(new Node(13));
48
        obj1.display();
49
        System.out.println();
50
51
        //----
        //----Test Delete Norm-----
53
        //----
54
        //Delete a middle node (32) from list an display
55
        System.out.println("Delete a Middle Node (32)");
56
        obj1.Delete(new Node(32));
        obj1.display();
        System.out.println();
59
60
61
        //-----Test Delete Tail-----
62
        //----
        //Delete the tail node '42' from list an display
64
        System.out.println("Delete tail: ");
65
        obj1.Delete(new Node(102));
        obj1.display();
67
        System.out.println();
68
69
70
71
     } }
```

```
/usr/lib/jvm/java-8-oracle/bin/java ...
Insert 3 items:
42, 32, 102,
Insert 12 after head node:
42, 12, 32, 102,
Prepend 13:
13, 42, 12, 32, 102,
Delete Head:
42, 12, 32, 102,
Delete a Middle Node (32)
```

42, 12, 102, Delete tail:

42, 12,