Sequential Auction Assignment

Richard Fox
Mathematics for Real-World Systems Centre for Doctoral Training
Warwick Mathematics Institute
University of Warwick
richard.fox@warwick.ac.uk

***Abstract*—Four games are considered, falling into two categories of auction, in which the aim is to create a bidding bot that succeeds in an unknown environment. First, analysis of simple strategies are performed, to create a viable policy. Passive reinforcement learning is then applied to these viable policies to optimise them. In extensive local testing this bot performed very well, and in collaborative testing still remained competitive. Its performance in a completely unknown environment is still yet to be seen. Future work could include analysis of critical regret matching, replicator dynamics or an alternative learning paradigm.**

I. Introduction

There are four auctions considered that fall within two types of auction, with two variations on each type, which are first to five and value accumulation. These require a bidder to accumulate a set number of items (in our case 5) of the same type after which the game ends, or to accumulate the highest number of points (sum of values of items owned) over the whole sequence of auctions, each auction is referred to as a round and there are R rounds in a sequence.

Formally, these are extensive-form finite games with imperfect information defined as a tuple,

$$\langle N, A, H, Z, \vec{i}, \vec{A}, \sigma, \mathbf{u} \rangle,$$

that contains:

- $N = \{1, \ldots, n\}$ the set of bidders
- A a set containing the possible actions
- R a set containing the choice nodes, with exactly one root r_0
- Z a set of leaf nodes (terminal states)
- $\vec{i} : R \to N$ a turn function, in this case turns are defined as each round of an auction
- $\vec{A} : R \to 2^A$ mapping the choice nodes to the power-set of actions thereby fixing all the allowed actions
- $\sigma : R \times A \to R \cup Z$ a successor function, i.e. where an action at a choice node leads, either another choice node or a terminal state that we require to be injective
- $\mathbf{u} = (u_1, \ldots, u_n) \quad u_i : Z \to \mathbb{R}$, u_i is a utility function and $\mathbf{u}$ is each bidders utility function

with an injective σ and a defined root r_0, can be expressed as a tree with branches for each choice node, the current choice node is indistinguishable from all others at the same level, and previous rounds are common knowledge[1], as is the item sold and price paid for it, and the next round only depends on the current round. All finite extensive-form games have at least one Nash equilibrium (NE), brute force will not work to find it in this case, instead the approach is to simulate very simple strategies, including considerations of degeneracy, and analytically find profitable deviations from this.

Equivalent bids are awarded randomly to one of the bidders who placed at bid at that price, first to five sequences end at r_{200} irrespective of if there is a winner by that point, value accumulation runs for a full 200 rounds and this allows for multiple players to have equivalent point totals. The bidders have a fixed sum of money to exchange for items, known as their budget $B_i \subseteq R$, where $B_i^{r_0} = 1000$, $\forall\, i \in N$. There is an assumption that the reader has a familiarity of the rules of the auctions examined in this assignment, but the details of which are explained in the Appendix otherwise.

Also, these games can be modelled as a Markov Decision Process (MDP), and, for completeness, more specifically as a Markov Reward Process (MRP). Unfortunately, the MDP state space is incredibly large [2] of order,

$$\sim O\big(\mid N \mid^{|R \times A|}\big) \sim 1 * 10^{100000}$$

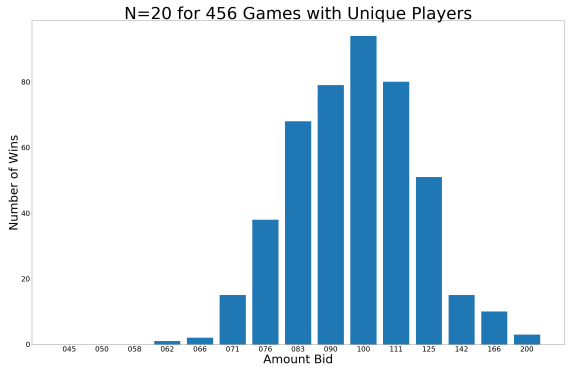
meaning that while theoretically possible to use tools like backward induction [1] or value iteration [1] , it is highly impracticable to compute even a small fraction of this state space effectively. Due to this, the approach taken is to use reinforcement learning on a viable policy, which effectively truncates the state space to exploration around this policy [2].

First, we must assess what is a 'viable' policy by evaluating simple pure strategies, to create a robust mixed strategy. Creating such a policy is not as straightforward as one may think due to integer bids in a finite range and item values being globally fixed, leading to bid degeneracy and therefore a dependence on $\mid N \mid$.

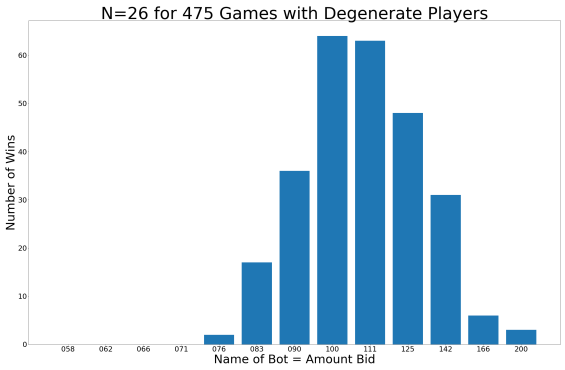
Secondly, passive reinforcement learning is applied, after examining the differences between variations in order to

[1] common knowledge is defined as knowledge that every individual knows, knows that every other individual knows, knows that they know that they and every other knows that they know this knowledge, ad infinitum

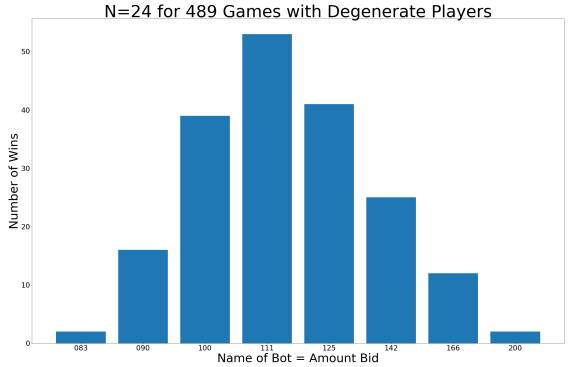
[2] calculated with an indicative value for $\mid N \mid = 10$, and assuming every bidder spends all their money linearly throughout the rounds, giving average $B_i = 500$ so taking $\mid R \times A \mid \approx 200^{500}$



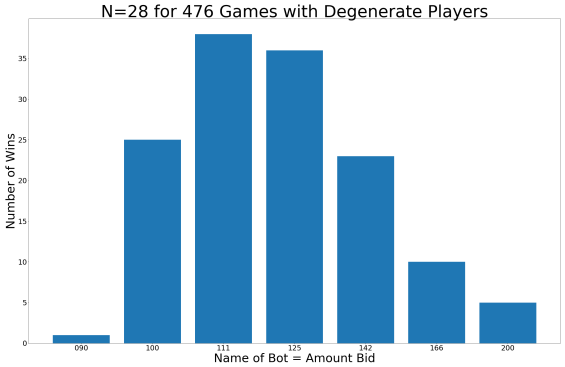
(a) No Degeneracy



(b) 2x Degeneracy



(c) 3x Degeneracy



(d) 4x Degeneracy

Fig. 1. Wins for each FlatBot$X$ over the course of $N = 500$ games, the x axis labels refer to the X value of the FlatBot. There are less than $N = 500$ games completed due to rare games where no one wins or errors are thrown during the game

exploit any particulars in each case. By using this technique it is hoped to gain insight into non-trivially analysed parameters, identified with the initial analytical approach. Whilst not as robust as it could be here, which is discussed later, there should either be significant improvements or a reasonable empirical proof of the viability of these strategies.

II. First to Five

A. Analysis of Strategies

Starting with bidders that bid the same amount each round is a simple strategy, known as 'FlatBot$X$' where X is the amount bid, serving as a initial testing ground. Although, some considerations of the auction itself are applied, namely that if you need 5 items to win, then you must be able to bid on that many items meaning that no FlatBot above $X = 200$ is considered. Conversely, as there are 4 types of item on offer there will be a maximum of 17 rounds before there are 5 of any type in a range, which relates to FlatBot58. Taking the maximum that can be bid consistently for n_{expect} rounds $X = \frac{B_0}{n_{expect}}$, where $n_{expect} \in [5\,..\,17]$, and ignoring bots that would not bid within 200 rounds, we populate our testing environment with the full range of these FlatBots, the increase in degeneracy shows a higher proficiency for higher X Fig.1.

In this regime, the first profitable deviation considered is to only bid on the required type with FlatBot200. Although, with $g > 4$ this is weak to smaller FlatBots as they can effectively 'spray and pray' for the type with the least amount of associated degeneracy. An increase in effectiveness can be achieved by looking for the next most frequent type, but this is now dependant on g, which is not known a priori and must be estimated, and when g is over estimated, next highest frequency is weak to the first highest frequency strategy. Further to this there is another deviation that can be performed, that is to bid on the next most (or least or next least) frequent type depending on the degeneracy in the system, or to perform switching based on observed degeneracy, although this switch may also be degenerate. Switching implies one cannot use the FlatBot200 as this can only secure 5 items, and you should not pay less than 58, as this introduces unnecessary risk of being out bid. On the other hand switching can fall pray to a 'sting in the tail' FlatBot if not careful, that is a FlatBot that bids everything it has left if it already has 4 of this type.

As the level of degeneracy is environment dependant, a behavioural strategy is more applicable as want to adapt as more of a particular environment is deduced. Therefore what we want to learn is a distribution for the probability of which strategy is taken given the amount of amount of items bought for less than 200 and given the estimated amount of degenerate strategies. Therefore creating a much smaller state space to learn over.

B. Variant Differences

In the first variant of first to five, the list of items, in the correct order, is published to all participants. This allows for the possibility of interrogating this list to find the first five of any type to bid on. Unlike all other auctions examined here, the second variant restricts the information available, by removing the published order of items.

The initial evaluation of FlatBots for the first to five auction does not take the item order into account, although the total number of each type of item is still available in both variants. Therefore, these two variant can be approached much the same, with the only difference being to evaluate frequency in a certain time period ahead, and to look for the most available item over all rounds. This approach is much weaker in the second variant as it cannot see unusually high entropy item order states, leading to a greater effectiveness of modified FlatBots.

This converges to almost the same states to learn over, with the first variant having the extra parameter of number of rounds in the future to evaluate.

III. Value Accumulation

A. Analysis of strategies

Here, we start with calculating the total number of points available throughout the auction v', by using the published list of items, and the published list of item values. Then, calculating the expected money per point value $v_{x_i} = \frac{B_{x_i}}{v'}$. Repeating this at the start of each round give a variable that scales with dwindling points available and when winning items and therefore accumulating points. Used with the items intrinsic value v_{item}, gives an individual their Value $V = v_{x_i} v_{item}$ for the item up for auction each round.

Now, this will obviously change contiguously for each bidder after each round Fig.2, and decrease if is winning, however aggression now comes into play. In first to five auction, maximal aggression, even so far has to bitterly stop rivals bids, is displayed at all times, with value accumulation the approach needs endurance as it always, as long as everyone doesn't run out of money, completes the full 200 rounds.

That said, aggression has its place here, as there is merit in paying a small amount more to secure points. The balance between aggression and endurance is very fine, and most if not all bidders will have varying levels of aggression, although there is plenty of information in the auction to try and fit to to extract this information, the approach is to wrap this up in stochastic noise.

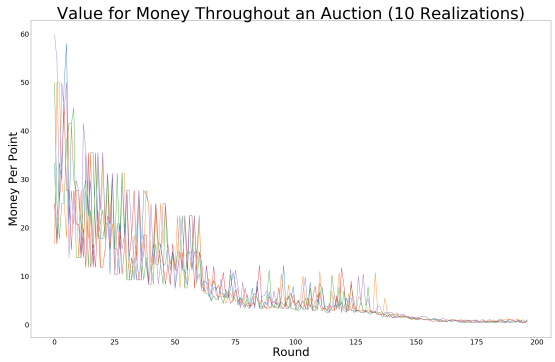


Fig. 2. Typical spread of how many points are gained for each money spent throughout a sequence of rounds.

Assume bids are approximated by a normal distribution around V as the mean, that we will update throughout the rounds, we want to be bidding in the $[\sigma, 2\sigma]$ range, where $\sigma =$ the standard deviation, region of the distributions probability density function. This means that a sufficient amount to be competitive in every round is bet, but not so much as to over pay and run out of money. As we are only in one realisation the statistic will be a sample, so each average taken will vary about the true average, with standard deviation σ_{SEM}. Thus, corrected by $\sigma = \sigma_{SEM}\sqrt{r_t}$ where $t =$ current round. As for own aggression this takes the simple form of a multiplicative constant that will be learned, here the level of aggression is proportional to consecutive wins/losses.

B. Variant Differences

The first variant is a sealed bid first price auction, as both variants of first to five are, which Vickrey showed in 1961 [3] to have a NE for an individual auction as $\frac{n-1}{n}V$. The second Variant is a second price auction, the winning bidder only pays what the second highest bidder bid, which is shown to have NE of V [3], i.e. to pay what the item is worth to you. All information provided in both auction variants is the exact same.

For both variants the NE from the literature is used as a baseline and apply the variance analysis and the aggression multiplier, with a view to having the RL also learn how to compensate for unreliable statistics before a sufficient number of rounds has pasted.

IV. Reinforcement Learning (RL)

A. Theory

By the use of the passive adjective, it is implied that there is a preexisting policy (defined in previous sections) upon which temporal difference leaning is applied (3a), comparing to deviations from this policy with a deviation rate of $(1-\epsilon)$, $\epsilon \in [0, 1]$.

$$V^{\pi}(s) = V^{\pi}(s) + \alpha[r(s) + \gamma V^{\pi}(s') - V^{\pi}(s)]$$

$$V^{\pi}(s) = \mathbb{E}_{Utility}\left[\sum_{t=0}^{R} \gamma^t r(s_t)\right]$$

where $V^{\pi}(s)$ is the value of the policy in state s, α is a weighing constant that determines learning rate based on our confidence that the new policy is defacto an improvement, γ a discounting term that controls the importance of future state values, $r(s)$ is the reward received in state s and $\mathbb{E}_{Utility}$ is the expected utility [4].

Values for terminal states is defined as well as rewards at each state, the value is then back-propagated in a Bellman Equation (value iteration) -esque manner. The optimal policy π_s^* is defined as the policy that maximises value in each state as such:

$$\pi_s^* = \arg\max_{\pi} V^{\pi}(s).$$

In addition to the selection of architecture for learning, there are different forms of input perception, defined as: Independent Learners, that look at their own actions and the associated rewards, Joint Action Learners, that also take into account all the other agents and Gradient based optimisation, that sits in the spectrum somewhere between the two. Here we take the Independent approach.

B. Implementation

A script to perform passive RL is created, also a file creating the policy that both the RL script reads from and updates and the actual played policy is drawn from. This take the form as a list of parameters that each have their own exploration rate ϵ, but all have the same weighting $\alpha = 0.5$. No parameter is used in every game and are therefore updated when in the appropriate auction, hence the necessity of individual exploration rates.

In the first to five regime, it is imperative to win as soon as possible and therefore if we set the value of negligible importance $\epsilon_{import} = 0.01$, then we want:

$$\gamma^{\mathbb{E}[r]} = \epsilon_{import}, \quad \gamma = \sqrt[\mathbb{E}[r]]{\epsilon_{import}}$$

where $\mathbb{E}[r]$, $r < R$, $r \in \mathbb{Z}^+$, is the expected number of rounds as discussed in section II.B, for example $\gamma \approx 0.76$ sets states > 17 rounds in the future as negligible. In value accumulation a win is only dependant on the outcome of all rounds, therefore $\gamma = 1$.

In the learning environment meant, there were a varying number of players of varying complexity, being randomly selected at the beginning of each round. In the submission of the work the value of parameters that is converged upon will be hard coded in, i.e. will not be dynamically leaning whilst playing, due to submission constraints.

V. Conclusions & Future Work

During individual testing this converged policy has performed very well, less so in collaborative testing, but is for the most part competitive. The work submitted as part of this assignment uses the best policy achieved, there is not formal proof that this policy has not fallen prey to optimising in a local minima. Indeed it is highly likely not to be, as the state space is vast, there is a requirement for very extensive testing and input from a cross-disciplinary team of researchers to be declared solved in any sense. That said, I believe what has been achieved is in some sense "good" or in some approximation "competitive".

Going forward the obvious step is to take a gradient based approach to learning. In a slightly different direction, taking other analytical polices and trying to optimise them, in a similar approach to performed here and comparing to this one. Further to this, evaluate replicator dynamics to find the evolutionary stable strategies (ESS) that are emergent from policies that are competitive, but may only have a small probabilistic advantage. Studies of such dynamical systems lead to the establishing of critical points of convergence, which may be a more appropriate way to truncate the state space.

A more game theoretic alternative would be to apply critical regret matching (CRM), where there is further punishment for not taking actions that would have won a round and are available to that bidder in a certain state. Learning the opponents strategy this way could lead to a much more adaptive strategy, for instance that can consider $\mid N \mid$ dependence, which is mentioned above but never taken into serious consideration.

References

[1] Y. Shoham and K. Leyton-Brown, Multiagent systems. Cambridge: Cambridge Univ. Press, 2012.

[2] Pineau and N. Roy, "Reinforcement learning with limited reinforcement: Using Bayes risk for active learning in POMDPs", Artificial Intelligence, vol. 187-188, pp. 115-132, 2012. Available: 10.1016/j.artint.2012.04.006.

[3] W. Vickrey, "Counterspeculation, Auctions, and Competitive Sealed Tenders", The Journal of Finance, vol. 16, no. 1, p. 8, 1961. Available: 10.2307/2977633.

[4] S. Russell and P. Norvig, Artificial intelligence.

Appendix
Assignment and Game details (on next page)

Auction Games

CS404 Agent-based Systems Coursework

1 Introduction

Imagine an auction of paintings by famous artists. There is an auction room, with an auctioneer that presents each piece to be sold, and all bidders then write their bids for the item onto a secret sealed note that is handed to the auctioneer. The auctioneer then declares the highest bidder the winner, takes their payment, and starts the next round with a new item. There might be a target number of paintings by the same artist to get, in which case the first bidder to get this wins; otherwise, the auction continues until everyone runs out of money or there are no more items to sell, and the bidder with the highest total value of paintings is the winner.

Your objective is to implement strategies for a Python 3 bidding bot that will participate in such an auction.

2 The Auction

The auction runs in the following way. A number of bidders connects to the AuctionServer, which then announces the winning condition, the item types (i.e, the artists) and the number of each being sold and value of each of these types, how much each money each bidder starts with, which bid the highest bidder will pay and optionally the sequence of the auction.

The AuctionServer will then announce an item type to be bid upon. Your AuctionClient will use the above information to determine an amount to bid and submit it to the AuctionServer. Once done, the AuctionServer will declare the highest bidder the winner, who will then be charged (but not necessarily the amount they bid, see below) and receives the item. If there are drawn positive bids, then the winner is chosen at random from the drawn bidders; if however the drawn bids are all at 0 (no-one bid), no-one wins, and the item is discarded.

The auction will continue until either there are no more items to sell, all bidders have run out of money, or if there is a set winning condition, e.g., a bidder managed to acquire N items of the same type, in which case they are declared the grand winner. If there is no set winning condition, the game will only end once one of the first two conditions is met, and then the bidder who ends with the highest total value in items is the grand winner.

Note that however whilst the highest bidder will always win, the auction may be set up so the highest bidder does not pay their own bid. It can be set up so that the highest bidder is only charged the second highest bid (see `winner_pays`).

You will write your strategies in AuctionClient, and then everyone's AuctionClients will all be logged into one AuctionServer, where a tournament will be played across four games:

- Game 1: First to buy 5 of any artist wins, the highest bidder pays own bid, and the auction order known.
- Game 2: First to buy 5 of any artist wins, the highest bidder pays own bid, but the auction order is not known.
- Game 3: Highest total value at the end wins, the highest bidder pays own bid, and the auction order known.
- Game 4: Highest total value at the end wins, the highest bidder pays the second highest bid, and the auction order known.

Note that for all games in this tournament, there will be four types, (Picasso, Van_Gogh, Rembrandt and Da_Vinci), the auction size will be 200, the starting budget 1000, and where the win condition is based on final value, each Picasso is worth 4, Van_Gogh worth 6, Rembrandt worth 8 and Da_Vinci worth 12.

3 Implementation

Provided to you are four Python 3 files: AuctionServer, AuctionClient, run_auction and run_client.

AuctionServer contains the definition for the AuctionServer class, which sets up and runs the auction. It has the following arguments:

- `host`: Where the server will be hosted, keep this at "localhost" for your own testing.
- `ports`: Either a list of port sockets that the clients will connect through, or if a single number, P, is given, it will use all ports from P to P + numbidders. If you get any error that a port is already in use or locked or something, just change this number to any other number above 1000 and try again.
- `numbidders`: The number of clients i.e., bidders, that will be playing.
- `neededtowin`: The number of the same type of painting a player needs to get to be the grand winner. If 0, then the total value once the auction has ended will be used to declare the grand winner instead.
- `itemtypes`: List of the different types of paintings, i.e, the different artists.
- `numitems`: A dict that can either be used to manually set how many items of each type there should be in the auction, or if unset i.e., "`numitems={}`", this is generated randomly according to `auction_size`.
- `auction_size`: If `numitems` is set, this should be 0. Otherwise, this is how many items are in the auction; items will be randomly generated until there are this many to sell.
- `budget`: The starting budget of all the players.
- `values`: A dict of the value for each type. This will be used to determine the total value for each player if `neededtowin == 0`
- `announce_order`: If True, the bidders are told the sequence of the items to be sold before auction begins. If False, they are not, and will have to use `numitems` and the past bidding record to guess which item will be sold next.
- `winner_pays`: An index of which bid the highest bidder pays. If 0, they pay their own bid, if 1 they pay the second highest, if 2 the third etc... Note however that the winner always pays at least 1, even if the second highest is 0.

The function `announce_auction` sets up and announces the auction details to all clients, whilst `run_auction` can then run the auction until it is completed.

AuctionClient is the code for a bidding client, and this is where you will implement your strategies. Note that an AuctionServer must be initialised first so it is listening on the given ports before an AuctionClient can be initialised and connect to the server. It has the following arguments:

- `host`: Where the AuctionServer to connect to is hosted, keep this at "localhost" for your own testing.
- `port`: The port to connect to the AuctionServer to. This must be unique to this AuctionClient (no other AuctionClient can use this port), and an initialised AuctionServer must be ready and listening to this port.
- `mybidderid`: The name for your bidder. If not given with initialisation, you will be asked to input this on the command line. Note that it must be unique to your bidder, and can contain alphanumeric or underscore characters.
- `verbose`: If set to True, then your client will print out statements detailing its input and progress.

The function `play_auction` runs the loop for receiving data from the AuctionServer and passing the bids back until the auction is done.

You will also see provisionally defined functions of `determinebid` and four `bidding_strategy` functions for each of the four games, and it is in these functions you should implement the algorithm for your strategies.

Currently, they just run the function `random_bid`, which returns a random number between 1 and amount of budget the bidder has left. You can use this random strategy to test your strategy against.

`determinebid` receives all the relevant data as input and has a comment block that explains what each argument is. If you wish, you are also free to access and create self defined variables, which may be useful if you want to store state between calls of `determinebid`. You may also add code to the `__init__` or `play_auction` functions to help achieve this, but be very careful not to change either the networking code or the auction parameters or you might either crash or confuse your bot! Also note that if your bidder tries to bid more than it has budget left (accessible by `standings[mybidderid]['money']`), the AuctionServer will cap the bid to its current budget.

run_auction is a script to initialise an AuctionServer with the given parameters and start it running and waiting for AuctionClients to connect. Once it is set up, you can manually run AuctionClients by initialising them with a port the AuctionServer is listening on, and once `numbidders` AuctionClients have connected to the server on their different ports, it will run the auction.

run_clients is a script that automatically initialises and runs not only an AuctionServer, but also `numbidders` number of AuctionClients, and is provided for convenient testing.

4 Submission

Your coursework submission will consist of a single compressed file (either .zip or .tgz) containing:

- An AuctionClient.py file, encoding the strategies for each game.
- A four pages Analysis.pdf file, with the analysis of each of the four corresponding strategies and the reasons of your design choices. The pdf should be written in IEEE two-column conference format.

The coursework file should be submitted through Tabula.

Each submission will be run as follows. We first set up the auction with the desired parameters for each tournament in the run_auction.py script, then run it. This will create an auction room with those parameters. Then we will run each student's client with the following three lines in a Python script:

import AuctionClient from [submission filename]

bidbot = AuctionClient(port=[Unique port], mybidderid=[student ID])

bidbot.play_auction()

Where [submission filename] is the filename of the student's submission, [Unique port] is one of the ports the AuctionSever is configured to listen on and [student ID] is the student's unique ID. Your submission file name should be named with your student ID, followed by the suitable extension.

Please make sure that your submission is suitable for this operation.

5 Evaluation

The coursework is worth 50% of the module credit. Its marking scheme is structured as follows:

- 30% strategy implementation (how you design and structure your strategies)
- 30% strategy performance (how each strategy performs in the games)
- 40% quality of the analysis (how you describe and analyse your strategies)

Each of the four strategies will be evaluated independently and the result averaged across them.