

# CMSC 303 Introduction to Theory of Computing

## Chapter 2 Practice

*Key*

1. Define language  $L = \{w \mid w \text{ starts and end with the same symbol and } w \in \{a, b\}^*\}$  Give a CFG generating  $L$ .

$$\begin{aligned} S &\rightarrow aTa \mid bTb \mid bLa \mid \epsilon \\ T &\rightarrow aT \mid bT \mid \epsilon \end{aligned}$$

2. Create a state diagram for a PDA recognizing the language as defined below.

$$Q = \{q_0, q_1, q_2, q_3\},$$

$$\Sigma = \{a, b\},$$

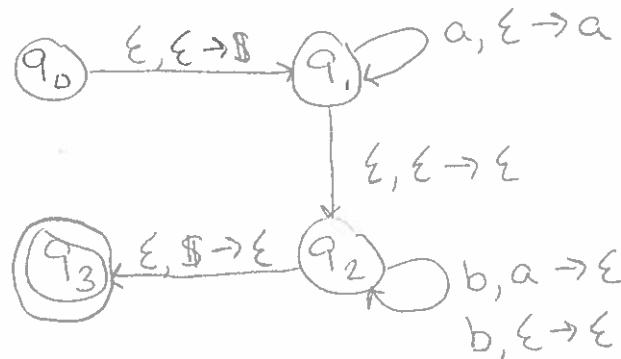
$$\Gamma = \{a, \$\},$$

$$q_0 = q_0,$$

$$F = \{q_3\}, \text{ and}$$

$\delta$  given by the following table, wherein blank entries signify  $\emptyset$ .

| $\delta$ | a                   | b                      | $\epsilon$             | a | b | $\epsilon$             |
|----------|---------------------|------------------------|------------------------|---|---|------------------------|
| pop      | a    \\$ $\epsilon$ | a    \\$ $\epsilon$    | a    \\$ $\epsilon$    |   |   |                        |
| $q_0$    |                     |                        |                        |   |   | ( $q_1$ , \$)          |
| $q_1$    |                     | ( $q_1$ , a)           |                        |   |   | ( $q_2$ , $\epsilon$ ) |
| $q_2$    |                     | ( $q_2$ , $\epsilon$ ) | ( $q_2$ , $\epsilon$ ) |   |   | ( $q_3$ , $\epsilon$ ) |
| $q_3$    |                     |                        |                        |   |   |                        |



3. Describe the strings this PDA recognizes.

$$L = \{a^n b^m \mid n \leq m\}$$

4. Use the following CFG to parse the string, aababb:

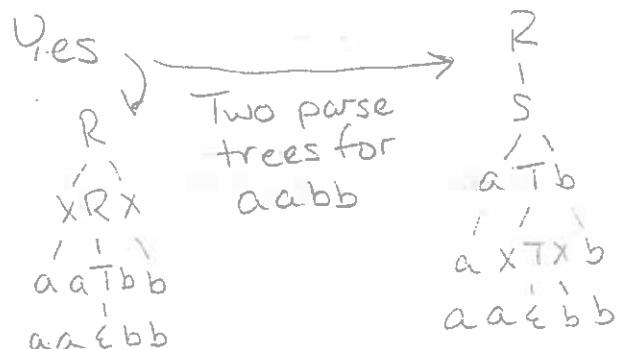
$$\begin{aligned} R &\rightarrow XRX \mid S \\ S &\rightarrow aTb \mid bTa \\ T &\rightarrow XTX \mid X \mid \epsilon \\ X &\rightarrow a \mid b \end{aligned}$$

one way

$$\begin{aligned} R \Rightarrow XRX &\Rightarrow aRX \Rightarrow aXRX \Rightarrow aaRXX \Rightarrow aaSXX \\ &\Rightarrow aabTaXX \Rightarrow aab\epsilon aXX \Rightarrow aabaXX \Rightarrow aabab \\ &\Rightarrow aabb \end{aligned}$$

5. Choose another string from this CFG to parse and show the parse tree. Is this grammar ambiguous?

Your choice.  
I pick aabb



6. Place the CFG below into Chomski Normal Form.  $\Sigma = \{a\}$

$$\begin{aligned} A &\rightarrow BAB \mid B \mid \epsilon \\ B &\rightarrow aa \mid \epsilon \end{aligned}$$

① New Start State

$$\begin{aligned} S &\rightarrow A \\ A &\rightarrow BAB \mid B \mid \epsilon \\ B &\rightarrow aa \mid \epsilon \end{aligned}$$

③ Get rid of  $A \rightarrow \epsilon$

$$\begin{aligned} S &\rightarrow A \mid \epsilon \\ A &\rightarrow BAB \mid B \mid A \mid BA \mid AB \mid BB \\ B &\rightarrow aa \end{aligned}$$

② Get rid of  $B \rightarrow \epsilon$

$$\begin{aligned} S &\rightarrow A \\ A &\rightarrow BAB \mid B \mid A \mid BA \mid AB \mid \epsilon \\ B &\rightarrow aa \end{aligned}$$

④ Get rid of  $A \rightarrow A + S \rightarrow A$

$$\begin{aligned} S &\rightarrow BAB \mid B \mid BA \mid AB \mid BB \mid \epsilon \\ A &\rightarrow BAB \mid B \mid BA \mid AB \mid BB \\ B &\rightarrow aa \end{aligned}$$

⑤ Get rid of  $S \rightarrow B + A \rightarrow B$

$$\begin{aligned} S &\rightarrow BAB \mid aa \mid BA \mid AB \mid BB \mid \epsilon \\ A &\rightarrow BAB \mid aa \mid BA \mid AB \mid BB \\ B &\rightarrow aa \end{aligned}$$

4. Use the following CFG to parse the string, aabbabb:

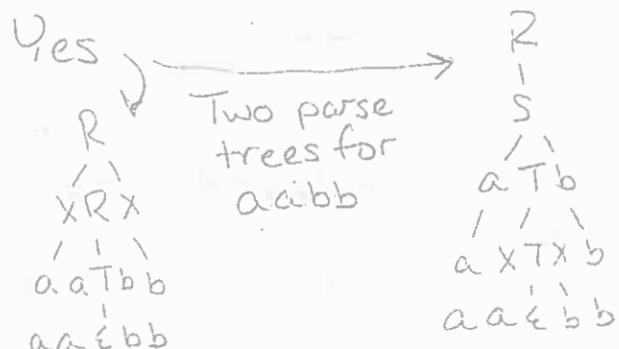
$$\begin{aligned} R &\rightarrow XRX \mid S \\ S &\rightarrow aTb \mid bTa \\ T &\rightarrow XTX \mid X \mid \epsilon \\ X &\rightarrow a \mid b \end{aligned}$$

one way

$$\begin{aligned} R &\Rightarrow XRX \Rightarrow aRx \Rightarrow aXRX \Rightarrow aaRXX \Rightarrow aaSXX \\ &\Rightarrow aabTaXX \Rightarrow aab\epsilon aXX \Rightarrow aabaXX \Rightarrow aabab \\ &\Rightarrow aabbabb \end{aligned}$$

5. Choose another string from this CFG to parse and show the parse tree. Is this grammar ambiguous?

Your choice.  
I pick aabb



6. Place the CFG below into Chomski Normal Form.  $\Sigma = \{a\}$

$$\begin{aligned} A &\rightarrow BAB \mid B \mid c \\ B &\rightarrow aa \mid c \end{aligned}$$

① New Start State      ③ Get rid of  $A \rightarrow \epsilon$

$$\begin{aligned} S &\rightarrow A \\ A &\rightarrow BAB \mid B \mid \epsilon \\ B &\rightarrow aa \mid \epsilon \end{aligned}$$

② Get rid of  $B \rightarrow \epsilon$

$$\begin{aligned} S &\rightarrow A \\ A &\rightarrow BAB \mid B \mid A \mid BA \mid AB \mid \epsilon \\ B &\rightarrow aa \end{aligned}$$

④ Get rid of  $A \rightarrow A + S \rightarrow A$

$$\begin{aligned} S &\rightarrow BAB \mid B \mid BA \mid AB \mid BB \\ A &\rightarrow BAB \mid B \mid BA \mid AB \mid BB \\ B &\rightarrow aa \end{aligned}$$

⑤ Get rid of  $S \rightarrow B + A \rightarrow B$

$$\begin{aligned} S &\rightarrow BAB \mid aa \mid BA \mid AB \mid BB \\ A &\rightarrow BAB \mid aa \mid BA \mid AB \mid BB \\ B &\rightarrow aa \end{aligned}$$

6 cont.  $S \rightarrow BAB|aa|BA|AB|BB|\epsilon$

$A \rightarrow BAB|aa|BA|AB|BB$

$B \rightarrow aa$

⑥ Place into correct form

$S \rightarrow XB|DD|BA|AB|BB|\epsilon$

$A \rightarrow XB|DD|BA|AB|BB$

$B \rightarrow DD$

$D \rightarrow a$

$X \rightarrow BA$

All in form  
 $A \rightarrow BC$   
 $A \rightarrow a$   
 $A \rightarrow \epsilon$