

# Polynomial Interpolation and Sums of Powers

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April 2022

## **Abstract**

Analysis of my investigation into the construction of an algorithm to find a formula for the sum of integers raised to some power.

## **Contents**

|          |                                 |          |
|----------|---------------------------------|----------|
| <b>1</b> | <b>Introduction</b>             | <b>2</b> |
| <b>2</b> | <b>Differences of Sequences</b> | <b>3</b> |
| <b>3</b> | <b>Polynomial Sequences</b>     | <b>5</b> |
| <b>4</b> | <b>Sums of Powers</b>           | <b>6</b> |

# 1 Introduction

This paper will discuss a method for computing a closed form formula for  $\sum_{r=1}^n (r^k)$ , where  $k \in \mathbb{N}$ .

The formulae for when  $k \in 1, 2, 3$  are generally well-known, but it begs to question how to generate the formula for large  $k$ .

$$\begin{aligned}\sum_{r=1}^n (r^1) &= \frac{1}{2}n(n+1) \\ \sum_{r=1}^n (r^2) &= \frac{1}{6}n(n+1)(2n+1) \\ \sum_{r=1}^n (r^3) &= \frac{1}{4}n^2(n+1)^2 \\ \sum_{r=1}^n (r^4) &= \frac{1}{30}n(n+1)(2n+1)(3n^2+3n-1)\end{aligned}$$

Perhaps the most well-known way of generating these polynomials is Faulhaber's formula, which gives an efficient and simple  $O(k^2)$  algorithm, however it depends on the generation of Bernoulli numbers, and so formulae can only be found for sufficiently small  $k$ .

$$\sum_{r=1}^n (r^k) = \frac{n^{p+1}}{p+1} + \frac{1}{2}n^p + \sum_{r=2}^k \left( \frac{B_r}{r!} k^{\overline{r-1}} n^{k-r+1} \right)$$

Where  $B_k$  denotes the  $k$ th Bernoulli number, and  $r^{\overline{k}}$  is the falling factorial.

Such methods have been thoroughly researched by mathematicians including Pierre de Fermat, Blaise Pascal and Jakob Bernoulli. I present a different method which follows the work of Leonhard Euler and his research into differences of sequences.

## 2 Differences of Sequences

As a build up to the polynomial interpolation outlined in the next chapter, we look at the  $n$ th difference of a sequence of numbers  $a_k$ , which is denoted by  $\Delta^n a_k$ . For example, consider the following sequence of numbers,

$$1, 1, 2, 3, 5, 8$$

Then the  $n$ th difference is calculated as,

$$1, 1, 2, 3, 5, 8$$

$$0, 1, 1, 2, 3$$

$$1, 0, 1, 1$$

$$-1, 1, 0$$

Where each element is the difference between the terms above it, a sort of "subtracting" Pascal's triangle.

This yields the relation,  $\Delta^{n+1}a_k = \Delta^n a_{k+1} - \Delta^n a_k$ . We now aim to find a formula for calculating  $\Delta^n a_k$  using this relation and Pascal's triangle.

Consider the sequence,  $a_1, a_2, a_3, a_4, a_5$ . The differences are shown below,

$$a_1, a_2, a_3, a_4, a_5$$

$$a_2 - a_1, a_3 - a_2, a_4 - a_3, a_5 - a_4$$

$$a_3 - 2a_2 + a_1, a_4 - 2a_3 + a_2, a_5 - 2a_4 + a_3$$

$$a_4 - 3a_3 + 3a_2 - a_1, a_5 - 3a_4 + 3a_3 - a_2$$

$$a_5 - 4a_4 + 6a_3 - 4a_2 + a_1$$

It seems as though the coefficients of  $a_k$  in the  $n$ th difference are the coefficients in the binomial expansion of  $(x-1)^n$ , which makes sense given its similarity to Pascal's triangle. From this, we give the following proposition,

**Proposition 2.1**  $\Delta^n a_k = \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} a_{k+i}$

Now we prove this formula by induction.

When  $n = 1$ ,  $\Delta^n a_k = a_{k+1} - a_k$  from our formula earlier. The formula gives,

$$\Delta^1 a_k = \sum_{i=0}^1 \binom{1}{i} (-1)^{1-i} a_{k+i} = \binom{1}{0} (-1)^1 a_k + \binom{1}{1} (-1)^0 a_{k+1} = a_{k+1} - a_k$$

So our formula holds when  $n = 1$ , which serves as our base case.

Now assume for our inductive hypothesis that,

$$\Delta^n a_k = \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} a_{k+i}$$

Now for the inductive step,

$$\begin{aligned} \Delta^{n+1} a_k &= \Delta^n a_{k+1} - \Delta^n a_k \\ \Delta^{n+1} a_k &= \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} a_{k+1+i} - \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} a_{k+i} \\ \Delta^{n+1} a_k &= \sum_{i=1}^{n+1} \binom{n}{i-1} (-1)^{n+1-i} a_{k+i} - \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} a_{k+i} \\ \Delta^{n+1} a_k &= \sum_{i=1}^n \binom{n}{i-1} (-1)^{n+1-i} a_{k+i} + \binom{n}{n} (-1)^0 a_{k+n+1} - \sum_{i=1}^n \binom{n}{i} (-1)^{n-i} a_{k+i} - \binom{n}{0} (-1)^n a_k \\ \Delta^{n+1} a_k &= \sum_{i=1}^n (-1)^{n+1-i} a_{k+i} \left( \binom{n}{i-1} + \binom{n}{i} \right) + \binom{n}{n} (-1)^0 a_{k+n+1} - \binom{n}{0} (-1)^n a_k \\ \Delta^{n+1} a_k &= \sum_{i=1}^n (-1)^{n+1-i} a_{k+i} \binom{n+1}{i} + a_{k+n+1} - (-1)^n a_k \\ \Delta^{n+1} a_k &= \sum_{i=0}^{n+1} (-1)^{n+1-i} a_{k+i} \binom{n+1}{i} \end{aligned}$$

Which has the same form as the formula for  $n + 1$ . We have shown that if the formula is true for  $n$  then it is true for  $n + 1$ . Since it holds for  $n = 1$ , the formula therefore holds for all positive integers  $n$ .

This approach was originally used by Leonhard Euler to prove that all numbers of the form  $4m + 1$  could be written as the sum of two squares, but we show in the next chapter that it gives us a formula for the  $n$ th difference of a polynomial sequence.

### 3 Polynomial Sequences

A polynomial sequence is any infinite sequence,  $f(1), f(2), f(3), \dots$  where,

$$f(x) = a_d x^d + a_{d-1} x^{d-1} + \dots + a_1 x + a_0$$

The goal of this to chapter is to find a method to find this polynomial, starting with the polynomial sequence. We now make the claim,

**Proposition 3.1** A polynomial of degree  $d$  is uniquely defined by  $d + 1$  points.

We prove this by contradiction. Assume that there exist two non-identical polynomials,  $f(x)$  and  $g(x)$ , both of degree  $d$ , which pass through the same  $d + 1$  points. Then  $f(x) - g(x)$  is a polynomial with degree  $\leq d$ , with  $d + 1$  roots. However a degree  $d$  polynomial has at most  $d$  roots. Contradiction. Hence  $f(x) - g(x)$  must be the zero polynomial, and so  $f(x)$  is identical to  $g(x)$ .

Hence, given a polynomial sequence of degree  $d$  and  $d + 1$  terms we know there exists a unique polynomial which produces the sequence. Consider the polynomial,

$$f(x) = a_d x^d + a_{d-1} x^{d-1} + \dots + a_1 x + a_0$$

Then we can create a polynomial sequence with  $f(x)$ . Let us only consider the leading coefficient of  $f(x)$  in this sequence. Since the coefficients are effectively independent of each other, then the coefficient of  $x^d$  in the  $n$ th difference is given by proposition 3.1 as,

$$\sum_{i=0}^n \binom{n}{i} (-1)^{n-i} (i + k)^d$$

The  $n$ th difference of the whole polynomial is simply the sum of the  $n$ th differences for each coefficient. The formula above tells us that the  $n + 1$ th difference of  $x^n$  is 0 because the two terms above are constant. And so there are  $d + 1 - n$  terms in the  $n$ th difference of a degree  $d$  polynomial. This tells us that in the case where  $k = 1$ ,

$$\Delta^n f(x) = \sum_{j=1}^{d+1-n} (a_{d+1-j} \sum_{i=0}^n \binom{n}{i} (-1)^{n-i} (i + 1)^{d+1-j})$$

Because  $k = 1$ , these values are equal to the leftmost column of the triangle of differences. If we know these values, we can calculate the values of the coefficients by back-substitution, since the corresponding matrix is in row-echelon form.

## 4 Sums of Powers

As mentioned in chapter 1, our aim is to find a formula for  $\sum_{r=1}^n (r^k)$ . Looking at the first 4 formulae suggests that a polynomial of degree  $k+1$  should suffice.

**Proposition 4.1**  $\sum_{r=1}^n (r^k) = f(n)$ , where  $\deg(f) = k+1$ .

<https://math.stackexchange.com/questions/1893693/elementary-proof-of-polynomial-degree-of-sum-of-pth-powers?rq=1>

Hence, our method of polynomial interpolation via proposition 3.1 can be used to find our polynomial formula. The value of the  $n$ th difference can be computed via calculating the first  $k+2$  terms since they are the sum of powers, and then applying the formula from proposition 3.1. This method doesn't seem to yield an explicit formula for  $f(n)$ , however it can be easily computed via back-substitution in  $O(n^3)$  by a computer algorithm. A proof of concept program is available on my GitHub page.

As an example, the formula for the sum of the first  $n$  powers of 12 is computed by the algorithm to be,

$$\sum_{r=1}^n (r^{12}) = \frac{1}{13}n^{13} + \frac{1}{2}n^{12} + n^{11} - \frac{11}{6}n^9 + \frac{22}{7}n^7 - \frac{33}{10}n^5 + \frac{5}{3}n^3 - \frac{691}{2730}n$$

Plotting the graph of  $(RunTime)^{1/4}$  against  $k$  verifies that the algorithm is  $O(k^4)$ .

