
Question:

Let B be the language of all palindromes over $\{0,1\}$ containing equal numbers of 0s and 1s. Show that B is not context free.

Answer:

----SETP1----

Let B be the language of all palindromes over $\{0, 1\}$ containing the equal numbers of 0's and 1's . To prove B is not a context free language by taking a contradiction.
Assume that B is a context free language.

Since B is a context free language, then by pumping lemma, there is a number p (the pumping length) where, if s is any string in B of length at least p , then s may be divided into five pieces $s = uvxyz$ satisfying the conditions

1. for each $i \geq 0, uv^i xy^i z \in A$,
2. $|vy| > 0$, and
3. $|vxy| \leq p$.

----SETP2----

Case 1:

Now select a string $s = 0^n 1^{2n} 0^n$.

Clearly s is a member of B of length at least p .

Assume the value of $n=2$ for the string s .

$$s = 0^2 1^4 0^2$$

$$s = 00111100$$

The string s can be divided into $uvxyz$ as follows:

$$\begin{array}{cccccc} 00 & 11 & 1 & 10 & 0 \\ u & v & x & y & z \end{array}$$

Now apply the first condition of the pumping lemma.

$$\text{for each } i \geq 0, uv^i xy^i z \in A$$

For $i = 2$:

$$\begin{array}{cccccc} 00 & \left(\frac{11}{v}\right)^2 & \frac{1}{x} & \left(\frac{10}{y}\right)^2 & \frac{0}{z} \\ u & v & x & y & z \end{array}$$
$$\begin{array}{cccccc} 00 & \left(\frac{1111}{v}\right) & \frac{1}{x} & \left(\frac{1010}{y}\right) & \frac{0}{z} \\ u & v & x & y & z \end{array}$$

Assume the obtained string $00 1111 1 1010 0$ as s' .

The obtained string s' is not a palindrome after applying the first condition of pumping lemma and $s' \notin B$. In the pumped string, the number of 0's and 1's is not equal.

So, the language B is not following the condition1 of the pumping lemma.

----SETP3----

Case 2:

The same string is selected as $s = 0^n 1^{2n} 0^n$

Clearly s is a member of B of length at least P .

Assume the value of $n=2$ for the string s .

$$s = 0^2 1^4 0^2$$

$$s = 00111100$$

The string s can be divided into $uvxyz$ as follows:

$$\begin{array}{c} 0 \quad 01 \quad 11 \quad 1 \quad 00 \\ \hline u \quad v \quad x \quad y \quad z \end{array}$$

Now apply the first condition of the pumping lemma.

for each $i \geq 0, uv^i xy^i z \in A$

For $i=2$:

$$\begin{array}{c} 0 \quad \left(\frac{01}{v}\right)^2 \quad 11 \quad \left(\frac{1}{y}\right)^2 \quad 00 \\ \hline u \quad \quad \quad x \quad \quad \quad z \end{array}$$

$$\begin{array}{c} 0 \quad \left(\frac{0101}{v}\right) \quad 11 \quad \left(\frac{11}{y}\right) \quad 00 \\ \hline u \quad \quad \quad x \quad \quad \quad z \end{array}$$

Assume the obtained string $0 \ 0101 \ 11 \ 11 \ 00$ as s' .

The obtained string s' is not a palindrome after applying the first condition of pumping lemma and $s' \notin B$. In the pumped string, the number of 0's and 1's is not equal.

Hence, the assumption B is a context free language is wrong.

Therefore, by the two cases it can be proved that B is not a context free language.