

Given:

A DFA $M = (Q, \Sigma, \delta, q_0, F)$ and suppose h is a state of DFA M known as its home state.

Proof:

The at most length of synchronizing sequences is $k^3/6$ for a k -state synchronizable DFA. In the year 1964, a Slovak scientist named Jan Cerny first tried to solve the problem of synchronizing automata in real time. This problem is sometimes referred as *Cerny's Conjecture*.

To prove the upper bound on the synchronizing sequence we try to device a greedy algorithm.

Algorithm:

1. Let a synchronizing DFA be $M = (Q, \Sigma, \delta, q_0, F)$. Initialize the synchronizing sequence $\omega \leftarrow \epsilon$ (empty word) and a set of states $P \leftarrow Q$.

2. **while** $|P| > 1$

a. Find a word v which belongs to Σ^* and it has minimum length $|\delta(P, v)| < |P|$.

b. If none exists, **return** failure.

c. $\omega \leftarrow \omega v$

$P \leftarrow \delta(P, v)$

3. **return** ω

• Now suppose that M is a k -state DFA, that is, $|Q| = k$ then clearly the main loop of the algorithm runs at most $k-1$ times. In order to get the length of the output word ω user has to estimate the length of each word v derived at each loop.

• Consider a generic step at which $|P| = n > 1$ and let $v = a_1 \cdots a_l$ with $a_i \in \Sigma$, $i = 1 \cdots l$. Then it is quite simple to see that the sets, $P_1 = P$, $P_2 = \delta(P_1, a_1)$, ..., $P_l = \delta(P_{l-1}, a_{l-1})$ are n -element subsets of Q .

• Furthermore, since $|\delta(P_i, a_i)| < |P_i|$, there exists two states $q_i, q'_i \in P_i$ such that $\delta(q_i, a_i) = \delta(q'_i, a_i)$.

• Now define two element subsets $R_i = \{q_i, q'_i\} \subseteq P_i$, $i = 1, \dots, l-1$. then the condition that v is a word that has minimum length $|\delta(P, v)| < |P|$ which implies that $R_i \not\subseteq P_j$ for $1 \leq j < i < l$. Now by the Peter Frankl inequality, the tight bound over l be $\binom{k-n+2}{2}$.

• Summing up these inequalities from $n = k$ to $n = 2$, user can get the upper bound over the synchronizing sequence $|\omega| \leq \frac{k^3 - k}{6}$.

Conclusion:

Therefore, at most length of synchronizing sequences is $k^3/6$ for a k -state synchronizable DFA.