获得的答案

(a) Languages are

 $L_1 = \{w \mid \text{the length of } w \text{ is at most 5} \} \text{ on } \Sigma = \{0,1\}$ 

And  $L_2 = \{w | \text{ every odd position of } w \text{ is } a1\} \text{ on } \Sigma = \{0,1\}$ 

 $M_1$  be the NFA that recognizes  $L_1$  and

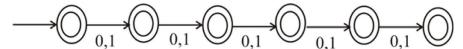
 ${\it M}_{\rm 2}$  be the NFA that recognizes  ${\it L}_{\rm 2}.$ 

Let  $L = L_1 0 L_2$ 

M be the NFA that recognizes L.

•  $L_1 = \{ w \mid \text{ the length of } w \text{ is at most 5} \}$ 

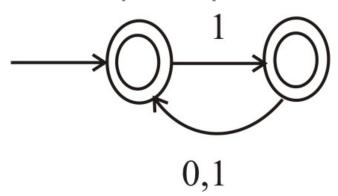
The state diagram of  $M_{\rm l}$  that recognizes  $L_{\rm l}$  is



•  $L_2 = \{w\}$  every odd position of w is a1

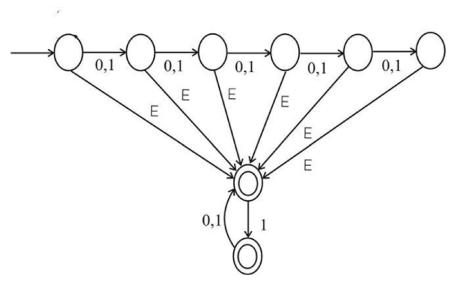
$$L_2 = (1\Sigma)^*$$

The state diagram of  $M_{\scriptscriptstyle 2}$  that recognizes  $L_{\scriptscriptstyle 2}$  is



 $\it L$  is concatenation of  $\it L_1$  and  $\it L_2$ 

So the state diagram of  $\boldsymbol{\mathit{M}}$  that recognizes  $\mathit{L}$  is described as follows



(b) Given Languages are

 $L_1 = \{ w \mid w \text{ contains at least three 1s} \} \text{ on } \Sigma = \{0,1\}$ 

浙ICP备16034203号-2

And  $L_2 = \{w \mid w \text{ is a empty set}\}\ \text{on } \Sigma = \{0,1\}$ 

 $\emph{M}_{\rm l}$  be the NFA that recognizes  $\emph{L}_{\rm l}$  and

 $M_{\rm 2}$  be the NFA that recognizes  $L_{\rm 2}.$ 

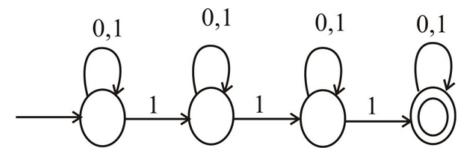
Let 
$$L = L_1 0 L_2$$

M be the NFA that recognizes L.

•  $L_1 = \{ w \mid w \text{ contains at least three1s} \}$ 

$$L_1 = (0,1)^* 1(0,1)^* 1(0,1)^* (0,1)^*$$

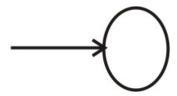
The state diagram of  $M_{\mathrm{I}}$  that recognizes  $L_{\mathrm{I}}$  is



•  $L_2 = \{ w \mid w \text{ is a empty set} \}$ 

$$L_2=\phi=\left\{\ \right\}$$

The state diagram of  $\boldsymbol{M_2}$  that recognizes  $\boldsymbol{L_2}$  is



 $\it L$  is concatenation of  $\it L_{\rm l}$  and  $\it L_{\rm 2}$ 

So the state diagram of M that recognizes L is described as follows

