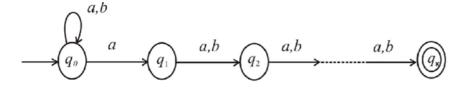
Given language is

$$C_{\scriptscriptstyle k} = \Sigma * a \; \Sigma^{k-1} \; \text{for each} \; \; K \geq 1 \text{, over the alphabet} \; \; \Sigma = \left\{a,b\right\}$$

 C_k is the language consisting of all strings that contains an 'a' exactly K places from the right – hand end.

Let N be the NFA with K+1 states that recognizes C_{K}

(i) The state diagram of NFA N is follows:



(ii) The formal description of NFA N is as follows:

$$N = (Q, \Sigma, \delta, q_0, F)$$

$$Q = \text{set of sates} = \{q_0, q_1, \dots, q_K\}$$

$$\Sigma = \text{set of alphabet } = \{a, b\}$$

$$q_0 = \text{start state } = \{q_0\}$$

$$F = \text{set of final states} = \{q_K\}$$

 δ =The transition function is given as follows:

$$\delta \left(q_{i},a\right) = \begin{cases} \left\{q_{0},q_{1}\right\} \text{ if } i=0\\ \left\{q_{i+1}\right\} & \text{if } 0 < i < k\\ \phi & \text{if } i=k \end{cases}$$

$$\begin{split} \delta\left(q_{i},b\right) = \begin{cases} \left\{q_{0}\right\} & \text{if } i = 0 \\ \left\{q_{i+1}\right\} & \text{if } 0 < i < k \\ \phi & \text{if } i = k \end{cases} \end{split}$$

$$\delta(q_i, \in) = \phi \forall i$$
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