
Question:

Let $D = \{xy \mid x, y \in \{0,1\}^* \text{ and } |x| = |y| \text{ but } x \neq y\}$. Show that D is a context-free language.

Answer:

----SETP1----

By the definition of **Context free language**, for showing that the language D is a CFL i.e. context free language, generate a context free grammar CFG G .

Consider the following grammar G :

$$\begin{aligned} S &\rightarrow AB \mid BA \\ A &\rightarrow 0 \mid 0A0 \mid 0A1 \mid 1A0 \mid 1A1 \\ B &\rightarrow 1 \mid 0B0 \mid 0B1 \mid 1B0 \mid 1B1 \end{aligned}$$

The given grammar $L(G)$ generates the language in the form $w_1 x w_2 v_1 y v_2$, where $|w_1| = |w_2| = k, |v_1| = |v_2| = l, x \neq y$ for $\Sigma = \{0,1\}^*$.

----SETP2----

- By the definition, any language which is generated by a context-free grammar is termed as a context-free language.
 - The grammar generated above is a Context Free Grammar. The language D can be generated using the above context free grammar G as follows:
 - A string is in D iff it can be written as xy with $|x| = |y|$ s.t. for some i , the i th character of x and y are different from one another. The above grammar can be used to obtain the required string by generating the i th characters and filling up with the remaining characters.
 - The generated language $w_1 x w_2 v_1 y v_2$ can be subjected to nested induction over k and l with case distinction over pairs (x, y) .
 - Now, w_2 and v_1 can exchange symbols because both carry symbols that are independent of the rest of the string.
 - Therefore, x and y in their respective half can have the same position, which implies $L(G) = D$ because G doesn't impose any restrictions on its language.
- Hence, the given language D is **context free language**.