Compute descriptive complexity of strings K(x) with an oracle for A_{TM} .

• For the given string x, start testing all the strings 'S' up to the length |x|+c

Where c = length of TM (Turing machine) that halts immediately upon starting.

- All the strings up to the length |x|+c are potential description of x.
- If S is well formed as $\langle M, w \rangle$ from all binary strings in lexicographic order, then we simulate M with input w and see if it halts with x on the tape.
- Here we do not know whether M will halt on input w or not.
- \bullet An oracle for ${\it A_{TM}}$ can determine this.

$$A_{TM} = \{\langle M, w \rangle | M \text{ is a } TM \text{ and } M \text{ accepts } w\}$$

- An oracle for A_{TM} will take $\left\langle M,w \right
 angle$ as input and determine whether M accepts w or not.
- If M doesn't halt we move on to the next string S, and so on.
- After that we will find lexicographically first string S among them.
- In this way shortest string will be determined and it is represented as minimal description d(x).
- From d(x), we find K(x) as

$$K(x) = |d(x)|$$

• By this procedure, we will calculate $\mathbf{K}(x)$ with an oracle for A_{TM} .

From the above procedure, K(x) is the descriptive complexity of strings and computed with an oracle A_{TM} .

Now, we have to know the definition of incompressible strings.

Incompressible strings:

Let x be a string. If x doesn't have any description shorter than itself then x is incompressible.

Let f be the function that is computable with an oracle for A_{TM} , for each n, calculate f(n) i.e., incompressible length of string n.

Create an oracle \emph{TM} (Turing machine) that does the following:

- 1. For a given number *n*, begin
- 2. Enumerate all strings x of length n.
- 3. For each x, calculate $\mathbf{K}(x)$ by using oracle for A_{TM} .
- 4. As soon as we find x = f(n) with K(x) >= |x| output x.

[K(x)] > |x| means getting a string which is greater than its description.]

5. Halt.