

获得的答案

Consider the language A is regular. Let $A_{\frac{1}{2}}$ be the set of all first halves of strings in A .

$$A_{\frac{1}{2}} = \{x \mid \text{for some } y, |x|=|y| \text{ and } xy \in A\}$$

Since A is a regular language, the DFA M recognizes the language A .

$$M = (Q, \Sigma, \delta, q_0, F)$$

where, Q is the set of states

Σ is the input alphabet

δ is the transition function

q_0 is the start state

F is the final state

The language is said to be regular if there exists an FA for it. In this case, construct an NFA N that recognizes $A_{\frac{1}{2}}$. Let x be the first part and choose y such that $|x|=|y|$. Here, $x \in A_{\frac{1}{2}}$. To prove the language $A_{\frac{1}{2}}$ is regular, run two DFAs at the same time one forward and the other backward. Run the DFA M on input x in forward direction and run the DFA M on input y in backward direction parallelly. The input string is accepted if both simulations reach the same state.

Construction of NFA N to recognize $A_{\frac{1}{2}}$:

Let $N = (Q', \Sigma, \delta', q_0', F')$ where,

(i) $Q' = Q \times Q \cup \{q_0'\}$ set of states contains the following:

- Special start state q_0' and
- A cross product $q \times q \times q$ where
- The first part tracks the performance of M on x
- The second part does the same thing for y .
- The third part tests whether the guess on M is consistent or not.

(ii) Σ = input alphabet

(iii) q_0' is the start state

(iv) $F' = \text{set of final states} = \{(q_i, q_j, q_k) \mid q_i, q_k, q_j \in Q\}$

(v) δ' = Rules of transition are as follows:

- There exists an ϵ move from the start state to the all the states in $\{(q_0, q_f) \mid q_f \in F\}$.
- Consider the states $q_i, q_j, q_k, q_l \in Q$. There exists a move from (q_i, q_j) to (q_k, q_l) on input symbol $a \in \Sigma$ if and only if $\delta(q_i, a) = q_k$ and $\delta(q_l, b) = q_j$.

The NFA N is constructed to recognize $A_{\frac{1}{2}}$. Thus, $A_{\frac{1}{2}}$ is regular.

Therefore, if A is regular then $A_{\frac{1}{2}}$ is also regular.