### ECE2800J

Programming and Introductory Data Structures

### **Abstract Data Types**

### **Learning Objectives:**

Understand what is an abstract data type (ADT)

Understand the usefulness of an ADT

Know how to define an ADT in C++

### Outline

- Introduction to Abstract Data Types
- Class in C++: A Trivial Example
- More Details on Class
- Another Class Example: a Mutable Set of Integers (IntSet)
- Improve the Efficiency of IntSet

- Suppose we want to build an abstraction that holds a **mutable** set of integers.
- This is a **<u>set</u>** in the mathematical sense:
  - A collection of zero or more integers, with **no duplicates**.
- The set is "mutable" because we can insert objects into and remove objects from the set.

- Suppose we want to build an abstraction that holds a **mutable** set of integers.
- There are four **operations** on this set that we will define:
  - 1. Insert a value into the set.
  - 2. Remove a value from the set.
  - 3. Query to see if a value is in the set.
  - 4. Count the number of elements in the set.

**Using Classes** 

• Here is an **incomplete** definition of a class implementing such an ADT: class IntSet { // OVERVIEW: a mutable set of integers public: void insert(int v); // MODIFIES: this // EFFECTS: this = this + {v} void remove(int v); // MODIFIES: this // EFFECTS: this = this - {v} bool query(int v); // EFFECTS: returns true if v is in this, false otherwise int size(); // EFFECTS: returns |this|.

```
class IntSet { // omitted OVERVIEW for space
  public:
    void insert(int v); // omitted RME for space
    void remove(int v); // omitted RME for space
    bool query(int v); // omitted RME for space
    int size(); // omitted RME for space
};
```

- The class is incomplete because we haven't chosen a representation for sets.
- Choosing a representation involves two things:
  - Deciding what **concrete** data elements will be used to **represent the values** of the set.
  - Providing an **implementation** for each **method**.

```
class IntSet { // omitted OVERVIEW for space
  public:
    void insert(int v); // omitted RME for space
    void remove(int v); // omitted RME for space
    bool query(int v); // omitted RME for space
    int size(); // omitted RME for space
};
```

- Despite not having a representation for a set, the (incomplete) definition above is all that a **customer** of the IntSet abstraction needs to know since it has:
  - The general overview of the ADT.
  - The specification of each method.

#### **Using Classes**

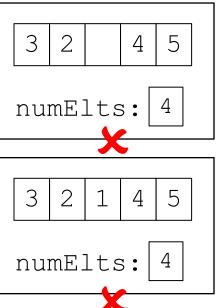
- Start with a representation for the set itself:
  - Use an array.
  - Represent a set of size N as an **unordered** array of integers with no duplicates, stored in the first N slots of the array.
  - int numElts: maintains the number of elements currently in the array.
- These last two statements are called **representation invariants** or **rep invariants** (more on this later).
- This invariant is a rule that the representation must obey both **immediately before** and **immediately after** any method's execution.

rep

invariant

**Using Classes** 

- Start with a representation for the set itself:
  - Use an array.
  - Represent a set of size N as an **unordered** array of integers with no duplicates, stored in the first N slots of the array.
  - int numElts: maintains the number of elements currently in the array.



rep

invariant

#### **Using Classes**

• Since this is an array, and arrays have maximum sizes, we have to choose a maximum size and modify the OVERVIEW:

```
// OVERVIEW: a mutable set of
// integers, |set| <= 100</pre>
```

• We also have to change the EFFECTS clause of insert:

```
// EFFECTS: this = this + {v} if
// room available, throws int
// 100 otherwise
```

```
const int MAXELTS = 100;
class IntSet {
    // OVERVIEW: a mutable set of integers ( |set | <= MAXELTS
              elts[MAXELTS]
    int
              numElts;
    int
                                   Use a global constant like we
 public:
    void insert(int v);
                                   have talked about.
      // MODIFIES: this
      // EFFECTS: this = this + {v} if room,
                 throws int MAXELTS otherwise
    void remove(int v);
      // MODIFIES: this
      // EFFECTS: this = this - {v}
   bool query(int v); // RME omitted for space
    int size();  // RME omitted for space
};
```

**Using Classes** 

Given this representation, and the representation invariants, we can write the methods.

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
   int      elts[MAXELTS];
   int      numElts;
   public:
      void insert(int v); // RME omitted for space
      void remove(int v); // RME omitted for space
      bool query(int v); // RME omitted for space
      int size(); // RME omitted for space
};
```

```
int IntSet::size() {
  return numElts;
}
```

Because our rep invariant says that numElts is always the size of the set, we can return it directly.

- Next, consider the three final routines:
  - query: search the array looking for a specific number.
  - remove: search the array for a number; if it exists, remove it.
  - insert: search the array for a number; if it doesn't exist, add it.
- All three of these have "search" in common.
- One might be tempted to just write insert and remove in terms of query, will this work?
  - Hint: think about remove.
- query only tells us **whether** the element exists, not **where** we need one more method...

**Using Classes** 

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
  int        elts[MAXELTS];
  int        numElts;

int indexOf(int v);

      // EFFECTS: returns the index of
       // v if it exists in the
       // array, MAXELTS otherwise.
```

```
public:
    void insert(int v);
    void remove(int v);
    bool query(int v);
    int size();
};
```

Note: This member function must be **private**. This is because it exposes details about the concrete representation. It is inappropriate to expose these details to a user of this class.

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
   int
           elts[MAXELTS];
   int
           numElts;
   int indexOf(int v); // RME omitted for space
 public:
   void insert(int v); void remove(int v); // RME omitted
   bool query(int v); int size();  // RME omitted
};
int IntSet::indexOf(int v) {
  for (int i = 0; i < numElts; i++) {
    if (elts[i] == v) return i;
  return MAXELTS;
```

**Using Classes** 

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
   int    elts[MAXELTS];
   int    numElts;
   int indexOf(int v); // RME omitted for space
   public:
    void insert(int v); void remove(int v); // RME omitted
   bool query(int v); int size(); // RME omitted
};
```

With indexOf, query is trivial...

```
bool IntSet::query(int v) {
  return (indexOf(v) != MAXELTS);
}
```



### How to Implement insert (v)?

Select all the correct answers.

- A. We can first search v to check if it is already there with indexOf (v)
- **B.** If v is not present, we then add v
- C. If we add v, it should be added as elts [numElts-1] before we increment numElts
- D. If v is added, we need to increment numElts



- The code for insert is not much more difficult than query:
  - First look for the indexOf the element to insert.
  - If it doesn't exist, we need to add this element to the **end** of the array.
  - What is the index of the current "end"?

- Place the element in the next slot and update numElts.
- The only exception to this is if numElts already equals MAXELTS.

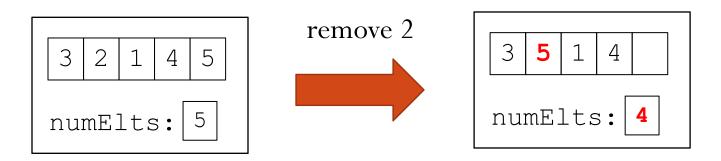
```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
   int.
         elts[MAXELTS];
   int numElts;
   int indexOf(int v); // RME omitted for space
 public:
   void insert(int v); void remove(int v); // RME omitted
   bool query(int v); int size();  // RME omitted
};
void IntSet::insert(int v) {
  if (indexOf(v) == MAXELTS) {
    if (numElts == MAXELTS) throw MAXELTS;
    elts[numElts++] = v;
```

### How about Remove?

- If the element (its index is called the victim) is in the array, we have to remove it leaving a "hole" in the array.
- What representation invariants are violated?
  - How can we fix them?

### How about Remove?

- Instead of moving each element after the victim to the left by one position, pick up the current "last" element and move it to the hole.
- This also breaks the invariant on numElts, so we must fix it.



```
void IntSet::remove(int v) {
  int victim = indexOf(v);
  if (victim != MAXELTS) {
    elts[victim] = elts[numElts-1];
    numElts--;
  }
}
```

**Using Classes** 

• <u>Question</u>: There is one problem with our implementation. What is it?

• <u>Hint</u>: Consider the newly-created set:

```
IntSet s;
```

What does the computer actually create when we declare S?

- <u>Question</u>: There is one problem with our implementation. What is it?
- Answer: On creation, S's data members are uninitialized!
- This means that the value of numElts could be a random value, but our representational invariant says it must be zero!
- How can we fix this?

**Automatically Initializing Classes** 

- Using constructor!
- The constructor (really, the **default** constructor) has the following type signature:

```
class IntSet { // OVERVIEW omitted for space
    ...
    public:
        IntSet();
        // EFFECTS: creates an empty IntSet
        ...
};
```

**Automatically Initializing Classes** 

```
IntSet();
  // EFFECTS: creates an empty IntSet
```

- The name of the function is the same as the name of the class.
- This function doesn't have a return type.
- It also does not take an argument in this case.
- It is guaranteed to be the **first** function called immediately after an object is created.
- It builds a "blank" uninitialized IntSet and makes it satisfy the rep invariant.

**Automatically Initializing Classes** 

```
IntSet();
  // EFFECTS: creates an empty IntSet
```

• Here's how it's written:

```
IntSet::IntSet(): numElts(0)
{
}
```

**Automatically Initializing Classes** 

```
IntSet::IntSet()
     : numElts(0)
{
}
```

```
Class_T::Class_T(): anInt(0),
     aDouble(1.2),
     aString("Yes")
{
}
```

- This syntax is called "initialization syntax".
- Each data member is initialized this way.
- <u>Note</u>: The order in which elements are initialized is the order they **appear in the definition**, NOT the order in the initialization list. It is a good practice to keep them in the same order to avoid confusion.

**Automatically Initializing Classes** 

• Alternatively, we could write this function as follows, but this is not considered as a good way!

```
IntSet::IntSet()
{
   numElts = 0;
}
```



A Benefit of Classes

• Now, instead of writing this:

```
void add_one (int a[], int elts);
```

and having to worry about the number of elements in the array, all we have to write is this:

```
void add_one (IntSet& set);
```

and we no longer have to worry about the array and its count being separated.

• A slight change to the class definition: const int MAXELTS = 100; class IntSet { int elts[MAXELTS]; int numElts; int indexOf(int v) const; public: void insert(int v); void remove(int v); bool query(int v) const; int size() const; **}**;

#### int size() const;

- Each member function of a class has an extra, implicit parameter named **this**.
  - "this" is a pointer to the current instance on which the function is invoked.
- **const** keyword modifies the implicit **this** pointer: **this** is now a pointer to a **const instance**.
  - <u>Means</u>: the member function **size()** cannot change the object on which **size()** is called.
  - By its definition, **size()** shouldn't change the object! Adding **const** keyword prevents any accidental change.
  - It is a good practice to add const keyword when possible!

```
• Implement size()
  int IntSet::size() const {
    return numElts;
}
The function body is the same as before.
```

• A **const** object can only call its **const** member functions!

```
const IntSet is;
cout << is.size(); ✓
is.insert(2); ✗</pre>
```

• If a const member function calls other member functions, they must be **const** too!

```
void A::g() const { f(); }
```





### Outline

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- Class in C++: A Trivial Example
- More Details on Class
- Another Class Example: a Mutable Set of Integers (IntSet)
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?

# How many elements of the array must indexOf examine?

• Suppose numElts >= 1.
int IntSet::indexOf(int v) {
 for (int i = 0; i < numElts; i++) {
 if (elts[i] == v) return i;
 }
 return MAXELTS;
}</pre>

- A. In the best case, 1 element
- **B.** In the worst case, numElts elements
- C. In the worst case, MAXELTS elements
- **D.** None of the above



- We say the time for indexOf grows **linearly** with the size of the set.
- If there are N elements in the set, we have to examine all N of them in the worst case. For large sets that perform lots of queries, this is too expensive!
- Luckily, we can replace this implementation with a different one that can be more efficient. The only change we need to make is to the **representation (implementation)** the abstraction can stay precisely the same.

- Still use an array to store the elements of the set and the values will still occupy the first numElts slots.
- However, now we'll keep the elements in **sorted** order.

# Question: Which Member Functions of the Class Should Be Changed?

```
const int MAXELTS = 100;
class IntSet {
    // OVERVIEW: a mutable set of integers
    int elts[MAXELTS];
    int numElts;
    int indexOf(int v) const;
 public:
    IntSet();
    void insert(int v);
    void remove(int v);
    bool query(int v) const;
    int size() const;
};
```

Improving Efficiency

• The constructor and size methods don't need to change at all since they just use the numElts field.

• query also doesn't need to change.

```
bool IntSet::query(int v) {
    return (indexOf(v) != MAXELTS);
}
```

- indexOf also doesn't need to change.
- However, insert and remove do need to change.

Improving Efficiency

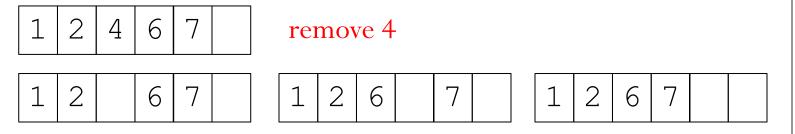
- We'll start with the easiest one: remove.
- Recall the old version that moved the last element from the end to somewhere in the middle, this will break the new "sorted" invariant.



• Instead of doing a swap, we have to "squish" the array together to cover up the hole.



- How are we going to do the "squish"?
  - Move the element next to the hole to the left leaving a new hole.
  - Keep moving elements until the hole is "off the end" of the elements.



- We'll reuse the variable victim as a loop variable.
- victim's invariant is that it always points at the hole in the array.

```
void IntSet::remove(int v) {
  int victim = indexOf(v);
  if (victim != MAXELTS) {
      // victim points at hole
    numElts--; // one less element
    while (victim < numElts) {</pre>
      // ..hole still in the array
      elts[victim] = elts[victim+1];
      victim++;
```

Improving Efficiency

• We also have to change insert since it currently just places the new element at the end of the array. This will also break the new "sorted" invariant.



- How are we going to do the insert?
  - Start by moving the last element to the right by one position.
  - Repeat this process until the correct location is found to insert the new element.
  - Stop if the start of the array is reached or the element is sorted.
  - We'll need a new loop variable called cand(idate) to track this movement.
  - Its invariant is that it always points to the next element that <a href="might"><u>might</u></a> have to move to the right.

```
void IntSet::insert(int v) {
  if (indexOf(v) == MAXELTS) { // duplicate not found
    if (numElts == MAXELTS) throw MAXELTS; // no room
    int cand = numElts-1; // last element
    while ((cand \geq 0) && elts[cand] \geq v) {
      elts[cand+1] = elts[cand];
      cand--;
    // Now, cand points to the left of the "gap".
    elts[cand+1] = v;
    numElts++; // repair invariant
    insert 5
                                                        4 | 5 |
                                   2
                                                      2
                                      4
                                              7
                                                   1
                       cand
                                       cand
                                                       cand
```

```
void IntSet::insert(int v) {
  if (indexOf(v) == MAXELTS) { // duplicate not found
    if (numElts == MAXELTS) throw MAXELTS; // no room
    int cand = numElts-1; // last element
    while ((cand >= 0)) && elts[cand] > v) {
      elts[cand+1] = elts[cand];
      cand--;
                                 Note: We are using the
    // Now, cand points to the "short-circuit" property
    elts[cand+1] = v;
                                 of &&. If cand is not
    numElts++; // repair invar;
                                 greater than or equal to
                                 zero, we never evaluate
                                 the right-hand clause.
```

```
void IntSet::insert(int v) {
  if (indexOf(v) == MAXELTS) { // duplicate not found
    if (numElts == MAXELTS) throw MAXELTS; // no room
    int cand = numElts-1; // largest (last) element
    while ((cand \geq 0) && elts[cand] \geq v) {
      elts[cand+1] = elts[cand];
      cand--;
    // Now, cand points to the left of the "gap".
    elts[cand+1] = v;
    numElts++; // repair invariant
         Question: What is the situation when the loop terminates due
         to cand < 0? Is our implementation correct?
```

Improving Efficiency

• **Question**: Do we have to change indexOf?

```
int IntSet::indexOf(int v) {
  for (int i = 0; i < numElts; i++) {
    if (elts[i] == v) return i;
  }
  return MAXELTS;
}</pre>
```

- Question: Do we have to change indexOf?
- <u>Answer</u>: No, but it can be made more efficient with the new representation.
- How? Using binary search! (The array is sorted)

```
int IntSet::indexOf(int v) {
  for (int i = 0; i < numElts; i++) {
    if (elts[i] == v) return i;
  }
  return MAXELTS;
}</pre>
```

Complexity

	Unsorted	Sorted
query	O(N)	?
insert	?	?
remove	?	?

Complexity

	Unsorted	Sorted
query	O(N)	O(log N)
insert	O(N)	O(N)
remove	O(N)	O(N)

insert and remove are still linear, because they may have to "swap" an element to the beginning/end of the array.

Complexity

	Unsorted	Sorted
query	O(N)	O(log N)
insert	O(N)	O(N)
remove	O(N)	O(N)

- If you are going to do more searching than inserting/removing, you should use the "sorted array" version, because query is faster there.
- However, if query is relatively rare, you may as well use the "unsorted" version. It's "about the same as" the sorted version for insert and remove, but it's MUCH simpler!

### References

- Problem Solving with C++ (8<sup>th</sup> Edition)
  - Chapter 10.3 Abstract Data Types
  - Chapter 10.2 Classes and constructors
- C++ Primer, 4th Edition
  - Chapter 7.7.1 const Member Function