

Binary Search

- Efficient searching method, which consumes less time
- The essential thing here is the array should be sorted first.
- Steps for binary search

Let the input array be $A[0 \dots n-1]$

Key is the element to be searched.

$A[m]$ is the mid element of array A .

Then,

- ① Check if $A[m] = \text{key}$
- ② If $A[m] > \text{key}$, then go to left sub list.
- ③ If $A[m] < \text{key}$, then go to right sub list.

Example

Input array $\rightarrow 0 \ 5 \ 10 \ 15 \ 20 \ 25 \ 30 \ 35 \ 40$

0	1	2	3	4	5	6	7	8
0	5	10	15	20	25	30	35	40
↑								↑
Low								High

Let $\text{key} = 5$ (No. to be searched)

Now find the middle element

$$\begin{aligned} m &= (low + high) / 2 \\ &= (0 + 8) / 2 \\ &= 4 \end{aligned}$$

$$\therefore A[m] = 20$$

Check if $A[m] = \text{key}$

Here, $A[m] = 20$

$\text{key} = 5$

$\therefore A[m] \neq \text{key}$

Now, $\text{key} < A[m]$

\therefore we have to search the left sub-list

0	1	2	3
0	5	10	15
↑			↑
Low			High

$$\begin{aligned} m &= (\text{Low} + \text{high}) / 2 \\ &= (0 + 3) / 2 \\ &= 1 \end{aligned}$$

Now $\text{key} = 5$
 $A[m] = 5$

$\therefore \text{key} = A[m]$

Thus, the no. is present in the list.

Algorithm for binary search

(1) Initialize array $A[0 \dots n-1]$

(2) $low \leftarrow 0$

(3) $high \leftarrow n-1$

(4) $Do((low+high)/2) \rightarrow m$

(5) if $(key = A[m])$ then

return m ;

(6) if $(key < A[m])$ then

$high \leftarrow m-1$ [Search left sub tree]

(7) if $(key > A[m])$ then

$low \leftarrow m+1$ [Search right sub tree]

(8) return -1 (if element is not present)