

Name - Rakesh Suthar

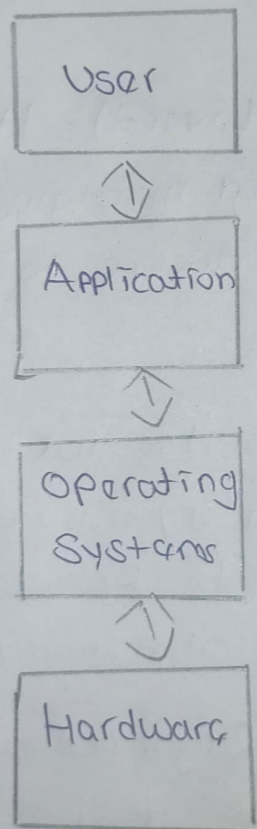
Registration No - 22BC010178

Subject - Operating Systems

Faculty - Dr. Chour Singh Rajpoot

### Mid Term Exam

#### 1) The Structure of Computer Systems



The layered approach is a design paradigm where an operating system is divided into distinct layers, each performing specific functions and relying on the services of the layers below it.

- \* **Hardware Layer (Layer 0)** - Directly interacts with the system physical components (CPU, memory, disk drives, etc). Provides basic services like memory allocation, device I/O, and interrupt handling.
- \* **Device Drivers Layer (Layer 1)** - Manages specific hardware devices, translating high level requests into low level commands. Hides device specific components.
- \* **I/O Management Layer (Layer 2)** - Handles input/output operations, buffering data and managing device queues. Coordinates with device drivers to ensure efficient data transfer.
- \* **Memory Management** - Allocates and deallocates memory to processes, optimizing memory usage. Implements virtual memory techniques to provide a larger address space than physical memory.
- \* **File System (Layer 4)** - Organizes and manages files and directories on storage devices. Provide services for creating, deleting, reading, and writing files.
- \* **Kernel Layer (Layer 5)** - Core of the Operating System. Responsible for process scheduling, inter process communication and system security. Manages system resources and handles system calls.



\* **User Interface Layer (Layer 6)** - Provides the user interface, including command line interfaces, graphical user interfaces, and application programming interfaces (APIs). Interacts with the user and translates user commands into system calls.

### Advantages -

- \* **Modularity** - Each layer is independent, making it easier to develop, test and modify.
- \* **Debugging** - Issues can be isolated to specific layers, simplifying troubleshooting.
- \* **Security** - Layers can be designed with varying levels of privilege, enhancing system security.
- \* **Flexibility** - New features can be added or modified by modifying layers.

### Disadvantages -

- \* **Performance Overhead** - Communication between layers can introduce overhead due to context switching and data copying.

2)

## (i) MultiProgramming v/s Multitasking

### Multi Programming -

- \* Multiple Programs are loaded into main memory simultaneously.
  - \* It increases CPU utilization by organizing jobs so that the CPU always has one job to execute.
  - \* Multiprogramming operating systems use the mechanism of job scheduling and CPU scheduling.
- Ex - A batch processing system where multiple jobs are executed sequentially, but the OS switches between them to prevent CPU idle time.

### Multitasking -

- \* A Single Program is divided into multiple tasks.
  - \* The OS switches between these tasks to give the illusion of concurrent execution.
- Ex - A user can ~~handle~~ run multiple applications simultaneously on single computer.



## Process V/s Program

### Program

Process -

- \* A static entity, a sequence of instructions stored in secondary memory.
- \* It doesn't consume system resources until it's executed
- \* It doesn't perform any action on its own

Ex - A .exe file or a Python Script

### Process -

- \* A dynamic entity, an instance of a program being executed.
- \* It consumes system resources (CPU, memory I/O devices) during its execution.
- \* Multiple processes can be instances of the same program running simultaneously

Ex - When you run a Python Script, the OS creates a process to execute its instructions

# Time Sharing v/s Parallel Processing System

## Time Sharing -

- \* Multiple users share a single computer system
- \* The OS allocates time slices to each user's process giving the illusion of simultaneous execution
- \* The time that each task gets to execute is called quantum

Ex - A mainframe computer used by multiple users for different tasks.

## Parallel Processing System -

Multiple processors or cores execute different parts of a program simultaneously.

This is used to speed up execution time for computationally intensive tasks.

Ex - A supercomputer with multiple processors solving complex scientific simulations.



3)

Inter Process Communication - It is a type of mechanism usually provided by the operating system. The main aim or goal of this mechanism is to provide communication in between several processes. This communication could involve a process letting another process know that some event has occurred or the transferring of data from one process to another.

Need of Inter Process -

- \* It helps to speed up modularity
- \* Computational
- \* Privilege Separation
- \* Convenience
- \* Helps Operating System communicate with each other synchronize their actions

# Approaches to Inter Process Communication

- \* Pipes
- \* Shared memory
- \* Message Passing
- \* Direct Communication
- \* Message Queue
- \* Indirect Communication
- \* FIFO
- \* Socket
- \* File
- \* Signal

**Mutual Exclusion** - It is generally required that only one process thread can enter the critical section at a time. This also helps in synchronization and creates a stable state to avoid the race condition.

**Race Condition** - A situation that occurs when two or more processes access and manipulate a shared resource concurrently, and the outcome of the execution depends on the specific timing of their actions. Race Conditions can lead to unpredictable and often incorrect results.



**Critical Section** - It is a code segment that accesses a shared resource and must be executed by only one process at a time to avoid race conditions.

The main goal of managing a critical section is to ensure data integrity by allowing only one process to execute in that section at a time, typically using mutual exclusion techniques like semaphores or mutexes.

### Key Points

- \* Entry Section
- \* Exit Section
- \* Remainder Section

#### 4) Necessary Conditions For Deadlock

- \* Mutual Exclusion - Atleast one resource can be held in a non-shareable mode, meaning only one process can use it at a time.
- \* Hold and wait - Atleast one process must be holding atleast one resource and waiting to acquire additional resources held by other processes.
- \* Non Preemption - Resources cannot be forcibly taken away from a process holding them; they must be released voluntarily.
- \* Circular Wait - A set of processes must exist each of which is waiting time for a resource held by the next process in set.

#### Resource Graph model -

A Resource graph model is a directed graph that depicts the current state of resource allocation among processes.

Process Nodes - Represent Process in System

Resource Nodes - Represent Resource in System.



- \* Request Edge - From a process node to a resource node, indicating that the process is requesting the resource.
- \* Assignment Edge - From a resource node to a process node, indicating that the resource has been allocated to the process.

### Safe and Unsafe States

A system is in a safe state if there exists a sequence of process executions that allows each process to finish without causing a ~~deadlock~~ ~~dead allocation~~ deadlock. In other words, there exists a safe sequence of resource allocation and deallocation that avoids feedback.

### Key Points -

- \* A System can be both safe and unsafe, depending on the order of process execution.
- \* Deadlock avoidance and detection techniques are employed to prevent or detect unsafe states.
- \* Resource allocation strategies, such as banker's algorithm, can be used to ensure system safety.

5)

Process	Arrival Time	Burst Time	Priority
P <sub>1</sub>	0	8	3
P <sub>2</sub>	1	4	2
P <sub>3</sub>	2	9	1
P <sub>4</sub>	3	5	4

Grant Chart

P <sub>1</sub>	P <sub>2</sub>	P <sub>4</sub>	26
0	8	12	17

$$TAT = CT - AT$$

$$WT = TAT - BT$$

Process	AT	BT	CT	TAT	WT
P <sub>1</sub>	0	8	8	8	0
P <sub>2</sub>	1	4	12	11	7
P <sub>3</sub>	2	5	17	14	9
P <sub>4</sub>	3	9	26	24	15

$$\text{Average Waiting Time} = (0 + 7 + 9 + 15) / 4$$

$$= 7.75$$

$$\text{Average TAT} = (8 + 11 + 14 + 24) / 4 = 14.25$$



Priority

Grant Chart

P <sub>1</sub>	P <sub>3</sub>	P <sub>2</sub>	P <sub>4</sub>
0	8	17	26

Process	AT	BT	CT	TAT	WT
P <sub>1</sub>	0	8	8	8	0
P <sub>3</sub>	2	9	17	15	6
P <sub>2</sub>	1	4	21	20	16
P <sub>4</sub>	3	5	26	23	18

Average Waiting Time =  $(0+6+16+18)/4 = 6$

Average Turn Around Time =  $(8+15+20+23)/4$   
 $= 16.5$