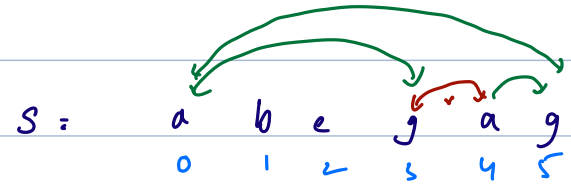
Question: Given a string of lowercase alphabets, return the count of pairs[indices] (i, j) such that i < j and s[i] = 'a' and s[j] = 'g'

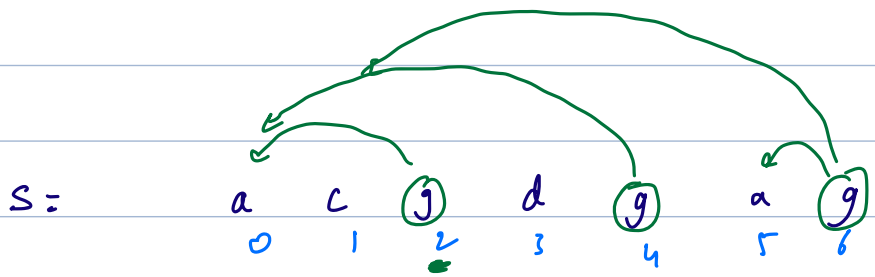
$len(s) \leq 10^6$



(0, 3)

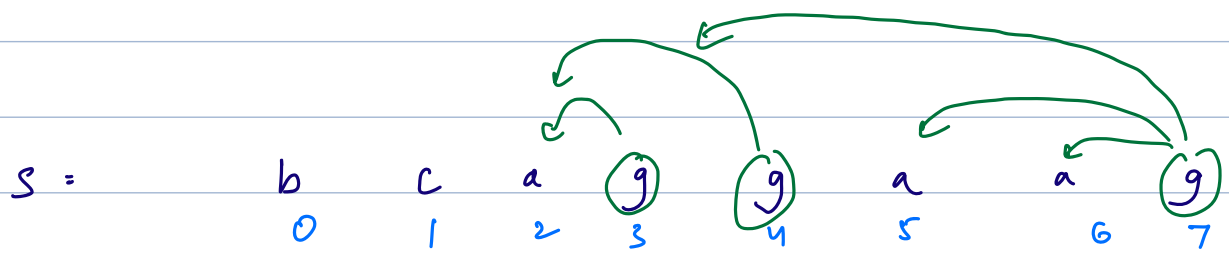
(0, 5)

(4, 5)



(0, 2) (0, 4) (0, 6) ⇒ 4

(5, 6)



(2, 3) (2, 4) (2, 7)

(5, 7) ⇒ 5 pairs

(6, 7)

Brute Force

Iterate the string, for every 'g' count the no. of pairs it forms by counting #a's to its left

```
count = 0;
for(int j = 0; j < N; j++){
      if(s[j] == 'g'){
            for(i = 0; i < j; i++){
                  if(s[i] == 'a'){
                        count++;
                  }
            }
      }
}
```

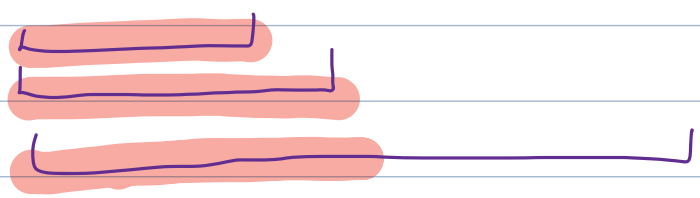
T.C: $O(N^2)$

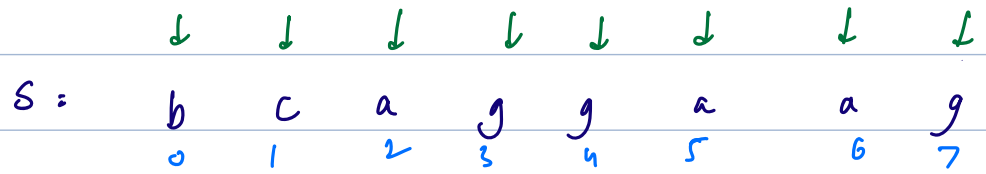
S.C: $O(1)$

Efficient:

S:

b	c	a	g	g	a	a	g
0	1	2	3	4	5	6	7





result = 0 + 1 + 1 + 3 = 5

count_a = 0 1 2 3

```
function count_ag(str) {
    result = 0;
    count_a = 0;
    for(i -> 0 to n-1) {
        if(str[i] == 'a') {
            count_a++;
        }
        else if(str[i] == 'g') {
            result += count_a;
        }
    }
    return result;
}
```

N

T.C : O(N)

S.C : O(1)

CARRY FORWARD

SUBARRAY

-> A contiguous part of an array
-> Single element can also be a subarray
-> Entire array is a subarray

A: 3 4 5 6 -2 8 10
0 1 2 3 4 5 6

[5, 6, -2] ✓ => (2, 4)
[3, 4, 6, -2] ✗
[3, 4, 5, 6, -2] ✓ => (0, 4)
[5, 4, 3] ✗
[5] ✓ (2, 2)
[3 4 5 6 -2 8 10] ✓ ( )

Quiz: 2 4 1 6 -3 7 9 4

How to represent a subarray?

(Start Index, end Index)

Quiz: [4, 2, 10, 3, 12, -2, 15]

0 1 2 3 4 5 6

start = 0

end : [0, 6] => 6 - 0 + 1 = 7

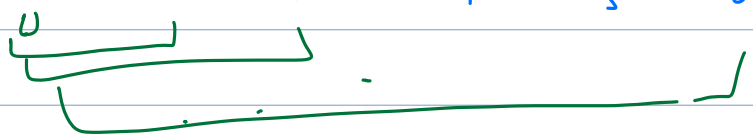
(0,0) (0,1) (0,2) (0,3) (0,4) (0,5) (0,6)

7 Subarrays

Quiz:

4 2 10 3 12 -2 15

0 1 2 3 4 5 6 N=7



start = 1

end = [1, N-1] => [1, 6] => 6 - 1 + 1 = 6

Question: Total No. of Subarrays

#subarrays Starting at index 0 => [0, N-1] (end) : N

+

#subarrays Starting at index 1 => [1, N-1] : N-1

+

#subarrays Starting at index 2 => [2, N-1] : N-2

+

#subarrays Starting at index 3 => [3, N-1] : N-3

+

⋮

#subarrays Starting at index N-1 [N-1, N-1] : 1

$$1+2+3+\ldots\ldots N-2+N-1+N = \boxed{\frac{N(N+1)}{2}} \Rightarrow$$

Question: Print a subarray

```
void printSubarray(int A[], int s, int e){
    for(i=s; i<=e; i++){
        print(A[i]);
    }
    print("\n");          // Go to new line
}
```

(Array)

T.C: O(N)

can we print an array is less than O(N)?

Question: Print all subarrays of an array

A = [1, 2, 3]
    0 1 2

N = 3

$\frac{N(N+1)}{2} = \frac{3 \cdot 4}{2} = 6$

s	e	Subarray
0	0	[1]
0	1	[1, 2]
0	2	[1, 2, 3]
1	1	[2]
1	2	[2, 3]
2	2	[3]

```
for(start = 0;  start < N;  start++) {
    for(end = start;  end < N;  end++) {
        => // (start, end) represent a subarray
        printSubarray(A, start, end);
    }
}
```

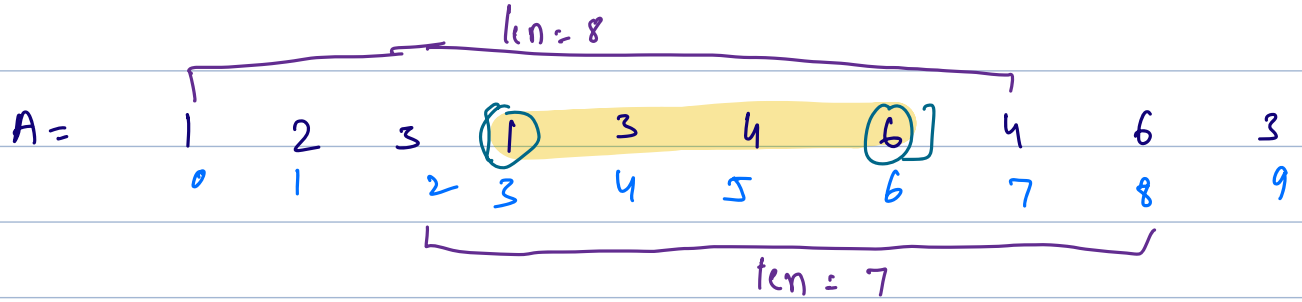
O(N) — printSubarray(A, start, end)

A — Array

T.C : $N^2 \times O(N) \Rightarrow O(N^3)$

S.C : $O(1)$

Question: Given an array, return the length of the smallest subarray (contiguous part of array) which contains both the min and max of the array



Max(A) = 6

Min(A) = 1

len = 4

8 : 32

A =	2	2	6	4	5	1	5	2	6	4	1
	0	1	2	3	4	5	6	7	8	9	10

Max(A) = 6

Min(A) = 1

len = 3

Brute Force

=> Find Max and Min => O(N)

=> Iterate over all the $\frac{N(N+1)}{2}$ arrays, for each subarray, check if it has max and min, if yes update your answer

T.C: $O(N) + \frac{N(N+1)}{2} \times O(N) = O(N^3)$

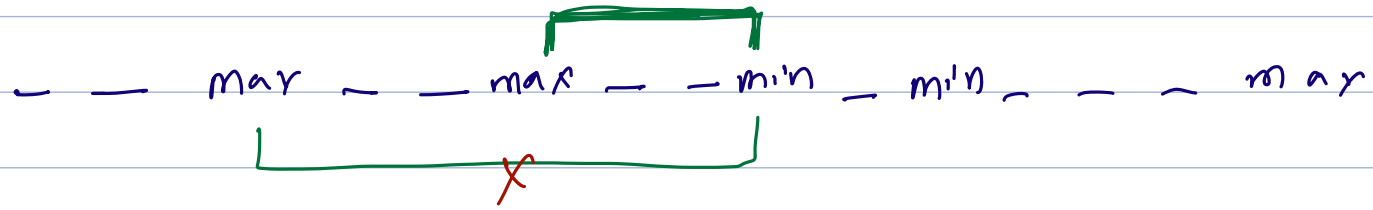
S.C: O(1)

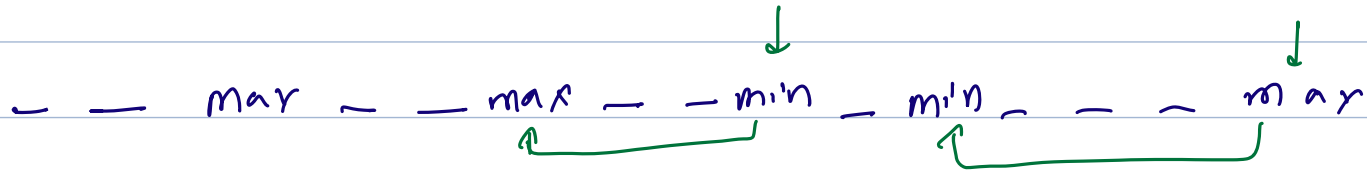
Approach:

Observations:

1) the corner elements of the answer subarray are max and min

2) Answer subarray will have exactly 1 max and 1 min





-) Carry foward index of latest max and latest min.

-) when you get a min element, use the index of latest max

-) when you get a max element, use the index of latest min

A:-	2	2	6	4	5	1	5	2	6	4	1
	0	1	2	3	4	5	6	7	8	9	10

Max(A) = 6

Min(A) = 1

latest_max = ~~-1~~ ~~7~~ 8

latest_min = ~~-1~~ ~~5~~ 10

Ans: ~~INT_MAX~~ ~~4~~ 3

(i - latest_max + 1)

A = 4 4 4

0 1 2

latest_max = -1 0 1 2

latest_min = -1 0 1 2

ans = INF 1

```
// Find max (maxm.)
// Find min (minm)
latest_max = -1;
latest_min = -1;
ans = INFINITY;
for(int i = 0; i < n; i++){
      if(A[i] == maxm){
            if(latest_min != -1){
                  ans = min(ans, i - latest_min + 1);
                  latest_max = i;
            }
            latest_max = i;
      }
      if(A[i] == minm){
            if(latest_max != -1){
                  ans = min(ans, i - latest_max + 1);
                  latest_min = i;
            }
            latest_min = i;
      }
}
return ans;
```

O(1)

O(1)

ans = min(5, 3) = 3

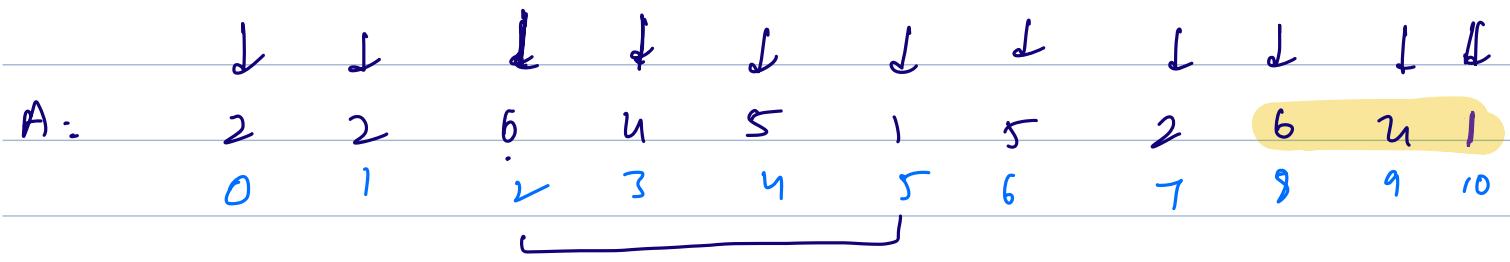
T.C : O(N)

S.C : O(1)

latest_min

len = i - latest_min + 1 =>

[a, b] = b - a + 1



Min(A) = 1

Max(A) = 6