CN Assignment 2
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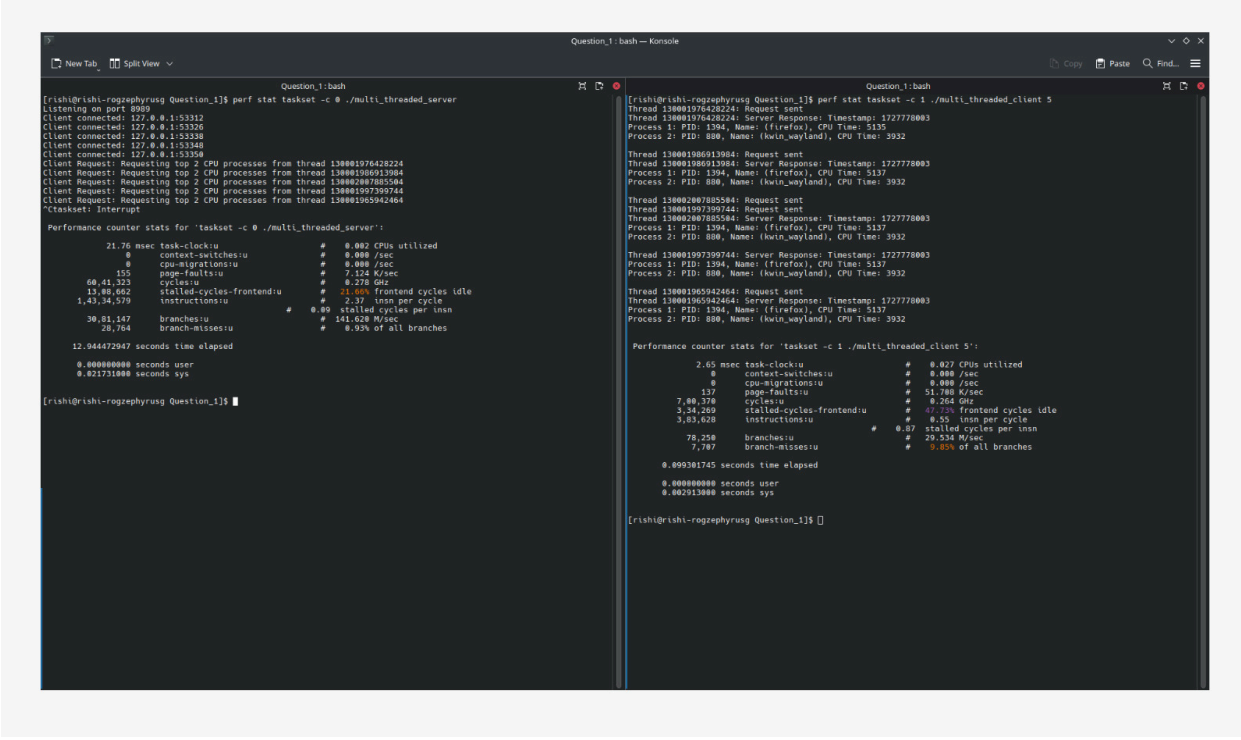
github

Q.1. Write a client-server socket program in C. The client process and the server process should run on separate VMs (or containers) and communicate with each other. Use "taskset" to pin the process to specific CPUs. This helps measure the performance. Your program should have the following features. [1+(1+2)+2+(1+2)+1+1]

1. The server sets up a TCP socket and listens for client connections.
2. The server should be able to handle multiple concurrent clients (multithreaded, concurrent server). The server accepts the client connection; hence, a new socket is created with 4 tuples (server IP, server listening port, client IP, client port). The server creates a new thread that continues to process the client connection (Hint: Use pthread library for multithreading). The original server socket continues to listen on the same listening port for newer incoming client connections.
3. The client creates a socket and initiates the TCP connection. Your client process should support initiating "n" concurrent client connection requests, where "n" is passed as a program argument.
4. After the client connection is established, the client sends a request to the server to get information about the server's top TWO CPU-consuming processes. The server finds out the top CPU-consuming process (user+kernel CPU time) and gathers information such as process name, pid (process id), and CPU usage of the process in user & kernel mode (you can report this time in clock ticks).
    - ○ Possible solution approach: The server should make use of the open() system call to read the proc" filesystem to get this information. You need to read /proc/[pid]/stat for all processes to parse the process name, pid, and CPU time (user+kernel). To understand the format of /proc/[pid]/stat file, refer the source code of show_stat function is available here for reference. More about "proc" file system is here.
5. Server sends the information collected in step (4) to the client.
6. The client prints this information & closes the connection.



Q.2. The socket programming source code that leverages “select” system call here. Modify the server code as per Q.1. Use the perf tool to analyze the performance of the following: [3+3+3]

(a) Single-threaded TCP client-server

(b) Concurrent TCP client-server

(c) TCP client-server using “select”

You can be creative for this analysis. You may take readings of various performance counters (CPU clocks, cache misses, context switches, etc.) across number of concurrent connections, etc. This is an OPEN question; you can be as comprehensive as you can.

Single-threaded TCP client-server

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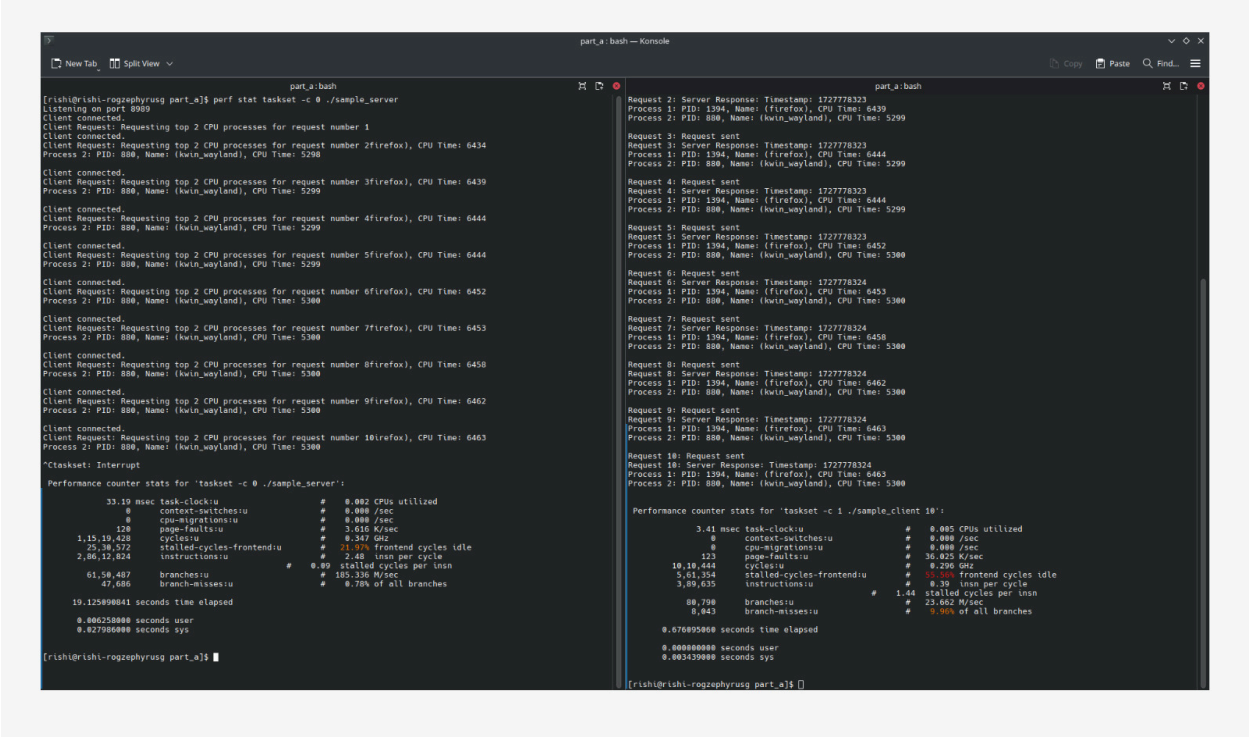
(a) Single-threaded TCP client-server



Server Analysis

Task-clock: 33.19 msec - longer processing time compared to client

Page-faults: 120 page faults - low in number - probably efficient memory access.

Cycles: 1,15,12,428 cycles - significant work on CPU as the server has to process multiple client requests, manage I/O etc.

Stalled-cycles-frontend: 21.93% - due to instruction cache misses or complex control flow - some inefficiency in fetching/decoding instructions.

Instructions per cycle (IPC): 2.48 - good efficiency of CPU

Branch misses: 0.78% is quite low - efficient branch prediction

Client Analysis

Task-clock: 3.41 msec- less time on CPU.

Page-faults: 123 indicates efficient memory handling.

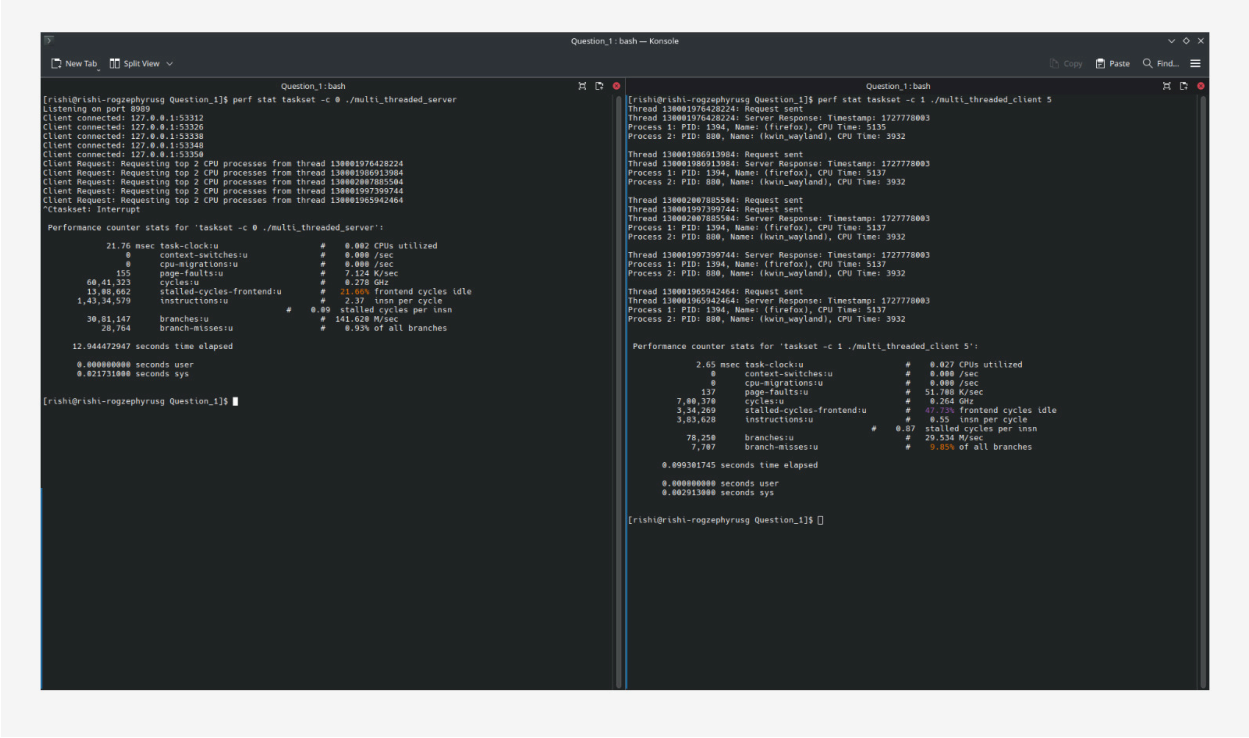
Cycles: 10,10,444 - client mainly sends and receives data.

Stalled-cycles-frontend: 55.56% client waits for instruction fetching

Instructions per cycle (IPC): 0.39 - quite low as client only sends and receives data

Branch misses: 9.96% is higher than the server' - may not be very efficient

(b) Concurrent TCP client-server



Multi-threaded Server Analysis

Task-clock: 21.76 msec indicates that the server was active for a significant duration, suggesting substantial workload handling multiple threads.
Page-faults: 155 page faults, which is still relatively low given that this is a multi-threaded server.
Cycles: 60,41,323 cycles show a higher workload compared to previous single-threaded observations, as expected in a multi-threaded setup.
Stalled-cycles-frontend: 21.66% suggests that the server efficiently fetches instructions, with fewer stalls compared to the single-threaded setup.
Instructions per cycle (IPC): 2.37 is quite efficient, indicating that the server can effectively handle multiple threads concurrently without becoming a bottleneck.
Branch misses: 0.93% is low, demonstrating efficient handling of control flow even in a multi-threaded environment.

Multi-threaded Client

Task-clock: 2.65 msec - relatively quick completion time
Page-faults: 137 - slightly higher than the single-threaded client but still within acceptable limits.
Cycles: 7,00,370 - lesser workload compared to the server; expected.
Stalled-cycles-frontend: 47.73% suggests that the client waits significantly
Instructions per cycle (IPC): 0.55 - less than server- client is not utilizing CPU efficiently.
Branch misses: 9.85% indicates that the client's branching logic is less efficient compared to the server

Comparisons Between Multithreaded and Single-threaded

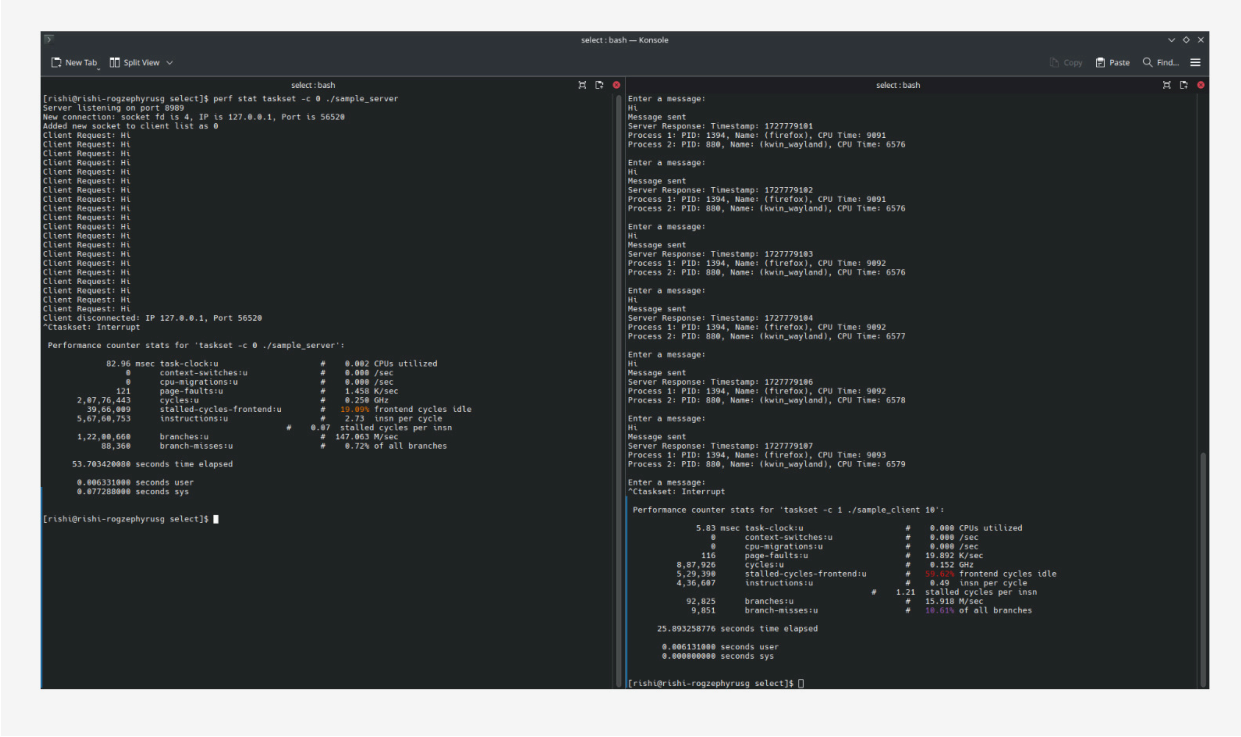
1. Server-side: The multi-threaded server handles a more significant workload efficiently, with a high IPC and low branch misprediction rate. This suggests good scalability in handling multiple threads and concurrent client requests.

  Instruction fetching is improved compared to the single-threaded setup, resulting in fewer stalled cycles.

2. Client-sider:

  The multi-threaded client doesn't perform as efficiently as the server, with a lower IPC and a high percentage of stalled cycles at the frontend, suggesting potential inefficiencies in its threading model or execution path.

(c) TCP client-server using "select"



Client

Task-clock: 5.83 msec - less time spent on CPU; expected.

Page-faults: 116 is acceptable and is not of major concern.

Cycles: 8,87,926 cycles indicate that the client node has low CPU utilization.

Stalled-cycles-frontend: 59.62 % i.e ~half of the cycles are waiting, suggesting a potential inefficiency in instruction fetching.

Instructions per cycle (IPC): 0.49 indicates suboptimal CPU usage. Ideally, this should be closer to 1 for better performance.

Branch misses: 10.61% indicates inefficiencies in branch prediction

Server

Task-clock: 82.96 msec is much longer than the client - more time spent handling requests.

Cycles: 20776443 shows heavier processing.

Stalled-cycles-frontend: 19.09% is lower than the client - more efficient instruction fetching but still a significant waiting period.

Instructions per cycle (IPC): 2.73 is quite efficient use of CPU

Branch misses: 0.72% is reasonably efficient

Comparison Between Threaded and Select Server-Client Models

Server:

1. Instruction Fetching & IPC:

The multi-threaded server exhibits a lower percentage of stalled-cycles-frontend & high IPC indicating efficient instruction fetching and concurrent request handling. The select-based server has even higher IPC, suggesting it is better optimized for handling multiple requests sequentially rather than creating new threads, resulting in fewer overheads.

2. Branch Misses:

The multi-threaded server has a low branch miss rate. The select server shows an even lower branch miss rate implying better optimization for linear, non-threaded processing.

Client:

1. CPU Utilization & IPC

The multi-threaded client shows moderate CPU utilization but has a lower IPC indicating suboptimal parallel execution. The high percentage of stalled cycles suggests instruction fetching. The select-based client has an even lower IPC and high stalled-cycles-frontend indicating that it spends more time waiting for instructions. However, the select model is simpler and may avoid overheads associated with threads.

2. Branch Misses

The threaded client struggles with efficient branch prediction. The select-based client performs worse in terms of branch prediction, maybe dynamic control flow leads to more mispredictions.