Keywood	Place of declaration	Litetime	Scope	Initial value	Place of Storage	Linkage
auto	Inside function/	Function/ block	Function/ block	Glasbage	Memosy (Stack)	None Vasiable
Register	m.		"	- " - 5	Regiotess	None
Static (local)		Program	30 11	Zeso	Memosy (data asea)	
Static (global)	Outside or inside fun (declaration of external variable)	Program/ Function	Declaration to end of file		Memosy (olata asea)	External
Extern	Outside Junction	Program	Definition end of	nto Zexo	Memory (data axea)	Internal

Storage Class Keyword:

· While declasing the vasiable we can mention stosage class

keyword.

· By obsessing the storage class compiler decides four important properties to a variable those are

1. Default value

2. Memory allocation 3. Scope

4. Life .

These are jour storage class keyword available in C, those are

1 · Auto

2. Static

3. Extern

Register

· In these four keyword all of them we can use while declasing the variable whereas static and extern we can use for function also.

We can declase a vasiable in our program in two places :

1) Above the main which is called as global variable.

2) Within the function which is called as local variable.

It we declare a local variable without mentioning any storage class type, compiler consider them as auto variable.

· Every auto variable is a local variable but every local may not be auto.

# include < stdio. h > auto int m ; Exxox i.e, we can't use auto variable as global vasiable Void main ()

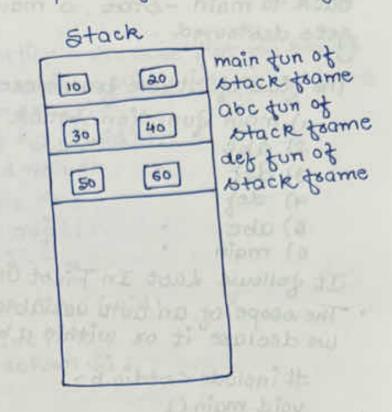
q inti; - auto auto int T - auto static int K; - Static

stack frame auto vasiable formal argument xeturn argument

- Default value of auto variable is a garbage or unknown value

Auto variable are stored in respected function stack frame

#include < stdio.h> void abc (void); void det (void); Void main () int i= 10, T= 20; abc(); void abc (void) 1 int K=30, L=40; def (); void det (void) int m=50, n=60;



Whenever our program is loaded into the RAM for execution operating system reserves 8 mb of stack for that particular

The stack is divided into small frames depends on the

function calls in our program

· In the above example - stat calls the main then main function stack frame gets created. In that stack frame if J variable memory gets allocated.

. Main is calling abo function, then abo function stack

grame gets created. In that stack frame Kfl

vasiable memosy gets exected.

· abc is calling det, then det function stack frame gets created. In that stack frame min variable memory gets exected.

· Once det junction execution completes, when it goes back to abo the det fun stack frame gets destroyed.

· After completion of abc when my function program execution goes back to main, abo function stack frame gets destroyed.

· After completion of main when my program execution goes back to main - stat, a main function stack frame gets destroyed. The flow of stack grame exection and destroy as follows: 1) main function stack frame created a) abc 3) det debtsoyed 4) det 5) abc 6) main It follows Last In Fixot Out. The scope of an auto vasiable is within a function where we declare it or within a block where we declare it # include < stdio. h> void main () 9 int i= 10; print; ("1) i= 1.dln",i); 1 int 1 = 20; point ("2) = 1.2 (n", i); print ("3) = /& (n", i); · Life of an auto vasiable stasts when the function stack frame is exected and life end when the function stack frame destroyed. · While designing the function we should not setus the address of auto variable which we declare inside that function. int \*abc (void); main () d int \*p; P = abc() i I int \*abc (void) 1 int 1 = 10; z setuxn 41;

Dangling Pointex: - It a pointex is holding life ended variable address.

Note: While designing any function make sure that you should not return the address of the variable whose life going to end after completion of that function execution.

```
uoid abc (int*);
void main ()
{ int i = 10;
abc (&i);
}
void abc (int *p)
{ *p = 20;
}
```

```
int *abc (void);

void main ()

fint *p;

P = abc ();

----

int *abc (void) X

fint i = 10;

xeturn fi;

}
```

Static Storage Class :-

- . Static variable can be declared local as well global but auto should be local.
- . The default value of static variable is 0.

· Static variables are stored in data segment.

Uninitialized static variables are stored in Bos (Block started by symbols) or (uninitialized data segment), whereas initialized static variables are stored in initialized data section.

```
# include < stdio.h >
static int i=10;
void main ()
{ print ("i=1.d ln", i);
}
L> 10
cc -c sl. c
size sl. 0
```

text data BBB dec hex filename
94 4 0 98 62 BL.0

#include < btdio.h >

ptatic int i=10;

void main ()

Estatic int i=20;

Printt("i=1/d In", i);

Ly 20

Note: We can declare a global variable as well as a local variable with same name.

In such case local is the higher priority.

There is no directly method to excess the global variables, it local is there with the same name but in C++, there is an operator called scope resolution operator (::) with that we can excess the global even local is these with the same name. Dope :- It we declare a static variable within a function the scope is inside that function only. If we declare a static variable outside the function (globally)

6

- the scope is within that file. Life: - Life of a static variable begins when program is loaded into the RAM for execution and life ends when the program execution ends.
- For static variable re-initialization will not happen #include < stdio.h> void abc (void); void main () tabe(): abc () ; abc(); Juoid abc (void)

d static int i=10; print; (" i= /d In", i);

9. How to access the global variable if the local variable is there with the same name .

# include < stdio. h > static int i = 10; int \*abc (void); void main () 2 static int i = 20; int \*p; printt (" i= /d In", i);

\* Multiple File compilation :-

m1.c ma.c

m1.s ma.c

m1.b m2.b

m1.o m2.o

Linkex Lib

· By default function are non-static (extern). Do, in one file if we declare one function, that function can be called from another file also.

Q. What is a static function?

=> When we defining the function if we put static keyword before the definition that function becomes static function, so that function is visible only inside that file we cannot call that function Isom another file.

Translator decides the scope (visibility) whereas linker decides who cannot access that.

Note: Without declasing the function if we call the function translator generates warning.

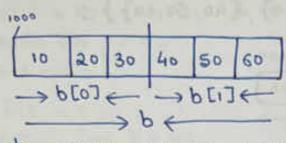
Without declasing the variables if we call that variable translator generates error.

	Extern inti	int i	int i= 10
Translator	Declaration	Definition	Initialization
		D+ M (Memor	A+m+A
Linker	× 16 to 2 to	allocate	a) a page a a
	1-60	Weak vasi	PHUAJCX

```
* auto & static vasiable axe
          int i = 10;
Linkage
                                  almost same
          Static int J = 20;
Internal-
                                  Local vasiable does not have
Linkage
          void main ()
                                  any linkage
          fint K = 30;
External
                                 · Register variable are stored
Linkage
           Static int & = 40;
                                   in CPU register and it this
16 not -
these
                                   otasop if youd as sstarps
           void abc(upid)
(No
 Linkage)
                                   Stack
                                 · We cannot take input from
                                  keyboard in register variable
                                   Register variable are used
                                   for fast processing
   2-D Axxay
 . 2D Axxay is a collection of I'D's axxay
  20 Assay is a assay of assays
   20 Array is also called as matrix
9. How to declare a 2D Array
                                             In each 1-D assay
                              Mo. of L-D
                                          number of elements
     datatype variable name [rows] [coloumn]
     # include < stdio.h >
     void main ()
     t int x, c;
      int b[2][3] = { [10,20,30], {40,50,60} };
      print; (" / & In ", sizeof (b));
                           (6[0]); 10 20
                           (b[o][o]);
                                          ← b[0] → 1← b[1] →
                   L> 24
       8 = 513eof (6) / 513eof (6[0]);
       c = bizeof (b[o] ) / bizeof (b[o][o]);
        print; ("x= 1.d c= 1.d ln", x,c);
               8 = 2
               C = 3
```

```
int b[2][3] = { {10,20,30}, {40,50,60}};
     1000 1004 1008 1 1012 1016 1020
       10 20 30 40 50 60
      ← [1]d → (- [0]d →
           →b<-
Q. What is the type of b.
=> Integer 2D array
Q. What is the type size of b.
=> 24 bytes
& How many elements are present in b.
=> 2 elements i.e., b[o], b[i]
9. What is the type of boo .
=> Integes L-D assay
Q. What is the size of b[o].
=> 12 bytes
9. How many elements are present
   in blo]
3 elements i.e., b[o][o], b[o][i], b[o][a]
8. What is the type of a blodlod.
=> Integer type
8. What is the size of b [0] [0].
=> 4 bytes
      fox ( i=0; i < x; i++)
        $08 (J=0; J<c; J++)
         print; (" /d", b[:][]);
        pxint ("In");
    int b[2][3] = {{10,20}, {40,50}};
        11 partial initialization
             L> 10 20 0
                40 50 0
```

-5



$$b+1 \rightarrow 1012$$
  
 $b[0]+1 \rightarrow 1004$   
 $b[0][0] \rightarrow 10+1$   
 $+1 = 11$   
 $4b+1 = 1024$   
 $4b[0]+1 = 1012$   
 $4b[0][0]+1 = 1004$ 

\* Difference between typedet and macro

## Macso

Typedet

- i) #define is a preprocessor directive 1) Uses defined datatype
- 2) # define is useful for communication 2) Typedef is useful for the with processor communication with translator
- 3) It is used for data types, symbol 3) Only for the data types. replacement with another symbol.
- 4) Replacement takes place
- 5) #define IPTR int \* IPTR P. 9; int \* P, 9;

Syntax: # define macroname microbody

- 4) No replacement.
- 5) typedet int \* INTPR P, 9; INTPR P.Q;

Syntax: Typedef oldname newname

\* Typedet: - It is one of the uses defined datatype which is used to provide another name to existing data type. With the help of this we can make understanding of variable or a data type .

Data	Common Data Structure Oper				workt			Space Complexity		
Structure	Access	Seasch :	Incestion	deletion .	Access	beasch	Insestion	Proposition and the same of th	Works	
-Assay	0(1)	0(n)	0(n)	0(n)	0(1)	o(n)	0(n)	0(n)	0(1	
Stack	O(n)	o(n)	0(1)	0(1)	o(n)	0(n)	0(1)	0(1)	0(1	
	o(n)	0(n)	0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(1	
Queue	o(n)	o(n)	0(1)	0(1)	o(n)	0(n)	0(1)	0(1)		
SLL			0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(1	
DLL	O(n)	O(n)		100 000	o(n)	0(n)	0(n)	0(n)	o (n log	
Skip List	O(log(n))	0(209(11)	((n)gos)0	21200		2.	"		O(n)	
Выт	"		-	E E		ar.	-			
asterian tree	NIA	H	#	*	N/A	Sec.	W =	248		
tash Table	. 14	0(1)	0(1)	0(1)	SHOTS	15 6		BIEL		
-Algorithm   Bent   Average		Lossot	woxat			Mostet				
Suickbost	O(n log(	277	log(n))	o(n'2)				0 (2091	(n))	
Nesperost				o(nlog(n))		O(n)				
Heapwoot		A STANSANTE		"		0(1)				
Bubble Soxt	o(n) o(n^2)		o(n^a)		0(1)					
neestion Sos		12.0	10	- 11						
relection box	The second second	1	W.	10						