Computing Gauss elimination on a 2-digit floating point computer.

$$\begin{bmatrix} 4 & 1 & 2 \\ 2 & 4 & -1 \\ 1 & 1 & -3 \end{bmatrix} \begin{bmatrix} \chi_1 \\ \gamma_2 \\ \gamma_3 \end{bmatrix} = \begin{bmatrix} 9 \\ -5 \\ -9 \end{bmatrix}$$

In augmented form

$$0.40\times10^{1}$$
 0.10×10^{1}
 0.20×10^{1}
 0.40×10^{1}
 0.40×10^{1}

Ri - ith row

$$R_2 \rightarrow R_2 - \frac{R_1 * \alpha_{21}}{4} \qquad R_3 \rightarrow R_3 - \frac{R_1 * \alpha_{31}}{4}$$

$$0.10 \times 10^{10}$$
 0.25×10^{10} 0.5×10^{10} 0.22×10^{10} 0.35×10^{10} -0.20×10^{10} -0.94×10^{10} 0.35×10^{10} -0.35×10^{10} -0.11×10^{10}

$$\begin{bmatrix}
1.0 & 0.25 \times 10^{\circ} & 0.5 \times 10^{\circ} & 0.22 \times 10^{\circ} \\
0 & 1.0 & -0.57 \times 10^{\circ} & -0.27 \times 10^{\circ} \\
0 & -0.30 \times 10^{\circ} & -0.9 \times 10^{\circ}
\end{bmatrix}$$

now back substitution
$$x_3 = \frac{-0.9}{-0.3} = 3$$

$$\chi_{2} + (-0.57)(\chi_{3}) = -2.7$$

$$\chi_{2} = -2.7 + 0.57 \times 3$$

$$= -2.7 + 0.71 = -2.7 + 2.7$$

$$= -2$$

$$= -2$$

$$\chi_{2} = 2.2 + (0.35 \times 1) - (0.5 \times 3)$$

$$= 2.2 + (0.35) - (1.5)$$

$$= 0.22 \times 10 + 0.02 \times 10^{2} - 0.15 \times 10^{2} = 0.9$$

So
$$\begin{bmatrix} \chi_1 \\ \chi_2 \end{bmatrix} = \begin{bmatrix} 0.9 \\ -1 \\ 3 \end{bmatrix}$$

Question 3

Computational complexity of of gours olimination of a tridiagonal matrix.

$$\begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 \\ 0 & a_{32} & a_{34} \end{bmatrix} = \begin{bmatrix} a_{11} \\ a_{21} \\ b_{22} \\ a_{33} \end{bmatrix}$$

Suppose the algorithm is at ith step then, operations are:

4 divisions at ith row to make an =1

then 4 multiplication 2 4 substraction to make

ain, it o. (Including operation on b)

No further subtraction are required because of triadiogonal nature of the matrix.

Total operations out ith row = 12 except at 1 row, At first row there will be 3 multiplicet 1

3 division 23 subtract "

Total operation =
$$9 + 12(n-2)$$

= $12n-3$

last now requires 1 division.

2 other rows require 1 multiplices, I subtrate a of division not required be cause we have 15 in diagonal exupt at first last row).

Total operation in back substitut? 1+ n(2)

Total operations to obtain the solut? of the system is

Total Operad" = 12m-3 + 2m + 1 = 14m-2 = 0(n)

Hence, solut of Ax=b, with A being tridlagonal, can be obtain in O(n) steps.

Question 5

The criteria for selecting library

1) I will check the timing measurements of the library to see which one is optimised.

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If they are using different algorithms then the package with less complexity will be selected.

- 2) How storage is managed in the also library and the storage requirements.
- 3) Cloually libraries are free but if not, cost will also play a role. Also the library which is popularly used and has produced reliable results for others

will be preferred. All of this is given considered given that the library can solve the physica problem I am interested in solving.

$$\frac{dy_1}{dx} = 32y_1 + 66y_2 + \frac{2x}{3} + \frac{2}{3} = f$$

$$\frac{dy_2}{dx} = -66y_1 - 123y_2 - \frac{1}{3}x - \frac{1}{3} = g$$

$$\left|\frac{\partial f}{\partial y_1}\right| = \frac{32}{\left|\frac{\partial f}{\partial y_2}\right|} = \left|\frac{-66}{=66}\right| = \frac{66}{66}$$

$$\left|\frac{\partial f}{\partial y_2}\right| = \frac{1-66}{66} = \frac{66}{66}$$

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- =) the propoderivates are finite and bounded : they satisfy Lipschitz condition
 - =) This problem is well posed and solution by numerical muthod is correct.

Also, solution seem to overlap with analytic solution.

Question 4

(e)
Sample is a function randomly generated between 0 & 1 whose average is \(\frac{1}{2} \) so DfT is a delta function and here power spectrum is a desired delta function at K=0.

(C)
$$k_2 = 0, \frac{2}{n\Delta}$$
 with $q = 1 \dots \frac{n-1}{2}$
 $k_3 = 0, \frac{2}{n\Delta}$ with $q = -\frac{n}{2}, \dots 1$

$$K_{may} = 0.499$$
 $K_{min} = -0.5$