

* Network layer → IP address

✓ Design issues

✓ IP addressing

✓ Sub-netting | Super-netting

* Network layer in Internet

✓ The IP protocol

✗ IPv4 header

✗ IPv6 header

✗ Routing algorithms

✗ Congestion algo

✓ Quality of service

— X — X — X —

Network layer → implements logical addressing

⇒ Third layer in the communication model.

⇒ Translates the logical address into physical address.

⇒ Determines the route from source to destination and manages traffic

⇒ move packets from sender to receiver.

⇒ Provides a logical connection betn different types of network.

*

Design issues

1) store and forward of packet switching :

→ So, host sends the packet to nearest router, now the

packet is stored in router till the entire packet arrives,

→ Once the entire packet has arrived, the checksum is done

and then the packet is sent to next router till

destination

Packet



router

(stored)

till complete

packet arrives

destination

2) Services Provided to transport layer:

* Services shld not be dependent on router.

• Network layer provides some services to transport layer;

Packet-independent entity.

→ connection less ⇒ routing and insertion of packets into subnets is done individually no extra setup.

→ connection oriented ⇒ all packets must be transmitted through single route or path.

3) Implementation Connection oriented:

⇒ Here a virtual circuit is setup in betⁿ host and sender, thus the path is already decided.

⇒ Each packet follows a certain virtual circuit.

4) connectionless :

⇒ here the packet is transmitted as an independent entity

⇒ each packet has its routing info. termed as datagram.

⇒ Each datagram follows his own route from source to destination.

e.g.: IP.

IP addressing (0.0.0.0 → 255.255.255.255) (32 bits)

⇒ Internet protocol address. → unique.

• helps internet to distinguish betⁿ diff. routers and components.

⇒ structure ⇒ set of 4 digits from 0.0.0.0 → 255.255.255.255

∴ x_1, x_2, x_3 = network ID

x_4 = host ID.

addressing
physical (logical)
(mac)
IP
VIA
public private

mac address \Rightarrow we cannot change
IP address \Rightarrow may change on restart.

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- * Network ID, identifies the location where device is located.
- * Host ID \Rightarrow identifies the specific device.
- * IPv4 \Rightarrow 32 bits, ~~42~~ hundred crores of IP, facing shortage
- * IPv6 \Rightarrow 128 bits, hexadecimal IP address, 3.4×10^{38} IP

IP address types: Public, private, Fixed & Dynamic

Public \Rightarrow web server, email etc which are accessed by the world will have a public IP.

Private \Rightarrow Router generates a unique IP for a device ⁱⁿ that network.

static \Rightarrow (fixed) \Rightarrow invalid IP, a dynamic IP will be provided at that time.

dynamic \Rightarrow changes over time. ISP provides Dynamic IP instead of static.

Classfull and Classless:

* classless \Rightarrow in classfull size of addressing is fixed and each address range has a default subnet mask.

\Rightarrow using classless, we can allocate IP more efficiently.

\Rightarrow used as a short soln to the demanding IP.

\Rightarrow Prefix data & Postfix data is attached to the IP.

\Rightarrow classless addressing is a specific instance of classfull.

* Classfull addressing:

- IP is divided into classes and each class has a specific subnet mask.

classes Range.

A 0.0.0.0 127.255.255.255

B 128.0.0.0 191.255.255.255

C 192.0.0.0 223.255.255.255

multicast D 224.0.0.0 239.255.255.255

reserved for future E 240.0.0.0 255.255.255.255

(2^{31}) • class A address can be used globally ($0000 \rightarrow 127.255.255.255$)

(1 reserved) Net ID = 8 bits Host ID = 24 bits.

Used in very large networks.

(2 reserved) • class B (2^{30})

Net ID = 16 bits Host ID = 16 bits

• medium size

(2^{29}) class C (2^{29})

(2 reserved) • Net ID = 24 host ID = 8

(2^{28}) • class D → multicast → 1st 4 bits reserved (1110)

(2^{28}) • class E → reserved → 1st 4 bits reserved (1111)

default
subnet mask

A \Rightarrow 255.0.0.0

B \Rightarrow 255.255.0.0

C \Rightarrow 255.255.255.0

D \Rightarrow 255

classfull vs classes

- classes → IP address are divided into 5 groups in classfull, in order to prevent depletion of IP address ⇒ classes is used.

- ⇒ classless is more beneficial than classfull.
 - ⇒ classfull netid & hostid are properly allotted.
but in classless, distinction of these does not exists.

- classless divides blocks into varying length

- prefix suffix is given

- Block of address must have Power of 2 addresses.

* slash notation for prefix length eg: 12.23.45.67 | 8.

167.199.170.82, 127

$$\begin{array}{r} 10101010 \\ 61010010 \\ \hline 00000 \end{array}$$

• keep 1st 2² bits same & change
rest to (1).

Subnetting Supernetting

* Cubnetting \Rightarrow

Subnet \Rightarrow logical division or partition of IP addressess network

• Segmented piece of larger network into small networks

⇒ The practice of dividing larger network in smaller = subnetting
small logical networks.

- * Computers belonging to a single subnet, addressed with identical most significant bit group.

$$\text{eg: } 192 - 168 = 14 \cdot 0 \quad \text{to} \quad 192 - 168 = 14 \cdot 255$$

(12 ... 244) = Subnets.

Subnet mask = 32 bit address used to distinguish.

bet'n netid & host id.

0000 = network address | 255255... = broadcast reserved

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- * ① identify class & default subnet mask.
- ② convert default subnet mask into binary.
- ③ note the no of host required per subnet ~~mask~~.
- ④ generate new subnet mask.
- ⑤ make network ranges.

* 255.255.0.0

These say that,

group of IP

should have

same value for
1st & 2nd pos!

↳ these say that the IP can have any

value, group of IP can have

any value at 3rd & 4th pos.

eg: 11.11.14.12. & 11.11.13.11

both belong to same group

* 216.21.5.0 in-to 30 host in each subnet.

Net 216.21 host id.

216.21.5.0 \Rightarrow class C.

∴ default subnet mask of C.

255.255.255.0.

$\Rightarrow 11111111.11111111.11111111.00000000$.

8 bits reserved

no. of host = 30 (11110) \rightarrow 5 bits

now reserve 5 bits & fill remaining with 1.

$\Rightarrow 11111111.11111111.11111111.1111000000$ & 8 reserved

subnet generator = 1st 1 from right

$$= 2^5 = 32$$

in octet 4

= 255.255.255.224

⇒ 216.21.5.0 + 32

⇒ 216.21.5.32. Keep adding 32 for 30 times.

216.21.5.0 - 216.21.5.31.

supernetting = Opposite of subnetting,

- ⇒ multiple networks are combined into 1 single network
 - ⇒ supernetting / supernet
- ⇒ used in route summarization.
- Where routes to multiple networks with same network prefixes are sent through some single route.
- ⇒ reduces traffic and size of routers.
- ⇒ aggregating small networks into one big.

Internet Protocol

- * main task of Ip is to deliver packets from source to destination, based on Ip address available in the packets.
- * provides a connectionless service accompanied by either TCP/IP or UDP/IP.
- * main function is also to provide addressing to hosts, and encapsulating data and routing data from source to destination.
- * Ip has its Packet ⇒ Ip header → data.
- * And Ip addressing ⇒ private & public.

- * IP address unique identifier assigned to comp. which is connected to Internet

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* Public address \Rightarrow external address, they are grouped under WAN. addresses used to access the internet, available on our comp

* with public address we can setup home server to access internet.

\Rightarrow scope is global.

* Private addresses \Rightarrow internal address, grouped under LAN

* used in communicating within the network.

* not routed on internet, thus no internet traffic.

* private address are those comp, printer that are placed in house.

\Rightarrow Scope is local, only, network communication.

* IP routing: process of determining the path of data.

* the path that each packet follows is done by routing algorithms.

* routing algo. considers various, constraints, size, header for efficient route

* router forwards the packets & it works on network layer.

* Constraints & measures in routing:

(1) Hop count \Rightarrow path with least hops, least routers is to be selected

(2) Delay \Rightarrow time taken by router to process data, hence path with less delay is used.

③ bandwidth \Rightarrow higher the capacity to transfer data measured in bits/sec. higher bandwidth more data transferred.

④ load ⑤ reliability

* static routing \Rightarrow give the route manually.

* default \Rightarrow used when router has to send data to single hop. (single exit point).

* dynamic \Rightarrow adaptive, discover new routes, can change routes.

* Routing Algorithms:

- To determine the route of transferring we use routing algo.
- We know, we have 2 points in IP. i.e. connections \oplus ^{WDP}_{IP} if connection oriented (\oplus TCP)
- Routing algo help in calculating least cost path.

* Types of Routing algo.

- ① Adaptive Routing algo.
- ② Non-adaptive routing algo.

* 1] Adaptive Routing algorithms:

- also are the dynamic routing algorithms, able to change the path and discover new path.
- Makes own decisions based on topology & traffic.
- Main parameters are hop count, distance, delay.

* They are of 3 types ① centralized

② isolated

(CIO)

③ distributive-decentralized.

(global
(local

Routing →

adaptive

non adaptive

-) • centralized
-) • isolated
- distributive.

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① centralized ⇒

- computes the least cost path by using the complete and global data about the network

- takes the data first and before performing the routing, and does the calc. of routing.

* Link state algorithm = centralized.

② isolation ⇒

- It takes in the information for routing from the local information or from nodes.

- ③ distributed ⇒ also called decentralized, computes the shortest path in iterative, repeated manner.
- here no node has knowledge about other nodes.
 - It has info about only its directly attached link. and from that it calculates least cost path.
 - never knows the complete path.

④ Nonadaptive algorithms:

- also known as static routing algorithms
- do not take routing decisions.

We have 2 types ① Flooding

② Random walks.

(node based)

a) Flooding \Rightarrow

- every incoming packet, sent to all other links connected or all other outgoing lines.
- several copies of some data may be generated.

b) Random walks \Rightarrow

- the received packet is sent to one of the neighbours randomly.
- It uses alternative routes very efficiently.

Distance vector algo (distributed) (Dynam)

- \hookrightarrow This is (1) distributed \Rightarrow each node receives info from one or more of directly attached nodes.
- (2) Iterative \Rightarrow process continues
- (3) Async \Rightarrow does not require all nodes.

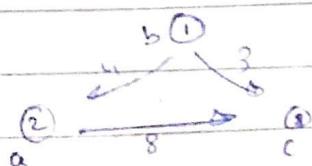
- * each router maintains a distance vector.
- * routing only to neighbours / connected to main.
- * cost for distance vector routing is based on hop count



- * Routing table \Rightarrow netid, cost, next hop.
Final destination \Rightarrow no. of hops \Rightarrow next router

\therefore all these data is stored in all routers, and that basis the distance vector is seen and data is transmitted to closer

- * after each transmission, we update data



$1 \rightarrow 2 = 4$	stable -
$2 \rightarrow 3 = 8$	
$1 \rightarrow 3 = 3$	
$\underline{\underline{8}}$	$1 \rightarrow 3 = 3$

$$3 \rightarrow 2 = 8.$$

Link State Routing: (Centralized)

Based on tech. in which each router shares knowledge about its neighbours to all other routers.

Flooding → * Instead of sending the routing table, router sends the information about its neighbours.

* Each router does send info to all other routers except neighbours. and now every router that receives the data, sends that to other routers.

* 2 states ~~3~~

1] Reliable Flooding

a) Initial stage ⇒ each node knows the cost of neighbour

b) Final stage ⇒ each node knows entire graph.

2] Route calc. ⇒

• uses Dijkstra to calc. node.

↳ iterative algo. known for finding least cost path.

* uses Dijkstra to find closest-path & share info.

* heavy traffic is created due to flooding.

Congestion \Rightarrow heavy traffic.

• delay in performance.

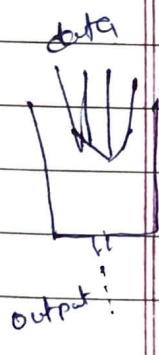
- Cong. control algo is a mechanism that controls entry of data packets in network
- Congestion avoidance algo is used in TCP.

Two types

- ① Leaky bucket
- ② Token bucket

① Leaky bucket

- allows to control, the rate at which data is given to network.
- Used in traffic shaping.
- Assume as a bucket
- Any input to the bucket is not affected but at a constant rate the packets are released.
- Basically bucket leaks some data at const. rate.
- means data input is higher than output.
- But this helps in making the data in a ordered sequence manner. \Rightarrow makes data regular flow.
- * leak rate is fix so bit inefficient, uses some buffer to control the data.
- * During large traffic, this causes issues.
- * in huge traffic into bigger than bucket may be lost



② Token bucket:

(bucket \Rightarrow buffer)

- When traffic is high, output is allowed to speed up. thus making it more efficient and flexible.
- Generally no data is lost.
- The output of bucket is based on the tokens data in the bucket
- Token bucket algo adds token to bucket at every 1/r sec
- If new token appears and previous bucket is full, then the new token is discarded.
- When not enough spaces available, then either the token is considered in non-conformant. they may be dropped
- * the space available in Token bucket only that no. of ~~as~~ packets will be allowed. \therefore no scope of overflow.

IPv4 vs IPv6

- IPv4
 - * current version - 32 bits.
 - * a set of 4, 8 bit octets.
 - * IPv4 is depleting, no more left, what to give

• IPv6

- next gen of IP. (hexadecimal + value),
- 128 bit address, has larger scopes & small heads

* transition strategy's IPV4 TO IPV6

- (1) Dual stacking \Rightarrow both version of V4 & V6 on same device
- (2) Tunneling \Rightarrow IPV6 communicates with V4.
- (3) Network addr. translation \Rightarrow allows communication with diff IP version

* V4 has 4 octets, V6 has 16 octets.

- numeric separated by dot
- class \rightarrow 5 classes are used
- supports virtual length subnet mask
- Manual & DHCP config
- end to end integrity is impossible.
- lower security
- alphanumeric separated by colon
- does not have any classes.
- Does not support VLSM
- manual, DHCP, auto config, renumbering.
- end to end integrity possible.
- Higher security.

- fragmentation is done by sender and routers
- uses checksum
- does not have encryption & decryption.
- fragmentation done by sender only
- has no checksum
- provides encryption & decryption.

Quality of service: set of tech. that work on network to give its best performance

- B • Bandwidth - high
 - D • delay - less
 - L • loss rate - less
 - J • error rate - med.
 - A • jitter - dk.
- } Based on router as we are checking the network layer data transmission.

L = irregular speed of packet resulting congestion.

- prioritize, queue, traffic marking.
- latency reduction, enhanced user exp.
- prevention