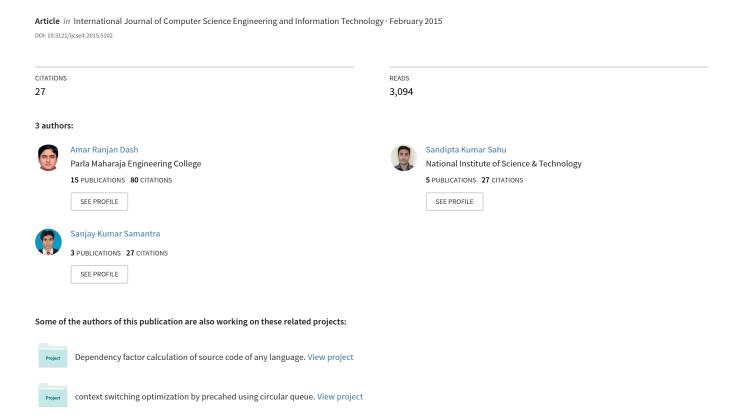
# An Optimized Round Robin CPU Scheduling Algorithm with Dynamic Time Quantum



# AN OPTIMIZED ROUND ROBIN CPU SCHEDULING ALGORITHM WITH DYNAMIC TIME QUANTUM

Amar Ranjan Dash<sup>1</sup>, Sandipta kumar Sahu<sup>2</sup> and Sanjay Kumar Samantra<sup>3</sup>

<sup>1</sup>Department of Computer Science, Berhampur University, Berhampur, India.

<sup>2</sup>Department of Computer Science, NIST, Berhampur, India.

<sup>3</sup>Department of Computer Science, NIST, Berhampur, India.

#### **ABSTRACT**

CPU scheduling is one of the most crucial operations performed by operating system. Different algorithms are available for CPU scheduling amongst them RR (Round Robin) is considered as optimal in time shared environment. The effectiveness of Round Robin completely depends on the choice of time quantum. In this paper a new CPU scheduling algorithm has been proposed, named as DABRR (Dynamic Average Burst Round Robin). That uses dynamic time quantum instead of static time quantum used in RR. The performance of the proposed algorithm is experimentally compared with traditional RR and some existing variants of RR. The results of our approach presented in this paper demonstrate improved performance in terms of average waiting time, average turnaround time, and context switching.

## **KEYWORDS**

CPU Scheduling, Round Robin, Response Time, Waiting Time, Turnaround Time

## 1. Introduction

Operating systems are resource managers. The resources managed by Operating systems are hardware, storage units, input devices, output devices and data. Operating systems perform many functions such as implementing user interface, sharing hardware among users, facilitating input/output, accounting for resource usage, organizing data, etc. Process scheduling is one of the functions performed by Operating systems. CPU scheduling is the task of selecting a process from the ready queue and allocating the CPU to it. Whenever CPU becomes idle, a waiting process from ready queue is selected and CPU is allocated to that. The performance of the scheduling algorithm mainly depends on CPU utilization, throughput, turnaround time, waiting time, response time, and context switch.

Different CPU scheduling algorithms described by Abraham Silberschatz et al. [1], viz. FCFS (First Come First Served), SJF (Shortest Job First), Priority and RR (Round Robin). Neetu Goel et al. [2] make a comparative analysis of CPU scheduling algorithms with the concept of schedulers. Jayashree S. Somani et al. [3] also make a similar analysis but with their characteristics and applications. In FCFS, the process that requests the CPU first is allocated the CPU first. In SJF, the CPU is allocated to the process with smallest burst time. When the CPU becomes available, it is assigned to the process that has the smallest next CPU burst. If the next CPU bursts of two processes are the same, FCFS scheduling is used to break the tie. In priority scheduling algorithm a priority is associated with each process, and the CPU is allocated to the process with the highest priority. Equal priority processes are scheduled in FCFS order. A major problem with priority scheduling is starvation. In this scheduling some low priority processes wait indefinitely to get the CPU.

DOI: 10.5121/ijcseit.2015.5102

In RR a small unit of time is used which is called Time Quantum or Time slice. The CPU scheduler goes around the Ready Queue allocating the CPU to each process for a time interval up to 1 time quantum. If a process's CPU burst exceeds 1 time quantum, that process is pre-empted and is put back in the ready queue .If a new process arrives then it is added to the tail of the circular queue. Out of the above discussed algorithms RR provides better performance as compared to the others in case of a time sharing operating system. The performance of a scheduling algorithm depends upon the scheduling criteria viz. Turnaround time, Waiting time, Response time, CPU utilization, and throughput.

Turnaround time is the time interval from the submission time of a process to the completion time of a process. Waiting time is the sum of periods spent waiting in the ready queue. The time from the submission of a process until the first response is called Response time. The CPU utilization is the percentage of time CPU remains busy. The number of processes completed per unit time is called Throughput. Context switch is the process of swap-out the pre-executed process from CPU and swap-in a new process to CPU. Context switch is the number of times the process switches to get execute. A scheduling algorithm can be optimized by minimizing response time, waiting time and turnaround time and by maximizing CPU utilization, throughput.

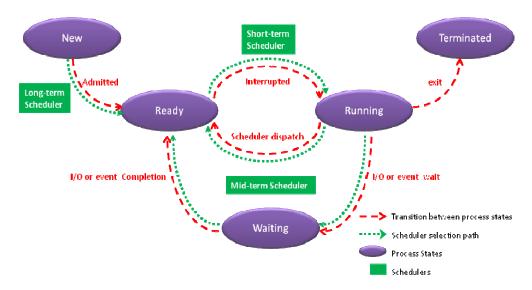


Figure 1: Different schedulers and Process states in CPU Scheduling.

The rest of the paper is organized as follows: In section 2, we describe the related works with a special emphasis on working procedure and dynamic time quantum selection procedure of different scheduling algorithms. Section 3 presents the proposed algorithm and its illustration. In section 4, we experimentally analyze the performance of seven scheduling algorithms, including our proposed algorithm, with six test cases. In section 5 we analyze the result obtained from our analysis. Section 6 provides the concluding remarks.

# 2. RELATED WORK

In recent time different approaches are used to increase the performance of CPU scheduling algorithm. Rakesh Kumar Yadav et al. [4] use the concept of SJF in RR algorithm. Ajit Singh et al. [5] combine the concept of SJF in RR algorithm. After each cycle they double the time quantum. Manish Kumar Mishra et al. [6] also merge the concept of Shortest Job First (SJF) with Round Robin (RR) to minimize the waiting time & turnaround time. After each complete cycle

they chose the burst time of shortest process as new time quantum. Rishi Verma [7] calculate the time quantum after every cycle by subtracting the minimum burst time from maximum burst time. Neetu Goel et al. [8] take two dynamic numbers K (as time quantum) and F. During execution if the remaining burst time of process in execution is less than time-quantum/F, then the process continues its execution otherwise the process stops its execution and goes to the end of ready queue. M. Ramakrishna et al. [9] add the concept of priority scheduling in RR scheduling to optimize the Round Robin scheduling.

Rami J. Matarneh [10] proposes an algorithm SARR to improve the performance of Round Robin. In SARR for each cycle the median of burst time of the processes is calculated and used as time quantum. H.S.Behera et al. [11] also use similar type of algorithm. But they again rearrange the process during their execution. It select the process with lowest burst time, then process with highest burst time, then process with second lowest burst time, and so on.

## 3. OUR PROPOSAL

# 3.1. DABRR Algorithm

```
TQ:
        Time Quantum
RQ:
        Ready Queue
        number of process
n:
        Process at ith index
Pi:
i, j:
        used as index of ready queue
TBT:
        Total Burst Time
[1]
        Arrange the processes in ascending order.
        n = number of processes in RQ
[2]
[3]
        i=0, TBT=0
[4]
        Repeat step 5 and 6 till i < n
[5]
          TBT += burst time of process Pi
[6]
          i++
        TQ = TBT/n
[7]
[8]
        j = 0
[9]
        Repeat from step 12 to 19 till j<n
[10]
          if (burst time of Pi) <= TQ
             Execute the process
[11]
[12]
             Take the process out of RQ
[13]
             n--
[14]
          Else
[15]
             Execute the process for a time interval up to 1 TQ
[16]
             Burst time of Pi = Burst time of Pi - TQ
[17]
             Add the process to ready queue for next round of execution
[18]
          j++
[19]
        If new process arrives
[20]
          goto step 1
[21]
        If RQ is not empty
[22]
          goto step 2
```

#### 3.2. Illustration:

In this section we analyzed the execution of the proposed algorithm. To demonstrate the proposed algorithm we have considered a ready queue with 5 processes p1, p2, p3, p4, p5. These processes are arrived at zero millisecond. The burst time of p1, p2, p3, p4, p5 are 15, 32, 102, 48, 29

milliseconds respectively. First the processes are arranged in ascending order of their burst time which provides the sequence p1, p5, p2, p4, p3. The time quantum is set equal to the mean of burst time of all 5 processes i.e. 45. After executing all processes for a time quantum of 45 millisecond execution of p1, p5, p2 get completed. So they are removed from the ready queue. After first cycle, the remaining burst time for p3 and p4 are 3 and 57 respectively. In the next cycle the new time quantum is set equal to the mean of the burst time of the processes in ready queue i.e. 30 and CPU is assigned to the processes for the new time quantum. After the second round the process p4 has finished execution and only process p5 remains in the ready queue with burst time 27 millisecond. As only one process is there in the ready queue so its burst time is directly chosen as time quantum and CPU is allocated to p5. According to our illustration the turnaround time for p1, p2, p3, p4, p5 are 15, 76, 226, 169, and 44 milliseconds respectively. The average turnaround time is 106. The waiting time for p1, p2, p3, p4, p5 are 0, 44, 124, 121, and 15 milliseconds respectively. The average waiting time is 60.8.

#### 4. EXPERIMENTAL ANALYSIS

First we divide the problems into two types based on arrival time of processes (processes with zero arrival time and processes without zero arrival time). We further divide each into 3 more types based on the burst time of processes (in ascending order, descending order, & random order). We analyzed all the algorithms based on six situations. In each we have considered five processes with their arrival time and burst time. We have taken 25 as static time quantum for Round Robin.

Ajit Singh et al. [5] have proposed an algorithm in which after each cycle the time quantum is doubled. They haven't given any specific name for that algorithm. So during our experimental analysis we have considered it as R.P-5.As that paper has been referred at fifth place in our reference.

#### 4.1. Assumptions

During analysis we have considered CPU bound processes only. In each test case 5 independent processes are analyzed in uni-processor environment. Corresponding burst time and arrival time of processes are known before execution. The context switch time of processes has been considered as zero. The time required for arranging the processes in ascending order also considered as zero.

#### 4.2. With ZERO arrival time

# **4.2.1.** Ascending Order (Case I)

Table 1: Processes with Zero arrival time and Burst time in increasing order.

Processes	Arrival time	Burst Time
P1	0	40
P2	0	55
P3	0	60
P4	0	90
P5	0	102

# RR:

<b>←</b>						2	25						$\rightarrow$
P1	P2	P3	P4	P5	P1	P2	P3	P4	P5	P2	P3	P4	P5
25	50	75	100	125	140	165	190	215	240	245	255	280	305

$$\begin{array}{c|cccc}
 & \leftarrow & 25 & \rightarrow \\
\hline
P4 & P5 & P5 \\
\hline
305 & 320 & 345 & 347
\end{array}$$

Average Turnaround time = (140+245+255+320+347)/5 = 1307/5 = 261.4

Average Waiting time

$$=((0+100)+(25+90+75)+(50+90+55)+(75+90+40+25)+(100+90+40+15+0))/5$$

$$= (100+190+195+230+245)/5 = 960/5 = 192$$

#### **DORRR:**

+		60		$\rightarrow$	← 3	6 →	← 6 →
P1	P5	P2	P4	P3	P5	P4	P5
40	100	155	215	275	311	341	347

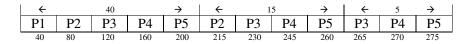
Average Turnaround time = (40+155+275+341+347)/5 = 1158/5 = 231.6

Average Waiting time

$$= (0 + 100 + 215 + (155+96) + (40+175+30))/5$$

$$= (0+100+215+251+245)/5 = 811/5 = 162.2$$

# **IRRVQ**:



$$\begin{array}{c|ccccc}
 & \leftarrow & 30 & \rightarrow & \leftarrow 12 \rightarrow \\
\hline
P4 & P5 & P5 & \\
\hline
275 & 305 & 335 & 347
\end{array}$$

Average Turnaround time = (40+215+265+305+347)/5 = 1172/5 = 234.4

Average Waiting time = (0 + (40+120) + (80+95+30) + (120+70+20+5) + (160+45+10+30))/5

$$= (0+160+205+215+245)/5 = 825/5 = 165$$

# **SARR:**

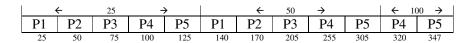
<b>←</b>		60		$\rightarrow$	← 3	6 →	←6→
P1	P2	P3	P4	P5	P4	P5	P5
40	95	155	215	275	305	341	347

Average Turnaround time = (40+95+155+305+347)/5 = 942/5 = 188.4

Average Waiting time = (0 + 40 + 95 + (155+60) + (215+30+0))/5

$$= (0+40+95+215+245)/5 = 595/5 = 119.0$$

## **R.P-5**:



Average Turnaround time = (140+170+205+320+347)/5 = 1182/5 = 236.4

Average Waiting time = ((0+100) + (25+90) + (50+95) + (75+105+50) + (100+130+15))/5

= (100+115+145+230+245)/5 = 835/5 = 167

# MRR:

<b>←</b>		62		$\rightarrow$	<b>←</b> 2	.5 →	← 2	25 →
P1	P2	P3	P4	P5	P4	P5	P4	25 → P5
40	95	155	217	279	304	329	332	347

Average Turnaround time = (40+95+155+332+347)/5 = 969/5 = 193.8Average Waiting time = (0+40+95+(155+62+25)+(217+25+3))/5= (0+40+95+242+245)/5 = 622/5 = 124.4

# **DABRR:**

Average Turnaround time = (40+95+155+314+347)/5 = 951/5 = 190.2Average Waiting time = (0+40+95+(155+69)+(224+21))/5= (0+40+95+224+245)/5 = 604/5 = 120.8

Table 2: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-I.

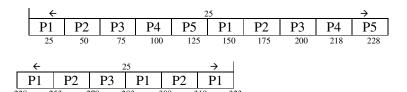
Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time	25	60,36,6	40,15,5,	60, 36, 6	25,50,	62, 25,	69,27,6
Quantum			30,12		100	25	
Context	16	7	14	7	11	8	7
Switch							
Average	192	162.2	165	119	167	124.4	120.8
Waiting time							
Average Turn	261.4	231.6	234.4	188.4	236.4	193.8	190.2
Around Time							

# 4.2.2. Descending Order (Case II)

Table 3: Processes with zero arrival time and Burst time in decreasing order.

Processes	Arrival time	Burst Time
P1	0	105
P2	0	85
P3	0	55
P4	0	43
P5	0	35

# RR:



Average Turnaround time = (323+318+283+218+228)/5 = 1370/5 = 274

Average Waiting time

$$=((0+100+78+30+10)+(25+100+78+30)+(50+100+78)+(75+100)+(100+93))/5$$

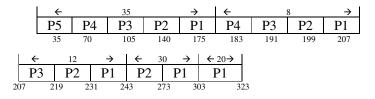
$$=(218+233+228+175+193)/5=1047/5=209.4$$

## **DQRRR**:

Average Turnaround time = (323+313+243+133+35)/5 = 1047/5 = 209.4

Average Waiting time = 
$$((35+153+30) + (133+95) + 188+90+0)/5$$
  
=  $(218+228+188+90+0)/5 = 724/5 = 144.8$ 

# **IRRVQ**:



Average Turnaround time = (323+273+219+183+35)/5 = 1033/5 = 206.6Average Waiting time

$$=((140+24+24+30+0)+(105+51+20+12)+(70+78+16)+(35+105)+0)/5$$

$$= (218+188+164+140+0)/5 = 710/5 = 142$$

#### **SARR:**

<b>←</b>		55		$\rightarrow$	<b>←</b> 4	0 >	←10→
P1	P2	P3	P4	P5	P1	P2	P1
55	110	165	208	243	283	313	323

Average Turnaround time = (323+313+165+208+243)/5 = 1252/5 = 250.4Average Waiting time = ((0+188+30) + (55+173) +110+165+208)/5= (218+228+110+165+208)/5 = 929/5 = 185.8

#### **R.P-5**:

<b>←</b>		25		$\rightarrow$		←	50	-3	<b>&gt;</b>	<b>←</b> 1	00 →
P1	P2	P3	P4	P5	P1	P2	P3	P4	P5	P1	P2
25	50	75	100	125	175	225	255	273	283	313	323

Average Turnaround time = (313+323+255+273+283)/5 = 1447/5 = 289.4Average Waiting time = ((0+100+108) + (25+125+88) + (50+150) + (75+155) + (100+148))/5 = (208+238+200+230+248)/5 = 1124/5 = 224.8

# MRR:

	,	<del>(</del>	70	-	<b>&gt;</b>	← 2	.5 →	← 25 →
	P5	P4	P3	P2	P1	P2	P1	P1
,	35	78	133	203	273	288	313	323

Average Turnaround time = (323+288+133+78+35)/5 = 857/5 = 171.4

Average Waiting time = 
$$((203+15+0) + (133+70) + 78 + 35 + 0))/5$$
  
=  $(218+203+78+35+0)/5 = 534/5 = 106.8$ 

# **DABRR:**

<b>←</b>		64		$\rightarrow$	<b>←</b> 3	1 →	←10→
P5	P4	P3	P2	P1	P2	P1	P1
35	78	133	197	261	282	313	323

Average Turnaround time = (323+282+133+78+35)/5 = 851/5 = 170.2Average Waiting time = ((197+21+0) + (133+64)+78+35+0)/5= (218+197+78+35+0)/5 = 528/5 = 105.6

Table 4: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-II.

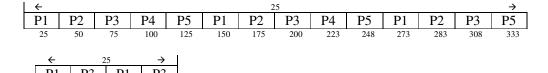
Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time	25	55,40,10	35,8,12,	55, 40, 10	25,50,	70, 25,	64,31,10
Quantum			30,20		100	25	
Context	15	7	14	7	11	7	7
Switch							
Average	209.4	144.8	142	185.8	224.8	106.8	105.6
Waiting time							
Average Turn	274	209.4	206.6	250.4	289.4	171.4	170.2
Around Time							

#### 4.2.3. Random Order (Case III)

Table 5: Processes with zero arrival time and Burst time in random order.

Processes	Arrival time	Burst Time
P1	0	105
P2	0	60
P3	0	120
P4	0	48
P5	0	75

# RR:



Average Turnaround time = (388+283+408+223+333)/5 = 1635/5 = 327Average Waiting time

=((0+100+98+60+25)+(25+100+98)+(50+100+83+50+5)+(75+100)+(100+98+60))/5

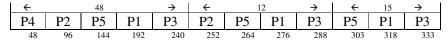
= (283+223+288+175+258)/5 = 1227/5 = 245.4

#### **DQRRR**:

1	<b>←</b>		75		$\rightarrow$	← 3	<b>←</b> 8 <b>→</b>	
Ī	P4	P3	P2	P1	P5	P3	P1	P3
	48	123	183	258	333	370	400	408

Average Turnaround time = (400+183+408+48+333)/5 = 1372/5 = 274.4Average Waiting time = ((183+112) + 123+ (48+210+30) + 0 + 258)/5 = (295+123+288+0+258)/5 = 964/5 = 192.8

# **IRRVQ:**



$$\begin{array}{c|cccc}
 & \leftarrow & 30 \rightarrow & \leftarrow 15 \rightarrow \\
\hline
P1 & P3 & P3 \\
\hline
333 & 363 & 393 & 408
\end{array}$$

Average Turnaround time = (363+252+408+48+303)/5 = 1374/5 = 274.8Average Waiting time

$$= ((144+72+27+15) + (48+144) + (192+36+30+30+0) + 0 + (96+108+24))/5$$

= (258+192+288+0+228)/5 = 966/5 = 193.2

# **SARR:**

Ĺ	<b>←</b>		120		$\rightarrow$
	P1	P2	P3	P4	P5
	105	165	285	333	408

Average Turnaround time = (105+165+285+333+408)/5 = 1296/5 = 259.2Average Waiting time = (0+105+165+285+333)/5 = 888/5 = 177.6

# **R.P-5**:

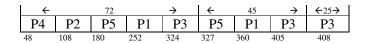
<b>←</b>		25		$\rightarrow$	<b>←</b>		50		$\rightarrow$	← 10	00 <b>→</b>
P1	P2	P3	P4	P5	P1	P2	P3	P4	P5	P1	P3
25	50	75	100	125	175	210	260	283	333	363	408

Average Turnaround time = (363+210+408+283+333)/5 = 1597/5 = 319.4Average Waiting time

$$= ((0+100+158) + (25+125) + (50+135+103) + (75+160) + (100+158))/5$$

= (258+150+288+235+258)/5 = 1189/5 = 237.8

#### MRR:



Average Turnaround time = (360+108+408+48+327)/5 = 1251/5 = 250.2Average Waiting time = ((180+75) + 48 + (252+36+0) + 0 + (108+144))/5= (255+48+288+0+252)/5 = 843/5 = 168.6

# **DABRR:**

<b>←</b>		81		$\rightarrow$	← 3	<b>←</b> 8 <b>→</b>	
P4	P2	P5	P1	P3	P1	P3	P3
48	108	183	264	345	369	400	408

Average Turnaround time = (369+108+408+48+183)/5 = 1116/5 = 223.2Average Waiting time = ((183+81) + 48 + (264+24+0) + 0 + 108)/5

= (264+48+288+0+108)/5 = 708/5 = 141.6

Table 6: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-III.

Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time Quantum	25	75,37,8	48,12,15,	120	25,50,	72, 45,	81,31,8
			30,15		100	25	
Context	17	7	14	4	11	8	7
Switch							
Average	245.4	192.8	193.2	177.6	237.8	168.6	141.6
Waiting time							
Average Turn	327	274.4	274.8	259.2	319.4	250.2	223.2
Around Time							

# 4.2.4. Comparison

This section provides the comparative analysis of seven algorithms on the basis of their resulted waiting time and turnaround time. Table 7 shows the performance analysis of seven algorithms by summarizing the waiting time and turnaround time resulted from case I, II, III.

Table 7: Comparison of six algorithms by aggregating their performance measure from case I, II, III.

	With 0 Arival Time											
	Co	ntext S	Switch		Waiting Time				Turnaround Time			
	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL
RR	16	15	17	48	192.00	209.40	245.40	646.80	261.40	274.00	327.00	862.40
DQRRR	7	7	7	21	162.20	144.80	192.80	499.80	231.60	209.40	274.40	715.40
IRRVQ	14	14	14	42	165.00	142.00	193.20	500.20	234.40	206.60	274.80	715.80
SARR	7	7	4	18	119.00	185.80	177.60	482.40	188.40	250.40	259.20	698.00
RP-5	11	11	11	33	167.00	224.80	237.80	629.60	236.40	289.40	319.40	845.20
MRR	8	7	8	23	124.40	106.80	168.60	399.80	193.80	171.40	250.20	615.40
DABRR	7	7	7	21	120.80	105.60	141.60	368.00	190.20	170.20	223.20	583.60

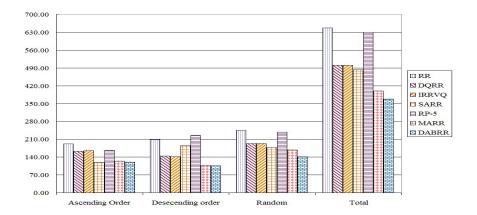


Figure 2: Waiting time for processes with 0 arrival time.

Figure 2 & 3 show the graphical analysis of the waiting time and turnaround time respectively of all algorithms based on case I, II, III. By evaluating these two graphs we conclude that our algorithm performs better than other six algorithms in case of processes arrived at zero arrival time.

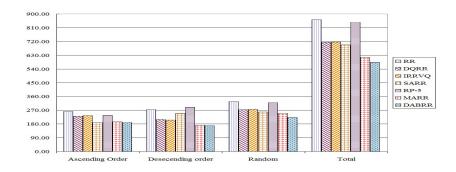


Figure 3: Turnaround time for processes with 0 arrival time.

# 4.3. Without ZERO arrival time

# 4.3.1. Ascending Order (Case IV)

Table 8: Processes without Zero arrival time and Burst time in increasing order.

Processes	Arrival time	Burst Time
P1	0	27
P2	3	32
P3	5	55
P4	7	82
P5	9	110

# RR:

1		<b>←</b>			2	5		$\rightarrow$			
	P1	P2	P3	P4	P5	P1	P2	P3	P4	P5	
	25	50	75	100	125	127	13/1	150	18/1	200	

Average Turnaround time = (127 + (134-3) + (214-5) + (271-7) + (306-9))/5 = 1028/5 = 205.6Average Waiting time

$$=((0+100)+(25+77-3)+(50+59+50-5)+(75+59+30+25-7)+(100+59+30+7+0-9))/5$$

$$=(100+99+154+182+187)/5 = 722/5 = 144.4$$

#### **DQRRR**:

← 27 <b>→</b>	<b>←</b>	6	8	$\rightarrow$	<b>←</b> 2	← 14 →	
P1	P2	P5	P3	P4	P5	P4	P5
27	59	127	182	250	278	292	306

Average Turnaround time = (27+(59-3)+(182-5)+(292-7)+(306-9))/5 = 842/5 = 168.4Average Waiting time

$$= (0 + (27-3) + (127-5) + (182+28-7) + (59+123+14-9))/5$$

$$= (0+24+122+203+187)/5 = 536/5 = 107.2$$

# **IRRVQ**:

<b>←</b> 27 <b>→</b>	<b>←</b>	3	2	$\rightarrow$	<b>←</b>	23	$\rightarrow$	<b>←</b> 2	7 <b>→</b>	←28→
P1	P2	P3	P4	P5	P3	P4	P5	P4	P5	P5
27	59	91 1	23	155 1	78 2	01 2	24 2	51 2	78	306

Average Turnaround time = (27 + (59-3) + (178-5) + (251-7) + (306-9))/5 = 797/5 = 159.4Average Waiting time

$$= (0 + (27-3) + (59+64-5) + (91+55+23-7) + (123+46+27+0-9))/5$$

$$= (0+24+118+162+187)/5 = 491/5 = 98.2$$

## **SARR:**

1	← 27 <b>→</b>	<b>←</b>	6	8	$\rightarrow$	<b>←</b> 2	8 >	← 14 →
	P1	P2	P3	P4	P5	P4	P5	P5
	27	59	114	182	250	264	292	306

Average Turnaround time = (27 + (59-3) + (114-5) + (264-7) + (306-9))/5 = 746/5 = 149.2Average Waiting time

$$= (0 + (27-3) + (59-5) + (114+68-7) + (182+14+0-9))/5$$

$$= (0+24+54+175+187)/5 = 440/5 = 88$$

# **R.P-5**:

	←25→	<b>←</b>		50		$\rightarrow$	+	100	$\rightarrow$
	P1	P1	P2	P3	P4	P5	P3	P4	P5
,	25	27	59	109	159	209	214	246	306

Average Turnaround time = (27 + (59-3) + (214-5) + (246-7) + (306-9))/5 = 828/5 = 165.6Average Waiting time

$$= (0 + (27-3) + (59+100-5) + (109+55-7) + (159+37-9))/5$$

$$= (0+24+154+157+187)/5 = 522/5 = 104.4$$

# MRR:

←27→	<b>←</b>	7	8	$\rightarrow$	← 2	28 →	←25→
P1	P2	P3	P4	P5	P4	P5	P5
27	59	114	192	270	274	302	306

Average Turnaround time = (27+(59-3)+(114-5)+(274-7)+(306-9))/5 = 756/5 = 151.2Average Waiting time

$$= (0 + (27-3) + (59-5) + (114+78-7) + (192+4+0-9))/5$$

$$= (0+24+54+185+187)/5 = 450/5 = 90$$

#### **DABRR:**

←27→		← 6	9 →		<b>←</b> 2	7 →	←14→
P1	P2	P3	P4	P5	P4	P5	P5
27	59	114	183	252	265	292	306

Average Turnaround time = (27 + (59-3) + (114-5) + (265-7) + (306-9))/5 = 747/5 = 149.4Average Waiting time

$$= (0+(27-3)+(59-5)+(114+69-7)+(183+13+0-9))/5$$

Table 9: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-IV.

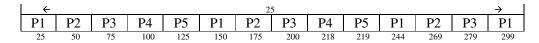
Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time Quantum	25	27,68,	27,32,23,	27, 68,	25,50,	27, 78,	27,69,27,
		28,14	27,28	28, 14	100	28, 25	14
Context	15	7	10	7	8	7	7
Switch							
Average	144.4	107.2	98.2	88	104.4	90	88.2
Waiting time							
Average Turn	205.6	168.4	159.4	149.2	165.6	151.2	149.4
Around Time							

# **4.3.2.** Descending Order (Case V)

Table 10: Processes without Zero arrival time and Burst time in decreasing order.

Processes	Arrival time	Burst Time
P1	0	95
P2	2	75
P3	4	60
P4	8	43
P5	16	26

# RR:



Average Turnaround time = (299 + (269-2) + (279-4) + (218-8) + (219-16))/5= 1254/5 = 250.8

Average Waiting time

<sup>= (0+24+54+176+187)/5 = 441/5 = 88.2</sup> 

$$= ((0+100+69+35) + (25+100+69-2) + (50+100+69-4) + (75+100-8) + (100+93-16))/5$$
  
=  $(204+192+215+167+177)/5 = 955/5 = 191$ 

#### **DORRR:**

Average Turnaround time = (95+(299-2)+(291-4)+(215-8)+(121-16))/5 = 991/5 = 198.2Average Waiting time

$$= (0 + (121+94+9-2) + (215+16-4) + (172-8) + (95-16))/5$$

$$= (0+222+227+164+79)/5 = 692/5 = 138.4$$

#### **IRRVQ:**

←95→	<b>←</b>	2	.6	$\rightarrow$	<b>←</b>	17	$\rightarrow$	← 1	.7 <b>→</b>	←15→
P1	P5	P4	P3	P2	P4	P3	P2	P3	P2	P2
95	121	147	173	199	216	233	250	267	284	299

Average Turnaround time = (95+(299-2)+(267-4)+(216-8)+(121-16))/5 = 968/5 = 193.6Average Waiting time

$$= (0 + (173+34+17+0-2) + (147+43+17-4) + (121+52-8) + (95-16))/5$$

$$= (0+222+203+165+79)/5 = 669/5 = 133.8$$

# **SARR:**

	←95→	<b>←</b>	51	1	$\rightarrow$	← 1	6 <b>→</b>	<b>←</b> 8 <b>→</b>
	P1	P2	P3	P4	P5	P2	P3	P2
,	95	146	197	240	266	282	291	299

Average Turnaround time = (95+(299-2)+(291-4)+(240-8)+(266-16))/5 = 1161/5 = 232.2Average Waiting time

$$= (0 + (95+120+9-2) + (146+85-4) + (197-8) + (240-16))/5$$

$$= (0+222+227+189+224)/5 = 862/5 = 172.4$$

# **R.P-5**:

←25→	<b>←</b>		50		$\rightarrow$	←	100	$\rightarrow$
P1	P1	P2	P3	P4	P5	P1	P2	P3
25	75	125	175	218	244	264	289	299

Average Turnaround time = (264 + (289-2) + (299-4) + (218-8) + (244-16))/5 = 1284/5 = 256.8Average Waiting time

$$= ((0+169) + (75+139-2) + (125+114-4) + (175-8) + (218-16))/5$$
  
=  $(169+212+235+167+202)/5 = 985/5 = 197$ 

# MRR:

I	←95→	←	4	.9	$\rightarrow$	<b>←</b> 2	.5 →	←25→
	P1	P5	P4	P3	P2	P3	P2	P2
	95	121	164	213	262	273	298	299

Average Turnaround time = (95+(299-2)+(273-4)+(164-8)+(121-16))/5 = 922/5 = 184.4Average Waiting time

$$= (0 + (213+11+0-2) + (164+49-4) + (121-8) + (95-16))/5$$

$$= (0+222+209+113+79)/5 = 623/5 = 124.6$$

# **DABRR:**

←95→		<b>←</b> 5	1 →		← 1	6 →	<b>←</b> 8 <b>→</b>
P1	P5	P4	P3	P2	P3	P2	P2
95	121	164	215	266	275	291	299

Average Turnaround time = (95+(299-2)+(275-4)+(164-8)+(121-16))/5 = 924/5 = 184.8Average Waiting time

= (0+(215+9+0-2)+(164+51-4)+(121-8)+(95-16))/5

= (0+222+211+113+79)/5 = 625/5 = 125

Table 11: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-V.

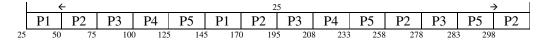
Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time Quantum	25	95,51,	95,26,17,	95,51,	25,50,	95, 49,	95,51,
		16,8	17,15	16,8	100	25, 25	16,8
Context	13	7	10	7	8	7	7
Switch							
Average	191	138.4	133.8	172.4	197	124.6	125
Waiting time							
Average Turn	250.8	198.2	193.6	232.2	256.8	184.4	184.8
Around Time							

# 4.3.3. Random Order (Case VI)

Table 12: Processes without Zero arrival time and Burst time in random order.

Processes	Arrival time	Burst Time
P1	0	45
P2	5	90
P3	8	70
P4	15	38
P5	20	55

# RR:



Average Turnaround time = (145+(298-5)+(278-8)+(208-15)+(283-20))/5= 1164/5 = 232.8

Average Waiting time

=((0+100)+(25+95+63+25-5)+(50+95+63-8)+(75+95-15)+(100+83+45-20))/5

= (100+203+200+155+208)/5 = 866/5 = 173.2

# **DQRRR**:

$$\begin{array}{|c|c|c|c|c|c|c|c|c|} \hline \leftarrow +45 \rightarrow & \leftarrow & 62 & \rightarrow & \leftarrow 18 \rightarrow & \leftarrow 10 \rightarrow \\ \hline P1 & P4 & P2 & P5 & P3 & P2 & P3 & P2 \\ \hline +45 & 83 & 145 & 200 & 262 & 280 & 288 & 298 \\ \hline \end{array}$$

Average Turnaround time = (45+(298-5)+(288-8)+(83-15)+(200-20))/5= 866/5 = 173.2

Average Waiting time = 
$$(0 + (83+117+8-5) + (200+18-8) + (45-15) + (145-20))/5$$
  
=  $(0+203+210+30+125)/5 = 568/5 = 113.6$ 

## **IRRVQ:**

Average Turnaround time = (45+(298-5)+(263-8)+(83-15)+(214-20))/5= 855/5 = 171

Average Waiting time = (0 + (159+34+15+0-5) + (121+55+17-8) + (45-15) + (83+76-20))/5= (0+203+185+30+139)/5 = 557/5 = 111.4

#### **SARR:**

Average Turnaround time = (45+(298-5)+(277-8)+(191-15)+(278-20))/5= 1041/5 = 208.2

Average Waiting time = (0 + (45+146+17-5) + (99+108-8) + (153-15) + (191+32-20))/5= (0+203+199+138+203)/5 = 743/5 = 148.6

# **R.P-5**:

←25→			←50→			←100→			
P1	P1	P2	P3	P4	P5	P2	P3	P5	
25	45	95	145	183	233	273	293	298	

Average Turnaround time = (45+(273-5)+(293-8)+(183-15)+(298-20))/5 = 1044/5 = 208.8Average Waiting time

$$= ((0+0) + (45+138-5) + (95+128-8) + (145-15) + (183+60-20))/5$$
  
=  $(0+178+215+130+223)/5 = 746/5 = 149.2$ 

# MRR:

←45→		←5	2→			←25→		
P1	P4	P5	P3	P2	P5	P3	P2	P2
45	83	135	187	239	242	260	295	298

Average Turnaround time = (45+(298-5)+(260-8)+(83-15)+(242-20))/5 = 880/5 = 176Average Waiting time

$$= (0 + (187+21+0-5) + (135+55-8) + (45-15) + (83+104-20))/5$$
  
=  $(0+203+182+30+167)/5 = 582/5 = 116.4$ 

# **DABRR:**

Average Turnaround time = (45+(298-5)+(271-8)+(83-15)+(138-20))/5= 787/5 = 157.4

Average Waiting time

$$= (0+(201+7+0-5)+(138+63-8)+(45-15)+(83-20))/5$$

= (0+203+193+30+63)/5 = 489/5 = 97.8

Table 13: Comparison between RR, DQRRR, IRRVQ, SARR, RP-5, MRR, DABRR for case-VI.

Algorithm	Round Robin	DQRRR	IRRVQ	SARR	RP-5	MRR	DABRR
Time Quantum	25	45, 62,	45,38, 17,	45,54,	25, 50,	45, 52,	45, 63, 17,
		18, 10	15, 20	16,20	100	35, 25	10
Context	13	7	10	8	8	8	7
Switch							
Average	173.2	113.6	111.4	148.6	149.2	116.4	97.8
Waiting time							
Average Turn	232.8	173.2	171	208.2	208.8	176	157.4
Around Time							

# 4.3.4. Comparison

This section provides the comparative analysis of seven algorithms on the basis of their resulted waiting time and turnaround time. Table 14 shows the performance analysis of seven algorithms by summarizing the waiting time and turnaround time resulted from case IV, V, VI.

Table 14: Comparison of six algorithms by aggregating their performance measure from case IV, V, VI.

	Without 0 Arival Time													
	Context Switch					Waiting Time				Turnaround Time				
	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL	Ascending Burst Time	Descending Burst Time	Random Burst Time	TOTAL		
RR	15	13	13	41	144.40	191.00	173.20	508.60	205.60	250.80	232.80	689.20		
DQRRR	7	7	7	21	107.20	138.40	113.60	359.20	168.40	198.20	173.20	539.80		
IRRVQ	10	10	10	30	98.20	133.80	111.40	343.40	159.40	193.60	171.00	524.00		
SARR	7	7	8	22	88.00	172.40	148.60	409.00	149.20	232.20	208.20	589.60		
RP-5	8	8	8	24	104.40	197.00	149.20	450.60	165.60	256.80	208.80	631.20		
MRR	7	7	8	22	90.00	124.60	116.40	331.00	151.20	184.40	176.00	511.60		
DABRR	7	7	7	21	88.20	125.00	97.80	311.00	149.40	184.80	157.40	491.60		

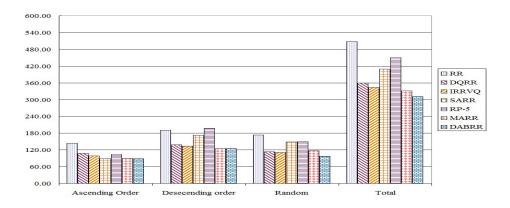


Figure 4: Waiting time for processes without 0 arrival time.

Figure 4 & 5 show the graphical analysis of the waiting time and turnaround time respectively of all algorithms based on case IV, V, VI. By evaluating these two graphs—we conclude that our algorithm performs better than other six algorithms in case of processes arrived with out zero arrival time.

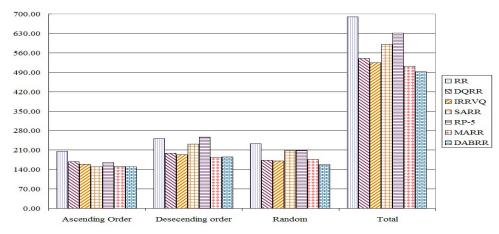


Figure 5: Turnaround time for processes without 0 arrival time.

# 5. RESULT ANALYSIS

From the analysis of all the algorithms it is concluded that our algorithm performs better than other compared algorithms in all cases.

Table 15: Percentage of waiting time reduced by each algorithm.

Waiting Time									
RR DQRRR IRRVQ SARR RP-5 MRR DABR							DABRR		
<b>Grand Total</b>	1,155.40	859.00	843.60	891.40	1,080.20	730.80	679.00		
Percentage Gain	0.00%	25.65%	26.99%	22.85%	6.51%	36.75%	41.23%		

Figure 6 depicts the percentage of waiting time saved by each algorithm than traditional Round Robin algorithm. Proposed algorithm "DABRR" saves 41% waiting time as compared to traditional Round Robin.

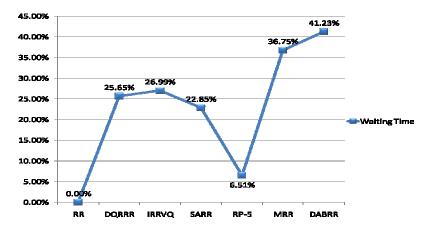


Figure 6: Percentage of waiting time reduced.

Table 16: Percentage of turnaround time reduced by each algorithm.

Turnaround Time									
	RR DQRRR IRRVQ SARR RP-5 MRR DABRR								
<b>Grand Total</b>	1,551.60	1,255.20	1,239.80	1,287.60	1,476.40	1,127.00	1,075.20		
Percentage Gain	0.00%	19.10%	20.10%	20.10%	4.85%	27.37%	30.70%		

Figure 7 depicts the percentage of turnaround time saved by each algorithm than traditional Round Robin algorithm. Proposed algorithm "DABRR" saves 31% turnaround time as compared to traditional waiting time.

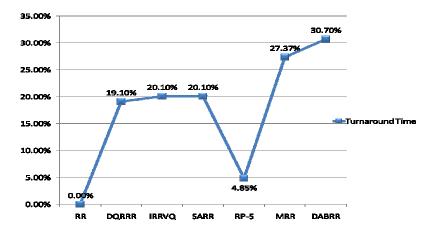


Figure 7: Percentage of turnaround time reduced.

# 6. CONCLUSIONS

This paper presents a variant of Round Robin scheduling algorithm. Comparative analysis of various algorithms like RR, DQRRR, IRRVQ, SARR, RP-5, MRR, and the proposed algorithm DABRR has been done. The proposed algorithm provides better performance metrics than the

above discussed algorithms by minimizing the average waiting time and average turnaround time. In future we want to improve this algorithm for multiprocessor environment.

#### REFERENCES

- [1] Abraham Silberschatz, Peter B. Galvin, Greg Gagne (2009) "Operating System Concepts", eighth edition, Wiley India.
- [2] Neetu Goel, R.B. Garg, (2012) "A Comparative Study of CPU Scheduling Algorithms", International Journal of Graphics & Image Processing Volume 2 issue 4, pp 245-251.
- [3] Jayashree S. Somani, Pooja K. Chhatwani, (2013) "Comparative Study of Different CPU Scheduling Algorithms", International Journal of Computer Science and Mobile Computing, PP 310-318.
- [4] Rakesh Kumar Yadav, Abhishek K Mishra, Navin Prakash, Himanshu Sharma (2010) "An Improved Round Robin Scheduling Algorithm for CPU Scheduling", International Journal on Computer Science and Engineering, pp 1064-1066.
- [5] Ajit Singh, Priyanka Goyal, Sahil Batra (2010) "An Optimized Round Robin Scheduling Algorithm for CPU Scheduling", International journal on Computer Science and Engineering, pp 2383-2385.
- [6] Manish Kumar Mishra, Dr. Faizur Rashid (2014) "An Improved Round Robin CPU Scheduling Algorithm with Varying Time Quantum", International Journal of Computer Science, Engineering and Applications (IJCSEA), pp 1-8.
- [7] Rishi Verma, Sunny Mittal, Vikram Singh (2014) "A Round Robin Algorithm using Mode Dispersion for Effective Measure", International Journal for Research in Applied Science and Engineering Technology (IJRASET), pp 166-174.
- [8] Neetu Goel, R.B. Garg (2013) "Simulation of an Optimum Multilevel Dynamic Round Robin Scheduling Algorithm", International Journal of Computer Applications, pp 42-46.
- [9] M. Ramakrishna, G. Pattabhi Rama Rao (2013) "Efficient Round Robin CPU Scheduling Algorithm for Operating Systems", International Journal of Innovative Technology and Research, pp 103-109.
- [10] Rami J. Matarneh (2009) "Self-Adjustment Time Quantum in Round Robin Algorithm Depending on Burst Time of the Now Running Processes", American Journal of Applied Sciences, pp 1831-1837.
- [11] H.S.Behera, R. Mohanty, Debashree Nayak, (2010) "A New Proposed Dynamic Quantum with Re-Adjusted Round Robin Scheduling Algorithm and Its Performance", International Journal of Computer Applications, pp 10-15.

#### Authors

Amar Ranjan Dash has achieved his B.Tech. degree from Biju Patnaik University of Technology, Odisha, India and M. Tech. degree from Berhampur University, Odisha, India. His research interests include CPU Scheduling, Web Accessibility, and Cloud Computing.



Sandipta Kumar Sahu has achieved his B.Tech. degree from Biju Patnaik University of Technology, Odisha, India and currently pursuing his M. Tech. degree in computer science and engineering at National Institute of Science And Technology, Odisha, India. His research interests include Operating System, Software engineering, and Computer Architecture.



**Sanjay Kumar Samantra** has achieved his MCA degree from Berhampur University, Odisha, India and currently pursuing his M. Tech. degree in computer science and engineering at National Institute of Science And Technology, Odisha, India. His research interests include CPU scheduling, Grid computing, and Cloud Computing.

