



Revision Notes: Software Engineering Concepts

These notes cover concepts related to subarrays, sum of odd indexed elements, special indices, and optimization strategies discussed in the class.

1. Subarrays and Their Properties

Printing All Possible Subarrays

- Given an array and specified starting (`s`) and ending (`e`) indices, subarrays can be printed using a double loop where the outer loop defines the start and the inner loop defines the end.
- Time Complexity:** $O(N^2)$ due to nested loops for each possible subarray.
- Space Complexity:** $O(1)$ constant space as no extra space is used beyond input storage.

2. Sum of Odd Indexed Elements

Problem Statement

- Calculate the sum of elements at odd indices within given subarray limits.

Approach

- Brute Force:** Iterate over each element in the given range and check if it is at an odd index. If yes, add to sum.
 - Complexity:**
 - Time: $O(n)$ per query where n is size of the subarray range.
 - Space: $O(1)$ since no additional space is needed.

2. Optimized Approach using Prefix Sum

- Build a prefix sum array that stores cumulative sums of elements at odd indices up to each index.



- **Time and Space Complexity:** Both reduced to $O(n)$ due to the prefix computation and constant-time queries .

3. Special Indices

Definition

- A "Special Index" is an index that when removed, makes the sum of elements at even indices equal to the sum of elements at odd indices in the resulting array.

Algorithm

1. Iterative Check:

- For each index, create a temporary array excluding this index.
- Check if the sums of even and odd indexed elements match.
- **Complexity:**
 - Time: $O(n^2)$ as for each index a new temp array is created and sums calculated.
 - Space: $O(n)$ for the temporary arrays .

2. Optimized Approach:

- Use two prefix sum arrays for odd and even indexed sums respectively.
- Calculate the sums using prefix arrays:
 - $\text{odd_sum}[i-1] + \text{even_sum}[n-1] - \text{even_sum}[i]$
 - $\text{even_sum}[i-1] + \text{odd_sum}[n-1] - \text{odd_sum}[i]$
- Check the sums after removal for equality .

4. Closest Min-Max Optimization

Problem Statement

- Given an array, find the smallest window which contains both the minimum and maximum value in the array.

Observations and Approach



and `max`.

- **Prefix Arrays:**

- `prefix_min[i]` : Last index $\leq i$ where `min_val` appears.
- `prefix_max[i]` : Last index $\leq i$ where `max_val` appears.

- **Calculation:**

- At each step, calculate the subarray length using these indices.
- Update the minimum subarray length accordingly.
- **Complexity:** Both Time and Space are reduced to $O(n)$.

Conclusion

These methods utilize both brute force and optimized algorithmic techniques such as prefix sums for efficient calculations, especially when handling multiple queries or operations on subarrays. Proper understanding and implementation of these methods are crucial in optimizing software applications dealing with large datasets or requiring real-time data processing.