

Recap : Regular expressions over an alphabet Σ

RE r	$L(r)$
\emptyset	\emptyset
ϵ	$\{\epsilon\}$
$a \ (a \in \Sigma)$	$\{a\}$

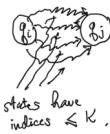
If r_1, r_2 are REs denoting sets R_1 and R_2
then

$(r_1 r_2)$	$R_1 \cup R_2$
(r_1^*)	R_1^*

Thm If $L = L(A)$ for some DFA, then there is a r.e.
 r s.t. $L = L(r)$.

Proof : Let $A = (Q, \Sigma, \delta, q_i, F)$

where $Q = \{q_1, \dots, q_n\}$, $F \subseteq Q$



Let R_{ij}^k denote the set of all strings w which take the DFA A from q_i to q_j w/o passing through any state q_ℓ whose index $\ell > k$.

$$L(A) = \bigcup_{q_i \in F} R_{i,j}^n$$



We define R_{ij}^k inductively as follows:

Base Case $R_{ij}^0 = \begin{cases} \{a \mid \delta(q_i, a) = q_j\} & \text{if } i \neq j \\ \{\epsilon\} & \text{if } i = j \end{cases}$

Inductive Case $R_{ij}^k = R_{ij}^{k-1} \cup R_{i,k}^{k-1} \cdot (R_{k,k}^{k-1})^* \cdot R_{k,j}^{k-1}$

We show that for each i, j and k there exists a r.e. r_{ij}^k which denotes the language R_{ij}^k by induction on k .

Basis ($k=0$) R_{ij}^0 is a finite set of strings each of which is either ϵ or a single symbol.

Then $r_{ij}^0 = a_1 + a_2 + \dots + a_p$ when $i \neq j$
 $R_{ij}^0 = \{a_1, \dots, a_p\}$

when $i=j$
 $R_{ij}^0 = \{a_1, a_2, \dots, a_p, \epsilon\}$
 $r_{ij}^0 = \epsilon$
i.e. for no $a \in \Sigma$, $q_i = \delta(q_i, a)$
 $r_{ij}^0 = \emptyset$

Inductive Step $r_{ij}^k = r_{ij}^{k-1} + r_{i,k}^{k-1} \cdot (r_{k,k}^{k-1})^* \cdot r_{k,j}^{k-1}$



$K=0$

$$\begin{aligned} r_{11}^0 &= 1 + \epsilon \\ r_{12}^0 &= \emptyset \\ r_{21}^0 &= \emptyset \\ r_{22}^0 &= 0 + 1 + \epsilon \end{aligned}$$

$K=1$

$$\begin{aligned} r_{11}^1 &= r_{11}^0 + r_{11}^0 \cdot (r_{11}^0)^* \cdot r_{11}^0 = (1 + \epsilon) + (1 + \epsilon)(1 + \epsilon)^* (1 + \epsilon) \\ r_{12}^1 &= r_{12}^0 + r_{11}^0 \cdot (r_{11}^0)^* \cdot r_{12}^0 \\ r_{21}^1 &= 0 + 0 \\ r_{22}^1 &= 0 + 0 \end{aligned}$$

RE for $L(A)$? $r = r_{12}^2$