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### Assignment 3

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1

$L$  is context free Language and by definition of  $L_1$ , it contains all prefixes of language  $L$ .

I am going to show  $L_1$  is also context free Language.

We know if  $L$  is context free then there exist a PDA for it. Let's call this PDA as  $M_1$ .

Now we take another copy of same PDA and call it  $M_2$  and remove input sign from starting state of  $M_2$  means starting state will be only of  $M_1$ .

In  $M_2$  we change the input symbol of all transitions to  $\epsilon$ .

Now we add transition from state of  $M_1$  to corresponding state in  $M_2$  as  $\epsilon, \epsilon \rightarrow \epsilon$ .

And only accepting state corresponding to  $M_1$  will be declared as accepting states in  $M_2$  and acceptance of  $M_1$  states will be removed.

This will be our whole PDA and start state will same remain as start state of  $M_1$ .

Let whole PDA is called as  $M$ .

We know that  $M$  initially takes input in  $M_1$  and then after getting all string as input we move to  $M_2$  using epsilon transitions and there we check if we can reach final state of  $M_2$ .

Hence all the Prefixes will be accepted by this PDA and this will generate all Prefix. Hence  $L_1$  is context-free Language.

2

$L$  is context free Language and by definition of  $L_1$ , it contains all suffix of language  $L$ .

I am going to show  $L_1$  is also context free Language.

We know if  $L$  is context free then there exist a PDA for it Let's call this as  $M_1$ .

Now we take another copy of same PDA and call it  $M_2$  and remove input sign of  $M_1$  means starting state will be only of  $M_2$ .

There will be no accepting state in  $M_2$ .

In  $M_2$  we change input symbol of all transitions to  $\epsilon$ .

Now we add transition from state of  $M_1$  to corresponding state in  $M_2$  as  $\epsilon, \epsilon \rightarrow \epsilon$ .