

## DFA Minimization

- Idea
- For any regular language  $L$ , the DFA  $M_L$  obtained from  $R_L$  is the unique (up to iso) minimum DFA that accepts  $L$ .
  - $R_M$  refines  $R_L$  for any DFA  $M$  that accepts  $L$ .

3. Idea behind minimizing a given DFA  $M$ :

— Merge states of  $M$  : merge  $q$  and  $q'$  iff

$$\delta(q, x) = q \text{ and } \delta(q', y) = q' \Rightarrow x R_L y$$

• Def Given a DFA  $M$  (where all states are reachable from  $q_0$ )  
Define equivalence reln.  $\equiv$  on the states of  $M$ :  
 $p \equiv q$  iff for each  $x \in \Sigma^*$ ,  $\delta(p, x) \in F \Leftrightarrow \delta(q, x) \in F$   
—  $p$  and  $q$  are equivalent if  $p \equiv q$   
distinguishable if  $p \not\equiv q$

- Need to show : Correctness of  $\equiv$  wrt which states can be merged  
 $p \equiv q \Leftrightarrow \forall x \forall y. \delta(q, x) = p \text{ and } \delta(q, y) = q \Rightarrow x R_L y$   
 $\Rightarrow$   
 $x R_L y$

Proof : Easy (Exercise)

Computing  $\equiv$  : How do we determine if  $p \equiv q$ ,  
where  $p, q \in Q$ ?

### Inductive Definition of $\neq$ (distinguishability)

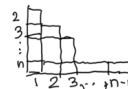
- If  $p \in F$  and  $q \notin F$ , or vice versa, then  $p \neq q$ .
  - If for some  $a$ ,  $\delta(p, a) \neq \delta(q, a)$  then  $p \neq q$ .
- Every other pair is equivalent.

Claim : This defn. is equivalent to the one given earlier

Proof : Exercise.

### Minimization Algorithm

- Eliminate all unreachable states from  $M$ .
- Distinguish  $\leftarrow \{ \langle p, q \rangle \mid p \in F \text{ and } q \notin F, \text{ or vice versa} \}$
- OldDistinguish  $\leftarrow \emptyset$  /\* initializations \*/
- while (Distinguish  $\neq$  OldDistinguish)
- OldDistinguish  $\leftarrow$  Distinguish
- for every symbol  $a$  and for every  $\langle p, q \rangle \notin$  Distinguish do
- if  $\langle \delta(p, a), \delta(q, a) \rangle \in$  Distinguish then  
Distinguish  $\leftarrow$  Distinguish  $\cup \{ \langle p, q \rangle \}$



### Time Complexity

# iteration of while loop :  $O(n)$  Why?

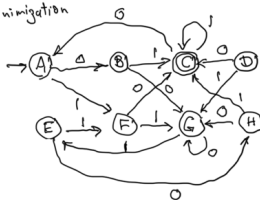
Time taken in each iteration :  $O(1 \leq n^2)$

Total time  $O(n^3)$

There is an algorithm that runs in time  $O(n \log n)$ .

### Ex DFA Minimization

DFA  $M$



### Table

	A	B	C	D	E	F	G
A							
B							
C							
D							
E							
F							
G							