

Assignment 2

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1

Language L consists of strings of length $2n$ where last n characters are all ones where $n \geq 0$.

Let there are 2 strings a, b where $|a| = n$ and $|b| = m$ where $n \neq m$.

We will show that relation R_L has infinite index. Taking $n < m$ and for all pairs of n, m appending 01^{n+1} at end of both strings.

we will show that $a01^{n+1} \in L$ and $b01^{n+1} \notin L$.

$a01^{n+1} \in L$ as taking w as $a0$ and it has length $n+1$ and string has $n+1$ 1's at end. So it belongs to L .

$b01^{n+1} \notin L$ as total length is $m+n+2$. So it must have $(m+n+2)/2$ 1's at end but it have only $n+1$ 1's which are less than required which can be shown as follows :-

$n < m$ and add $n+2$ both sides gives

$2n+2 < m+n+2$ and dividing by 2 gives desired result.

So a and b must be in different classes for all n and m .

Hence number of classes of R_L are infinite and by Myhill-Nerode theorem L is non-regular.

2

Language L consists of strings of length $2n$ where last n characters are complement of first n characters where $n \geq 0$.

Let there are two strings a and b where $a=0^n$ and $b=0^m$ where $n \neq m$.

We will show that relation R_L has infinite index. Taking $n < m$ and for all pairs of n, m appending 1^n at end of both strings.

we will show that $0^n 1^n \in L$ and $0^m 1^n \notin L$.

As $0^n 1^n \in L$ is trivial as taking w as 0^n gives required string. Hence $0^n 1^n \in L$.

Also $0^m 1^n \notin L$ is trivial as $n \neq m$ So, it can never be in L .

So a and b are in different classes for all n and m .

Hence number of classes of R_L are infinite and by Myhill-Nerode theorem L is non-regular.