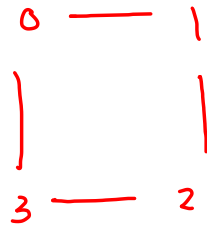
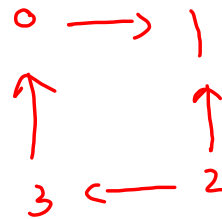


## Kosaraju Algorithm

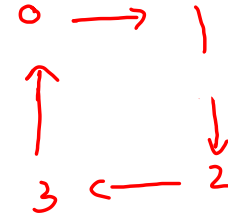
: Strongly connected comp  
└──────────────────┘  
    ↓  
directed comp



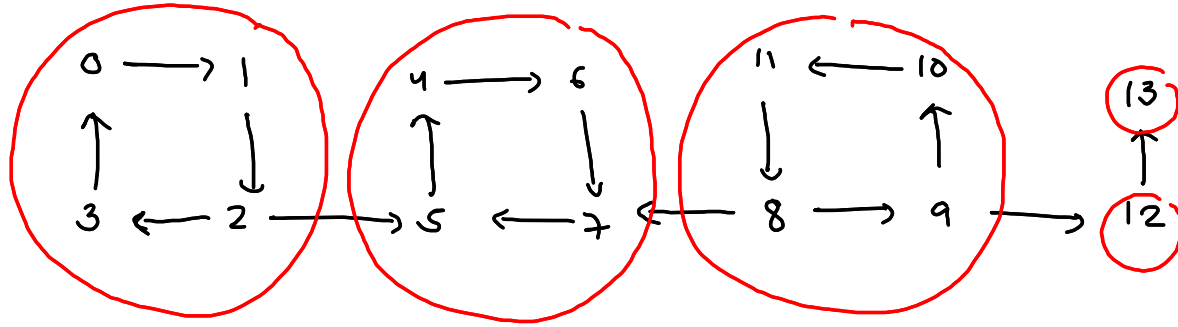
connected comp



not a  
strongly conn.  
comp



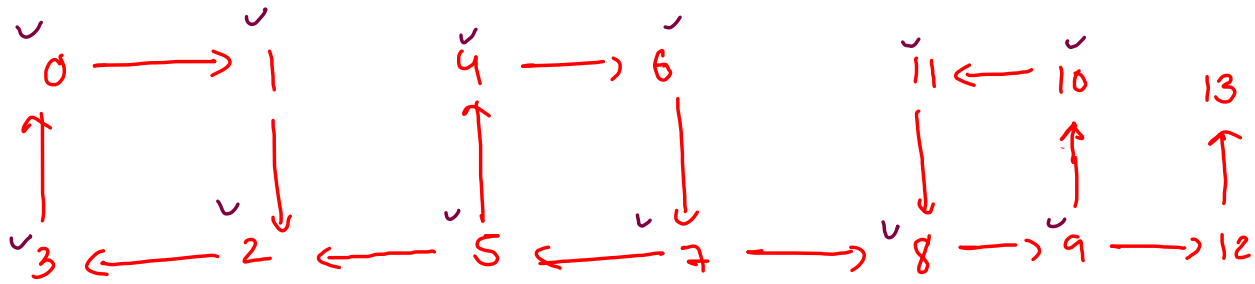
Strongly connected  
comp.



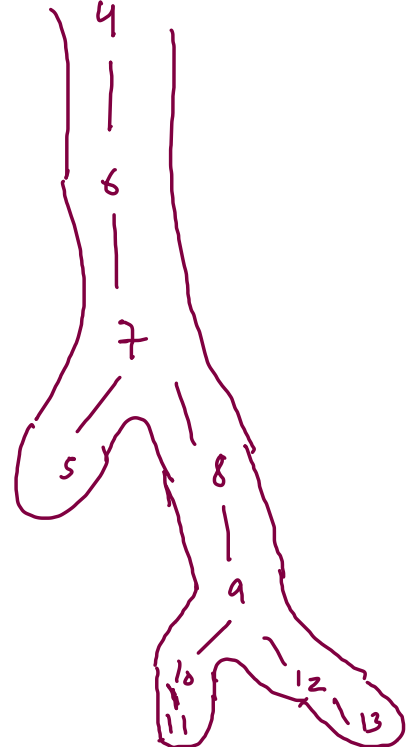
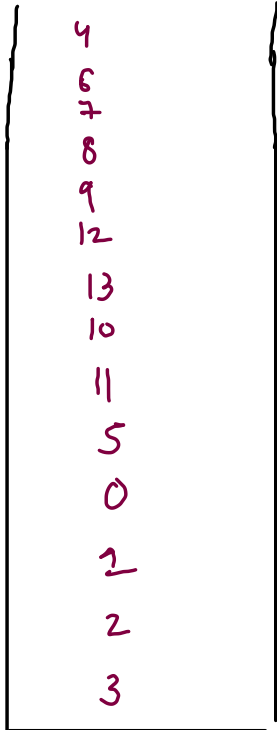
5 SCC

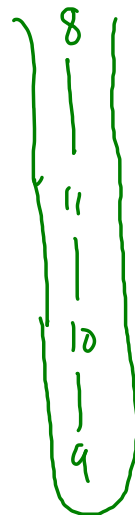
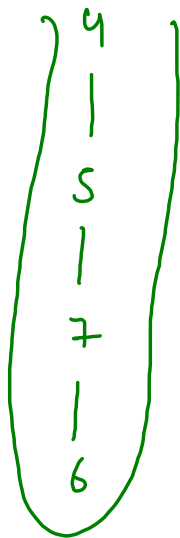
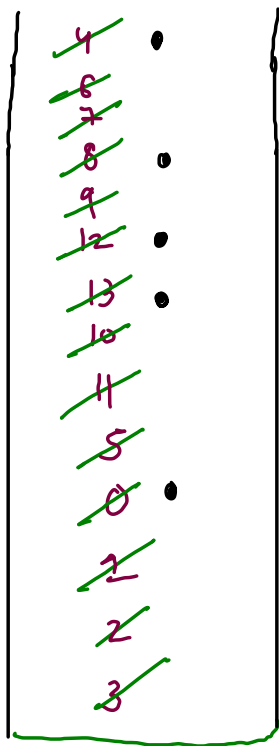
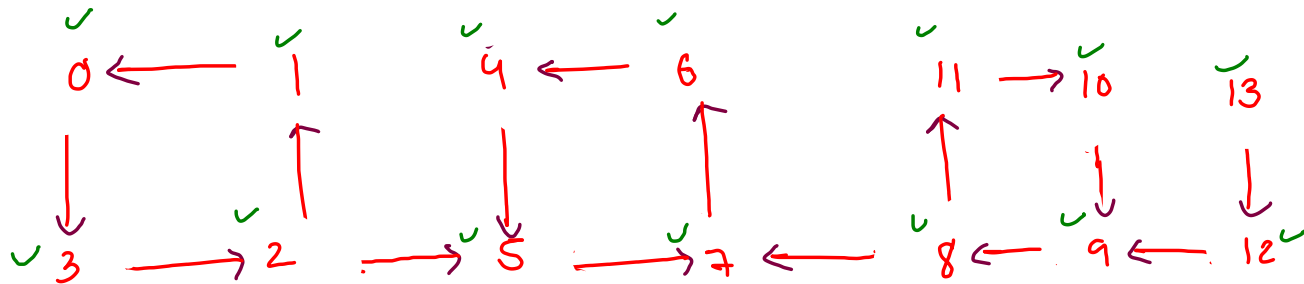
Strongly connected comps.

no of cycles  $\neq$  no. of SCC



①. DFS

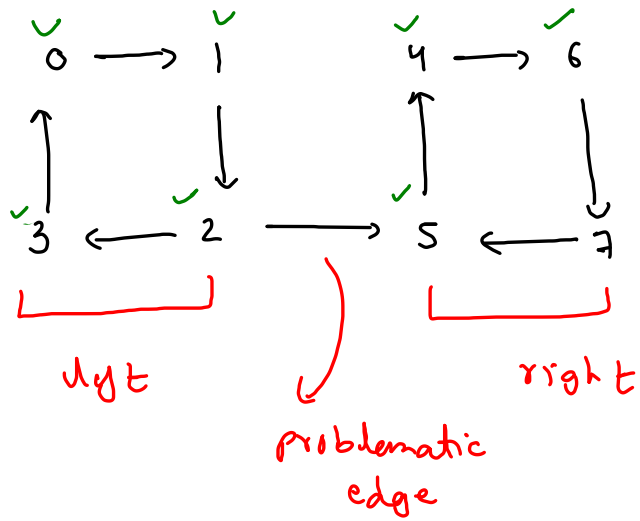




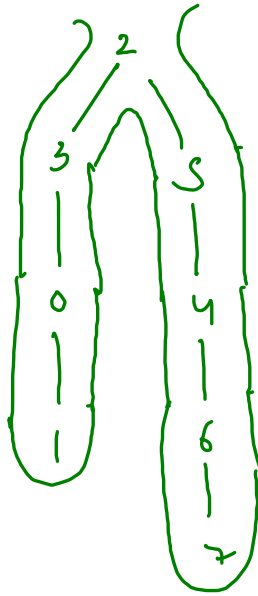
12

13

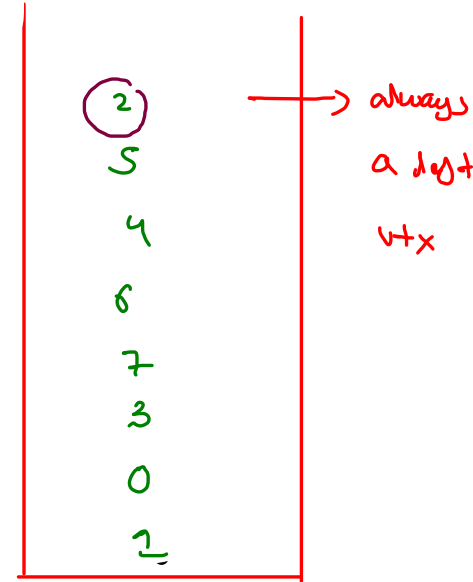


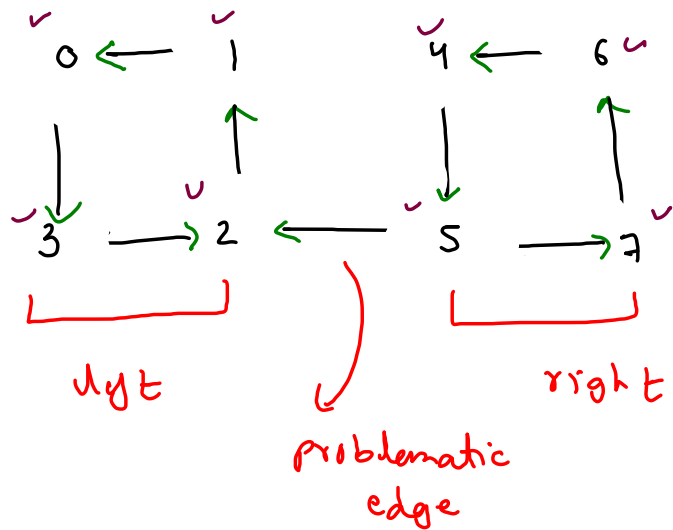


why?

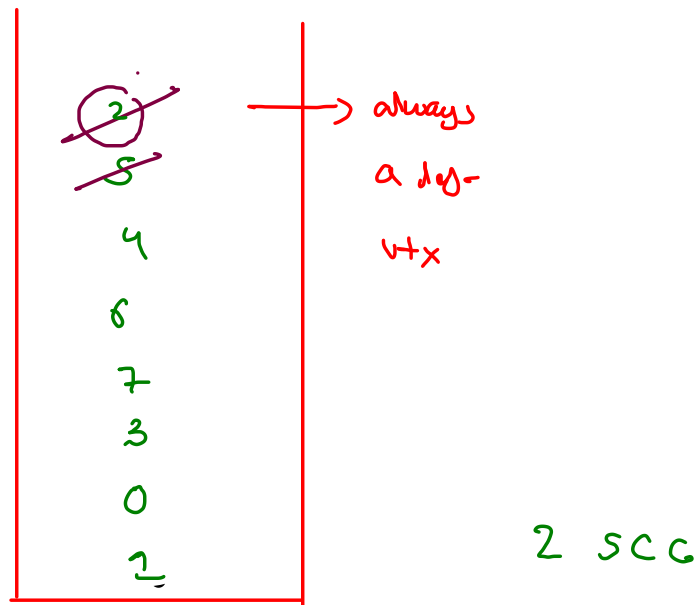
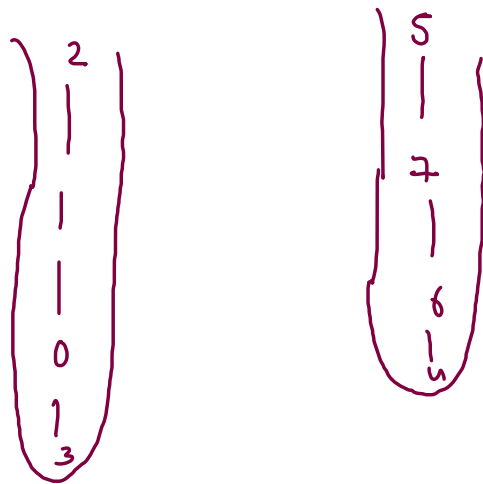


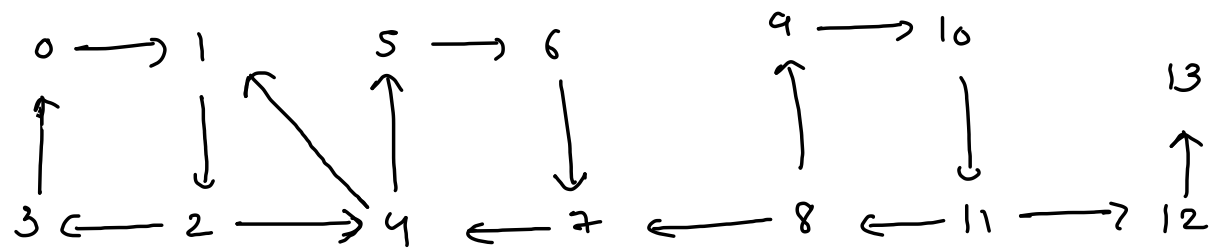
- ①. performs <sup>complete</sup> dfs and add vertices in stack in postorder.
- ②. reverse this graph.
- ③. now travel dfs using stack (no. of dfs == scc)



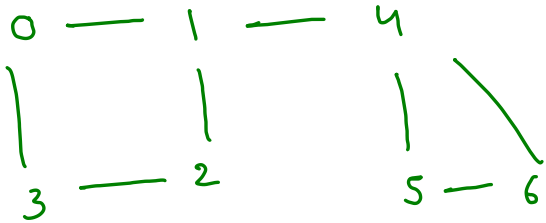


2





## Articulation Point



1 & 4 are

articulation points.

brute force :

remove  
each  
vtx



no. of  
connected comps





discovery time &  
low time

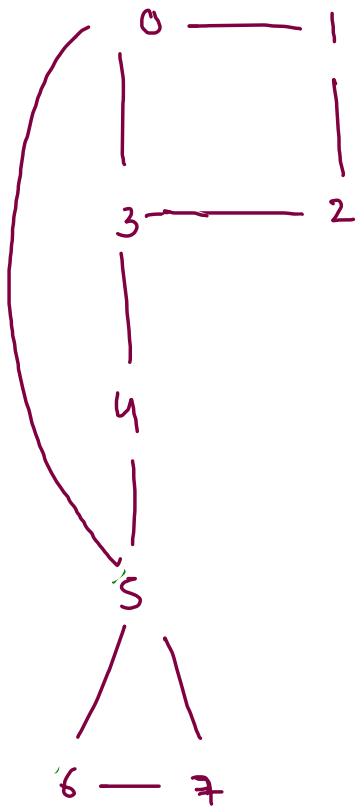
5 is ap

4 is ap

3 is ap

discovery time : the no. at  
which you are  
travelled in dfs.

low time : except your parents,  
the lowest vertex you know  
(discovered earlier)



if (nbr == par) ?

// ignore

}

else if (nbr == unvisited) {

①. go to nbr.

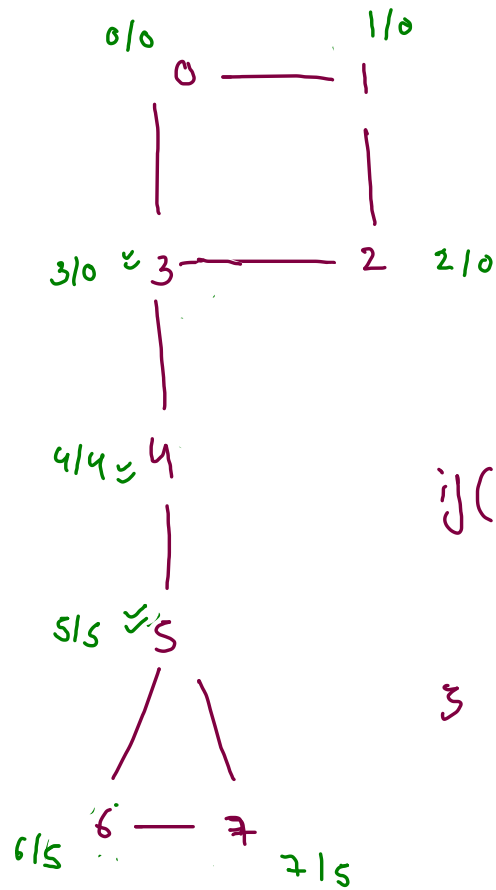
②.  $\text{low}[\text{src}] = \text{math.min}(\text{low}[\text{nbr}], \text{low}[\text{src}])$

}

else if (nbr == visited) {

$\text{low}[\text{src}] = \text{math.min}(\text{low}[\text{src}], \text{dis}[\text{nbr}])$

}



src  $\rightarrow$  manage differently

if (nbr == par) ?

// ignore

{

else if (nbr == unvisited) {

① - go to nbr.

② -  $\text{low}[\text{src}] = \text{math.min}(\text{low}[\text{nbr}], \text{low}[\text{src}])$

}

else if (nbr == visited) {

$\text{low}[\text{src}] = \text{math.min}(\text{low}[\text{src}], \text{disc}[\text{nbr}])$

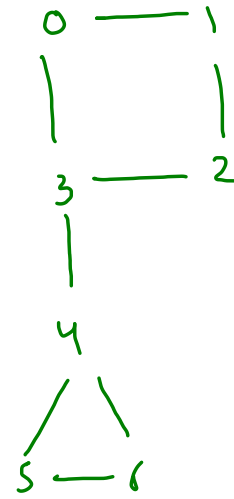
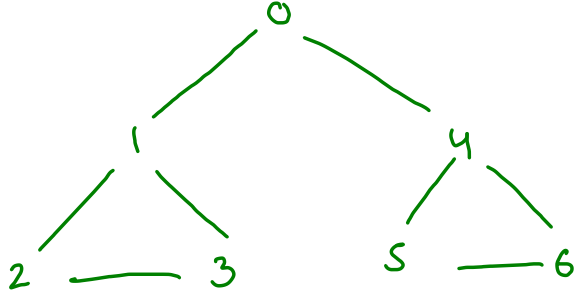
}

if ( $\text{disc}[\text{src}] \leq \text{low}[\text{nbr}]$ ) {

src is an ap;

}

0 is ap



0 is not ap