# Project 3: Designing a 32-bit CPU

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### 1 Introduction

This goal of this project is to design a stripped down version of the MIPS processor. The processor will be a 32-bit version of the processor discussed in class and the text book, however, its instruction set will be a small subset of the MIPS processor's full capability.

### 1.1 Background Information

#### 1.1.1 MIPS

MIPS is a reduced instruction set computer (RISC) instruction set architecture (ISA). It defines three types of instruction types: R (register), I (Immediate), and J (Jump). For the implementation that this project is focused on, only R and I instructions will be executed. R type instructions are the most common form of instructions. The format for an R-type instruction is:

Bits[31:26]	Bits[25:21]	Bits[20:16]	Bits[15:11]	Bits[10:6]	Bits[5:0]
opcode	Rs	Rt	Rd	shamt	funct

For this instruction, the opcode field is always  $000000_2$ , while the function code funct is used to determine which instruction is to be carried out. Rs and Rt are the two registers in which the operation reads and Rd is the destination of the result. Some instructions require a shift amount (shamt), so it is specified explicitly.

The I type instruction involves an immediate value, so the instruction format must accommodate this. The format of this type of instruction is:

Bits[31:26]	Bits[25:21]	Bits[20:16]	Bits[15:0]
opcode	Rs	Rt	immediate

For this instruction, the op code field is used to define the specific instruction, Rs is the register in which the operation acts on along with the immediate value as the other operand. Rt is the destination register in which the result is stored.

#### 1.1.2 Datapath and Control

A datapath is a collection of functional units that perform data processing operations. It includes units such as a program counter, a register file, instruction memory, an ALU, data memory, and a control unit. Figure 1 shows a high level overview of a simple datapath with control.

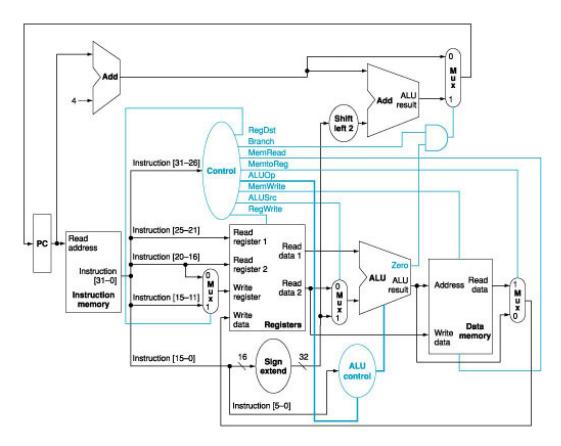


Figure 1: Datapath Overview

# 2 Design

#### 2.1 Instruction Set

Table 1 shows the instructions that were chosen to be implemented in the CPU with the respective OpCode and Function Field for each instruction.

OpCode[31:26]	Function Field [5:0]	Instruction	Example Operation
$100011_2$		lw	lw \$t3, 200(\$s2)
$101011_2$		SW	sw \$t4, 100(\$t3)
$000000_2$	$100000_2$	add	add \$s3, \$t2, \$s2
$000000_2$	$110000_2$	sub	sub \$s3, \$t2, \$s2
$000100_2$		beq	beq \$s5, \$s2, 500
$000000_2$	$000001_2$	nand	nand \$s5, \$s1, \$s2
$000010_2$		andi	andi \$s6, \$s2, 0x00FF
$000000_2$	$000010_2$	or	or \$s8, \$s1, \$s2
$000011_2$		ori	ori \$s7, \$s1, 0x00FF

Table 1: CPU Instruction Set

Because it was only required to implement 9 instructions and the MIPS instruction set format requires 6 bits for op code and function field, it was an easy decision to choose these values for the implemented instructions. For all R-type instructions, the functions fields were chosen to be vastly different from one another to make debugging easier for the team. Likewise, the same approach was taken for the op code decisions for the I-type instructions.

### 2.2 Memory

For this project, it seemed unnecessary to implement memory of 4GB  $(2^{32})$ . It was chosen to use an array of 256 words instead. If need be, this memory size could be upgraded easily, so this choice does not hinder performance on the actual design of the CPU.

### 2.3 Datapath

Because of the simplicity of this design, the implemented datapath did not need to be modified by much from Figure 1. Therefore, the design of a single cycle datapath from the textbook acted as the skeletal structure of the final

implementation. Because an ALU and Register File were previously implemented in earlier projects, it was necessary to extend their functionality to be able to handle 32-bit words. Once this was complete, this left the data memory entity to be completed so that it could be included in the processor entity. As mentioned earlier in Section 2.2, this entity contains an array of 256 words, and allows for reading and writing.

The processor entity combines all of the components into the desired datapath. It synchronizes the clock of the instruction memory, data memory, and register file so that the entire system is in sync with an external clock signal. The program counter is updated during the rising edge of the clock, and all writes happen on the falling edge. The processor relies on the control unit to carry out the instruction read from memory. Figure 2 shows the overview block diagram of the implemented datapath for the CPU.

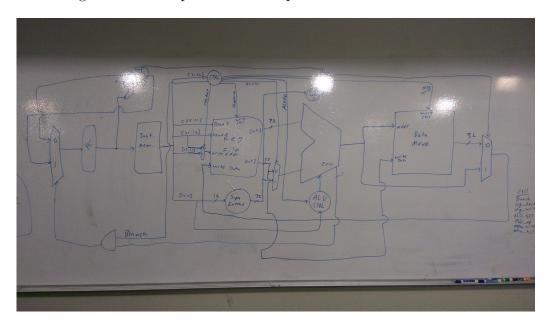


Figure 2: Implemented Data Path Overview

#### 2.4 Control

The control lines can also be seen in Figure 2. It is a simple design of several signals acting as the sel line for a series of multiplexers. Based on the op codes and function field read from the instruction memory, the signals are

asserted accordingly to relay the correct signals into the Register File, ALU, and Data memory. This unit is what determines which units will read/write, and what operations the ALU should perform.

# 3 Analysis

While this processor was optimized to be able to fully accomplish the tasks specified in the business requirements document, it could still be improved. In its current stage, it can be considered a bare bones prototype. To transform the current design into a processor on par with the current industry standard, a complete instruction set would have to be implemented. Furthermore, pipelining is a necessity to add. Any processor that doesn't implement pipelining is not making efficient use of its own components. After pipelining is implemented, hazard controls would need to coexist. This would allow for cool features of the processor to exist such as forwarding, making it a truly efficient piece of hardware.

#### 4 Simulation Results

Once the processor was completely designed, it was necessary to write some test bench code. To test each instruction, data had to first be written to memory, along with the program being loaded onto the CPU. Due to the amount of signals involved in the CPU, not all will be shown in the simulation. The clock, contents of the registers, data memory, and program counter will only be shown. Data Memory addresses  $0x00000001 \rightarrow 0x00000005$  were initialized with starting data. For simplicity, register numbers  $1 \rightarrow 8$  are s registers. This is not the convention in a usual MIPS implemented processor, however, for testing the instructions, this assignment is arbitrary. Please refer to the Test Bench Code in the Appendix for a detailed view of the testing procedure. The execution of the program begins at 38ns. The tested instructions are:

- 1. lw \$s1, 1(\$zero)
- 2. sw \$s1, 6(\$zero)
- 3. lw \$s2, 2(\$zero)

- 4. add \$s3, \$s1, \$s2
- 5. sub \$s4, \$s2, \$s1
- 6. beq \$s1, \$s2, 100
- 7. lw \$2, 4(\$zero)
- 8. nand \$s5, \$s1, \$s2
- 9. andi \$s6, \$s2, 0x00FF
- 10. ori \$s7, \$s1, 0x00FF
- 11. or \$s8, \$s1, \$s2
- 12. beq \$s1, \$s1, -0x000B

Figure 3 shows the first instruction being executed. The data memory at address 0x00000001 holds the value 0xAAAAAAA and register \$s1 is subsequently loaded with the data 0xAAAAAAA.

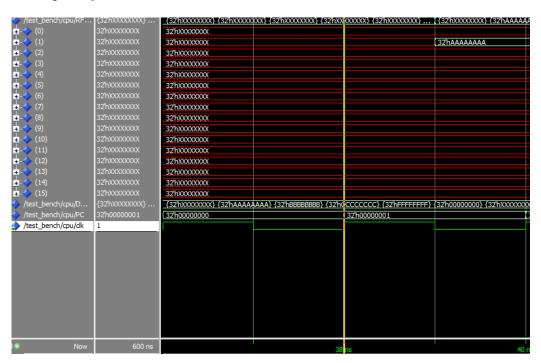


Figure 3: lw \$s1, 1(\$zero)

Figure 4 shows the second instruction. The value of 0xAAAAAAA in register \$s1 is successfully stored into the data memory at address 0x00000006.

/test_bench/cpu/RF	{32'hXXXXXXXXXX}	[32]hXXXXXXXXX [32]hAAAA	ΔΔΔ} {32'hXXXXXXXX} {32'	nXXXXXXXXXX {32'hXXXXXXXXXX {3
/test_bench/cpu/D	{32'hXXXXXXXXX}			[{32'hXXXXXXXXX} {32'hAAAAA
	32'hXXXXXXXXX	32'hXXXXXXXXX	,	
- (1)	32'hAAAAAAAA	32'hAAAAAAAA		
- (2)	32'hBBBBBBBB	32'hBBBBBBBB		
- (3)	32'hCCCCCCCC	32'hCCCCCCCC		
	32'hFFFFFFFF	32'hFFFFFFF		
	32'h00000000	32'h00000000		
- <del>-</del> (6)	32'hXXXXXXXX	32'hXXXXXXXX		32'hAAAAAAAA
	32'hXXXXXXXX	32'hXXXXXXXX		
J- <b>(</b> 8)	32'hXXXXXXXX	32'hXXXXXXXX		
)- <del>(</del> (9)	32'hXXXXXXXX	32'hXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXXX		
J- <b>(</b> 14)	32'hXXXXXXXX	32'hXXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXXX		
J- <b>(</b> 21)	32'hXXXXXXXX	32'hXXXXXXXX		
J- <del>(</del> (22)	32'hXXXXXXXX	32'hXXXXXXXX		
J- <b>(</b> 23)	32'hXXXXXXXX	32'hXXXXXXXXX		
J- <del>-</del> (24)	32'hXXXXXXXX	32'hXXXXXXXX		
J <del>-</del> (25)	32'hXXXXXXXX	32'hXXXXXXXX		
J <del>-</del> (26)	32'hXXXXXXXX	32'hXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXX		
Now	600 ps			

Figure 4: sw \$s1, 6(\$zero)

Figure 5 shows the third instruction. The value of 0xBBBBBBB at address 0x00000002 is successfully loaded into register \$s2.

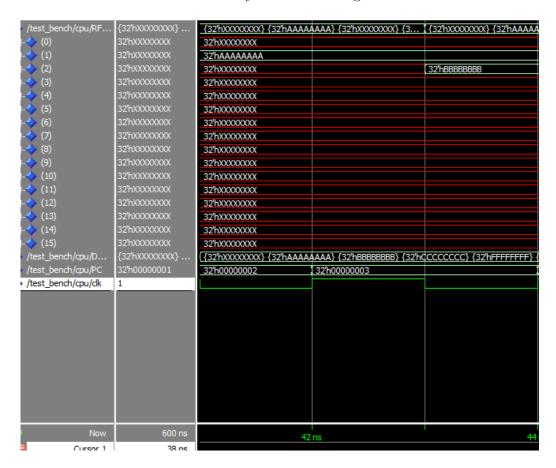


Figure 5: lw \$s2, 2(\$zero)

Figure 6 shows the fourth instruction. The values of 0xAAAAAAAA and 0xBBBBBBB are successfully added with the correct result of 0x66666665 being written into register \$s3.

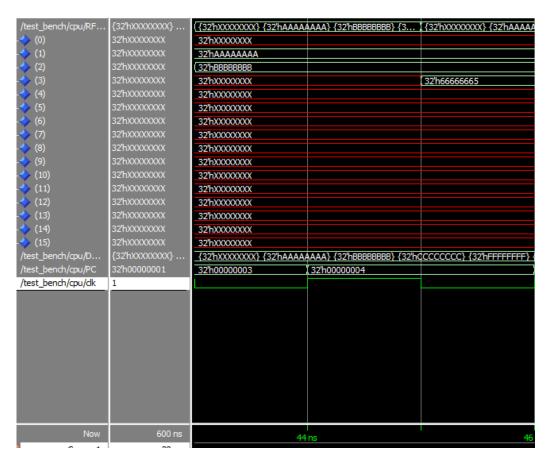


Figure 6: add \$s3, \$s1, \$s2

Figure 7 shows the fifth instruction. The value of 0xAAAAAAA in register \$s1 is subtracted from register \$s2. The result is 0x11111111 and is written back into register \$s4.

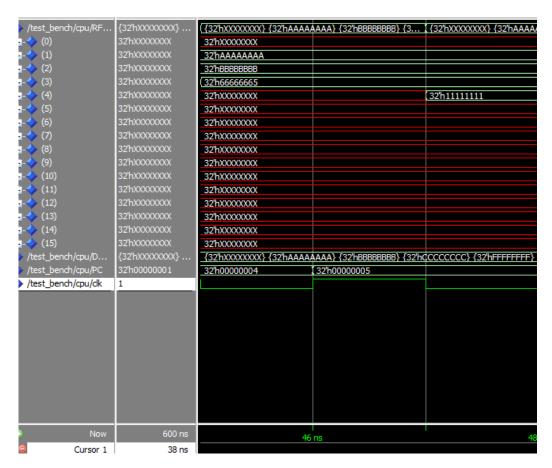


Figure 7: sub \$s4, \$s2, \$s1

Figure 8 shows the sixth instruction. The value of 0xBBBBBBBB in register \$s2 is compared to 0xAAAAAAA in register \$s1 for equivalence. Since they are not, the program counter increments to the next instruction.

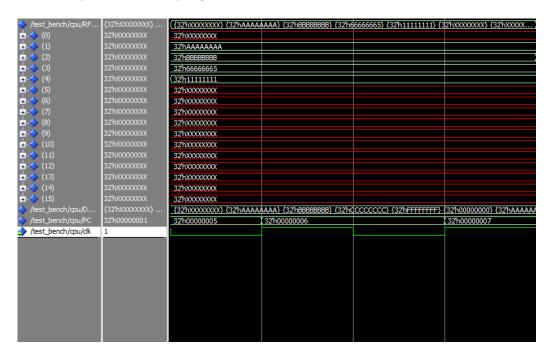


Figure 8: beq \$s1, \$s2, 100

/test_bench/cpu/RF	{32'hXXXXXXXXX}	{32'hXXXXXXXXX} {32'hAAAAA	AAA} {32'hBBBBBBBB} {3	{32'hXXXXXXXX} {32'hAAAAA
(0)	32'hXXXXXXXXX	32'hXXXXXXXX		
- <del>-</del> (1)	32'hXXXXXXXXX	32'hAAAAAAA		
- <del>-</del> (2)	32'hXXXXXXXX	32'hBBBBBBBB		32'hFFFFFFFF
	32'hXXXXXXXX	32'h66666665		
( <del>4</del> )	32'hXXXXXXXX	32h11111111		
	32'hXXXXXXXX	32'hXXXXXXXXX		
- <del>-</del> (6)	32'hXXXXXXXXX	32'hXXXXXXXXX		
· <del></del> (7)	32'hXXXXXXXX	32'hXXXXXXXX		
- <del>-</del> (8)	32'hXXXXXXXX	32'hXXXXXXXXX		
(9)	32'hXXXXXXXX	32'hXXXXXXXX		
- <del>-</del> (10)	32'hXXXXXXXXX	32'hXXXXXXXXX		
	32'hXXXXXXXXX	32'hXXXXXXXXX		
- <del>-</del> (12)	32'hXXXXXXXX	32'hXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXX		
- <del>-</del> (14)	32'hXXXXXXXX	32'hXXXXXXXXX		
> (15)	32'hXXXXXXXX	32'hXXXXXXXX		
/test_bench/cpu/D	{32'hXXXXXXXXX}	{32'hXXXXXXXX} {32'hAAAAA	AAA} {32'hBBBBBBBB} {32'h0	CCCCCCC} {32'hFFFFFFF}}
· <del></del> (0)	32'hXXXXXXXX	32'hXXXXXXXXX		
> (1)	32'hAAAAAAAA	32'hAAAAAAAA		
- <del>-</del> (2)	32'hBBBBBBBB	32'hBBBBBBBB		
(3)	32'hCCCCCCCC	32'hCCCCCCCC		
(4)	32'hFFFFFFFF	32'hFFFFFFF		
(5)	32'h00000000	32'h00000000		
- <del>-</del> (6)	32'hXXXXXXXX	32'hAAAAAAAA		
· <del></del> (7)	32'hXXXXXXXX	32'hXXXXXXXX		
- <del>(</del> 8)	32'hXXXXXXXX	32'hXXXXXXXX		
I- <b>(</b> 9)	32'hXXXXXXXX	32'hXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXX		
	32'hXXXXXXXX	32'hXXXXXXXX		
Now	600 ns	50 i		F2
	20	501	15	52

Figure 9: 1w \$s2, 4(\$zero)

Figure 10 shows the eighth instruction. The value of 0xFFFFFFFF in register \$s2 and 0xAAAAAAA in register are nanded and the result of 0x55555555 is both accurate and stored in register \$s5.

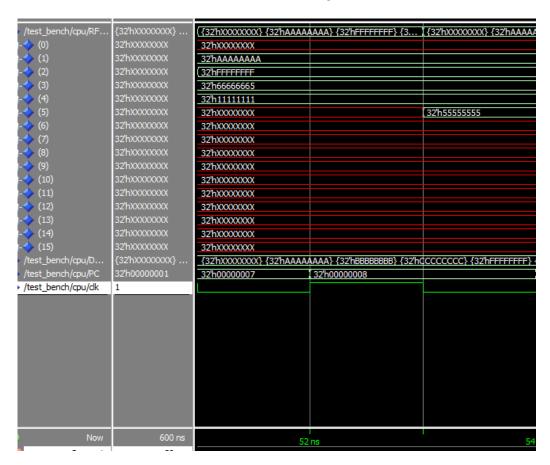


Figure 10: nand \$s5, \$s1, \$s2

Figure 11 shows the ninth instruction. The immediate value of 0x0000000FF is anded with the value of 0xFFFFFFFF in register \$s2. The correct result of 0x000000FF is calculated and written into register \$s6.

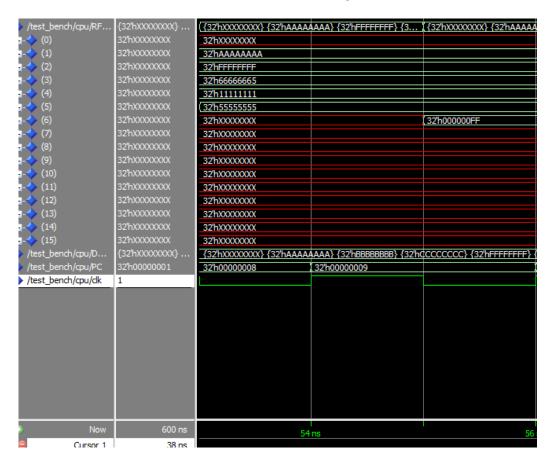


Figure 11: andi \$s6, \$s2, 0x00FF

Figure 12 shows the tenth instruction. The immediate value of 0x0000000FF is or'ed with the value of 0xAAAAAAA in register \$s1. The correct result of 0xAAAAAAFF is written into register \$s7.

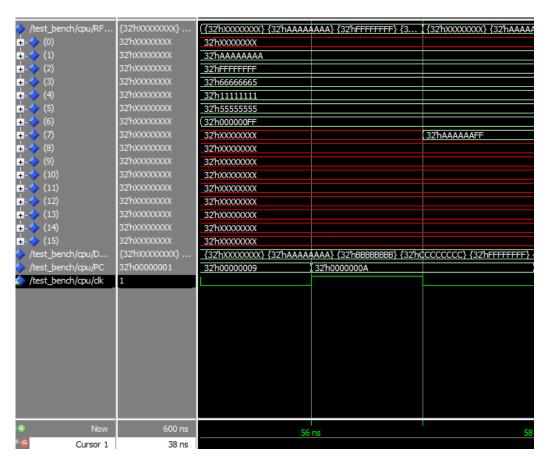


Figure 12: ori \$s7, \$s1, 0x00FF

Figure 13 shows the eleventh instruction. The value of 0xFFFFFFFF in register \$s2 and 0xAAAAAAA in register \$s1 are or'ed. The correct result of 0xFFFFFFFF is written into register \$s8.

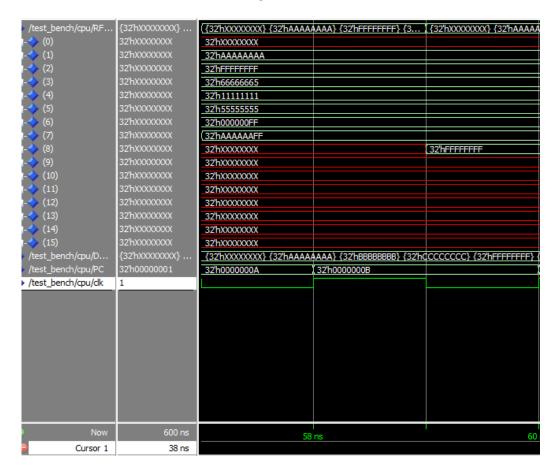


Figure 13: or \$s8, \$s1, \$s2

Figure 14 shows the twelfth instruction. The value in register \$s1 is checked for equivalence with itself. Since it is equal, the program counter branches backwards successfully back to 0x00000001. Thus the test program infinitely loops.

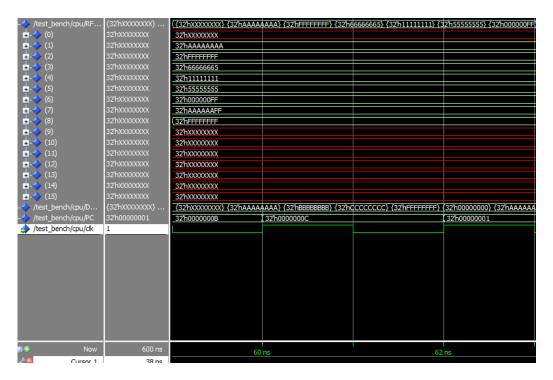


Figure 14: beq \$s1, \$s1, -0x000B

### 5 Conclusion

The design and implementation of a 32-bit CPU was a success. A set of 9 instructions were successfully implemented and verified with test bench code. All requested functionality was achieved. This 32-bit CPU can now be used in further projects and can be expanded upon to become a more efficient piece of hardware.

# **Appendix**

Listing 1: CPU Code

```
1 library ieee;
2 use ieee.std_logic_1164.all;
3 use ieee.numeric_std.all;
5 entity regFile is
    port (
6
7
    regA : out std_logic_vector(31 downto 0);
    regB : out std_logic_vector(31 downto 0);
9
    selA : in std_logic_vector(3 downto 0);
    selB : in std_logic_vector(3 downto 0);
10
    wData : in std_logic_vector(31 downto 0);
11
    registerWrite : in std_logic;
12
13
    selW : in std_logic_vector(3 downto 0);
    clk : in std_logic);
15 end regFile;
16
17 architecture behavioral of regFile is
18 type reg_arr is array (0 to 15) of std_logic_vector (31
     downto 0);
19 signal rData : reg_arr;
20 begin
21
    with selA
22
      select regA \le x"00000000" when b"0000",
      rData(to_integer(unsigned(selA))) when others;
23
    with selB
24
25
       select regB \le x"00000000" when b"0000",
      rData(to_integer(unsigned(selB))) when others;
26
27
28
    wrProc: process(clk) is
29
    begin
30
      if falling_edge(clk) then
       if(registerWrite = '1') then
31
32
         rData(to_integer(unsigned(selW))) <= wData;
      end if;
33
34
      end if;
35
    end process;
```

```
36 end behavioral;
37
38 ---
39
40 library ieee;
41 use ieee.std_logic_1164.all;
42 use ieee.numeric_std.all;
43
44 entity control is
45
     port (
46
       inst_in : in std_logic_vector(5 downto 0);
       func : in std_logic_vector(5 downto 0);
47
48
       stall : in std_logic;
       branch : out std_logic;
49
       reg_dest : out std_logic;
50
51
       reg_write : out std_logic;
52
       ALU_src : out std_logic;
53
      ALU_op : out std_logic_vector(2 downto 0);
54
      mem_write : out std_logic;
       mem_to_reg : out std_logic
55
56
     );
57 end control;
59 architecture behavioral of control is
     signal branch_o, reg_dest_o, reg_write_o, ALU_src_o,
     mem_write_o, mem_to_reg_o : std_logic;
     signal ALU_op_o : std_logic_vector(2 downto 0);
61
62
     signal branch_f, reg_dest_f, reg_write_f, ALU_src_f,
     mem_write_f, mem_to_reg_f : std_logic;
     signal ALU_op_f : std_logic_vector(2 downto 0);
63
64 begin
65
    -- set intermediate signals incase of r-type
     instruction
     with func select
66
       branch_f \le '0' \text{ when "100000"}, --add
67
         '0' when "110000", —sub
68
         '0' when "000001", --nand
69
         0, when 000010, —or
70
71
         '0' when others;
72
     with func select
```

```
reg_dest_f \le '1' \text{ when "100000", } --add
 73
          '1' when "110000", —sub
 74
          '1' when "000001", —nand
 75
          '1' when "000010", --or
 76
          'Z' when others;
 77
 78
      with func select
        reg_write_f <= '1' \text{ when "} 100000", --add
 79
          '1' when "110000", — \operatorname{sub}
 80
          '1' when "000001", --nand
 81
          '1' when "000010", —or
 82
 83
          'Z' when others;
      with func select
 84
 85
        ALU_src_f \le '0' \text{ when "100000", } --add
          '0' when "110000", —sub
 86
          '0' when "000001", —nand
 87
          '0' when "000010", ---or
 88
          'Z' when others;
 89
 90
      with func select
 91
        ALU_{op_f} \le "000" \text{ when } "100000", --add
          "001" when "110000", —sub
 92
          "010" when "000001", --nand
 93
          "100" when "000010", —or
 94
          "ZZZ" when others;
 95
 96
      with func select
        mem_write_f \ll 0 when "100000", --add
 97
          ^{\prime}0\,^{\prime} when ^{"}110000\,^{"} , --\mathrm{sub}
98
          '0' when "000001", —nand
99
           '0' when "000010", —or
100
101
          'Z' when others;
102
      with func select
        mem_to_reg_f <= '1' \text{ when "100000"}, --add
103
          '1' when "110000", —sub
104
          '1' when "000001", —nand
105
          '1' when "000010", --or
106
          'Z' when others;
107
108
     -- set intermediate signals incase of non r-type
109
       instruction
110
      with inst_in select
        branch_o <= '0' when "100011", --lw
111
```

```
^{\prime}0 ' when "101011" , \, —sw
112
113
          '1' when "000100", —beq
114
          '0' when "000010",
                              ---andi
115
          '0' when "000011",
                              --- o r i
116
          '0' when others;
      with inst_in select
117
118
        reg_dest_o \le '0' \text{ when "100011"}, --lw
          '0' when "101011", —sw
119
          '0' when "000100", —beq
120
          '0' when "000010", --andi
121
122
          '0' when "000011",
                              —— o r i
123
          'Z' when others;
124
      with inst_in select
125
        reg_write_o \ll 1' when "100011", —-lw
          '0' when "101011", ---sw
126
          '0' when "000100", —beq
127
128
          '1' when "000010", —andi
                              --- o r i
129
          '1' when "000011",
          'Z' when others;
130
131
      with inst_in select
        ALU\_src\_o \le '1' when "100011", --lw
132
          '1' when "101011", --sw
133
          '0' when "000100", —beq
134
          '1' when "000010", —andi
135
          '1' when "000011", --ori
136
          'Z' when others;
137
138
      with inst_in select
        ALU_{op_o} \le "000" when "100011", --lw
139
          "000" when "101011", --sw
140
          "001" when "000100", —beq
141
          "011" \ \ when \ "000010" \ , \quad --andi
142
143
          "100" when "000011", —ori
          "ZZZ" when others;
144
145
      with inst_in select
        mem_write_o <= '0' when "100011", --lw
146
          '1' when "101011", —sw
147
          '0' when "000100", —beq
148
          '0' when "000010", —andi
149
          '0' when "000011", —ori
150
          'Z' when others;
151
```

```
with inst_in select
152
153
        mem_to_reg_o \le '0' \text{ when "100011"}, --lw
          '1' when "101011", —sw
154
          '1' when "000100", —beq
155
          '1' when "000010", —andi
156
          '1' when "000011", —ori
157
          'Z' when others;
158
159
     -- select from intermediate signals
160
     with inst_in select
161
        branch <= branch_f when "000000",
162
          branch_o when others;
163
164
     with inst_in select
        reg_dest \le reg_dest_f when "000000",
165
          reg_dest_o when others;
166
167
     with inst_in select
168
       reg_write <= reg_write_f when "000000",
          reg_write_o when others;
169
170
     with inst_in select
        ALU_src \le ALU_src_f when "000000",
171
172
          ALU_src_o when others;
     with inst_in select
173
       ALU_{op} \le ALU_{op_f} when "000000",
174
          ALU_op_o when others;
175
     with inst_in select
176
177
       mem_write <= mem_write_f when "000000",
          mem_write_o when others;
178
179
     with inst_in select
180
        mem_to_reg <= mem_to_reg_f when "000000",
181
          mem_to_reg_o when others;
182 end behavioral;
183
184 ----
185
186 library ieee;
187 use ieee.std_logic_1164.all;
188 use ieee.numeric_std.all;
189
190 entity dataMem is
     port (
```

```
192
     data : out std_logic_vector(31 downto 0);
193
     sel : in std_logic_vector(31 downto 0);
     wData : in std_logic_vector(31 downto 0);
194
195
     memWrite : in std_logic;
196
     clk : in std_logic);
197 end dataMem;
198
199 architecture behavioral of dataMem is
200 type mem_arr is array(0 to 255) of std_logic_vector(31
      downto 0);
201 signal mData : mem_arr;
202 begin
203
     data <= mData(to_integer(resize(unsigned(sel),8)));
204
     wrProc: process(clk) is
205
206
     begin
       if falling_edge(clk) then
207
        if (memWrite = '1') then
208
209
              mData(to_integer(resize(unsigned(sel),8))) <=
      wData;
210
       end if;
211
       end if;
212
     end process;
213 end behavioral;
214
215
216
217 ---
218 library ieee;
219 use ieee.std_logic_1164.all;
220 use ieee.numeric_std.all;
221
222 entity ALU is
223
       port (
224
       inA : in std_logic_vector(31 downto 0);
225
       inB : in std_logic_vector(31 downto 0);
226
       ctl: in std_logic_vector(2 downto 0);
227
       res : out std_logic_vector(31 downto 0));
228 end ALU;
229
```

```
230 architecture behavioral of ALU is
231 signal add : std_logic_vector(31 downto 0);
232 signal sub : std_logic_vector(31 downto 0);
233 signal andres: std_logic_vector(31 downto 0);
234 signal nandres : std_logic_vector(31 downto 0);
235 signal orres : std_logic_vector(31 downto 0);
236
       begin
237
       add <= std_logic_vector(signed(inA)+signed(inB));
238
       sub <= std_logic_vector(signed(inA)-signed(inB));
239
       andres <= std_logic_vector(unsigned(inA) and
      unsigned (inB));
240
       nandres <= std_logic_vector(not(unsigned(inA) and
      unsigned (inB)));
241
       orres <= std_logic_vector(unsigned(inA) or unsigned(
      inB));
242
243
     -- Multiplexer
244
     with ctl select
245
       res \ll add when "000",
         sub when "001",
246
          nandres when "010",
247
          andres when "011",
248
249
          orres when "100",
         "00000000000000000000000000000000" when others;
250
251 end behavioral;
252
253 --
254
255 library ieee;
256 use ieee.std_logic_1164.all;
257 use ieee.numeric_std.all;
258 entity processor is
259
       port (
260
       extPC : in std_logic_vector(31 downto 0);
261
       IMdata: in std_logic_vector(31 downto 0);
       DMdata: in std_logic_vector(31 downto 0);
262
263
       IMwrite : in std_logic;
264
       DMwrite: in std_logic;
       DMaddr: in std_logic_vector(31 downto 0);
265
266
       stall: in std_logic;
```

```
267
       clk: in std_logic
268);
269 end processor;
270
271 architecture behavioral of processor is
272 signal im_wrEn, im_clk : std_logic;
273 signal im_data,im_addr,im_wData : std_logic_vector(31
      downto 0);
274 signal dm_wrEn, dm_clk : std_logic;
275 signal dm_data, dm_addr, dm_wData : std_logic_vector(31
      downto 0);
276 signal PC: std_logic_vector(31 downto 0);
277 signal regA, regB, wData: std_logic_vector(31 downto 0);
278 signal selA, selB, selW: std_logic_vector(3 downto 0);
279 signal aluCtl : std_logic_vector(2 downto 0);
280 signal regWrite, regDest, regClk, dm_write, aluSrc,
      memtoreg : std_logic;
281 signal aluA, aluB, aluRes : std_logic_vector(31 downto 0);
282 signal branch, branchI, zero : std_logic := '0';
283 signal braAddr : std_logic_vector(15 downto 0);
284 signal op_code, func : std_logic_vector(5 downto 0);
285
286 begin
287
     IM : entity work.dataMem port map(im_data,im_addr,
      im_wData, im_wrEn, im_clk);
     DM: entity work.dataMem port map(dm_data,dm_addr,
288
      dm_wData, dm_wrEn, dm_clk);
289
       RF: entity work.regFile port map(regA, regB, selA,
      selB, wData, regWrite, selW, regClk);
       ALU: entity work.ALU port map(aluA, aluB, aluCtl,
290
      aluRes):
291
     CTRL: entity work.control port map(op_code, func,
      stall, branchI, regDest, regWrite, aluSrc, aluCtl,
      dm_write, memtoreg);
292
293
       —all clocks synced
294
       im_clk \ll clk;
295
       dm_{clk} \ll clk;
       regClk <= clk;
296
297
```

```
298
       im_wData <= IMData:
       im_wrEn <= IMWrite;
299
       -- allow testbench to initialize
300
301
        process (clk)
302
        begin
303
        if (rising_edge(clk)) then
            if(stall = '1') then
304
                PC \le extPC;
305
            elsif(branch = '1') then
306
307
                PC \le std_logic_vector(unsigned(PC) + (
       unsigned (resize (signed (braAddr), 32)));
308
            else
309
                PC <= std_logic_vector(unsigned(PC) + x"1");
310
            end if;
311
        end if;
312
        end process;
313
        braAddr <= im_data(15 downto 0);
314
        im_addr \ll PC;
315
        aluA \ll regA;
316
        with aluSrc
            select aluB <= regB when '0',
317
            std_logic_vector(unsigned(resize(signed(im_data
318
       (15 downto 0)), 32))) when '1',
            x"00000000" when others;
319
320
        with regDest
            select selW <= im_data(19 downto 16) when '0',
321
322
            im_data(14 downto 11) when '1',
323
            "ZZZZ" when others;
        with stall
324
325
            select dm_addr <= DMaddr when '1',
326
            aluRes when others;
327
        with stall
            select dm_wData <= DMdata when '1',
328
            regB when others;
329
330
        with stall
            select dm_wrEn <= DMWrite when '1',
331
332
        dm_write when others;
333
334
        with memtoreg
            select wData <= aluRes when '1',
335
```

```
336
            dm_data when others;
337
338
        with aluRes
339
            select zero <= '1' when x"00000000",
            '0' when others;
340
341
        branch <= branchI and zero;</pre>
342
343
344
        op_code <= im_data(31 downto 26);
345
        func \le im_data(5 downto 0);
346
347
        selA \le im_data(24 downto 21);
348
        selB \le im_data(19 downto 16);
349 end behavioral;
                    Listing 2: Test Bench Code
 1 library ieee;
 2 use ieee.std_logic_1164.all;
 3 use ieee.numeric_std.all;
 5 entity test_bench is
 6 end test_bench;
 8 architecture behavioral of test_bench is
 9
      signal clk : std_logic;
      signal extPC, IMdata, DMdata, DMaddr : std_logic_vector
 10
       (31 \text{ downto } 0) := x"00000000";
11
      signal IMwrite, DMwrite, stall : std_logic := '0';
 12
 13 begin
 14
     cpu: entity work.processor port map(extPC, IMdata,
      DMdata, IMwrite, DMwrite, DMaddr, stall, clk);
15
 16
     -- clk process
 17
     clkgen: process
 18
     begin
 19
        clk <= '1';
 20
        wait for 1 ns;
 21
        clk \ll 0;
 22
        wait for 1 ns;
```

```
23
    end process;
24
25
    tester: process
26
    begin
      -- init values
27
       stall \ll '1';
28
29
            IMwrite <= '0';
            DMwrite \leq '1';
30
      — put some data into the DM
31
      DMdata \le x"AAAAAAA";
32
      DMwrite <= '1';
33
      DMaddr \le x"00000001";
34
35
      wait for 2 ns;
      DMdata <= x"BBBBBBBB";
36
37
      DMwrite \ll '1';
      DMaddr \le x"00000002";
38
      wait for 2 ns;
39
      DMdata \le x"CCCCCCCC";
40
      DMwrite <= '1';
41
      DMaddr \le x"00000003";
42
43
      wait for 2 ns;
44
      DMdata <= x"FFFFFFFF;;
      DMwrite <= '1';
45
      DMaddr \le x"00000004";
46
      wait for 2 ns;
47
48
      DMdata \le x"00000000";
      DMwrite <= '1';
49
      DMaddr \le x"00000005";
50
51
      wait for 2 ns;
52
      -- Now load program, start from address 1
53
      DMwrite <= '0';
54
      IMwrite <= '1';
55
      -- lw $1, 1($zero)
56
      extPC \le x"00000001";
57
      58
      wait for 2 ns;
59
      -- sw $1, 6($zero)
60
61
      extPC \le x"00000002";
      IMdata <= b"10101100000000010000000000000110";
62
```

```
wait for 2 ns;
63
64
      — lw $2, 2($zero)
       extPC \le x"00000003";
65
       66
       wait for 2 ns;
67
      — add $3, $1, $2
68
       extPC \le x"00000004";
69
70
       IMdata <= b"000000000010001100001100000100000";
       wait for 2 ns;
71
72
      — sub $4, $2, $1
       extPC \le x"00000005";
73
74
       IMdata <= b"00000000010000010010000000110000";
75
       wait for 2 ns;
      -- beg $1, $2, 100
76
77
       extPC \le x"00000006";
       IMdata <= b"000100000010001000000001100100";
78
79
       wait for 2 ns;
80
81
      -- lw $2, 4($zero)
       extPC \le x"00000007";
82
83
       84
       wait for 2 ns;
      — nand $5, $1, $2
85
       extPC \le x"00000008";
86
       IMdata <= b"00000000001000100010100000000001";
87
88
       wait for 2 ns;
      — andi $6, $2, 00FF
89
90
       extPC \le x"00000009";
91
       IMdata <= b"0000100001000110000000001111111111";
92
       wait for 2 ns;
      — ori $7, $1, 00FF
93
       extPC \le x"0000000A";
94
       IMdata <= b"0000110000100111000000001111111111";
95
       wait for 2 ns;
96
      -- or $8, $1, $2
97
       extPC \le x"0000000B";
98
      IMdata <= b"000000000010001001000000000000010";
99
       wait for 2 ns;
100
101
      -- beq $1, $1 -0x000B
       extPC \le x"0000000C";
102
```

```
103
      wait for 2 ns;
104
105
106 — Begin execution here
      wait for 2 ns;
107
      IMwrite <= '0';
108
      extPC \le x"00000000";
109
      wait for 2 ns;
110
111
      stall \ll 0;
112
113 — allow enough time for processor to execute
     instructions
      wait for 100 ns;
114
115
116
    end process;
117 end behavioral;
```