Project 3: Designing a 32-bit CPU

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1 Introduction

This goal of this project is to design a stripped down version of the MIPS processor. The processor will be a 32-bit version of the processor discussed in class and the text book, however, its instruction set will be a small subset of the MIPS processor's full capability.

1.1 Background Information

1.1.1 MIPS

MIPS is a reduced instruction set computer (RISC) instruction set architecture (ISA). It defines three types of instruction types: R (register), I (Immediate), and J (Jump). For the implementation that this project is focused on, only R and I instructions will be executed. R type instructions are the most common form of instructions. The format for an r-type instruction is:

Bits[31:26]	Bits[25:21]	Bits[20:16]	Bits[15:11]	Bits[10:6]	Bits[5:0]
opcode	Rs	Rt	Rd	shamt	funct

For this instruction, the opcode field is always 000000_2 , while the function code funct is used to determine which instruction is to be carried out. Rs and Rt are the two registers in which the operation reads and Rd is the destination of the result. Some instructions require a shift amount (shamt), so it is specified explicitly.

The I type instruction involves an immediate value, so the instruction format must accommodate this. The format of this type of instruction is:

Bits[31:26]	Bits[25:21]	Bits[20:16]	Bits[15:0]
opcode	Rs	Rt	immediate

For this instruction, the op code field is used to define the specific instruction, Rs is the register in which the operation acts on along with the immediate value as the other operand. Rt is the destination register in which the result is stored.

1.1.2 Datapath and Control

A datapath is a collection of functional units that perform data processing operations. It includes units such as a program counter, a register file, instruction memory, an ALU, data memory, and a control unit. Figure 1 shows a high level overview of a simple datapath with control.

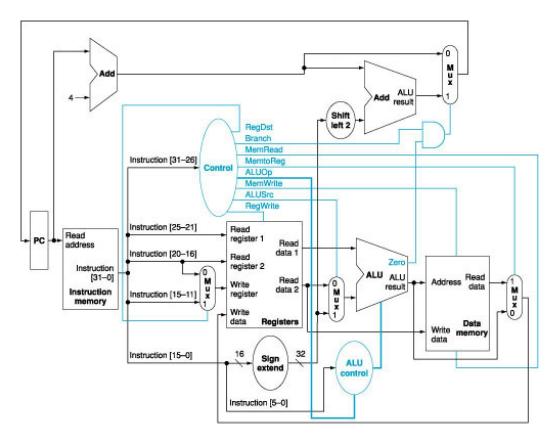


Figure 1: Datapath Overview

2 Design

2.1 Instruction Set

Table 1 shows the instructions that were chosen to be implemented in the CPU with the respective OpCode and Function Field for each instruction.

OpCode[31:26]	Function Field [5:0]	Instruction	Example Operation
100011_2		lw	lw \$t3, 200(\$s2)
101011_2		SW	sw \$t4, 100(\$t3)
000000_2	100000_2	add	add \$s3, \$t2, \$s2
000000_2	110000_2	sub	sub \$s3, \$t2, \$s2
000100_2		beq	beq \$s5, \$s2, 500
000000_2	000001_2	nand	nand \$5, \$1, \$2
000010_2		andi	andi \$6, \$2, 00FF
000000_2	000010_2	or	or \$8, \$1, \$2
000011_2		ori	ori \$7, \$1, 00FF

Table 1: CPU Instruction Set

Because it was only required to implement 9 instructions and the MIPS instruction set format requires 6 bits for op code and function field, it was an easy decision to choose these values for the implemented instructions. For all R-type instructions, the functions fields were chosen to be vastly different from one another to make debugging easier for the team. Likewise, the same approach was taken for the op code decisions for the I-type instructions.

2.2 Memory

For this project, it seemed unnecessary to implement memory of 4GB (2^{32}) . It was chosen to use an array of 256 words instead. If need be, this memory size could be upgraded easily, so this choice does not hinder performance on the actual design of the CPU.

2.3 Datapath

Because of the simplicity of this design, the implemented datapath did not need to be modified by much from Figure 1. Therefore, the design of a single cycle datapath from the textbook acted as the skeletal structure of the final implementation. Because an ALU and Register file were previously implemented in earlier projects, it was necessary to extend their functionality to be able to handle 32-bit words. Once this was complete, this left the data memory entity to be completed so that it could be included in the processor entity. As mentioned earlier in Section 2.2, this entity contains an array of

256 words, and allows for reading and writing.

The processor entity combines all of the components into the desired datapath. It synchronizes the clock of the instruction memory, data memory, and register file so that the entire system is in sync with an external clock signal. The program counter is updated during the rising edge of the clock, and all writes happen on the falling edge. The processor relies on the control unit to carry out the instruction read from memory. Figure 2 shows the overview block diagram of the implemented datapath for the CPU.

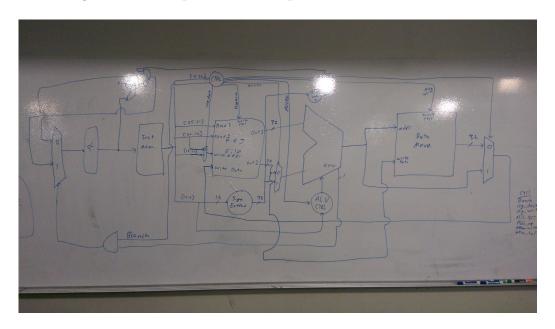


Figure 2: Implemented Data Path Overview

2.4 Control

The control lines can also be seen in Figure 2. It is a simple design of several signals acting as the sel line of a series of multiplexers. Based on the op codes and function field read from the instruction memory, the signals are asserted accordingly to relay the correct signals into the Register File, ALU, and Data memory. This unit is what determines which units will read/write, and what operations the ALU should perform.

3 Analysis

While this processor was optimized to be able to fully accomplish the tasks specified in the business requirements document, it could still be improved. In its current stage, it can be considered a bare bones prototype. To transform the current design into a processor on par with the current industry standard, a complete instruction set would have to be implemented. Furthermore, pipelining is a necessity to add. Any processor that doesn't implement pipelining is not making efficient use of its own components. After pipelining is implemented, hazard controls would need to coexist. This would allow for cool features of the processor to exist such as forwarding, making it a truly efficient piece of hardware.

4 Simulation Results

Once the processor was completely designed, it was necessary to write some test bench code to ensure the functionality it provided was desired. To test each instruction, data had to first be written to memory, along with the program being loaded onto the CPU. Due to the amount of signals involved in the CPU, not all will be shown in the simulation. The clock, contents of the registers, data memory, and program counter will only be shown. Data Memory addresses $00000001_{16} \rightarrow 00000005_{16}$ were initialized with starting data. Please refer to the Test Bench in the Appendix for a detailed view of the testing procedure. The execution of the program begins at 38ns. The tested instructions are:

- 1. lw \$1, 1(\$zero)
- 2. sw \$1, 6(\$zero)
- 3. lw \$2, 2(\$zero)
- 4. add \$3, \$1, \$2
- 5. sub \$4, \$2, \$1
- 6. beq \$1, \$2, 100
- 7. lw \$2, 4(\$zero)

- 8. nand \$5, \$1, \$2
- 9. andi \$6, \$2, 0x00FF
- 10. ori \$7, \$1, 0x00FF
- 11. or \$8, \$1, \$2
- 12. beq \$1, \$1, -0x000B

Figure 3 shows the first instruction being executed. The data memory at address 0x00000001 houses the data holds the value 0xAAAAAA and register \$1 is loaded with the data 0xAAAAAAA.

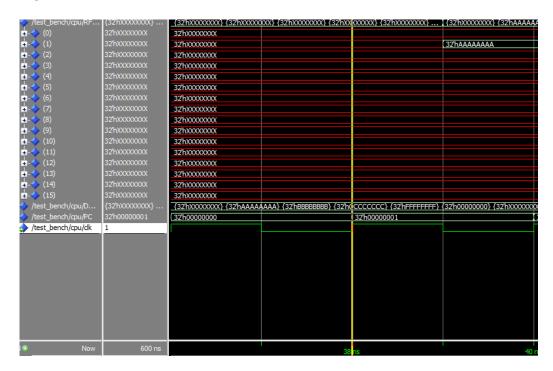


Figure 3: Instruction 1

5 Conclusion

The design and implementation of a 32-bit CPU was a success. a set of 9 instructions were successfully implemented and verified with test bench code.

All requested functionality was achieved. This 32-bit CPU can now be used in further projects.

Appendix

Listing 1: CPU Code

```
1 library ieee;
2 use ieee.std_logic_1164.all;
3 use ieee.numeric_std.all;
4
5 entity regFile is
    port (
    regA : out std_logic_vector(31 downto 0);
    regB : out std_logic_vector(31 downto 0);
9
    selA : in std_logic_vector(3 downto 0);
    selB : in std_logic_vector(3 downto 0);
10
11
    wData : in std_logic_vector(31 downto 0);
12
    registerWrite : in std_logic;
13
    selW : in std_logic_vector(3 downto 0);
    clk : in std_logic);
14
15 end regFile;
17 architecture behavioral of regFile is
18 type reg_arr is array (0 to 15) of std_logic_vector (31
     downto 0);
19 signal rData : reg_arr;
20 begin
21
    with selA
      select regA \le x"00000000" when b"0000",
22
      rData(to_integer(unsigned(selA))) when others;
23
24
    with selB
      select regB \le x"00000000" when b"0000",
25
26
      rData(to_integer(unsigned(selB))) when others;
27
    wrProc: process(clk) is
28
29
    begin
      if falling_edge(clk) then
30
      if (registerWrite = '1') then
31
```

```
32
        rData(to_integer(unsigned(selW))) <= wData;
33
      end if;
34
      end if;
35
    end process;
36 end behavioral;
37
38 ----
39
40 library ieee;
41 use ieee.std_logic_1164.all;
42 use ieee.numeric_std.all;
43
44 entity control is
45
    port (
46
       inst_in : in std_logic_vector(5 downto 0);
47
      func : in std_logic_vector(5 downto 0);
48
       stall : in std_logic;
      branch : out std_logic;
49
50
      reg_dest : out std_logic;
      reg_write : out std_logic;
51
52
      ALU_src : out std_logic;
      ALU_op : out std_logic_vector(2 downto 0);
53
      mem_write : out std_logic;
54
55
      mem_to_reg : out std_logic
56
    );
57 end control;
58
59 architecture behavioral of control is
     signal branch_o, reg_dest_o, reg_write_o, ALU_src_o,
     mem_write_o, mem_to_reg_o : std_logic;
     signal ALU_op_o : std_logic_vector(2 downto 0);
61
     signal branch_f , reg_dest_f , reg_write_f , ALU_src_f ,
62
     mem_write_f, mem_to_reg_f : std_logic;
     signal ALU_op_f : std_logic_vector(2 downto 0);
63
64 begin
65
    -- set intermediate signals incase of r-type
     instruction
    with func select
66
      branch_f <= '0' when "100000", --add
67
         '0' when "110000", —sub
68
```

```
'0' when "000001", —nand
69
 70
          '0' when "000010", —or
 71
         '0' when others;
72
     with func select
       reg_dest_f \le '1' \text{ when "100000"}, --add
73
         '1' when "110000", —sub
74
         '1' when "000001", --nand
 75
         '1' when "000010",
76
 77
         'Z' when others;
 78
     with func select
79
       reg_write_f <= '1' when "100000", --add
         '1' when "110000", —sub
80
         '1' when "000001", —nand
81
82
         '1' when "000010", —or
83
         'Z' when others;
84
     with func select
85
       ALU_src_f \le '0' \text{ when "100000", } --add
         ^{\prime}0 ' when "110000", —sub
86
         '0' when "000001", —nand
87
         '0' when "000010", —or
88
89
         'Z' when others;
90
     with func select
       91
         "001" when "110000", —sub
92
         "010" when "000001", —nand
93
         "100" when "000010", —or
94
         "ZZZ" when others;
95
96
     with func select
97
       mem_write_f <= '0' when "100000", --add
         '0' when "110000", —sub
98
         '0' when "000001", —nand
99
          '0' when "000010", --or
100
101
         'Z' when others;
     with func select
102
       mem\_to\_reg\_f \le '1' when "100000", --add
103
         '1' when "110000", —sub
104
         '1' when "000001", —nand
105
         '1' when "000010", —or
106
         'Z' when others;
107
108
```

```
109
     -- set intermediate signals incase of non r-type
       instruction
      with inst_in select
110
111
        branch_o <= '0' when "100011", --lw
           '0' when "101011", —sw
112
          '1' when "000100", —beq
113
          '0' when "000010", —andi
114
          '0' when "000011",
115
                               --- o r i
           '0' when others;
116
      with inst_in select
117
118
        reg_dest_o \ll 0' \text{ when "100011"}, --lw
          o, when "101011", —sw
119
          ^{\prime}0^{\,\prime} when ^{\prime\prime}000100^{\prime\prime} , ^{\,} —beq
120
           '0' when "000010", —andi
121
           '0' when "000011",
122
                               —- o r i
123
           'Z' when others;
124
      with inst_in select
        reg_write_o <= '1' when "100011", --lw
125
           '0' when "101011", —sw
126
127
          '0' when "000100", —beq
          '1' when "000010", —andi
128
          '1' when "000011",
129
          'Z' when others;
130
131
      with inst_in select
        ALU_{src_o} \ll '1' \text{ when "} 100011", --lw
132
           '1' when "101011", --sw
133
          '0' when "000100", —beq
134
          '1' when "000010", --andi
135
          '1' when "000011",
136
                               --- o r i
          'Z' when others;
137
      with inst_in select
138
        ALU_{op_o} \le "000" \text{ when } "100011", --lw
139
          "000" when "101011", —sw
140
          "001" when "000100",
                                 ---beq
141
          "011" when "000010", --andi
142
          "100" when "000011",
143
          "ZZZ" when others;
144
145
      with inst_in select
        mem_write_o <= '0' when "100011", --lw
146
           '1' when "101011", --sw
147
```

```
'0' when "000100", —beq
148
           ^{\prime}0^{\,\prime} when ^{\prime\prime}000010^{\prime\prime} , \, —andi
149
           '0' when "000011", --ori
150
           'Z' when others;
151
      with inst_in select
152
153
        mem_to_reg_o \ll 0' when "100011", --lw
           '1' when "101011", --sw
154
          '1' when "000100",
                               --beg
155
          '1' when "000010", —andi
156
           '1' when "000011", —ori
157
           'Z' when others;
158
159
160
     -- select from intermediate signals
      with inst_in select
161
162
        branch <= branch_f when "000000",
163
          branch_o when others;
164
      with inst_in select
        reg_dest \le reg_dest_f \text{ when "000000"},
165
166
          reg_dest_o when others;
      with inst_in select
167
        reg_write <= reg_write_f when "000000",
168
          reg_write_o when others;
169
      with inst_in select
170
        ALU_src \le ALU_src_f \text{ when "000000"},
171
172
          ALU_src_o when others;
173
      with inst_in select
        ALU_{op} \leftarrow ALU_{op_f} \text{ when "000000"},
174
175
          ALU_op_o when others;
176
      with inst_in select
177
        mem_write <= mem_write_f when "000000",
          mem_write_o when others;
178
      with inst_in select
179
180
        mem_to_reg \le mem_to_reg_f when "000000",
          mem_to_reg_o when others;
181
182 end behavioral;
183
184 -
185
186 library ieee;
187 use ieee.std_logic_1164.all;
```

```
188 use ieee.numeric_std.all;
190 entity dataMem is
191
     port (
192
     data : out std_logic_vector(31 downto 0);
193
     sel: in std_logic_vector(31 downto 0);
194
     wData: in std_logic_vector(31 downto 0);
195
     memWrite : in std_logic;
196
     clk : in std_logic);
197 end dataMem;
198
199 architecture behavioral of dataMem is
200 type mem_arr is array (0 to 255) of std_logic_vector (31
       downto 0);
201 signal mData : mem_arr;
202 begin
     data <= mData(to_integer(resize(unsigned(sel),8)));
203
204
205
     wrProc: process(clk) is
206
     begin
207
       if falling_edge(clk) then
208
        if (memWrite = '1') then
209
              mData(to_integer(resize(unsigned(sel),8))) <=
      wData;
210
       end if;
211
       end if;
212
     end process;
213 end behavioral;
214
215
216
217 ---
218 library ieee;
219 use ieee.std_logic_1164.all;
220 use ieee.numeric_std.all;
221
222 entity ALU is
223
        port (
224
       inA : in std_logic_vector(31 downto 0);
       inB : in std_logic_vector(31 downto 0);
225
```

```
226
        ctl: in std_logic_vector(2 downto 0);
227
       res : out std_logic_vector(31 downto 0));
228 end ALU;
229
230 architecture behavioral of ALU is
231 signal add : std_logic_vector(31 downto 0);
232 signal sub : std_logic_vector(31 downto 0);
233 signal andres: std_logic_vector(31 downto 0);
234 signal nandres : std_logic_vector(31 downto 0);
235 signal orres : std_logic_vector(31 downto 0);
236
       begin
237
       add <= std_logic_vector(signed(inA)+signed(inB));
238
       sub <= std_logic_vector(signed(inA)-signed(inB));
239
       andres <= std_logic_vector(unsigned(inA) and
       unsigned (inB));
240
       nandres <= std_logic_vector(not(unsigned(inA) and
      unsigned (inB)));
241
       orres <= std_logic_vector(unsigned(inA) or unsigned(
      inB));
242
     -- Multiplexer
243
244
     with ctl select
245
       res \ll add when "000",
         sub when "001",
246
          nandres when "010",
247
          andres when "011",
248
          orres when "100",
249
         "00000000000000000000000000000000" when others;
250
251 end behavioral;
252
253 - -
254
255 library ieee;
256 use ieee.std_logic_1164.all;
257 use ieee.numeric_std.all;
258 entity processor is
259
       port (
260
       extPC: in std_logic_vector(31 downto 0);
       IMdata: in std_logic_vector(31 downto 0);
261
       DMdata: in std_logic_vector(31 downto 0);
262
```

```
263
       IMwrite: in std_logic;
       DMwrite : in std_logic;
264
       DMaddr: in std_logic_vector(31 downto 0);
265
266
       stall : in std_logic;
267
       clk : in std_logic
268);
269 end processor;
270
271 architecture behavioral of processor is
272 signal im_wrEn, im_clk : std_logic;
273 signal im_data,im_addr,im_wData : std_logic_vector(31
      downto 0);
274 signal dm_wrEn, dm_clk : std_logic;
275 signal dm_data, dm_addr, dm_wData : std_logic_vector(31
      downto 0);
276 signal PC: std_logic_vector(31 downto 0);
277 signal regA, regB, wData: std_logic_vector(31 downto 0);
278 signal selA, selB, selW: std_logic_vector(3 downto 0);
279 signal aluCtl : std_logic_vector(2 downto 0);
280 signal regWrite, regDest, regClk, dm_write, aluSrc,
      memtoreg : std_logic;
281 signal aluA, aluB, aluRes : std_logic_vector(31 downto 0);
282 signal branch, branchI, zero : std_logic := '0';
283 signal braAddr : std_logic_vector(15 downto 0);
284 signal op_code, func : std_logic_vector(5 downto 0);
285
286 begin
     IM : entity work.dataMem port map(im_data,im_addr,
287
      im_wData, im_wrEn, im_clk);
288
     DM: entity work.dataMem port map(dm_data,dm_addr,
      dm_wData, dm_wrEn, dm_clk);
289
       RF: entity work.regFile port map(regA, regB, selA,
      selB, wData, regWrite, selW, regClk);
290
       ALU: entity work.ALU port map(aluA, aluB, aluCtl,
      aluRes);
291
     CTRL: entity work.control port map(op_code, func,
      stall, branchI, regDest, regWrite, aluSrc, aluCtl,
      dm_write, memtoreg);
292
293
       —all clocks synced
```

```
294
        im_clk \ll clk;
        dm_{clk} \ll clk;
295
        regClk <= clk;
296
297
298
       im_wData <= IMData;
299
       im_wrEn <= IMWrite;
       -- allow testbench to initialize
300
301
        process (clk)
302
        begin
303
        if (rising_edge(clk)) then
304
            if(stall = '1') then
                PC \le extPC;
305
            elsif(branch = '1') then
306
                PC <= std_logic_vector(unsigned(PC) + (
307
       unsigned (resize (signed (braAddr), 32)));
308
            else
                PC <= std_logic_vector(unsigned(PC) + x"1");
309
            end if;
310
311
        end if;
312
        end process;
        braAddr <= im_data(15 downto 0);
313
        im_addr \ll PC;
314
        aluA \le regA;
315
316
        with aluSrc
317
            select aluB <= regB when '0',
            std_logic_vector(unsigned(resize(signed(im_data
318
       (15 downto 0)), 32))) when '1',
319
            x"00000000" when others;
320
        with regDest
321
            select selW <= im_data(19 downto 16) when '0',
322
            im_data(14 downto 11) when '1',
323
            "ZZZZ" when others;
324
        with stall
            select dm_addr <= DMaddr when '1',
325
326
            aluRes when others;
327
        with stall
328
            select dm_wData <= DMdata when '1',
            regB when others;
329
330
        with stall
            select dm_wrEn <= DMWrite when '1',
331
```

```
332
        dm_write when others;
333
334
        with memtoreg
335
            select wData <= aluRes when '1',
            dm_data when others;
336
337
338
        with aluRes
            select zero \ll '1' when x"00000000",
339
            '0' when others;
340
341
342
        branch <= branchI and zero;</pre>
343
344
        op_code <= im_data(31 downto 26);
        func <= im_data(5 downto 0);
345
346
        selA <= im_data(24 downto 21);
347
348
        selB \le im_data(19 downto 16);
349 end behavioral;
                    Listing 2: Test Bench Code
 1 library ieee;
 2 use ieee.std_logic_1164.all;
 3 use ieee.numeric_std.all;
 4
 5 entity test_bench is
 6 end test_bench;
 7
 8 architecture behavioral of test_bench is
      signal clk : std_logic;
 9
10
      signal extPC, IMdata, DMdata, DMaddr : std_logic_vector
      (31 \text{ downto } 0) := x"00000000";
11
12
      signal IMwrite, DMwrite, stall : std_logic := '0';
13 begin
14
     cpu : entity work.processor port map(extPC, IMdata,
      DMdata, IMwrite, DMwrite, DMaddr, stall, clk);
15
16
     -- clk process
     clkgen: process
17
18
     begin
```

```
19
      clk <= '1';
20
      wait for 1 ns;
21
      clk <= '0';
22
      wait for 1 ns;
23
    end process;
24
25
    tester: process
    begin
26
      -- init values
27
28
      stall \ll '1';
29
            IMwrite \ll '0';
            DMwrite \leq '1';
30
31
      — put some data into the DM
      DMdata \le x"AAAAAAAA";
32
33
      DMwrite \leq '1';
      DMaddr \le x"00000001";
34
35
      wait for 2 ns;
      36
37
      DMwrite <= '1';
      DMaddr \le x"00000002";
38
39
      wait for 2 ns;
      DMdata \le x"CCCCCCCC";
40
      DMwrite <= '1';
41
42
      DMaddr \le x"00000003";
      wait for 2 ns;
43
44
      DMdata \le x"FFFFFFFF";
      DMwrite <= '1';
45
46
      DMaddr \le x"00000004";
47
      wait for 2 ns;
      DMdata \le x"000000000";
48
      DMwrite <= '1';
49
      DMaddr \le x"00000005";
50
      wait for 2 ns;
51
52
      -- Now load program, start from address 1
53
      DMwrite \leq '0';
54
      IMwrite <= '1';
55
      -- lw $1, 1($zero)
56
57
      extPC \le x"00000001";
      58
```

```
wait for 2 ns;
59
      -- sw $1, 6($zero)
60
      extPC \le x"00000002";
61
      IMdata <= b"10101100000000010000000000000110";
62
      wait for 2 ns;
63
      -- lw $2, 2($zero)
64
      extPC \le x"00000003";
65
66
      wait for 2 ns;
67
      — add $3, $1, $2
68
      extPC \le x"00000004";
69
70
      IMdata <= b"000000000010001100001100000100000";
71
      wait for 2 ns;
      — sub $4, $2, $1
72
73
      extPC \le x"00000005";
      IMdata <= b"0000000001000001001000000110000";
74
      wait for 2 ns;
75
      — beg $1, $2, 100
76
77
      extPC \le x"00000006";
      IMdata <= b"0001000000100010000000001100100";
78
79
      wait for 2 ns;
80
      -- lw $2, 4($zero)
81
      extPC \le x"00000007";
82
      83
84
      wait for 2 ns;
      — nand $5, $1, $2
85
86
      extPC \le x"00000008";
      IMdata <= b"00000000001000100010100000000001";
87
88
      wait for 2 ns;
89
      — andi $6, $2, 00FF
      extPC \le x"00000009";
90
      IMdata <= b"000010000100011000000001111111111";
91
      wait for 2 ns;
92
      — ori $7, $1, 00FF
93
      extPC \le x"0000000A";
94
      IMdata <= b"0000110000100111000000001111111111";
95
      wait for 2 ns;
96
97
      — or $8, $1, $2
      extPC \le x"0000000B";
98
```

```
IMdata <= b"00000000001000100100000000000010";
99
      wait for 2 ns;
100
      -- beq $1, $1 -0x000B
101
      extPC <= x"0000000C";
102
      103
      wait for 2 ns;
104
105
106 — Begin execution here
      wait for 2 ns;
107
      IMwrite <= '0';
108
      extPC \le x"00000000";
109
      wait for 2 ns;
110
111
       stall \ll '0';
112
113 — allow enough time for processor to execute
      instructions
      wait for 100 ns;
114
115
116
     end process;
117 end behavioral;
```