

# COMP3121 21T2 Assignment 3 Q1

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In order to count the number of ways to place brackets in the expression, so that the evaluation will return true, two sub-problems need to be considered. Assuming there is total  $n$  symbols and  $n - 1$  operations in the expression.

1. Number of ways to place bracket from  $l_{th}$  symbol to  $r_{th}$  symbol, resulting the expression as true.
2. Number of ways to place bracket from  $l_{th}$  symbol to  $r_{th}$  symbol, resulting the expression as false.

Hence  $T(l_{th}, r_{th})$  can be defined as number of ways to place bracket from  $l_{th}$  symbol to  $r_{th}$  symbol, resulting the expression as true, for example, *true* AND *false* OR *true* will result  $T(1, 2)$  to  $1(\text{true}) \times 0(\text{false}) = 0$  true expression by evaluating *true* AND *false*. Therefore the problem can be solved by utilising the strategy of dividing large problem into smaller sets until it cannot be smaller. In the previous example, *true* AND *false* OR *true* can be divided into (*true* AND *false*) OR (*true*) by separating OR operator, then evaluate the smaller problem sets on each side, and eventually combine all the expressions. In general, this strategy can be standardised as:

$$T(l, r) = \sum_{m=1}^{r-1} TSplit(l, m, r) \quad (1)$$

$$T(l, r) = \sum_{m=1}^{r-1} FSplit(l, m, r) \quad (2)$$

Split functions perform differently based on different operators, the summations are concluded below:

Table 1: The relationship between operator and TSplit function

Opeator\Split function	$TSplit(l, m, r)$
AND	$T(l, m) \times T(m + 1, r)$
OR	$T(l, m) \times F(m + 1, r) + T(l, m) \times T(m + 1, r) + F(1, m) \times T(m + 1, r)$
NAND	$T(l, m) \times F(m + 1, r) + F(l, m) \times F(m + 1, r) + F(1, m) \times T(m + 1, r)$
NOR	$F(l, m) \times F(m + 1, r)$

Table 2: The relationship between operator and FSplit function

Opeator\Split function	$FSplit(l, m, r)$
AND	$T(l, m) \times F(m + 1, r) + F(l, m) \times F(m + 1, r) + F(1, m) \times T(m + 1, r)$
OR	$F(l, m) \times F(m + 1, r)$
NAND	$T(l, m) \times T(m + 1, r)$
NOR	$T(l, m) \times F(m + 1, r) + T(l, m) \times T(m + 1, r) + F(1, m) \times T(m + 1, r)$

Overall, the complexity is  $O(n^3)$ , since there are  $O(n^2)$  ranges as well as the situation  $O(n)$  in Tsplit and Fsplit need to be covered.