

PRACTICAL No. 4

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Topic: Parsing

Platform: Windows or Linux

Language to be used: Python or Java (based on the companies targeted for placement)

Aim: (A) Write a program to validate a natural language sentence. Design a natural language grammar, compute and input the LL (1) table. Validate if the given sentence is valid or not based on the grammar.

Input: NLP grammar and LL (1) parsing table (from file)

Implementation: String parsing rules

Output: Each step-in string parsing and whether the input string is valid or invalid.

CODE :

```
import pandas as pd

def util(ll1):
    tab = pd.DataFrame(
        ll1,
        columns=[
            "championship",
            "ball",
            "toss",
            "is",
            "want",
            "won",
            "Played",
            "me",
            "I",
            "you",
            "India",
            "Australia",
            "Steve",
            "John",
            "the",
            "a",
            "an",
        ],
    )
    tab["Nonterm"] = ["S", "NP", "VP", "N", "V", "P", "PN", "D"]
    tab.set_index("Nonterm", inplace=True)
    return tab

def validator(input):
```

```
from beautifultable import BeautifulTable
```

[illegible]

```

        "Australia",
        "Steve",
        "John",
        "",
        "",
        "",
    ],
    ["", "", "", "", "", "", "", "", "", "", "", "", "", "", "the",
"a", "an"],
]
tab = util(ll1)
table = BeautifulTable()
table.column_headers = ["Buffer", "Stack"]
buffer = input.split(" ")
buffer.reverse()
stack = ["S"]
table.append_row([buffer.copy(), stack.copy()])
while buffer != [] and stack != []:
    index = stack.pop(0)
    key = buffer[-1]
    if key not in tab.columns:
        print("Invalid input")
        table.append_row([buffer.copy(), stack.copy()])
        print(table)
        return

    rule = tab.loc[index][key].split(" ")

    if "" in rule:
        print("Invalid input")
        table.append_row([buffer.copy(), stack.copy()])
        print(table)
        return

    stack = rule + stack
    table.append_row([buffer.copy(), stack.copy()])

    if key in rule:
        buffer.remove(key)
        stack.remove(key)
        table.append_row([buffer.copy(), stack.copy()])
print(table)
print("Valid input")

input = "India won the championship"
validator(input)
input = "championship India won"
validator(input)

```

OUTPUT:

| Buffer | Stack |
|---|-------|
| ['championship', 'the', 'won', 'India'] | ['S'] |

| | | |
|---|------------------|--|
| +-----+-----+ | | |
| ['championship', 'the', 'won', 'India'] | ['NP', 'VP'] | |
| +-----+-----+ | | |
| ['championship', 'the', 'won', 'India'] | ['PN', 'VP'] | |
| +-----+-----+ | | |
| ['championship', 'the', 'won', 'India'] | ['India', 'VP'] | |
| +-----+-----+ | | |
| ['championship', 'the', 'won'] | ['VP'] | |
| +-----+-----+ | | |
| ['championship', 'the', 'won'] | ['V', 'NP'] | |
| +-----+-----+ | | |
| ['championship', 'the', 'won'] | ['won', 'NP'] | |
| +-----+-----+ | | |
| ['championship', 'the'] | ['NP'] | |
| +-----+-----+ | | |
| ['championship', 'the'] | ['D', 'N'] | |
| +-----+-----+ | | |
| ['championship', 'the'] | ['the', 'N'] | |
| +-----+-----+ | | |
| ['championship'] | ['N'] | |
| +-----+-----+ | | |
| ['championship'] | ['championship'] | |
| +-----+-----+ | | |
| [] | [] | |
| +-----+-----+ | | |

Valid input

Invalid input

| | | |
|----------------------------------|-------|--|
| +-----+-----+ | | |
| Buffer | Stack | |
| +-----+-----+ | | |
| ['won', 'India', 'championship'] | ['S'] | |
| +-----+-----+ | | |
| ['won', 'India', 'championship'] | [] | |
| +-----+-----+ | | |

(B) Use Virtual Lab on LL1 parser to validate the string and verify your string validation using simulation.

Link for Virtual Lab: http://vlabs.iitb.ac.in/vlabs-dev/vlab_bootcamp/bootcamp/system_deligators/labs/exp2/index.php

Output: Validation from Virtual lab simulator

Details:

PART A:

- Construct and consider a natural language grammar that can validate an English sentence.
- Solve the NLP grammar by hand for LL(1) parser and create parsing table
- Input the above parsing table and grammar using a file.
- Write program for performing string validation

PART B:

- Go to Virtual lab: Go through all the tabs, paste screen shots for all steps (including tests), validate your string parsing with the simulator (screen shot expected).

← → ↻ Not secure | vlabs.iitb.ac.in/vlabs-dev/vlab_bootcamp/bootcamp/system_deligators/labs/exp2/pretest.php

IoT Based Fire Alar... Control Lights with... https://www.astaga... (DOC) Fall of Indian... Heart Disease with...

☐ code generation

4) $E = E + T$ is correct grammar to parse the string without back tracing?

☐ Yes

☒ No

5) Grammar is ambiguous if _____

☐ It has 2 parse trees+

☐ It has two left most derivations

☐ It has two right most derivations

☒ all of the above

6) How to generate parse table?

☒ Using First and Follow

☐ Using RSA algo

7) The concept of grammar is much used in this part of the compiler

☐ lexical analysis

☒ parser

☐ code generation

☐ code optimization

8) Examine the given input string is syntactically correct or not id ** id

☐ Correct

☒ Incorrect

Evaluate

1) Correct

2) Correct

3) Correct

4) Correct

5) Correct

6) Correct

7) Correct

8) Correct

Lab contributed by KS School of Bussiness Management, Gujarat University

← → ↻ Not secure | vlabs.iitb.ac.in/vlabs-dev/vlab_bootcamp/bootcamp/system_deligators/labs/exp2/posttest.php

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Simulation

Post Test

References

1) Once you land on the simulator page what is the default condition grammar for LL(1) parser?

☒ Left recursive eliminated grammar

☐ Left recursive grammar

☐ Right recursive grammar

☐ Any grammar

2) What result do you expect when top down parsing will perform using right recursive grammar?

☐ infinite loop

☒ Top down parsing with back tracking

☐ Parse the given input

☐ Syntax analysis

3) What importance of parser table to parse the given input string?

☐ Top down parsing with back tracking

☐ Bottom up parsing

☐ A formal grammar that contains left recursion cannot be parsed by a LL(k) parser

☒ Automatic Predictive top down parsing of the input string without backtracking

4) What do you observe when left recursive grammar is used without rewriting?

☐ A formal grammar that contains left recursion can be parsed by a LL(k)-parser

☐ Top down parsing with back tracking

☒ It leads them into infinite recursion

☐ Bottom up parsing

5) What is the use of first and follow algorithm?

☐ LL1 parsing of the given input string

☐ To generate left recursive eliminated grammar

☐ To generate left recursive grammar

☒ predictive parse table construction for LL1 parsing

6) Generate LL(1) parsing for the grammar given in Task 2 for given input string (LD4) (LD5) id + id ** (id + id)

id + * id + (id * id)

7) analyze

id*+id

Evaluate

1) Correct

2) Correct

3) Correct

4) Correct

5) Correct

LL(1) Parser Visualization

Write your own context-free grammar and see an LL(1) parser in action!

Written by Zak Kincaid and Shaowei Zhu based on
<http://jgmachines.sourceforge.net/machines/ll1.html>

1. Write your LL(1) grammar (empty string "" represents ε):

```
S :-> NP VP
NP :-> P
NP :-> PN
NP :-> D N
VP :-> V NP
N :-> championship
N :-> ball
N :-> toss
V :-> is
V :-> want
V :-> won
V :-> played
P :-> me
P :-> I
P :-> you
PN :-> India
PN :-> Australia
PN :-> Steve
PN :-> John
D :-> the
D :-> a
D :-> an
```

Generate tables

Valid LL(1) Grammars

For any production $S \rightarrow A \mid B$, it must be the case that:

- For no terminal t could A and B derive strings beginning with t
- At most one of A and B can derive the empty string
- If B can derive the empty string, then A does not derive any string beginning with a terminal in $\text{Follow}(A)$

Formatting Instructions

- The non-terminal on the left-hand-side of the first rule is the start non-terminal
- Write each production rule in a separate line (see example to the left)
- Separate each token using whitespace
- $\$$ is reserved as the end-of-input symbol, and S is reserved as an artificial start symbol. The grammar is automatically augmented with the rule $S \rightarrow \text{start } S$

Debugging

- More information about the parser construction is printed on the console
- The source code follows the pseudocode in lecture. In particular, see `computeFollow`, `computeFirst`, `computeFollow`, and `computeL1Tables`

2. Nullable/First/Follow Table and Transition Table

| Nonterminal | Nullable? | First | Follow |
|-------------|-----------|---|---|
| S | X | me, I, you, India, Australia, Steve, John, the, a, an | \$ |
| NP | X | me, I, you, India, Australia, Steve, John, the, a, an | is, want, won, played, \$ |
| VP | X | is, want, won, played | \$ |
| N | X | championship, ball, toss | is, want, won, played, \$ |
| V | X | is, want, won, played | me, I, you, India, Australia, Steve, John, the, a, an |
| P | X | me, I, you | is, want, won, played, \$ |
| PN | X | India, Australia, Steve, John | is, want, won, played, \$ |
| D | X | the, a, an | championship, ball, toss |

| | \$ | championship | ball | toss | is | want | won | played | me | I | you | India | Australia | Steve | John | the | a | an |
|----|----|-------------------------------------|-----------------------------|-----------------------------|---------------------------|-----------------------------|----------------------------|-------------------------------|---|---|---|---|---|---|---|---|---|---|
| S | | | | | | | | | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ | $S \rightarrow NP VP$ $S \rightarrow \$$ |
| NP | | | | | | | | | $NP \rightarrow P$ | $NP \rightarrow P$ | $NP \rightarrow P$ | $NP \rightarrow PN$ | $NP \rightarrow PN$ | $NP \rightarrow PN$ | $NP \rightarrow PN$ | $NP \rightarrow D N$ | $NP \rightarrow D N$ | $NP \rightarrow D N$ |
| VP | | | | | $VP \rightarrow V NP$ | $VP \rightarrow V NP$ | $VP \rightarrow V NP$ | $VP \rightarrow V NP$ | | | | | | | | | | |
| N | | $N \rightarrow \text{championship}$ | $N \rightarrow \text{ball}$ | $N \rightarrow \text{toss}$ | | | | | | | | | | | | | | |
| V | | | | | $V \rightarrow \text{is}$ | $V \rightarrow \text{want}$ | $V \rightarrow \text{won}$ | $V \rightarrow \text{played}$ | | | | | | | | | | |
| P | | | | | | | | | $P \rightarrow \text{me}$ | $P \rightarrow \text{I}$ | $P \rightarrow \text{you}$ | | | | | | | |
| PN | | | | | | | | | | | | $PN \rightarrow \text{India}$ | $PN \rightarrow \text{Australia}$ | $PN \rightarrow \text{Steve}$ | $PN \rightarrow \text{John}$ | | | |
| D | | | | | | | | | | | | | | | | $D \rightarrow \text{the}$ | $D \rightarrow \text{a}$ | $D \rightarrow \text{an}$ |

3. Parsing

Token stream separated by spaces:

Start/Reset

Step Forward

Stack

\$

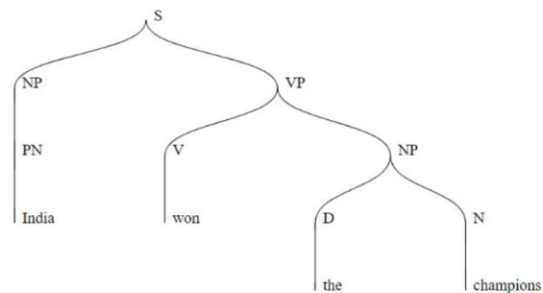
Remaining Input

\$

Rule

Match championship

Partial Parse Tree



Championship

ball

loss

is

want

won

Played

me

I

you

and

Am

Steve

John

the a an

S

NP

VP

N

V

P

PN

D

N → Chess

N → ball

N → loss

VP →
V NP

VP →
V NP

NP →
V NP

VP →
V NP

NP →

NP →

NP →

NP →

NP →

NP →

NP →

NP →

NP →

NP →

N → is V → want V → won V → played

P → me P → you

PN → and PN → Am

PN → Steve John

the a an