CS 211 RECITATIONS WEEK 9

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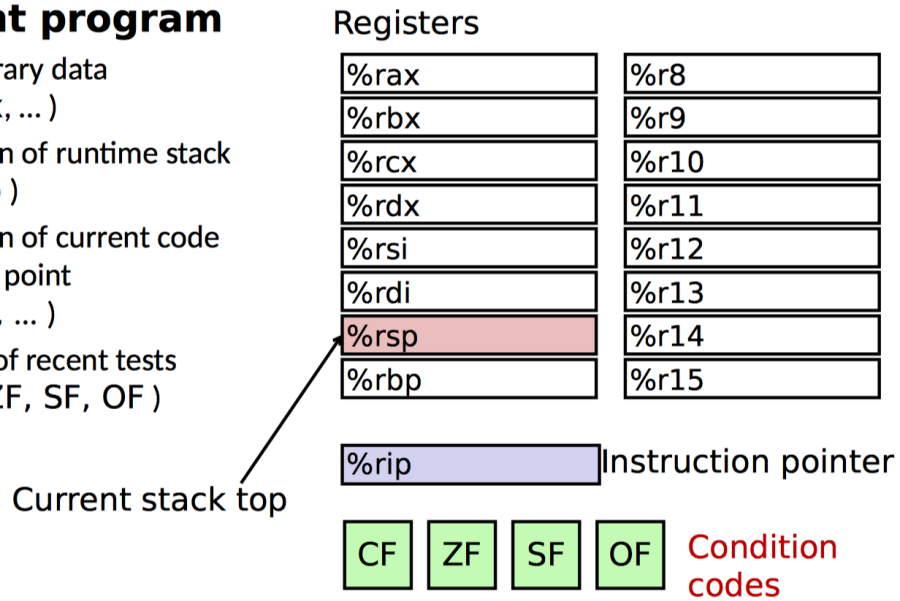
Content

- Condition codes & branches
- Loops
- Switch statements
- Functions
- Recursions

Processor State

- **Info about current program**
  - Temporary data (%rax, …)
  - Location of runtime stack (%rsp)
  - Location of current code control point (%rip, …)
  - Status of recent tests (CF, ZF, SF, OF)



Explicit Setting by Compare Instruction

- cmpq Src2, Src1
- cmpq b,a like computing a-b without setting destination

- CF set if carry out from most significant bit (used for unsigned comparisons)
- ZF set if a == b
- SF set if (a-b) < 0 (as signed)
- OF set if two's-complement (signed) overflow

  (a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)

Implicitly set (think of it as side effect) by arithmetic operations

Example: addq Src,Dest ↔ t = a+b

CF set if carry out from most significant bit (unsigned overflow)

ZF set if t == 0

SF set if t < 0 (as signed)

OF set if two's-complement (signed) overflow

(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)

Single bit registers

- CF Carry Flag (for unsigned)
- ZF Zero Flag
- SF Sign Flag (for signed)
- OF Overflow Flag (for signed)

Explicit Setting by Test instruction

- testq Src2, Src1
  - testq b,a like computing a&b without setting destination
- Sets condition codes based on value of Src1 & Src2
- Useful to have one of the operands be a mask
- ZF set when a&b == 0
- SF set when a&b < 0

SetX Instructions

- Set low-order byte of destination to 0 or 1 based on combinations of condition codes
- Does not alter remaining 7 bytes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~(SF^OF)&~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF)|ZF	Less or Equal (Signed)
seta	~CF&~ZF	Above (unsigned)
setb	CF	Below (unsigned)

- **SetX Instructions:**
  - Set single byte based on combination of condition codes
- **One of addressable byte registers**
  - Does not alter remaining bytes
  - Typically use movzbl to finish job

```
int gt (long x, long y)
{
  return x > y;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
cmpq   %rsi, %rdi    # Compare x:y
setg   %al           # Set when >
movzbl %al, %eax     # Zero rest of %rax
ret
```

jX Instructions

- Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF)&~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF)|ZF	Less or Equal (Signed)
ja	~CF&~ZF	Above (unsigned)
jb	CF	Below (unsigned)

Conditional Branch Example

- Generation

```
$ gcc -Og -S -fno-if-conversion control.c
```

```
long absdiff
  (long x, long y)
{
  long result;
  if (x > y)
    result = x-y;
  else
    result = y-x;
  return result;
}
```

```
absdiff:
   cmpq    %rsi, %rdi  # x:y
   jle     .L4
   movq    %rdi, %rax
   subq    %rsi, %rax
   ret
.L4:       # x <= y
   movq    %rsi, %rax
   subq    %rdi, %rax
   ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

“Do-While” Loop Example

C Code

```
long pcount_do
  (unsigned long x) {
  long result = 0;
  do {
    result += x & 0x1;
    x >>= 1;
  } while (x);
  return result;
}
```

“Do-While” Loop Compilation

Goto Version

```
long pcount_goto
  (unsigned long x) {
  long result = 0;
 loop:
  result += x & 0x1;
  x >>= 1;
  if(x) goto loop;
  return result;
}
```

Register	Use(s)
%rdi	Argument x
%rax	result

```
   movl    $0, %eax    #  result = 0
.L2:                   # loop:
   movq    %rdi, %rdx
   andl    $1, %edx    #  t = x & 0x1
   addq    %rdx, %rax  #  result += t
   shrq    %rdi        #  x >>= 1
   jne     .L2         #  if (x) goto loop
   ret
```

General “While” Translation #1

- “Jump-to-middle” translation
- Used with **-Og**

While version

```
while (Test)
  Body
```

Goto Version

```
  goto test;
loop:
  Body
test:
  if (Test)
    goto loop;
done:
```

While Loop Example #1

C Code

```
long pcount_while
  (unsigned long x) {
  long result = 0;
  while (x) {
    result += x & 0x1;
    x >>= 1;
  }
  return result;
}
```

Jump to

```
long pcount_goto_jtm
  (unsigned long x) {
  long result = 0;
  goto test;
 loop:
  result += x & 0x1;
  x >>= 1;
 test:
  if(x) goto loop;
  return result;
}
```

- **Compare to do-while version of function**
- **Initial goto starts loop at test**

“For” Loop Do-While Conversion

```
#define WSIZE 8*sizeof(int)
```

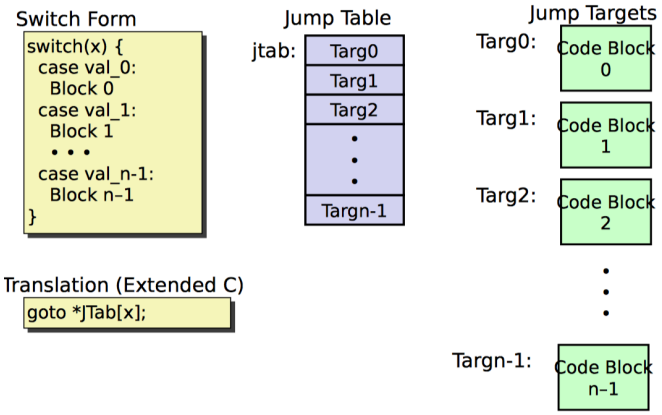
```
long pcount_for
  (unsigned long x)
{
  size_t i;
  long result = 0;
  for (i = 0; i < WSIZE; i++)
  {
    unsigned bit =
      (x >> i) & 0x1;
    result += bit;
  }
  return result;
}
```

- Initial test can be optimized away

```
long pcount_for_goto_dw
  (unsigned long x) {
  size_t i;
  long result = 0;
  i = 0;                          Init
  if (!(i < WSIZE))
    goto done;                    !Test
 loop:
  {
    unsigned bit =
      (x >> i) & 0x1;             Body
    result += bit;
  }
  i++;                            Update
  if (i < WSIZE)                  Test
    goto loop;
 done:
  return result;
}
```

```
long switch_eg
   (long x, long y, long z)
{
    long w = 1;
    switch(x) {
    case 1:
        w = y*z;
        break;
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
    case 5:
    case 6:
        w -= z;
        break;
    default:
        w = 2;
    }
    return w;
}
```

Jump Table Structure



Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
      . . .
    }
    return w;
}
```

Jump table

```
.section    .rodata
  .align 8
.L4:
  .quad     .L8 # x = 0
  .quad     .L3 # x = 1
  .quad     .L5 # x = 2
  .quad     .L9 # x = 3
  .quad     .L8 # x = 4
  .quad     .L7 # x = 5
  .quad     .L7 # x = 6
```

Setup:

```
    switch_eg:
        movq    %rdx, %rcx
        cmpq    $6, %rdi        # x:6
        ja      .L8             # Use default
        jmp     *.L4(,%rdi,8) # goto *JTab[x]
```

Indirect jump

Assembly Setup Explanation

- Table Structure
  - Each target requires 8 bytes
  - Base address at .L4

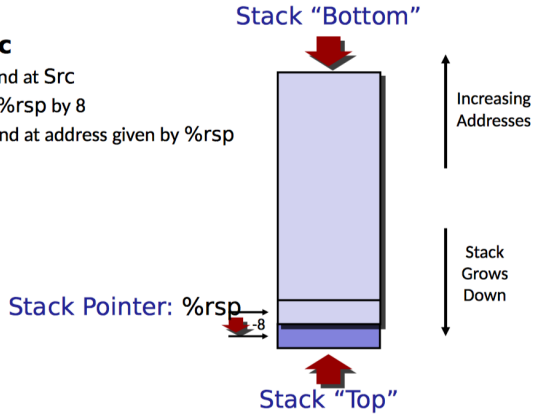
Jump table

```
.section    .rodata
  .align 8
.L4:
  .quad     .L8 # x = 0
  .quad     .L3 # x = 1
  .quad     .L5 # x = 2
  .quad     .L9 # x = 3
  .quad     .L8 # x = 4
  .quad     .L7 # x = 5
  .quad     .L7 # x = 6
```

- Jumping
  - Direct: `jmp .L8`
  - Jump target is denoted by label .L8
  - Indirect: `jmp *.L4(,%rdi,8)`
  - Start of jump table: .L4
  - Must scale by factor of 8 (addresses are 8 bytes)
  - Fetch target from effective Address .L4 + x*8
    - Only for $0 \le x \le 6$

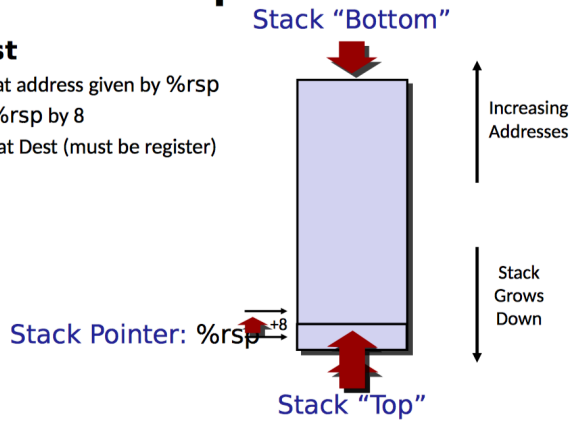
x86-64 Stack: Push

- **pushq Src**
  - Fetch operand at Src
  - Decrement %rsp by 8
  - Write operand at address given by %rsp



x86-64 Stack: Pop

- **popq Dest**
  - Read value at address given by %rsp
  - Increment %rsp by 8
  - Store value at Dest (must be register)



Procedure Control Flow

- **Use stack to support function call and return**
- **Procedure call:** `call label`
  - Push return address on stack
  - Jump to label
- **Return address:**
  - Address of the next instruction right after call
  - Example from disassembly
- **Procedure return:** `ret`
  - Pop address from stack
  - Jump to address

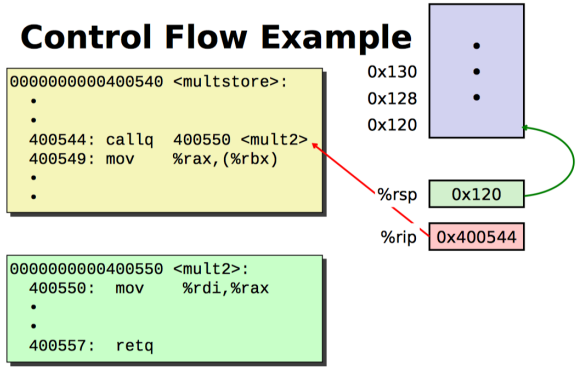
```
void multstore(long x,
   long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
0000000000400540 <multstore>:
  push   %rbx            # Save %rbx
  mov    %rdx,%rbx       # Save dest
  callq  400550 <mult2>  # mult2(x,y)
  mov    %rax,(%rbx)     # Save at dest
  pop    %rbx            # Restore %rbx
  retq                   # Return
```

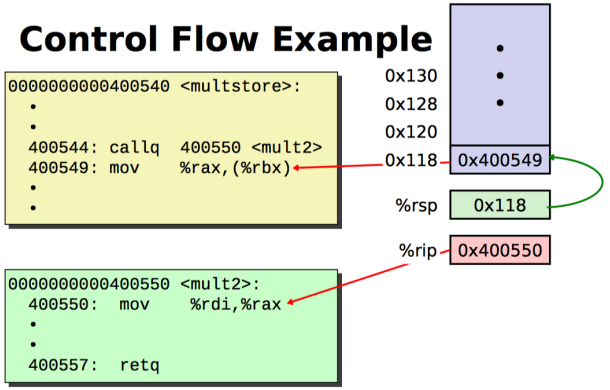
```
long mult2
  (long a, long b)
{
  long s = a * b;
  return s;
}
```

```
0000000000400550 <mult2>:
  mov    %rdi,%rax  # a
  imul   %rsi,%rax  # a * b
  retq              # Return
```

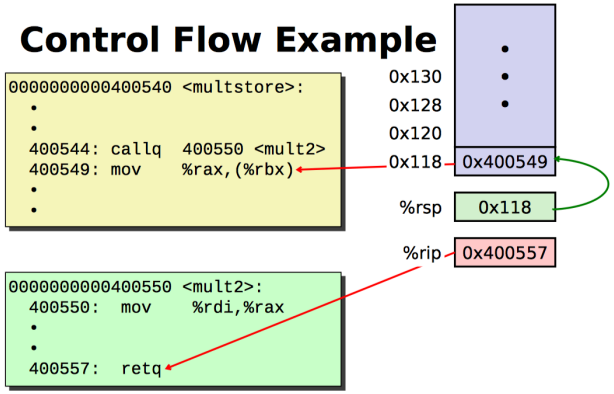
Control Flow Example



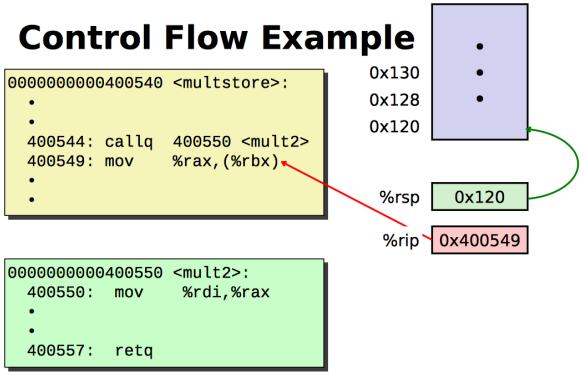
Control Flow Example



Control Flow Example



Control Flow Example



Procedure Data Flow

Registers

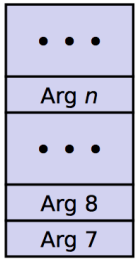
- **First 6 arguments**

%rdi
%rsi
%rdx
%rcx
%r8
%r9

- **Return value**

%rax

Stack



- **Only allocate stack space when needed**

Data Flow Examples

```
void multstore
 (long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

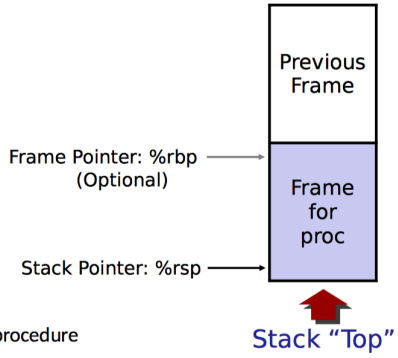
```
0000000000400540 <multstore>:
  # x in %rdi, y in %rsi, dest in %rdx
  • • •
  400541: mov    %rdx,%rbx        # Save dest
  400544: callq  400550 <mult2>   # mult2(x,y)
  # t in %rax
  400549: mov    %rax,(%rbx)      # Save at dest
  • • •
```

```
long mult2
  (long a, long b)
{
  long s = a * b;
  return s;
}
```

```
0000000000400550 <mult2>:
  # a in %rdi, b in %rsi
  400550:  mov    %rdi,%rax   # a
  400553:  imul   %rsi,%rax   # a * b
  # s in %rax
  400557:  retq               # Return
```

Stack Frames

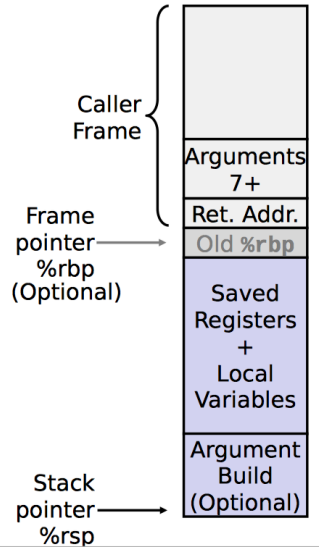
- **Contents**
  - Return information
  - Local storage (if needed)
  - Temporary space (if needed)



- **Management**
  - Space allocated when enter procedure
    - "Set-up" code
    - Includes push by **`call`** instruction
  - Deallocated when return
    - "Finish" code
    - Includes pop by **`ret`** instruction

- **Current Stack Frame ("Top" to Bottom)**
  - "Argument build:" Parameters for function about to call
  - Local variables If can't keep in registers
  - Saved register context
  - Old frame pointer (optional)

- **Caller Stack Frame**
  - Return address
    - Pushed by call instruction
  - Arguments for this call



Observations About Recursion

- **Handled Without Special Consideration**
  - Stack frames mean that each function call has private storage
    - Saved registers & local variables
    - Saved return pointer
  - Register saving conventions prevent one function call from corrupting another's data
    - Unless the C code explicitly does so (e.g., buffer overflow in Lecture 9)
  - Stack discipline follows call / return pattern
    - If P calls Q, then Q returns before P
    - Last-In, First-Out
- **Also works for mutual recursion**
  - P calls Q; Q calls P

x86-64 Procedure Summary

- **Important Points**
  - Stack is the right data structure for procedure call / return
    - If P calls Q, then Q returns before P
- **Recursion (& mutual recursion) handled by normal calling conventions**
  - Can safely store values in local stack frame and in callee-saved registers
  - Put function arguments at top of stack
  - Result return in %rax
- Pointers are addresses of values
  - On stack or global