Lecture 9: Optimization II

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Meaning of the Multiplier

$$\begin{aligned} \max_{\{x,y\}} & f(x,y) \\ \text{s.t} & g(x,y) = a \end{aligned}$$

We'll consider a to be a parameter which varies from problem to problem.

For any fixed value of a, write $(x^*(a), y^*(a))$ be a solution for the problem, and write $\lambda^*(a)$ for the multiplier which corresponds to the solution.

 $f(x^*(a), y^*(a))$ will be the corresponding optimal value of the objective function.

Meaning of the Multiplier

Theorem

Let f and g be C^1 functions of two variables. For any fixed value of the parameter a, let $(x^*(a), y^*(a))$ be a solution to the above problem with corresponding multiplier $\lambda^*(a)$. Suppose that x^* , y^* , and λ^* are all C^1 functions of a, and that the NDCQ holds. Then,

$$\lambda^*(a) = \frac{\mathrm{d}}{\mathrm{d}a} f(x^*(a), y^*(a))$$

The multiplier tells us how our maximized utility changes when we change the amount we're constrained by.

Envelope Theorems - Unconstrained

Theorem

Let $f(\mathbf{x}; a)$ be a C^1 funtion of $\mathbf{x} \in \mathbb{R}^n$ and a scalar a. For each choice of parameter a, consider the unconstrained maximisation problem

$$\max_{\mathbf{x}} \quad f(\mathbf{x}; a)$$

Let \mathbf{x}^* be a solution of this problem. Suppose that \mathbf{x}^* is a C^1 function of a. Then

$$\frac{\mathrm{d}}{\mathrm{d}a}f(\mathbf{x}^*(a);a) = \frac{\partial}{\partial a}f(\mathbf{x}^*(a);a)$$

The proof comes straight from the chain rule

$$\frac{\mathrm{d}}{\mathrm{d}a}f(\mathbf{x}^*(a);a) = \sum_{i} \frac{\partial f}{\partial x_i}(\mathbf{x}^*(a);a) \frac{\mathrm{d}x_i}{\mathrm{d}a}(a) + \frac{\partial f}{\partial a}(\mathbf{x}^*(a);a)$$

Envelope Theorems - Unconstrained

Exercise 9.1

A firm produces y microchips at a cost c(y), where c'(y) > 0, c''(y) < 0. Of the chips it produces, $1 - \alpha$ are defective and cannot be sold. The non-defective chips can be sold at price p. How will an increase in the production quality affect the firm's profit?

Envelope Theorems - Constrained

Theorem

Let $f, h_1, ..., h_k : \mathbb{R}^n \times \mathbb{R}^1 \to \mathbb{R}$ be C^1 functions. Let $\mathbf{x}^*(a) = (x_1^*(a), ..., x_n^*(a))$ denote a solution to the problem of maximising f on the constraint set

$$h_1(\mathbf{x};a) = 0, \dots, h_k(\mathbf{x},;) = 0$$

for any fixed choice of the parameter a. Suppose that $\mathbf{x}^*(a)$ and the Lagrange multipliers $\lambda_1(a), \ldots, \lambda_k(a)$ ar \mathcal{C}^1 functions of a and that the NDCQ holds. Then

$$\frac{\mathrm{d}}{\mathrm{d}a}f(\mathbf{x}^*(a);a) = \frac{\partial \mathcal{L}}{\partial a}(\mathbf{x}^a(a),\lambda(a);a)$$

where \mathcal{L} is the Lagrangian for the problem.

Envelope Theorems - Constrained

Exercise 9.2

- 1. Verify the meaning of the multiplier.
- 2. In the infinite horizon consumption-savings problem, what is the effect on optimum lifetime utility of a change in the interest rate r?

Inequality Constrained Optimization

Consider the problem with one inequality constraint

$$\max_{\{x_1, x_2\}} f(x_1, x_2)$$
s.t $g(x_1, x_2) \le h$

As before, it'll be necessary that at a local max/min, \mathbf{x}^*

$$\nabla f(\mathbf{x}^*) = \mu \nabla g(\mathbf{x}^*)$$

But we need to impose some further restrictions. We've still got constraints, but they're not as constraining as before.

Now we can not only attain an optimum on the boundary, but there's also the possibility of attaining an optimum on the interior of the constraint set.

Inequality Constrained Optimization

Previously, λ was a critical point of the Lagrangian. Taking the partial derivative with respect to λ and setting it to zero would result in the budget constraint being retrieved. However, if we were to structure the Lagrangian in the usual way

$$\mathcal{L} = f(\mathbf{x}) + \mu \left[h - g(\mathbf{x}) \right]$$

Taking this would yield exactly the same result.

$$\frac{\partial \mathcal{L}}{\partial \mu} = 0 \Rightarrow g(\mathbf{x}) = h$$

This isn't the result we want, because we wish to allow our constraint to be slack. This derivative requires that the constraint bind. So for inequality constraints, we replace this condition with two conditions.

Inequality Constrained Optimization

Non-negativity: $\mu > 0$

Complementary slackness: $\mu[h - g(\mathbf{x})^*] = 0$

Kuhn-Tucker Theorem

Theorem (Kuhn-Tucker)

Suppose that f, g_1, \ldots, g_k are C^1 functions of n variables. Suppose that $\mathbf{x}^* \in \mathbb{R}^n$ is a local maximiser of f on the constraint set defined by the k inequalities

$$g_1(\mathbf{x}) \le b_1, \dots, g_k(\mathbf{x}) \le b_k$$

Assume that at \mathbf{x}^* , the first k_0 constraints are binding at \mathbf{x}^* and the last $k-k_0$ constraints are slack. Suppose that the following NDCQ is satisfied at \mathbf{x}^*

The rank of the Jacobian matrix of the binding constraints

$$\begin{pmatrix} \frac{\partial g_1}{\partial x_1}(\mathbf{x}^*) & \dots & \frac{\partial g_1}{\partial x_n}(\mathbf{x}^*) \\ \vdots & \ddots & \vdots \\ \frac{\partial g_{k_0}}{\partial x_1}(\mathbf{x}^*) & \dots & \frac{\partial g_{k_0}}{\partial x_n}(\mathbf{x}^*) \end{pmatrix}$$

is k_0 .

Kuhn-Tucker Theorem

Theorem (Kuhn-Tucker)

From the Lagrangian

$$\mathcal{L} = f(\mathbf{x}) + \mu_1[b_1 - g_1(\mathbf{x})] + \dots + \mu_1[b_1 - g_1(\mathbf{x})]$$

there exist multipliers, μ_1^*, \ldots, μ_n^* , such that

(a)
$$\frac{\partial \mathcal{L}}{\partial x_1}(\mathbf{x}^*, \mu^*) = 0, \dots, \frac{\partial \mathcal{L}}{\partial x_1}(\mathbf{x}^*, \mu^*) = 0,$$

(b)
$$\mu_1^*[g_1(\mathbf{x}^*) - b_1] = 0, \dots, \mu_n^*[g_n(\mathbf{x}^*) - b_n] = 0$$

(c)
$$\mu_1^* \ge 0, \dots, \mu_n^* \ge 0$$

(d)
$$g_1(\mathbf{x}^*) \leq b_1, \dots, g_n(\mathbf{x}^*) \leq b_n$$

Exercises

Exercise 9.3 Bliss Point

An agent receives utility from consuming a good, c. The agent's utility function is characterised by $U(c) = -c^2 + 10c + 1$. The agent may consume up to their income y.

Set up and solve the agent's problem.

Exercises

Exercise 9.4 Consumption with Free Disposal of Wealth

An agent lives for two periods. The agent receives utility from consumption, with period utility satisfying U'(c) > 0, U''(c) < 0. Second period utility is discounted by a factor $\beta \in (0,1)$. At the beginning of time, the agent is endowed with wealth w, which they use to finance consumption in every period. The agent is not required to spend all of their income on consumption. The price of the consumption good in period 1 is equal to one unit of wealth. The price of the consumption good in the second period is p. The agent chooses consumption in each period to maximise lifetime utility. Solve the agents problem.

Exercises

Exercise 9.5 Inada Conditions

An agent lives for two periods. The agent receives utility from consumption, with the period utility function, U(c), being strictly increasing, strictly concave and satisfying $\lim_{c\to 0} U'(c) = \infty$, $\lim_{c\to\infty} U'(c) = 0$ (called the Inada conditions). Second period utility is discounted by a factor $\beta \in (0,1)$. At the beginning of time, the agent is endowed with wealth w, which they use to finance consumption in every period. The agent is not required to spend all of their income on consumption. The price of the consumption good in each period is equal to one unit of wealth. Consumption in each period must be non-negative. The agent chooses consumption in each period to maximise lifetime utility.

Set up and solve the agents problem.

Adding in Randomness

$$\max \mathbb{E}_0 \sum_{t=0}^{\infty} \beta^t U(c_t)$$
s.t. $c_t + k_{t+1} = A_t k_t^{\alpha} + (1 - \delta) k_t$

$$\ln A_t = \rho \ln A_{t-1} + \varepsilon_t$$

$$\varepsilon \sim N(0, \sigma^2)$$

$$\mathcal{L} = \mathbb{E}_0 \sum_{t=0}^{\infty} \left\{ \beta^t U(c_t) + \lambda_t [A_t k_t^{\alpha} + (1-\delta)k_t - c_t - k_{t+1}] \right\}$$

Adding in Randomness

A solution satisfying **rational expectations** uses the mathematical expectation conditioned on all information available to the agent at the time.

The usual Lagrange method gives the following characterization of a solution:

$$U'(c_t) = \beta \mathbb{E}_t [U'(c_{t+1})(A_{t+1}k_{t+1}^{\alpha} + 1 - \delta)]$$

$$c_t + k_{t+1} = A_t k_t^{\alpha} + (1 - \delta)k_t$$

We'll come back to this model in more detail after Dynamic Programming.

Learning Outcomes

You should be able to:

- Apply an envelope theorem.
- Solve an optimization problem using KKT conditions.