

# CSE 210A - HW 3 - Induction

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## Exercise 1

Given:

A:  $t := x; x := y; y := t$   
B:  $t := y; y := x; x := t$

Proof by counter example. To show that this assertion is false, find an example where this assertion does not hold. In other words, show an example where A and B are not the same. Choose

$x = 1$   
 $y = 2$   
 $t = 3$

Then,

A:  $t := x = 1; x := y = 2; y := t = 1$   
A:  $t = 1; x = 2; y = 1$   
B:  $t := y = 2; y := x = 1; x := t = 2$   
B:  $t = 2; x = 1; y = 2$   
A  $\neq$  B

## Exercise 2

If:

$\langle \text{while } b \text{ do } y := y - x, s \rangle \Downarrow s'$

then there exists  $k$  such that:

$$s(y) = s'(y) + k \cdot s(x)$$

Case: while-false

$$\frac{\langle b, s \rangle \Downarrow False}{\langle w, s \rangle \Downarrow s}$$

$$s' = s$$

$$k = 0 \text{ so } s(y) = s'(y) + 0$$

Therefore, inequality holds for the while-false case

Case: while-true

$$\frac{\langle b, s \rangle \Downarrow True \quad \langle y := s(y) - s(x), s \rangle \Downarrow s''}{\langle w, s \rangle \Downarrow s'}$$

This above tree is the initial execution of the while loop. It then calls this tree where the input state is now  $s''$

$$\frac{\langle b, s'' \rangle \Downarrow True \quad \langle y := s''(y) - s(x), s'' \rangle \Downarrow s'}{\langle w, s'' \rangle \Downarrow s'}$$

By our inductive hypothesis, we know that  $\langle w, s'' \rangle \Downarrow s'$  i.e.  $s''(y) = s'(y) + k * s''(x)$   
 Need to show:  $s(y) = s'(y) + k * s(x)$

$$\begin{aligned} \text{Have: } s'' &= s[y \mapsto s(y) - s(x)] \\ \text{Hence: } s(y) &= s(y) + k * s(x) \text{ for some } k \\ s(y) + k * s(x) &= s''(y) \text{ for } k = 1 \\ s(y) &= s''(x) \\ s(y) &= s''(x) + k * s(x) \end{aligned}$$

We have arrived back at our inductive hypothesis. We have shown that this is true for both the while-true and while-false cases. By showing it is true for all cases, we have proved that this is true.