

# Homework 8

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## Problem 13:

## Problem 14:

The goal of this problem to to find a solution to the polynomial

$$\left( \sum_{i=1}^n x_i v_i \right) \bmod n = L \bmod n$$

given a series of positive values for  $v$  and  $L$ . Also,  $x$  can be 0 or 1.

This problem is similar to the subset sum problem, except that values larger than  $L$  can not be immediately pruned. Instead, at a given level, if two nodes have the same value *modulo*  $n$ , one can be arbitrarily pruned.

Additionally, since the goal is to find a solution to the polynomial, a bit string can be created to consider the possibilities in which a value was included or discluded in the sum.

Enumerating all possibilities and then considering the above pruning rule gives way to the following algorithm.

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**Function:** *Modular Subset Sum*

**Input:** *Positive Integers  $v_1, \dots, v_n, L$*

**for**  $i$  *to*  $n$  **do**

**for**  $s$  *to*  $L \bmod n$  **do**

**if**  $A[k, s]$  *is defined* **then**

$/* ::$  *is the concatenation operator.*  $a :: c \rightarrow ac$   $*/$

$A[k + 1, s] = A[k, s] :: 0$

$A[k + 1, (s + v_{k+1} \bmod n) \bmod n] = A[k, s] :: 1$

**Return:**  $A[n, L \bmod n]$

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When this algorithm completes, the answer, if it exists, will be the bit string at  $A[n, L \bmod n]$ . This bit string starting, read from left to right, will represent the coefficient of each  $v$  in the polynomial.

**Problem 15:**

**Problem 16:**