

## Course 8

**Matrix of a list of vectors,  
matrix of a linear map**



Prof. dr. Septimiu Crivei

# Chapter 3. Matrices and Linear Systems

- 1 Elementary operations
- 2 Applications of elementary operations
- 3 The matrix of a list of vectors
- 4 The matrix of a linear map

# Applications: Hill cipher, image transformations

We present Hill cipher, which was the first application of linear algebra to cryptography, based on operations with invertible matrices.

We analyze image transformations: rotation of 2D-images with an angle around the origin.

## Definition

Let  $V$  be a vector space over  $K$ ,  $B = (v_1, \dots, v_n)$  a basis of  $V$  and  $X = (u_1, \dots, u_m)$  a list of vectors in  $V$ . Let

[illegible]

be the unique writings of the vectors in  $X$  as linear combinations of vectors of the basis  $B$ , for some  $a_{ij} \in K$ .

## Definition

The *matrix of the list of vectors*  $X$  in the basis  $B$  is the matrix having as its rows the coordinates of the vectors in  $X$  in the basis  $B$ , that is,

$$[X]_B = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix}.$$

# Example

Consider the canonical basis  $B = (e_1, e_2, e_3, e_4)$  and the list  $X = (u_1, u_2, u_3)$  in the canonical real vector space  $\mathbb{R}^4$ , where

$$\begin{cases} u_1 &= (1, 2, 3, 4) \\ u_2 &= (5, 6, 7, 8) \\ u_3 &= (9, 10, 11, 12) \end{cases}.$$

Since the coordinates of a vector in the canonical basis are just its components, we get

$$[X]_B = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 5 & 6 & 7 & 8 \\ 9 & 10 & 11 & 12 \end{pmatrix}.$$

## Theorem

*Let  $V$  be a vector space over  $K$ ,  $B = (v_1, \dots, v_n)$  a basis of  $V$  and  $X = (u_1, \dots, u_m)$  a list of vectors in  $V$  having the matrix  $A$  in the basis  $B$ . Then:*

- (i)  $\dim \langle X \rangle = \text{rank}(A)$ .*
- (ii) A basis of  $\langle X \rangle$  is the list of non-zero row-vectors  $(c_1, \dots, c_r)$  of an echelon form  $C$  equivalent to  $A$ .*

## Example

Let us determine the dimensions of the subspaces  $S$ ,  $T$ ,  $S + T$  and  $S \cap T$  of the canonical real vector space  $\mathbb{R}^4$ , where

$$S = \langle (-3, 5, -1, 1), (-1, 1, 0, 1), (1, 1, -1, -3) \rangle,$$

$$T = \langle (1, 0, 2, 0), (2, 1, -1, 2) \rangle.$$

We have  $\dim S = \dim T = 2$ . Furthermore,  $S + T = \langle S \cup T \rangle$ .

We write the matrix of  $S \cup T$  in the canonical basis and we have

$$\begin{pmatrix} -3 & 5 & -1 & 1 \\ -1 & 1 & 0 & 1 \\ 1 & 1 & -1 & -3 \\ 1 & 0 & 2 & 0 \\ 2 & 1 & -1 & 2 \end{pmatrix} \sim \dots \sim \begin{pmatrix} 1 & 1 & -1 & -3 \\ 0 & -1 & 3 & 3 \\ 0 & 0 & 5 & 4 \\ 0 & 0 & 0 & 33 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$$

Then  $\dim(S + T) = 4$  and a basis of  $S + T$  consists of the non-zero row-vectors from the echelon form, that is,

$((1, 1, -1, -3), (0, -1, 3, 3), (0, 0, 5, 4), (0, 0, 0, 33))$ . Finally,

$\dim(S \cap T) = \dim S + \dim T - \dim(S + T) = 2 + 2 - 4 = 0$ .



# Matrix of a vector

It is more convenient to define it as a column-matrix in order to avoid formulas involving transposes.

## Definition

Let  $V$  be a vector space over  $K$ ,  $v \in V$  and  $B = (v_1, \dots, v_n)$  a basis of  $V$ . If  $v = k_1 v_1 + \dots + k_n v_n$  ( $k_1, \dots, k_n \in K$ ) is the unique writing of  $v$  as a linear combination of the vectors of the basis  $B$ , then the *matrix* of  $v$  in the basis  $B$  is  $[v]_B = \begin{pmatrix} k_1 \\ \vdots \\ k_n \end{pmatrix}$ .

## Example

Consider the vector  $v = (1, 2, 3)$  in the canonical real vector space  $\mathbb{R}^3$ , and let  $E$  be the canonical basis. Then  $[v]_E = \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix}$ .

# Matrix of a linear map I

## Definition

Let  $f : V \rightarrow V'$  be a  $K$ -linear map,  $B = (v_1, \dots, v_n)$  a basis of  $V$  and  $B' = (v'_1, \dots, v'_m)$  a basis of  $V'$ . Then we can uniquely write the vectors in  $f(B)$  as linear combinations of the vectors of the basis  $B'$ , say

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for some  $a_{ij} \in K$ .

# Matrix of a linear map II

## Definition

Then the *matrix of the  $K$ -linear map  $f$*  in the bases  $B$  and  $B'$  is the matrix having as its columns the coordinates of the vectors of  $f(B)$  in the basis  $B'$ , that is,

$$[f]_{BB'} = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix}.$$

If  $V = V'$  and  $B = B'$ , then we simply denote  $[f]_B = [f]_{BB'}$ .

We have to emphasize that we put the coordinates on the columns of the matrix of a linear map and not on the rows as we did for the matrix of a list of vectors.

# Example

Consider the  $\mathbb{R}$ -linear map  $f : \mathbb{R}^4 \rightarrow \mathbb{R}^3$  defined by

$$f(x, y, z, t) = (x + y + z, y + z + t, z + t + x), \quad \forall (x, y, z, t) \in \mathbb{R}^4.$$

Let  $E = (e_1, e_2, e_3, e_4)$  and  $E' = (e'_1, e'_2, e'_3)$  be the canonical bases in  $\mathbb{R}^4$  and  $\mathbb{R}^3$  respectively. Since

$$\begin{cases} f(e_1) = f(1, 0, 0, 0) = (1, 0, 1) = e'_1 + e'_3 \\ f(e_2) = f(0, 1, 0, 0) = (1, 1, 0) = e'_1 + e'_2 \\ f(e_3) = f(0, 0, 1, 0) = (1, 1, 1) = e'_1 + e'_2 + e'_3 \\ f(e_4) = f(0, 0, 0, 1) = (0, 1, 1) = e'_2 + e'_3 \end{cases}$$

it follows that the matrix of  $f$  in the bases  $E$  and  $E'$  is

$$[f]_{EE'} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{pmatrix}.$$

# Property of the matrix of a linear map I

## Theorem

Let  $f : V \rightarrow V'$  be a  $K$ -linear map,  $B = (v_1, \dots, v_n)$  a basis of  $V$ ,  $B' = (v'_1, \dots, v'_m)$  a basis of  $V'$  and  $v \in V$ . Then

$$[f(v)]_{B'} = [f]_{BB'} \cdot [v]_B.$$

*Proof.* Let  $[f]_{BB'} = (a_{ij}) \in M_{m,n}(K)$ . Let

$$v = \sum_{j=1}^n k_j v_j,$$

$$f(v) = \sum_{i=1}^m k'_i v'_i$$

for some  $k_i, k'_i \in K$ .

# Property of the matrix of a linear map II

On the other hand, using the definition of the matrix of  $f$  in the bases  $B$  and  $B'$ , we have

$$\begin{aligned} f(v) &= f\left(\sum_{j=1}^n k_j v_j\right) = \sum_{j=1}^n k_j f(v_j) \\ &= \sum_{j=1}^n k_j \left(\sum_{i=1}^m a_{ij} v'_i\right) = \sum_{i=1}^m \left(\sum_{j=1}^n a_{ij} k_j\right) v'_i. \end{aligned}$$

But the writing of  $f(v)$  as a linear combination of the vectors of the basis  $B'$  is unique, hence we must have  $k'_i = \sum_{j=1}^n a_{ij} k_j$  for every  $i \in \{1, \dots, m\}$ . Therefore,  $[f(v)]_{B'} = [f]_{BB'} \cdot [v]_B$ .  $\square$

# Rank of a linear map

## Theorem

*Let  $f : V \rightarrow V'$  be a  $K$ -linear map. Then  $\text{rank}(f) = \text{rank}([f]_{BB'})$ , where  $B$  and  $B'$  are any bases of  $V$  and  $V'$  respectively.*

*Proof.* Let  $B = (v_1, \dots, v_n)$  and  $[f]_{BB'} = A$ . We have:

$$\begin{aligned}\text{rank}(f) &= \dim(\text{Im}f) = \dim f(V) = \dim f(\langle v_1, \dots, v_n \rangle) \\ &= \dim \langle f(v_1), \dots, f(v_n) \rangle = \text{rank}(A^T) = \text{rank}(A) = \text{rank}([f]_{BB'})\end{aligned}$$

Now take some other bases  $B_1 = (u_1, \dots, u_n)$  of  $V$  and  $B'_1$  of  $V'$  and denote  $[f]_{B_1 B'_1} = A_1$ . Then

$$\begin{aligned}\text{rank}([f]_{B_1 B'_1}) &= \text{rank}(A_1) = \text{rank}(A_1^T) = \dim \langle f(u_1), \dots, f(u_n) \rangle \\ &= \dim(\text{Im}f) = \dim \langle f(v_1), \dots, f(v_n) \rangle = \text{rank}([f]_{BB'}) .\end{aligned}$$

# Example

Consider the  $\mathbb{R}$ -linear map  $f : \mathbb{R}^4 \rightarrow \mathbb{R}^3$  defined by

$$f(x, y, z, t) = (x + y + z, y + z + t, z + t + x), \quad \forall (x, y, z, t) \in \mathbb{R}^4.$$

Let  $E = (e_1, e_2, e_3, e_4)$  and  $E' = (e'_1, e'_2, e'_3)$  be the canonical bases in  $\mathbb{R}^4$  and  $\mathbb{R}^3$  respectively. By a previous example it follows that

$$[f]_{EE'} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{pmatrix} \sim \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & -1 & 0 & 1 \end{pmatrix} \sim \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 2 \end{pmatrix}.$$

Now it follows that  $\text{rank}(f) = \text{rank}([f]_{EE'}) = 3$ .



# Connection between linear maps and matrices

## Theorem

Let  $V$ ,  $V'$  and  $V''$  be vector spaces over  $K$  with  $\dim V = n$ ,  $\dim V' = m$  and  $\dim V'' = p$  and let  $B = (v_1, \dots, v_n)$ ,  $B' = (v'_1, \dots, v'_m)$  and  $B'' = (v''_1, \dots, v''_p)$  be bases of  $V$ ,  $V'$  and  $V''$  respectively. Then  $\forall f, g \in \text{Hom}_K(V, V')$ ,  $\forall h \in \text{Hom}_K(V', V'')$  and  $\forall k \in K$ , we have

$$[f + g]_{BB'} = [f]_{BB'} + [g]_{BB'} ,$$

$$[kf]_{BB'} = k \cdot [f]_{BB'} ,$$

$$[h \circ f]_{BB''} = [h]_{B'B''} \cdot [f]_{BB'} .$$

*Proof.* Let  $[f]_{BB'} = (a_{ij}) \in M_{m,n}(K)$ ,  $[g]_{BB'} = (b_{ij}) \in M_{m,n}(K)$  and  $[h]_{B'B''} = (c_{ki}) \in M_{pm}(K)$ . Then

$$f(v_j) = \sum_{i=1}^m a_{ij} v'_i, \quad g(v_j) = \sum_{i=1}^m b_{ij} v'_i, \quad h(v'_i) = \sum_{k=1}^p c_{ki} v''_k$$

$\forall j \in \{1, \dots, n\}$  and  $\forall i \in \{1, \dots, m\}$  [...].

## Theorem

*Let  $V$  and  $V'$  be vector spaces over  $K$  with  $\dim V = n$  and  $\dim V' = m$ , and let  $B$  and  $B'$  be bases of  $V$  and  $V'$  respectively. Then the map*

$$\varphi : \operatorname{Hom}_K(V, V') \rightarrow M_{m,n}(K), \quad \varphi(f) = [f]_{BB'}, \quad \forall f \in \operatorname{Hom}_K(V, V')$$

*is an isomorphism of vector spaces.*

*Proof. [...]*

This isomorphism allows us to work with matrices instead of linear maps, which is much easier from a computational point of view.

Under this isomorphism, the kernel and the image of a linear map  $f : V \rightarrow V'$ , where  $V$  and  $V'$  are vector spaces over  $K$  with  $\dim(V) = n$  and  $\dim(V') = m$ , and bases  $B$  and  $B'$  respectively, correspond to the so-called *null space* and to the *column space* of its associated matrix  $A = [f]_{BB'} \in M_{m,n}(K)$  respectively.

Thus, the *null space* of  $A$  consists of vectors  $x \in K^n$  such that  $Ax = 0$ , while the *column space* of  $A$  consists of all linear combinations of the columns of  $A$ .

A vector  $b \in K^m$  belongs to the column space of  $A$  if and only if the system  $Ax = b$  has a solution.

By the First Dimension Theorem it follows that the sum of the dimensions of the null space and the column space of  $A$  equals  $n$ .

## Theorem

Let  $V$  be a vector space over  $K$  with  $\dim V = n$ , and let  $B$  be a basis of  $V$ . Then the map

$$\varphi : \text{End}_K(V) \rightarrow M_n(K), \quad \varphi(f) = [f]_B, \quad \forall f \in \text{End}_K(V)$$

is an isomorphism of vector spaces and of rings.

*Proof.* Note that  $(\text{End}_K(V), +, \circ)$  and  $(M_n(K), +, \cdot)$  are rings. The required isomorphisms follow by the above theorem.  $\square$

## Corollary

Let  $f \in \text{End}_K(V)$ . Then  $f \in \text{Aut}_K(V) \iff \det([f]_B) \neq 0$ , where  $B$  is any basis of  $V$ .

*Proof.*  $f \in \text{Aut}_K(V) \iff f$  is invertible in  $(\text{End}_K(V), +, \circ) \iff [f]_B$  is invertible in  $(M_n(K), +, \cdot) \iff \det([f]_B) \neq 0$ .  $\square$

## Extra: Hill cipher I

Let  $n \in \mathbb{N}^*$  and consider the canonical vector space  $V = \mathbb{Z}_2^n$  over  $\mathbb{Z}_2$  with canonical basis  $E$ . The vectors of  $V$  may be identified with  $n$ -bit binary strings. Suppose that Alice needs to send an  $n$ -bit plaintext  $p \in \mathbb{Z}_2^n$  to Bob.

*Hill cipher:*

- 1 (Key establishment) Alice and Bob randomly choose an invertible matrix  $K \in M_n(\mathbb{Z}_2)$  as a key, and compute its inverse.
- 2 (Encryption) Alice computes the ciphertext  $c$  according to the formula  $[c]_E^T = [p]_E^T \cdot K$ .
- 3 (Decryption) Bob computes the plaintext  $p$  according to the formula  $[p]_E^T = [c]_E^T \cdot K^{-1}$ .

The Hill cipher is nowadays insecure.

## Extra: Hill cipher II

Alice wants to send the message  $p = (1, 0, 1) \in \mathbb{Z}_2^3$  to Bob.

Alice and Bob agree on the matrix  $K$  and its inverse:

$$K = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 0 \end{pmatrix} \in M_3(\mathbb{Z}_2), \quad K^{-1} = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix} \in M_3(\mathbb{Z}_2).$$

Alice encrypts the message by computing the ciphertext  $c$  as:

$$[c]_E^T = [p]_E^T \cdot K = (1 \ 0 \ 1) \cdot \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 0 \end{pmatrix} = (0 \ 1 \ 1).$$

Bob decrypts the message by computing the plaintext  $p$  as:

$$[p]_E^T = [c]_E^T \cdot K^{-1} = (0 \ 1 \ 1) \cdot \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix} = (1 \ 0 \ 1).$$

## Extra: Image transformations I

Suppose that we have a 2D-image that we want to rotate counterclockwise with  $\theta$  degrees around the origin. By such a rotation, the point of coordinates  $(1, 0)$  becomes the point of coordinates  $(\cos \theta, \sin \theta)$ , while the point of coordinates  $(0, 1)$  becomes the point of coordinates  $(-\sin \theta, \cos \theta)$ .

We look for an  $\mathbb{R}$ -linear map  $f : \mathbb{R}^2 \rightarrow \mathbb{R}^2$  satisfying the following conditions:

$$\begin{aligned}f(1, 0) &= (\cos \theta, \sin \theta), \\f(0, 1) &= (-\sin \theta, \cos \theta).\end{aligned}$$

Recall that every linear map is determined by its values at the elements of a basis (the canonical basis in our case). Hence the matrix of the linear map  $f$  in the canonical basis  $E$  of the canonical real vector space  $\mathbb{R}^2$  is:

$$[f]_E = \begin{pmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{pmatrix}.$$

## Extra: Image transformations II

For any point  $v = (x, y) \in \mathbb{R}^2$  of a 2D-image, its corresponding point in the rotated image is computed as  $f(v) = (x', y') \in \mathbb{R}^2$ , where

$$\begin{pmatrix} x' \\ y' \end{pmatrix} = [f(v)]_E = [f]_E \cdot [v]_E = \begin{pmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{pmatrix} \cdot \begin{pmatrix} x \\ y \end{pmatrix}.$$

For instance, for a counterclockwise rotation of  $90^\circ$  around the origin one has the matrix:

$$[f]_E = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}.$$