```
// P_i
     do {
        flag[i] = true;
        while (flag[j]) {
            if (turn == j) {
                 flag[i] = false;
6
7
                 while (turn == j)
8
                 ; /* do nothing */
9
                 flag[i] = true;
10
11
        }
12
         /* critical section */
13
        turn = j;
        flag[i] = false;
14
         /* remainder section */
15
    } while (true);
```

```
// P j
2
     do {
3
        flag[j] = true;
        while (flag[i]) {
            if (turn == i) {
                 flag[j] = false;
7
                 while (turn == i)
                 ; /* do nothing */
9
                 flag[j] = true;
10
11
        }
12
         /* critical section */
13
        turn = i;
        flag[j] = false;
        /* remainder section */
    } while (true);
```

首先說明共享變數flag及turn代表的意義

共享變數:flag[2] of Boolean。(初值皆為 False)

共享變數:turn:int。(初值皆為0)

flag[i] = True P_i 有意願進入C.S; flag[i] = False P_i 無意願進入C.S。 Turn = i or j 代表哪個process目前擁有可以進入C.S的權限。 接著證明符合critical-section problem 三大特性:

1. Mutual Exclusion

假設P_i&P_j皆有意願進入C.S且i≠j,由於turn只會是i or j不可能同時是i & j,因此只有與turn對應的processs能夠進入C.S,而與turn不對應的process只能等待直到另一個process結束C.S將turn權限轉換才能進入C.S,符合Mutual Exclusion假設確保同一時間只有一個process能夠進入C.S。

2. Progress

- 情況一: 假設只有P_i有意願進入且P_j無意願進入C.S不管turn = i or j皆會跳出第三行while (flag[j]) 直接讓P_i進入C.S, 符合 Progress假設。
- 情況二:假設P_i&P_j皆有意願進入C.S, 只有與turn對應的process 才擁有進入C.S的權限, 在演算法中第四行if (turn == j) 會檢查turn 權限是否在對方, 如果是在對方會透過第五行flag[i] = false將自身狀態改成無意願進入同時進入第六行while (turn == j) do nothing 直到timer結束切換回擁有turn權限的process, 此時因為已經將沒有turn權限的process狀態改成無意願進入就會像情況一可以跳出第三行while (flag[j]) 直接進入C.S, 符合Progress假設確保擁有turn權限且有意願進入的process一定能夠優先進入C.S。

3. Bounded-waiting

● Dekker's演算法透過交替轉換turn值確保Bounded-waiting, 假設P_i&P_j皆有意願進入C.S且turn = i, 藉由Progress特性擁有turn權限且有意願進入的process一定能夠優先進入C.S, 在結束的同時會轉換turn值turn = j, 給正在等待的process, 這邊舉一個比較極端的例子在turn轉換且進入下一個迴圈flag[i] = true此時也會因為turn已經轉換透過if (turn == j) 檢查是否有權限並透過flag[i] = false; 將自身狀態改成無意願進入同時進入第六行while (turn == j) donothing直到timer結束, 這樣的設計機制使得 P_i 無法再度早於 P_j 進入 C.S., 一定是 P_j 先進入 C.S., 所以 P_j 至多等待一次後即可進入 C.S., 防止一個process能夠無限期地阻塞另一個process。

Q2

1. include相關library、定義worker_thread數量&每個word_thread所要產生的隨機點數量、Pthreads Mutex Locks和統計總共落在圓內點數量兩個global variable。

2. 選定一個幸運數字做為隨機種子及初始化5個worker thread和Mutex。

```
srand(20000154);

pthread_t threads[NUM_THREADS];
int i;

pthread_mutex_init(&mutex, NULL);
```

3. 創建5個worker_thread分配任務計算蒙地卡羅近似圓周率, 同時使用pthread_join() api 讓main thread等待worker_thread結束再繼續執行並回收其資源, 而為了能夠達到 multithread平行運行將pthread_create及pthread_join分成不同loop進行。

```
for (i = 0; i < NUM_THREADS; i++) {
    pthread_create(&threads[i], NULL, monteCarlo, NULL);
}

for (i = 0; i < NUM_THREADS; i++) {
    pthread_join(threads[i], NULL);
}</pre>
```

4. 接著頗析每個worker_thread執行計算蒙地卡羅近似圓周率的函式,進行助教要求 1000次生成隨機點,這裡確保生成的隨機點能落在(-1,1)之間同時透過計算半徑判斷是 否落在圓內,若有落在圓內則將數量加到前面提到統計落在圓內總數的global variable,值得注意的是為了確保global variable在各個worker_thread不會inconsistency使用 mutex_lock確保這個加的過程是atomic的。

```
void *monteCarlo(void *params) {
   double x, y, radius;
   int incircle = 0;
   for (i = 0; i < NUM_POINTS_PER_THREAD; i++) {
       x = ((double)rand() / RAND_MAX) * 2 - 1;
       y = ((double)rand() / RAND_MAX) * 2 - 1;
       radius = sqrt(x * x + y * y);
       if (radius <= 1.0) {
           incircle++;
   pthread_mutex_lock(&mutex);
   printf("%d points fall inside the circle\n", incircle);
   printf("Estimated value of pi per thread: %f\n", 4.0 * incircle / NUM_POINTS_PER_THREAD);
   total_incircle += incircle;
   total_point += NUM_POINTS_PER_THREAD;
   printf("Accumulate pi estimates from %d points: %f\n", total_point, 4.0 * total_incircle / total_point)
   pthread_mutex_unlock(&mutex);
   pthread_exit(NULL);
```

5. 等待所有worker_thread完成統計後釋放mutex佔用的資源使其無效最後再由 main_thread完成圓周率的估計。

Next, estimate π by performing the following calculation:

 $\pi = 4 \times$ (number of points in circle) / (total number of points)

```
pthread_mutex_destroy(&mutex);

double pi_estimate = 4.0 * total_incircle / (NUM_THREADS * NUM_POINTS_PER_THREAD);
printf("Estimated value of pi: %f\n", pi_estimate);

return 0;
```

6. 螢幕截圖結果, 上面是將每個worker_thread生成隨機點增加至100000, 下面是助教規定1000個點, 可以發現生成隨機點越多圓周率估計越精準。

```
(base) robert@LAPTOP-46RPPSGN:/mnt/c/users/rober/desktop/111-20S/hw3$ gcc Q2.C -o Q2 -lm (base) robert@LAPTOP-46RPPSGN:/mnt/c/users/rober/desktop/111-20S/hw3$ ./Q2
78547 points fall inside the circle
Estimated value of pi per thread: 3.141880
Accumulate pi estimates from 100000 points: 3.141880
78518 points fall inside the circle
Estimated value of pi per thread: 3.140720
Accumulate pi estimates from 2000000 points: 3.141300
78463 points fall inside the circle
Estimated value of pi per thread: 3.138520
Accumulate pi estimates from 300000 points: 3.140373
78794 points fall inside the circle
Estimated value of pi per thread: 3.151760
Accumulate pi estimates from 400000 points: 3.143220
78317 points fall inside the circle
Estimated value of pi per thread: 3.132680
Accumulate pi estimates from 5000000 points: 3.141112
Estimated value of pi: 3.141112
```

```
(base) robert@LAPTOP-46RPPSGN:/mnt/c/users/rober/desktop/111-20S/hw3$ gcc Q2.C -o Q2 -lm (base) robert@LAPTOP-46RPPSGN:/mnt/c/users/rober/desktop/111-20S/hw3$ ./Q2
789 points fall inside the circle
Estimated value of pi per thread: 3.156000
Accumulate pi estimates from 1000 points: 3.156000
779 points fall inside the circle
Estimated value of pi per thread: 3.116000
Accumulate pi estimates from 2000 points: 3.136000
774 points fall inside the circle
Estimated value of pi per thread: 3.096000
Accumulate pi estimates from 3000 points: 3.122667
797 points fall inside the circle
Estimated value of pi per thread: 3.188000
Accumulate pi estimates from 4000 points: 3.139000
797 points fall inside the circle
Estimated value of pi per thread: 3.188000
Accumulate pi estimates from 5000 points: 3.148800
Estimated value of pi: 3.148800
```