

From Rational Number Reconstruction to Set Reconciliation

Antoine Amarilli, Fabrice Ben Hamouda, Florian Bourse,
Robin Morisset, David Naccache, and Pablo Rauzy

École normale supérieure, Département d’informatique
45, rue d’Ulm, F-75230, Paris Cedex 05, France.
`surname.name@ens.fr` (except for `fabrice.ben.hamouda@ens.fr`)

Abstract. This work revisits *set reconciliation*, a problem consisting in synchronizing two multisets of fixed-size values while minimizing the amount of data transmitted. We propose a new number theoretic reconciliation protocol called “Divide & Factor” (D&F) which achieves optimal asymptotic transmission complexity like prior proposals. We then study the problem of synchronizing sets of variable-size files, and describe how constant-factor improvements can be achieved through the use of hashing with a carefully chosen hash size (balancing the quantity of data transferred and the risk of collisions). We show how this process can be applied to synchronize file hierarchies, taking into account the location of files. We describe `btrsync`, our open-source implementation of the protocol, and benchmark it against the popular software `rsync` to demonstrate that `btrsync` uses more CPU time but transmits less data.

1 Introduction

This work revisits *set reconciliation*, a problem consisting in synchronizing two multisets while minimizing the amount of data transmitted. Set reconciliation arises in many practical situations, the most typical of which is certainly incremental backups performed over a slow network link.

Several efficient and elegant solutions are known to achieve set reconciliation of multisets containing atomic elements of a fixed size. For instance, [8] manages to perform set reconciliation using a bandwidth which is linear in the size of the symmetric difference of the multisets multiplied by the size of the elements, which is optimal in this setting. We refer the reader to [8,9,5] (to quote a few references) for more on this problem’s history and its existing solutions.

However, in the case where the elements to be synchronized can be very large (e.g., files during a backup), we must use checksums to identify the differing files before transferring them, and the question of the size of the checksum to use is non-trivial. In this article, we propose a new reconciliation protocol called “Divide & Factor” (D&F) based on number theory. In terms of asymptotic transmission complexity, the proposed procedure reaches optimality as well. In addition, the new protocols offer a very interesting gamut of parameter trade-offs. We provide an analysis of the protocol’s complexity in terms of transmission and computation, as well as a probabilistic analysis of the possible choices of checksum sizes; we also provide an implementation of the protocol and experimental results.

*Si cet abstract
vous convient,
il faudrait en
reprendre des
bouts dans
l'introduction.*

*Fabrice:
j'écirai
quelque chose
de plus positif
comme:
“btrsync
transmits
much less
data at the
expense of a
small compu-
tationnal
overhead”
c'est vrai
qu'on apporte
quelque chose
de nouveau
dans le cas où
c'est des
fichiers de
taille non fixe
? — Fabrice:
je ne sais pas,
il faudrait que
l'on réanalyse
la complexité
propre de la
dernière
variante
Fabrice@Antoine:
Je ne suis pas
tout à fait
d'accord avec
ce
paragraphe,
ton abstract
est beaucoup
mieux...*

This paper is structured as follows: Section 2.2 presents a basic version of the proposed protocol. This basic version suffers from two limitations: it works only if the number of differences to reconcile is bound and it may fail leave the synchronized party in an erroneous state. Failure avoidance is overcome in section 2.3 and an extension to arbitrary numbers of differences is given in section 2.4. The protocol’s transmission complexity is treated in section 3. Section 3 also introduces two transmission optimizations and analyzes them in detail. Section 4 analyzes the computational complexities of the proposed protocols and 5 reports practical experiments and benchmarks against the popular software `rsync`.

2 “Divide & Factor” Set Reconciliation

2.1 Problem Definition and Notations

Oscar possesses an old version of a directory \mathcal{D} that he wishes to update. Neil has the new, up-to-date version \mathcal{D}' : \mathcal{D} and \mathcal{D}' can differ both in their files and in their tree structures. Oscar wishes to obtain \mathcal{D}' but *exchange as little data as possible* during the synchronization process.

To tackle this problem we separate the “*what*” from the “*where*” by considering files as a tuple of their location and content. In other words, we will first synchronize all the file contents and then move files to the adequate location. We consider that \mathcal{D} is a multiset of files which we denote as $\mathfrak{F} = \{F_0, \dots, F_n\}$, and likewise represent \mathcal{D}' as $\mathfrak{F}' = \{F'_0, \dots, F'_{n'}\}$.

Let t_0 be the number of discrepancies between \mathfrak{F} and \mathfrak{F}' that Oscar wishes to learn, i.e. the symmetric difference of \mathfrak{F} and \mathfrak{F}' :

$$t_0 = \#\mathfrak{F} + \#\mathfrak{F}' - 2\#(\mathfrak{F} \cap \mathfrak{F}') = \#(\mathfrak{F} \cup \mathfrak{F}') - \#(\mathfrak{F} \cap \mathfrak{F}')$$

Given a file F , we denote by $\text{Hash}(F)$ its image by a collision-resistant hash function such as SHA-1. Let $\text{HashPrime}(F)^1$ be a function hashing files (uniformly) into primes smaller than 2^u for some $u \in \mathbb{N}$. Define the shorthand notations: $h_i = \text{HashPrime}(F_i)$ and $h'_i = \text{HashPrime}(F'_i)$.

2.2 Description of the Basic Exchanges

The number of differences t_0 is unknown to Oscar and Neil. However, for the time being, we will assume that t_0 is smaller than some t and attempt to perform synchronization. If $t_0 \leq t$, synchronization will succeed; if $t_0 > t$ the parties will transmit more information later to complete the synchronization, as explained in section 2.4.

We generate a prime p such that:

$$2^{2ut+1} \leq p < 2^{2ut+2} \tag{1}$$

¹ The design of `HashPrime` is addressed in Appendix C.

Fabrice:
NON, ce n'est
pas vrai: on
considère
vraiment
qu'un fichier
est "path +
content"

Fabrice:
Qu'est-ce que
 \mathcal{D} ? Est-ce \mathfrak{F}
? Dans tous
les cas, il
s'agit d'un
set, vu la
représentation
indiquée

Fabrice: `Hash`
can be
introduced
latter when
needed
(section 2.3)

Given \mathfrak{F} , Oscar generates and sends to Neil the redundancy:

$$c = \prod_{F_i \in \mathfrak{F}} \text{HashPrime}(F_i) = \prod_{i=1}^n h_i \bmod p$$

Neil computes:

$$c' = \prod_{F'_i \in \mathfrak{F}'} \text{HashPrime}(F'_i) = \prod_{i=1}^{n'} h'_i \bmod p \quad \text{and} \quad s = \frac{c'}{c} \bmod p$$

Using [13] the integer s can be written as:

$$s = \frac{a}{b} \bmod p \text{ where the } G_i \text{ denote files and } \begin{cases} a = \prod_{G_i \in \mathfrak{F}' \wedge G_i \notin \mathfrak{F}} \text{HashPrime}(G_i) \\ b = \prod_{G_i \notin \mathfrak{F}' \wedge G_i \in \mathfrak{F}} \text{HashPrime}(G_i) \end{cases}$$

Note that if our assumption $t_0 \leq t$ is correct, \mathfrak{F} and \mathfrak{F}' differ by at most t elements and a and b are strictly less than 2^{ut} . The problem of recovering a and b from s efficiently is known as *Rational Number Reconstruction* [10,14]; theorem 1 (see [6]) guarantees that it can be solved in this setting.

Theorem 1. *Let $a, b \in \mathbb{Z}$ such that $-A \leq a \leq A$ and $0 < b \leq B$. Let $p > 2AB$ be a prime and $s = ab^{-1} \bmod p$. Then a, b can be recovered from A, B, s, p in polynomial time.*

Taking $A = B = 2^{ut} - 1$, Equation (1) implies that $2AB < p$. Moreover, $0 \leq a \leq A$ and $0 < b \leq B$. Thus Oscar can recover a and b from s in polynomial time: a possible option is to use Gauss algorithm for finding the shortest vector in a bi-dimensional lattice [13]. By testing the divisibility of a and b by the h_i and the h'_i , Neil and Oscar can attempt to identify the discrepancies between \mathfrak{F} and \mathfrak{F}' and settle them.

The formal description of the protocol is given in Figure 1. The “output $\perp_{\text{bandwidth}, \square}$ ” protocol interruptions will:

- never occur if the assumption $t_0 \leq t$ holds.
- occur with high probability if $t_0 > t$. Indeed, for a potential $\perp_{\text{bandwidth}, 1}$ to be overlooked, the ut -bit number a must perfectly factor over a set of n primes of size u . If we assume that a is “random”, the probability γ that a is divisible by some h_i is essentially $\gamma \sim 1/h_i \sim 2^{-u}$, the probability that a is divisible by exactly t digests is:

$$\alpha = \binom{n}{t} \gamma^t (1 - \gamma)^{n-t} \sim \binom{n}{t} 2^{-ut} (1 - 2^{-u})^{n-t}$$

and the probability that the protocol does not terminate by a $\perp_{\text{bandwidth}, \square}$ when $t_0 > t$ is $\sim \alpha^2$.

Fabrice: Quel est l'intérêt de cette citation ???

Fabrice: je ne comprends pas le “where the G_i denote files...”

Fabrice: certes, mais on utilise directement un Euclide étendu tronqué, et en plus je ne suis pas sûr que le lecteur comprenne le lien si facilement que cela avec les réseaux... en plus l'article ne donne pas de complexité temporelle... we are not interested in the fact a is divisible by “exactly” t digests but the fact that a can be factorized over of the basis of h_i ...

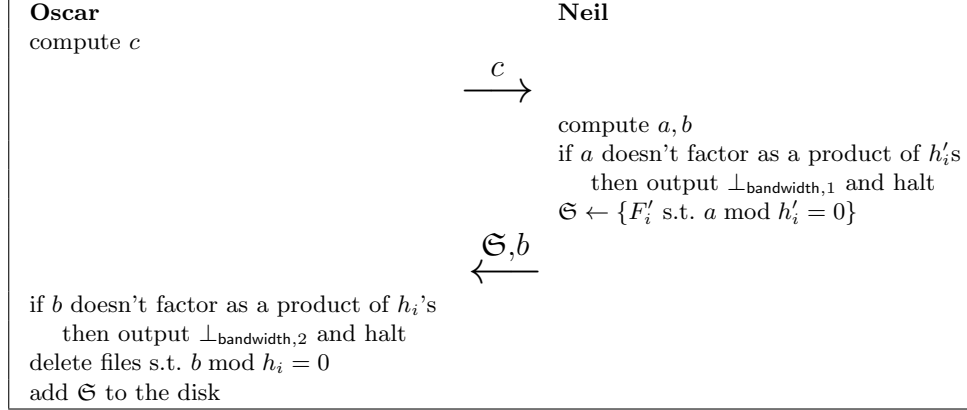


Fig. 1. Basic Protocol.

The very existence of $\perp_{\text{bandwidth},\square}$'s is annoying for two reasons:

- A file synchronization procedure that works *only* for a limited number of differences is not really useful in practice. Thus, section 2.4 explains how to extend the protocol to perform the synchronization even when the number of differences t_0 exceeds the initial estimation t .
- If, by sheer bad luck, both $\perp_{\text{bandwidth},\square}$'s went undetected (double accidental factorization) the Basic Protocol (Fig. 1) may leave Oscar in an inconsistent state.

Double accidental factorization is not only possible source of inconsistent states: as we did not specifically require **HashPrime** to be collision-resistant, the events

$$\perp_{\text{collision},1} = \begin{cases} h'_i = h'_j \text{ for } i \neq j \\ a \bmod h_i = 0 \end{cases} \quad \text{and/or} \quad \perp_{\text{collision},2} = \begin{cases} h_i = h_j \text{ for } i \neq j \\ b \bmod h'_i = 0 \end{cases}$$

will cause Neil to send wrong files in \mathfrak{S} ($\perp_{\text{collision},1}$) and/or have Oscar unduely delete files owned by Neil ($\perp_{\text{collision},2}$).

Inconsistent states may hence stem from three events:

- accidental double factorization of a and/or b when $t_0 > t$ (probability α^2)
- $\perp_{\text{collision},1}$ = collisions within the set $\{h'_i\}$
- $\perp_{\text{collision},2}$ = collisions within the set $\{h_i\}$

*Fabrice: ce
qui est un peu
tordu, c'est
que les
collisions qui
nous
intéressent
sont les
collisions
entre \mathfrak{S} et
 $\mathfrak{F} \cup \mathfrak{F}' \dots$ il
faut qu'on en
discute*

Section 2.3 explains how protect the protocol from all inconsistent events at once.

where $\lambda = t/t_0$ is the number of rounds required to complete the protocol. As we have no control over t , decreasing u is the main natural optimization option. We will get back to this later on in this paper (section 3.2).

3.1 Probabilistic Decoding: Reducing p

Generate a prime p about twice shorter than the p recommended in section 2.2, namely:

$$2^{u(t+1)} < p \leq 2^{u(t+1)+1} \quad (2)$$

Let $\eta = \max(n, n')$. The new redundancy c is calculated as previously and is hence also approximately twice smaller. Namely:

$$s = \frac{a}{b} \bmod p \text{ and } \begin{cases} a = \prod_{G_i \in \mathfrak{F}' \wedge G_i \notin \mathfrak{F}} \text{HashPrime}(G_i) \\ b = \prod_{G_i \notin \mathfrak{F}' \wedge G_i \in \mathfrak{F}} \text{HashPrime}(G_i) \end{cases}$$

and since there are at most t differences, we must have:

$$ab \leq 2^{ut} \quad (3)$$

By opposition to section 2.2 we do not have a fixed bound for a and b anymore; Equation (3) only provides a bound for the product ab . Therefore, we define a sequence of $t+1$ couples of bounds:

$$(A_i, B_i) = (2^{ui}, 2^{u(t-i)}) \forall i \in \{0, \dots, t\}$$

Equations (2) and (3) imply that there must exist at least one index i such that $0 \leq a \leq A_i$ and $0 < b \leq B_i$. Then using Theorem 1, since $2A_iB_i = 2^{u(t+1)} < p$, given (A_i, B_i, p, s) one can recover (a, b) , and hence the difference between \mathfrak{F} and \mathfrak{F}' .

The problem is that (unlike section 2.2) we have no guarantee that such an (a, b) is unique. Namely, we could (in theory) stumble over an $(a', b') \neq (a, b)$ satisfying (3) for some index $i' \neq i$. We conjecture that such failures happen with negligible probability (that we do not try to estimate here) when u is large enough, but this makes the modified protocol heuristic only. If failures never occur, this variant will roughly halve the quantity of transmitted bits with respect to section 2.2.

3.2 The File Laundry: Reducing u

What happens if we brutally shorten u in the basic Divide & Factor protocol? As expected by the birthday paradox, we should start seeing collisions. Let us analyze the statistics governing the appearance of collisions.

Fabrice:
 “heuristic”
 → not
 really, because
 of the final
 hash
 verification,
 though we
 cannot
 compute
 exactly the
 complexity...
 Fabrice:

maybe say it
 in another
 way, since we
 already do
 not suppose
 HashPrime is
 collision-
 resistant
 Fabrice: What
 is η exactly ?
 $n, n', n + n',$
 \dots

Consider **HashPrime** as a random function from $\{0, 1\}^*$ to $\{0, \dots, 2^u - 1\}$. Let X_i be the random variable:

$$X_i = \begin{cases} 1 & \text{if file } F_i \text{ collides with another file.} \\ 0 & \text{otherwise.} \end{cases}$$

Clearly, we have $\Pr[X_i = 1] \leq \frac{\eta-1}{2^u}$. The average number of colliding files is hence:

$$\mathbb{E} \left[\sum_{i=0}^{\eta-1} X_i \right] \leq \sum_{i=0}^{\eta-1} \frac{\eta-1}{2^u} = \frac{\eta(\eta-1)}{2^u}$$

For instance, for $\eta = 10^6$ files and 32-bit digests, the expected number of colliding files is less than 233.

However, it is important to note that a collision can only yield a *false positive*, and never a *false negative*. In other words, while a collision may oblivate a difference² a collision will never create a nonexistent difference *ex nihilo*.

Thus, it suffices to replace **HashPrime**(F) by a diversified $\tilde{h}_k(F) = \mathbf{HashPrime}(k|F)$ to quickly filter-out file differences by repeating the protocol for $k = 1, 2, \dots$. At each iteration the parties will detect new files and new deletions, fix these and “launder” again the remaining multisets.

Assume that the diversified $\tilde{h}_k(F)$ ’s are random and independent. To understand why the probability that a stubborn file persists colliding decreases exponentially with the number of iterations k , assume that η remains invariant between iterations and define the following random variables:

$$X_i^\ell = \begin{cases} 1 & \text{if file } F_i \text{ collides with another file during iteration } \ell. \\ 0 & \text{otherwise.} \end{cases}$$

$$Y_i = \prod_{\ell=1}^k X_i^\ell = \begin{cases} 1 & \text{if file } F_i \text{ collides with another file during all the } k \text{ first} \\ & \text{protocol iterations.} \\ 0 & \text{otherwise.} \end{cases}$$

By independence, we have:

$$\Pr[Y_i = 1] = \prod_{\ell=1}^k \Pr[X_i^\ell = 1] = \Pr[X_i^1 = 1] \dots \Pr[X_i^k = 1] \leq \left(\frac{\eta-1}{2^u} \right)^k$$

² e.g. make the parties blind to the difference between `index.htm` and `iexplore.exe`.

Fabrice:
difference or
discrepancy ?

Fabrice:
already said
in Section
2.3... and
maybe it is
better to say
that, in this
first analysis,
we suppose
we do not
filter out files,
because this
only improves
the algo...

*Fabrice je ne
vois pas la
difference
entre X_i^ℓ dans
cette section
et Z_i^ℓ dans la
suivante.
Peux-tu
preciser STP*

Therefore the average number of colliding files is:

$$\mathbb{E} \left[\sum_{i=0}^{\eta-1} Y_i \right] \leq \sum_{i=0}^{\eta-1} \left(\frac{\eta-1}{2^u} \right)^k = \eta \left(\frac{\eta-1}{2^u} \right)^k$$

And the probability that at least one false positive will survive k rounds is:

$$\epsilon_k \leq \eta \left(\frac{\eta-1}{2^u} \right)^k$$

For the previously considered instance³ we get $\epsilon_2 \leq 5.43\%$ and $\epsilon_3 \leq 2 \cdot 10^{-3}\%$.

A more refined (but somewhat technical) analysis. As mentioned previously, the parties can remove the files confirmed as different during iteration k and work during iteration $k+1$ only with common and colliding files. Now, the only collisions that can fool round k , are the collisions of file-pairs (F_i, F_j) such that F_i and F_j have both already collided during *all the previous iterations*⁴. We call such collisions “masquerade balls” (cf. Figure 4). Define the two random variables:

$$Z_i^\ell = \begin{cases} 1 & \text{if } F_i \text{ participated in masquerade balls during all the } \ell \\ & \text{first protocol iterations.} \\ 0 & \text{otherwise.} \end{cases}$$

$$X_{i,j}^\ell = \begin{cases} 1 & \text{if files } F_i \text{ and } F_j \text{ collide during iteration } \ell. \\ 0 & \text{otherwise.} \end{cases}$$

Set $Z_i^0 = 1$ and write $p_\ell = \Pr \left[Z_i^\ell = 1 \text{ and } Z_j^\ell = 1 \right]$ for all ℓ and $i \neq j$. For $k \geq 1$, we have:

$$\begin{aligned} \Pr \left[Z_i^k = 1 \right] &= \Pr \left[\exists j \neq i, X_{i,j}^k = 1, Z_i^{k-1} = 1 \text{ and } Z_j^{k-1} = 1 \right] \\ &\leq \sum_{j=0, j \neq i}^{\eta-1} \Pr \left[X_{i,j}^k = 1 \right] \Pr \left[Z_i^{k-1} = 1 \text{ and } Z_j^{k-1} = 1 \right] \\ &\leq \frac{\eta-1}{2^u} p_{k-1} \end{aligned}$$

Furthermore $p_0 = 1$ and

³ $\eta = 10^6, u = 32$.

⁴ Note that we do not require that F_i and F_j repeatedly collide *with each other*. e.g. we may witness during the first round $h_1(F_1) = h_1(F_2)$, $h_1(F_3) = h_1(F_4)$ and $h_1(F_5) = h_1(F_6)$ while during the second round $h_2(F_1) = h_2(F_2)$, $h_2(F_3) = h_1(F_6)$ and $h_2(F_5) = h_2(F_4)$ as shown in Figure 4.

*Fabrice:
maybe put
this in
appendix and
add a few
comments...*

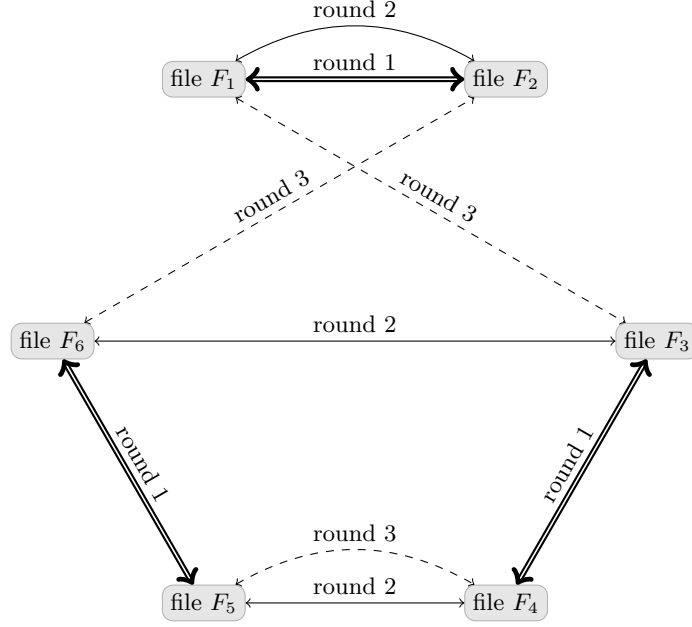


Fig. 4. Illustration of three masquerade balls. Each protocol round is materialized by a different type of arrow. Arrows denotes collisions.

$$\begin{aligned}
p_\ell &= \Pr \left[X_{0,1}^\ell = 1, Z_0^\ell = 1 \text{ and } Z_1^\ell = 1 \right] + \Pr \left[X_{0,1}^\ell = 0, Z_0^\ell = 1 \text{ and } Z_1^\ell = 1 \right] \\
&\leq \Pr \left[X_{0,1}^\ell = 1, Z_0^{\ell-1} = 1 \text{ and } Z_1^{\ell-1} = 1 \right] \\
&\quad + \sum_{i \geq 2, j \geq 2} \Pr \left[X_{0,i}^\ell = 1, X_{1,j}^\ell = 1, Z_0^{\ell-1} = 1 \text{ and } Z_1^{\ell-1} = 1 \right] \\
&= \Pr \left[X_0^\ell = X_1^\ell \right] \Pr \left[Z_0^{\ell-1} = 1 \text{ and } Z_1^{\ell-1} = 1 \right] \\
&\quad + \sum_{i \geq 2, j \geq 2} \Pr \left[X_{0,i}^\ell = 1 \right] \Pr \left[X_{1,j}^\ell = 1 \right] \Pr \left[Z_0^{\ell-1} = 1 \text{ and } Z_1^{\ell-1} = 1 \right] \\
&\leq \frac{1}{2^u} p_{\ell-1} + \frac{(\eta-2)^2}{2^{2u}} p_{\ell-1} = p_{\ell-1} \left(\frac{1}{2^u} + \frac{(\eta-2)^2}{2^{2u}} \right)
\end{aligned}$$

hence:

$$p_\ell \leq \left(\frac{1}{2^u} + \frac{(\eta-2)^2}{2^{2u}} \right)^\ell,$$

and

$$\Pr \left[Z_i^\ell = 1 \right] \leq \frac{\eta-1}{2^u} \left(\frac{1}{2^u} + \frac{(\eta-2)^2}{2^{2u}} \right)^{k-1}$$

And finally, the survival probability of at least one false positive after k iterations satisfies:

$$\epsilon'_k \leq \frac{\eta(\eta-1)}{2^u} \left(\frac{1}{2^u} + \frac{(\eta-2)^2}{2^{2u}} \right)^{k-1}$$

For $(\eta = 10^6, u = 32, k = 2)$, we get $\epsilon'_2 \leq 0.013\%$.

How to select u ? For a fixed k , ϵ'_k decreases as u grows. For a fixed u , ϵ'_k also decreases as k grows. Transmission, however, grows with both u (bigger digests) and k (more iterations). We write for the sake of clarity: $\epsilon'_k = \epsilon'_{k,u,\eta}$.

Fix η . Note that the number of bits transmitted per iteration ($\simeq 3ut$), is proportional to u . This yields an expected transmission complexity bound $T_{u,\eta}$ such that:

$$T_{u,\eta} \propto u \sum_{k=1}^{\infty} k \cdot \epsilon'_{k,u,\eta} = \frac{u\eta(\eta-1)}{2^u} \sum_{k=1}^{\infty} k \left(\frac{1}{2^u} + \frac{(\eta-2)^2}{2^{2u}} \right)^{k-1} = \frac{u\eta(\eta-1)8^u}{(2^u - 4^u + (\eta-2)^2)^2}$$

Dropping the proportionality factor $\eta(\eta-1)$, neglecting $2^u \ll 2^{2u}$ and approximating $(\eta-2) \simeq \eta$, we can optimize the function:

$$\phi_{\eta}(u) = \frac{u \cdot 8^u}{(4^u - \eta^2)^2}$$

$\phi_{10^6}(u)$ admits an optimum for $u = 19$.

Note: The previous analysis is incomplete because of the following approximations:

- We consider u -bit prime digests while u -bit strings contain only about $2^u/u$ primes.
- We used a fixed u in all rounds. Nothing forbids using a different u_k at each iteration, or even fine-tuning the u_k s adaptively as a function of the laundry's effect on the progressively reconciliated multisets..
- Our analysis treats t as a constant, but large t values increase p and hence the number of potential files detected as different per iteration - an effect disregarded *supra*.

A different approach is to optimize t and u experimentally, e.g. using the open source D&F program `btrsyc` developed by the authors (*cf.* section 5).

3.3 How to Stop a Probabilistic Washing Machine?

We now combine both optimizations and assume that ℓ laundry rounds are necessary for completing some given reconciliation task using a half-sized p . By opposition to section 2.2, confirming correct protocol termination is now non-trivial.

We say that a *round failure* occurs whenever a round results in an $(a', b') \neq (a, b)$ satisfying Equation (3). Let the round failure probability be some function $\zeta(u)$ (that we did not estimate). If u is kept small (for efficiency reasons), the probability $(1 - \zeta(u))^\ell$ that the protocol will properly terminate may dangerously drift away from one.

If v of $\ell + v$ rounds fail, Oscar needs to solve a problem called *Chinese Remaindering With Errors* [2]:

Problem 1. (Chinese Remaindering With Errors Problem: CRWEP). Given as input integers v , B and $\ell + v$ points $(s_1, p_1), \dots, (s_{\ell+v}, p_{\ell+v}) \in \mathbb{N}^2$ where the p_i 's are coprime, output all numbers $0 \leq s < B$ such that $s \equiv s_i \pmod{p_i}$ for at least ℓ values of i .

We refer the reader to [2] for more on this problem, which is beyond our scope. Boneh [3] provides a polynomial-time algorithm for solving the CRWEP under certain conditions satisfied in our setting.

But how can we confirm the solution? As mentioned in section 2.3, Neil will send to Oscar $H = \text{Hash}(\mathfrak{F}')$ as the interaction starts. As long as Oscar's CRWEP resolution will not yield a state matching H , the parties will continue the interaction.

4 Computational Complexity

Let $\mu(k)$ be the time required to multiply two k -bit numbers⁵. For naïve (i.e. convolutive) algorithms $\mu(k) = O(k^2)$, but using FFT multiplication [11], $\mu(k) = \tilde{O}(k)$. FFT is experimentally faster than convolutive methods starting at $k \sim 10^6$. The modular division of two k -bit numbers and the reduction of $2k$ -bit number modulo a k -bit number are also known to cost $\tilde{O}(\mu(k))$ [4]. Indeed, in packages such as **gmp**, division and modular reduction run in $\tilde{O}(k)$ for sufficiently large k .

Given that $p \sim 2^{ut}$, we obtain the following complexity analysis:

Entity	Computation	Complexity expressed in \tilde{O} of			
Both	generate primes	nu^4	Rabin-Miller	nu^3	Rabin-Miller
Both	compute redundancies c and c'	$n \cdot \mu(ut)$	naïve product	nut	using FFT
Oscar	compute $s = c'/c \pmod{p}$	$\mu(ut)$	naïve inversion	ut	using FFT
Oscar	find a, b such that $s = a/b \pmod{p}$	$(ut)^2$	naïve ext. GCD	ut	using [10,14]
Both	factor a (resp. b) by modular reductions	$n \cdot \mu(ut)$	naïve reduction	nut	using FFT
Overwhelming complexity:		$\max(nu^4, n \cdot \mu(ut), (ut)^2)$		$nu \max(t, u^3)$	

Table 1. Global Protocol Complexity. In practice $\max((ut)^2, n \cdot \mu(ut)) = n(ut)^2$ because $(ut)^2 \ll n \cdot \mu(ut) = n(ut)^2$. This boils down to $\max(n(ut)^2, nu^4) = n(u \max(t, u))^2$.

⁵ We assume that $\forall k, k', \mu(k + k') \geq \mu(k) + \mu(k')$.

Je ne comprends pas quelle la différence entre les colonnes 3 4 et les colonnes 5 6 de ce tableau.

Fabrice: oups, pas sûr de tout comprendre ici, il faudrait que l'on en rediscute... En particulier, le CRT n'a pas d'erreur a priori, sauf erreur de ma part. La définition de $\zeta(u)$ est assez gratuite vu que ça n'intervient pas dans la suite.

4.1 Improvements Using Product Trees

commIndiquer clairement à quelles lignes du tableau ça se réfère, et si la complexité donnée dans le tableau prend ces optims en compte. The non-overwhelming (but nonetheless important) complexities of the computations of (c, c') and of the factorizations can be even reduced to $\tilde{O}(\frac{n}{t}\mu(ut))$ using convolutive methods. These complexities can be reduced to $\tilde{O}(nu)$ with FFT [11]. To simplify the presentation, assume that $t = 2^\tau$ is a power of two dividing n .

The idea is the following: group h_i 's by subsets of t elements and compute the product of each such subset in \mathbb{N} .

$$H_j = \prod_{i=jt}^{jt+t-1} h_i \in \mathbb{N}$$

Each H_j can be computed in $\tilde{O}(\mu(ut))$ using the standard product tree method described in Algorithm 1 (for $j = 0$) illustrated in Figure 5. Thus, all these $\frac{n}{t}$ products can be computed in $\tilde{O}(\frac{n}{t}\mu(ut))$. We can then compute c by multiplying the H_j modulo p , which costs $\tilde{O}(\frac{n}{t}\mu(ut))$.

Algorithm 1 Product Tree Algorithm

Require: the set h_i

Ensure: $\pi = \pi_1 = \prod_{i=0}^{t-1} h_i$, and π_i for $i \in \{1, \dots, 2t-1\}$ as in Figure 5

```

1:  $\pi \leftarrow$  array of size  $t$ 
2: function PRODTREE( $i, \text{start}, \text{end}$ )
3:   if  $\text{start} = \text{end}$  then
4:     return 1
5:   else if  $\text{start} + 1 = \text{end}$  then
6:     return  $h_{\text{start}}$ 
7:   else
8:      $\text{mid} \leftarrow \lfloor \frac{\text{start} + \text{end}}{2} \rfloor$ 
9:      $\pi_{2i} \leftarrow$  PRODTREE( $2i, \text{start}, \text{mid}$ )
10:     $\pi_{2i+1} \leftarrow$  PRODTREE( $2i+1, \text{mid}, \text{end}$ )
11:    return  $\pi_{2i} \times \pi_{2i+1}$ 
12:  $\pi_1 \leftarrow$  PRODTREE( $1, 0, t$ )
```

The same technique applies to factorization⁶, but with a slight caveat.

After computing the tree product, we can compute the residues of a modulo H_0 . Then we can compute the residues of $a \bmod H_0$ modulo the two children π_2 and π_3 of $H_0 = \pi_1$ in the product tree (depicted in Figure 5), and so on. Intuitively, we descend the product tree doing modulo reduction. At the end (i.e., as we reach the leaves), we obtain the residues of a modulo each of the h_i ($i \in \{0, \dots, t-1\}$). This is described in Algorithm 6 and illustrated

*What is the
caveat?*

⁶ We explain the process with a , this is applicable *ne variatur* to b as well.

Does it make sense to nest \tilde{O} 's?

in Figure 6. We can use the same method for the tree product associated to any H_j , and the residues of a modulo each of the h_i ($i \in \{jt, \dots, jt+t-1\}$) for any j , i.e., a modulo each of the h_i for any i . Complexity is $\tilde{O}(\mu(ut))$ for each j , which amounts to a total complexity of $\tilde{O}(\frac{n}{t}\tilde{O}(\mu(ut)))$.

Algorithm 2 Division Using a Product Tree

Require: $a \in \mathbb{N}$, π the product tree of Algorithm 1

Ensure: $A[i] = a \bmod \pi_i$ for $i \in \{1, \dots, 2t-1\}$, computed as in Figure 2

```

1:  $A \leftarrow$  array of size  $t$ 
2: function MODTREE( $i$ )
3:   if  $i < 2t$  then
4:      $A[i] \leftarrow A[\lfloor i/2 \rfloor] \bmod \pi_i$ 
5:     MODTREE( $2i$ )
6:     MODTREE( $2i+1$ )
7:  $A[1] \leftarrow a \bmod \pi_1$ 
8: MODTREE( $2$ )
9: MODTREE( $3$ )
```

4.2 Adapting p

Let $\text{Prime}[i]$ denote the i -th prime⁷. Besides conditions on size, the *only* property required from p is to be co-prime with all the h_i and all the h'_i . We can hence consider the following variants:

Variant 1: Smooth p_i :

$$p_i = \prod_{j=r_i}^{r_{i+1}-1} \text{Prime}[j]$$

Where the bounds r_i are chosen to ensure that each p_i has the proper size.

Variant 2: $p_i = \text{Prime}[i]^{r_i}$ where the exponents r_i are chosen to ensure that each p_i has the proper size.

Variant 3: Progressively work modulo a growing power of two. This variant, is probably the most efficient of all, but somewhat complex to explain. We hence describe it in detail in Appendix B. This variant is compatible with the use of FFT-multiplication, hence asymptotic complexity is preserved. In addition, it avoids all modular reductions and all CRT re-combinations and hence offers considerable constant-factor accelerations.

⁷ with $\text{Prime}[1] = 2$

What is the goal of the following variants? Is it to reduce the complexity of the generation of the primes?

Have this written by Fabrice...

```
(les modulus etant  $2^{\text{ut}}$ ,  $2^{2\text{ut}}$ , ...):
- calcul des  $H_j$  comme dans 4.1 par product tree
- produit des  $H_j$  de la facon suivante pour round 1:
  -  $\text{prod} = 1$  (de taille  $\text{ut}$ )
  - pour  $j=1, \dots$ 
     $\text{carry}_j \mid \text{prod} = \text{prod} \times H_j$ 
  - retourner prod, et garder carry
- produit des  $H_j$  pour round suivants:
  - prod = 0
  - pour  $j=1, \dots$ 
     $\text{carry}_j \mid \text{prod} = \text{prod} \times H_j + \text{carry}_j$ 
```

L'idée est la suivante (invariant): au round i et au tour j , prod contient les bits $(i-1)\text{ut}$, carry_j contient le carry genere par ces bits...

Remarks Fabrice:

- we should separate move resolution algo from implementation
- we should implement the non-prime version stuff... but difficult to find the correct subset of h_i which factorizes...

5 Implementation

We implemented and benchmarked the Divide & Factor protocol described in the previous sections. The implementation is called `btrfsync`, its source code is available from [1].

The program performs unidirectional synchronization, which is simpler to understand. The code is divided into two subprograms: a shell script and a Python program:

The Shell Script sets up two instances of the Python program on Oscar and Neil and establishes a bidirectional communication channel between them using two Unix pipes between their standard inputs and outputs.

The Python Program uses `gmp` to perform all the number theory operations and performs the actual synchronization. It proceeds in two phases:

Finding Different Files

1. Hash of files' contents concatenated with their paths, types (folder/file), and permissions (not supported yet).

*on parle de
"in the
previous
sections"
mais de
quelle variate
s'agit-il?
~~Fabrice: in~~
the paper, it
is unidirec-
tional too
!*

2. Implement the protocol proposed in Section ?? with input data coming from `stdin` and output data going to `stdout`.

More precisely:

*sur "a first
prime number
 p_1 ", l'implem
marche avec
 p_i ou des 2^u ?*

- Oscar sends it product of hashes modulo a first prime number p_1 .
- Neil receives the product, divides by its own product of hashes, reconstructs the fraction modulo p_1 and checks if he can factor the denominator using his hashes base. If he can, he stops and sends the numerator and the list of tuples (path, type, hash of content of the file) corresponding to the denominator's factors. Otherwise he sends "None" [is this the ASCII string "None"? if not what does he send precisely?].
- If Neil sent "None", Oscar computes the product of hashes modulo another prime p_2 , sends it... CRT mechanism... [can we elaborate more on what happens here? which functions in GMP are used to do the CRT?]
- If Neil sent the numerator and a list of tuples, then Oscar factors the numerator over his own hash values. Now each party (Neil, Oscar) knows precisely the list of files (path + type + hash of content) that differs from the other party.

[please structure the following:]

2. synchronize all the files. This part is not completely optimized.

We just remove all folders Oscar should not have and create new folders.

Then we remove all files Oscar should not have and synchronize using `rsync` the last files.

We could check for move (since we have the list of hash of contents of files) and do moves locally.

We can even try to detect moves of complete subtrees...

*encore une
fois
"Implement
the protocol
proposed in
Section ??"
qu'a t on
implemente
au juste?*

5.1 Move Resolution Algorithm

To reproduce the structure of Oscar on Neil's disk, we need to perform a sequence of file moves. Sadly, moves are not straightforward to apply the moves, because, if we take a file to move, its destination might be blocked, either because a file already exists (we want to move a to b , but b already exists), or because a folder cannot be created (we want to move a to b/c , but b already exists as a file and not as a folder). Note that for a move operation $a \rightarrow b$, there is at most one file blocking the location b : we will call it the *blocker*.

If the blocker absent on Oscar, then we can just delete the blocker. However, if a blocker exists, then we might need to move it somewhere else before we solve the move we are interested in. This move itself might have a blocker, and so on. It seems that we just need to continue until we reach a move which has no blocker or whose blocker can be deleted, but we can get caught in a cycle: if we must move a to b , b to c and c to a , then we will not be able to perform the operations without using a temporary location.

How can we perform the moves? A simple way would be to move each file to a unique temporary location and then rearrange files to our liking: however, this performs many unnecessary moves and could lead to problems if the program is interrupted. We can do something more clever by performing a decomposition in Strongly Connected Components (SCC) of the *move graph* (with one vertex per file and one edge per move operation going from the file to its blocker or to its destination if no blocker exists). The computation of the SCC decomposition is simplified by the observation that because two files being moved to the same destination must be equal, we can only keep one arbitrary in-edge per node, and look at the graph pruned in this fashion: its nodes have in-degree at most one, so the strongly connected components are either single nodes or cycles. Once the SCC decomposition is known, the moves can be applied by applying each SCC in a bottom-up fashion, an SCC's moves being solved either trivially (for single files) or using one intermediate location (for cycles).

The detailed algorithm is implemented as two mutually recursive functions and presented as Algorithm 3.

An optimization implemented by **btrfsync** over the algorithm described here is to move files instead of copying them and then remove the original file. Moves are faster than copies on most filesystems as the OS does not need to copy the actual file contents to perform moves.

5.2 Experimental Comparison to **rsync**

We compared **rsync**⁸ and our Divide & Factor implementation (called **btrfsync**) under the following experimental conditions:

Test Directories: The directories used for transmission and time comparisons are described in Table 3.

Command-Line Options: **rsync** was called with the following options, for the reasons below:

- ▶ **--delete** to delete existing files on Oscar which do not exist on Neil like **btrfsync** does.

- ▶ **-I** to ensure that **rsync** did not cheat by looking at file modification times (which **btrfsync** does not do).

- ▶ **--chmod="a=rx,u+w"** in an attempt to disable the transfer of file permissions (which **btrfsync** does not transfer). Although these settings ensure that **rsync** does not need to transfer permissions, verbose logging suggests that it does transfer them anyway, so **rsync** must lose a few bytes per file as compared to **btrfsync** for this reason.

⁸ **rsync** version 3.0.9, used both as a competitor to benchmark against and as an underlying call in our own code.

Algorithm 3 Perform Moves

Require: \mathfrak{D} is a dictionary where $\mathfrak{D}[f]$ denotes the intended destinations of f

```

1:  $M \leftarrow []$ 
2:  $T \leftarrow []$ 
3: for  $f$  in  $\mathfrak{D}$ 's keys do
4:    $M[f] \leftarrow \text{not\_done}$ 
5:   function UNBLOCK_COPY( $f, t$ )
6:     if  $t$  is blocked by some  $b$  then
7:       if  $b$  is not in  $\mathfrak{D}$ 's keys then
8:          $\text{unlink}(b)$  ▷ We don't need  $b$ 
9:       else
10:         $\text{RESOLVE}(b)$  ▷ Take care of  $b$  and make it go away
11:     if  $T[f]$  was set then
12:        $f \leftarrow T[f]$ 
13:      $\text{copy}(f, d)$ 
14:   function RESOLVE( $f$ )
15:     if  $M[f] = \text{done}$  then
16:       return ▷ Already managed by another in-edge
17:     if  $M[f] = \text{doing}$  then
18:        $T[f] \leftarrow \text{mktemp}()$ 
19:        $\text{move}(f, T[f])$ 
20:        $M[f] \leftarrow \text{done}$ 
21:     return ▷ We found a loop, moved  $f$  out of the way
22:    $M[f] \leftarrow \text{doing}$ 
23:   for  $d \in \mathfrak{D}[f]$  do
24:     if  $d \neq f$  then
25:        $\text{unblock\_copy}(f, d)$  ▷ Perform all the moves
26:   if  $f \notin \mathfrak{D}[f]$  and  $T[f]$  was not set then
27:      $\text{unlink}(f)$ 
28:   if  $T[f]$  was set then
29:      $\text{unlink}(T[f])$ 
30: for  $f$  in  $\mathfrak{D}$ 's keys do
31:    $\text{RESOLVE}(f)$ 

```

► **-v** Transmission accounting was performed by calling `rsync` with the `-v` flag (which reports the number of sent and received bytes). For `btrfsync` we added a piece of code counting the amount of data transmitted during `btrfsync`'s own negotiations.

Network Configuration: Experiments were performed without any network transfer, by synchronizing two folders on the same host. Hence, time measurements should mostly represent the CPU cost of the synchronization.

Results: Results are given in Table 2. In general, `btrfsync` spent more time than `rsync` on computation (especially when the number of files is large, which is typically seen in the experiments involving `synthetic`). Transmission results, however, turn out to be favorable to `btrfsync`.

In the trivial experiments where either Oscar or Neil have no data at all, `rsync` outperforms `btrfsync`. This is especially visible when Neil has no data: `rsync` immediately notices that there is nothing to transfer, but `btrfsync` engages in information transfers to determine the symmetric difference.

On non-trivial tasks, however, `btrfsync` outperforms `rsync`. This is the case of the `synthetic` datasets, where `btrfsync` does not have to transfer information about all unmodified files, and even more so in the case where there are no modifications at all. For Firefox source code datasets, `btrfsync` saves a very small amount of bandwidth, presumably because of unmodified files. For the `btrfsync` source code dataset, we notice that `btrfsync`, unlike `rsync`, was able to detect the move and avoid retransferring the moved folder.

6 Conclusion and Further Improvements

The main contributions of our work are:

- We present the novel “Divide & Factor” protocol for set reconciliation, which is based on number theory and is optimal with respect to transfer size.
- We study the problem of set reconciliation of directories of files. We discuss the optimal size of message digests in this setting, as well as a move resolution algorithm to reproduce a directory structure.
- We present `btrfsync`, an open source implementation of the “Divide & Factor” protocol.
- We demonstrate the usability of this implementation through benchmarks on synthetic and real-world tasks, and show that `btrfsync` exchanges less data than the popular software `rsync`.
- The optimizations presented in this paper apply to [] as well.

Many fine questions of the probabilistic discussions in the paper are left as future work. Another further line of research would be to pursue development of `btrfsync` to make it suitable for end users.

*Vérifier les
claims de
cette liste, et
en parler dès
l'intro.*

*TODO
résoudre ce
point
Be more
specific!*

7 Acknowledgment

The authors acknowledge Guillaïn Potron for his early involvement in this research work.

8 ToDo

- Pourquoi on a deux biblios?! Laquelle est la bonne ?
- Merge two reference files rsynch and wagner.
- @Fabrice: Fig. 1 pas clair, je pense qu'il faut mieux créer une notation μ gcd car en fait la FFT ne modifie que le μ ... Et on pourrait plutôt indiquer dans ce tableau la complexité (avec μ et μ gcd) en utilisant l'optimisation décrite ou sans utiliser l'optimisation décrite). N'oublions pas de préciser que les temps des expériences d'Antoine comptent aussi le tirage des nombres premiers (qui doit être négligeable peut-être dans notre cas, je ne me rappelle plus...)
- @Fabrice: Pour éviter le cas empty \rightarrow source trop gros, on pourrait imaginer l'astuce suivante: si jamais Neil la taille de c est plus petite que la taille du produit des nombres premiers $p_1 \dots p_n$ utilisés, Neil envoie un message pour l'indiquer, et on arrête là le protocole. Et Oscar peut directement factoriser ce nombre envoyé...
- @Fabrice: il faut discuter de la taille de la taille maximale des "petits" premiers utilisés pour les variantes de p et montrer que cela n'enlève pas trop d'entropie pour les h_i . Encore une fois, je m'en occupe la semaine prochaine si besoin.
- Refaire une dernière fois les expériences, vu que Fabrice a significativement amélioré les perfs.
- Faire clarifier par Fabrice l'histoire du doublement...
- Fabrice: comparer nos perfs avec <http://ipsit.bu.edu/programs/reconcile/> ? ou pas ?
- Prendre en compte les remq suivants de Fabrice

- je pense que l'événement `bottom_bandwidth` est ou bien l'événement "on n'a pas réussi à synchroniser les fichiers"

- je pense que tu essayes de traiter 2 cas ensembles: `modulus` trop petit ($t_0 > t$) et `modulus` trop grand ($t_0 < t$)

- dans l'intro, au lieu de "This article proposes a new reconciliation protocol based on asymptotic transmission complexity, the proposed procedure reaches optimality as $n \rightarrow \infty$ ", j'aurai dit quelque chose comme:

"In this article, we first propose a new reconciliation protocol based on number theory. Then, we show how to apply this set reconciliation protocol to efficiently synchronize files. We show that the naive method of hashing the files and synchronizing these hashes can be efficiently optimized. First, we show that the size of the hashes can be carefully optimized to reduce the bandwidth. Then, we show that a protocol to detect moves of files and carefully use this information can be designed. Finally, we run some practical experiments and compare ... rsync ..."

References

1. Amarilli, A., Ben Hamouda, F., Bourse, F., Morisset, R., Naccache, D., Rauzy, P.: <https://github.com/RobinMorisset/Btrsnc>
2. Bleichenbacher, D., Nguyen, P.Q.: Noisy polynomial interpolation and noisy chinese remaindering. In: *Advances in Cryptology – Proceedings of Eurocrypt 2000*. pp. 53–69. Springer-Verlag (2000)
3. Boneh, D.: Finding smooth integers in short intervals using crt decoding. In: *Proceedings of the 32nd Annual ACM Symposium on Theory of Computing*. pp. 265–272 (2000)
4. Burnikel, C., Ziegler, J., Stadtwald, I., D-Saarbrücken: Fast recursive division (1998)
5. Eppstein, D., Goodrich, M., Uyeda, F., Varghese, G.: What’s the difference?: efficient set reconciliation without prior context. In: *ACM SIGCOMM Computer Communication Review*. vol. 41, pp. 218–229. ACM (2011)
6. Fouque, P.A., Stern, J., Wackers, J.G.: Cryptocomputing with rationals. In: Blaze, M. (ed.) *Financial Cryptography*. Lecture Notes in Computer Science, vol. 2357, pp. 136–146. Springer (2002)
7. Hofmann, M., Pierce, B.C., Wagner, D.: Edit lenses. In: Field, J., Hicks, M. (eds.) *POPL*. pp. 495–508. ACM (2012)
8. Minsky, Y., Trachtenberg, A.: Scalable set reconciliation. In: *40th Annual Allerton Conference on Communications, Control and Computing* (October 2002), a full version can be downloaded from <http://ipsit.bu.edu/documents/BUTR2002-01.ps>
9. Minsky, Y., Trachtenberg, A., Zippel, R.: Set reconciliation with nearly optimal communication complexity. *Information Theory, IEEE Transactions on* 49(9), 2213–2218 (2003)
10. Pan, V., Wang, X.: On rational number reconstruction and approximation. *SIAM Journal on Computing* 33, 502 (2004)
11. Schönhage, A., Strassen, V.: Schnelle multiplikation grosser zahlen. *Computing* 7(3), 281–292 (1971)
12. Tridgell, A.: Efficient algorithms for sorting and synchronization. Ph.D. thesis, PhD thesis, The Australian National University (1999)
13. Vallée, B.: Gauss’ algorithm revisited. *J. Algorithms* 12(4), 556–572 (1991)
14. Wang, X., Pan, V.: Acceleration of Euclidean algorithm and rational number reconstruction. *SIAM Journal on Computing* 32(2), 548 (2003)

A Extended Protocol

First phase during which Neil amasses modular information on the difference	
Oscar	Neil
	start the protocol with p_1
	$\xrightarrow{c_1}$
	computes a, b using p_1 if a factors properly then go to Final Phase else perform the protocol with p_2
	$\xrightarrow{c_2}$
	computes $c \bmod p_1 p_2 = \text{CRT}_{p_1, p_2}(c_1, c_2)$ computes a, b using $p_1 p_2$ if a factors properly then go to Final Phase else perform the protocol with p_3
	$\xrightarrow{c_3}$
	computes $c \bmod p_1 p_2 p_3 = \text{CRT}_{p_1, p_2, p_3}(c_1, c_2, c_3)$ computes a, b using $p_1 p_2 p_3$ if a factors properly then go to Final Phase else perform the protocol with p_4
	\vdots
Final Phase	
	Let $\mathfrak{S} = \{F'_i \text{ s.t. } a \bmod h'_i = 0\}$
	$\xleftarrow{\mathfrak{S}, b}$
deletes files s.t. $b \bmod h_i = 0$ adds \mathfrak{S} to the disk	

Note that parties do not need to store the p_i 's in full. Indeed, the p_i s could be subsequent primes sharing their most significant bits. This reduces storage per prime to a very small additive constant $\cong \ln(p_i) \cong \ln(2^{2tu+2}) \cong 1.39(tu + 1)$ of about $\log_2(tu)$ bits.

B Power of Two Protocol

In this variant Oscar computes c in \mathbb{N} :

$$c = \prod_{F_i \in \mathfrak{F}} \text{HashPrime}(F_i) = \prod_{i=1}^n h_i \in \mathbb{N}$$

and considers $c = \bar{c}_{n-1}|\dots|\bar{c}_2|\bar{c}_0$ as the concatenation of n successive u -bit strings. Again, we omit the treatment of \perp s for the sake of clarity.

First phase during which Neil amasses modular information on the difference	
Oscar computes $c \in \mathbb{N}$	Neil
$\xrightarrow{\bar{c}_0}$	computes a, b modulo 2^u if a factors properly then go to Final Phase else request next chunk \bar{c}_1
$\xrightarrow{\bar{c}_1}$	construct $c \bmod 2^{2u} = \bar{c}_1 \bar{c}_0$ computes a, b modulo 2^{2u} if a factors properly then go to Final Phase else request next chunk \bar{c}_2
$\xrightarrow{\bar{c}_2}$	construct $c \bmod 2^{3u} = \bar{c}_2 \bar{c}_1 \bar{c}_0$ computes a, b modulo 2^{3u} if a factors properly then go to Final Phase else request next chunk \bar{c}_3
\vdots (for $2t$ rounds)	
Final Phase	
deletes files s.t. $b \bmod 2^{2tu} = 0$ adds \mathfrak{S} to the disk	Let $\mathfrak{S} = \{F'_i \text{ s.t. } a \bmod 2^{2tu} = 0\}$ $\xleftarrow{\mathfrak{S}, b}$

C Hashing Into Primes

Hashing into primes is frequently needed in cryptography. A recommended implementation of $\text{HashPrime}(F)$ is given in Algorithm 5. If u is large enough (e.g. 160) one might sacrifice uniformity to avoid repeated file hashings by defining $\text{HashPrime}(F) = \text{NextPrime}(\text{Hash}(F))$. Yet another acceleration (that further destroys uniformity) consists in replacing NextPrime by Algorithm 4 where $\alpha = 2 \times 3 \times 5 \times \dots \times \text{Prime}[d]$ is the product of the first primes until some rank d .

analyse
ci-contre

Fabrice: Cela accélère un peu, car il y a environ $\frac{n}{\log(n)\varphi(\alpha)}$ nombres premiers $\leq n$ et congrus à 1 modulo α , contre $\frac{n}{\log(n)}$ nombres premiers $\leq n$ (voir http://fr.wikipedia.org/wiki/Th%C3%A9or%C3%A8me_de_la_progression_arithm%C3%A9tique#Version_quantitative).

Dans l'algo 4, h est donc premier avec proba $\frac{\frac{n}{\log(n)\varphi(\alpha)}}{\frac{n}{\alpha}} = \frac{\alpha}{\log(n)\varphi(\alpha)}$, tandis que dans l'algo 5, h est premier avec proba $\frac{1}{\log(n)}$. On a alors une accélération d'environ 10 si on prend $d = 60$.

Algorithm 4 Fast Nonuniform Hashing Into Primes

```

1:  $h = \alpha \left\lfloor \frac{\text{Hash}(F)}{\alpha} \right\rfloor + 1$ 
2: while  $h$  is composite do
3:    $h = h - \alpha$ 
4: return  $h$ 
```

Algorithm 5 Possible Implementation of HashPrime(F)

```

1:  $i = 0$ 
2: repeat
3:    $h = 2 \cdot \text{Hash}(F|i) + 1$ 
4:    $i = i + 1$ 
5: until  $h$  is prime
6: return  $h$ 
```

*Antoine sur
quelle
plate-formes
ont été
obtenus les
résultats ex-
perimentaux?
quelles
vitesses de
processeur
etc?*

Entities and Datasets		Transmission (Bytes)						Time (s)	
Neil's \mathfrak{F}'	Oscar's \mathfrak{F}	TX_{rs}	RX_{rs}	TX_{bt}	RX_{bt}	$\delta_{rs} - \delta_{bt}$	$\frac{\delta_{bt}}{\delta_{rs}}$	time_{rs}	time_{bt}
source	empty	1613	778353	1846	788357	10237	1.0	0.2	7.7
empty	source	11	29	12436	6120	18516	463.9	0.1	5.5
empty	empty	11	29	19	28	7	1.2	0.1	0.3
synthetic	synthetic_shuffled	24891	51019	3638	4147	-68125	0.1	0.2	26.8
synthetic_shuffled	synthetic	24701	50625	3443	3477	-68406	0.9	0.2	26.6
synthetic	synthetic	25011	50918	327	28	-75574	0.0	0.1	25.7
firefox-13.0.1	firefox-13.0	90598	28003573	80895	27995969	-17307	1.0	2.6	4.2
source_moved	source	2456	694003	1603	1974	-692882	0.0	0.2	2.5

Table 2. Experimental results. **rs** and **bt** subscripts respectively denote **rsync** and **btrfsync**. The two first columns indicate the datasets. Synchronization is performed *from* Neil to Oscar. **RX** and **TX** denote the quantity of received and sent bytes and $\delta_{\square} = \text{TX}_{\square} + \text{RX}_{\square}$. $\delta_{rs} - \delta_{bt}$ and δ_{bt}/δ_{rs} express the absolute and the relative differences in transmission between the two programs. The last two columns show timing results.

Directory	Description
synthetic	A directory containing 1000 very small files containing the numbers 1, 2, ..., 1000.
synthetic_shuffled	synthetic with: 10 deleted files 10 renamed files 10 modified files
source	A snapshot of btrfsync 's own source tree
source_moved	source with one big folder (a few megabits) renamed.
firefox-13.0	The source archive of Mozilla Firefox 13.0.
firefox-13.0.1	The source archive of Mozilla Firefox 13.0.1
empty	An empty folder.

Table 3. Test Directories.

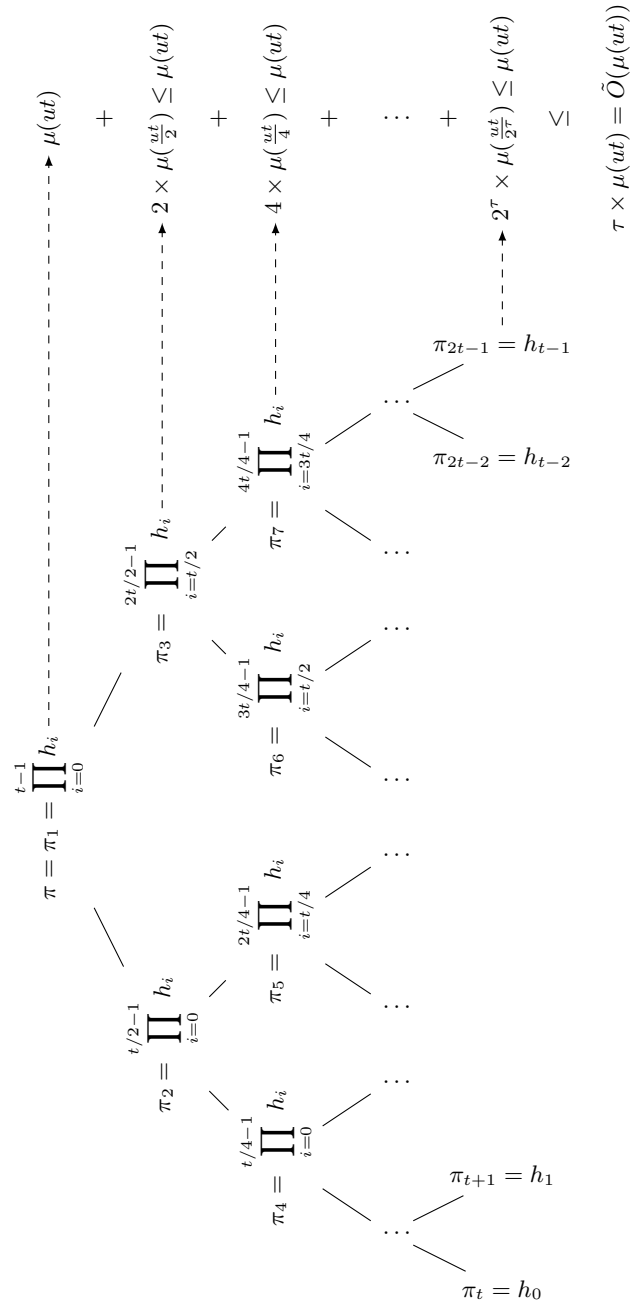


Fig. 5. Product Tree

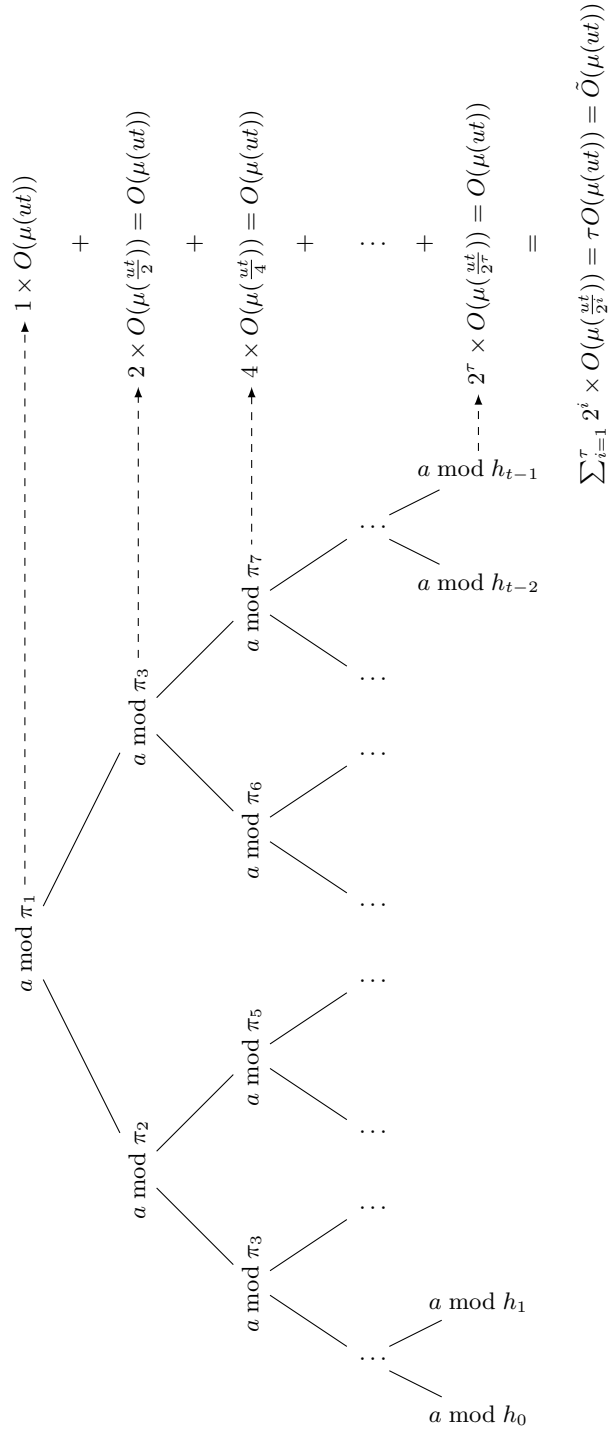


Fig. 6. Modular Reduction From Product Tree