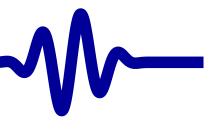


Danie

6. Transmitting
Data Words

Gisselquist Technology, LLC

Daniel E. Gisselquist, Ph.D.





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Debugging is one of the hardest parts of digital logic design

- You can't see what's happening inside the FPGA
- LED's are one solution
  - FPGA's operate 50MHz+
  - Your eye operates at < 60Hz
- The serial port can be a second solution

Let's learn to send data through our serial port! Objectives

- Transform Hello World into a debugging output
- Learn about formal abstraction
- Experiment with using neurses with Verilator
- Extract internal design variables from within Verilator



### **Data Transmitter**

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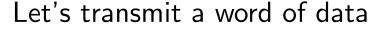
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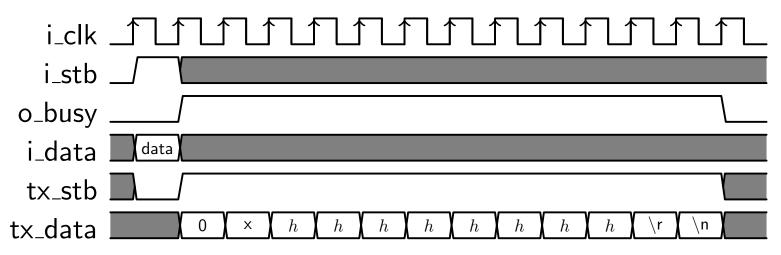
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Each word will . . .

- Start with 0x
- Contain the number sent, but in hexadecimal

this is much easier than doing decimal!

Four bits can be encoded at a time

End with a carriage return / line-feed pair



### **Data Transmitter**

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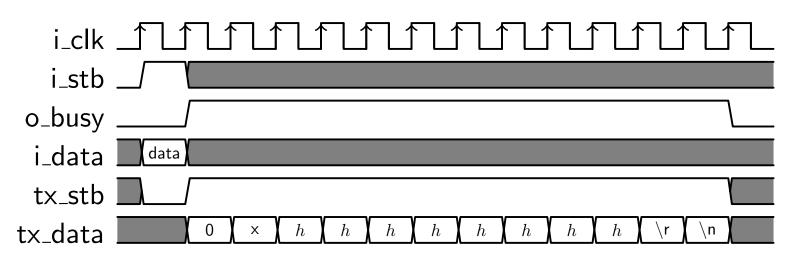
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You should know how to build this design already



Remember how we've built state machines before

- In this case, you have two triggers
  - One trigger, i\_stb, starts the process
  - A busy line from the serial port, tx\_busy (not shown), controls the movement from one character to the next
- This design will be the focus of this lesson



### **Data Transmitter**

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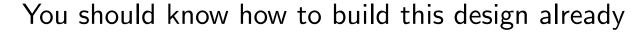
Verilator data

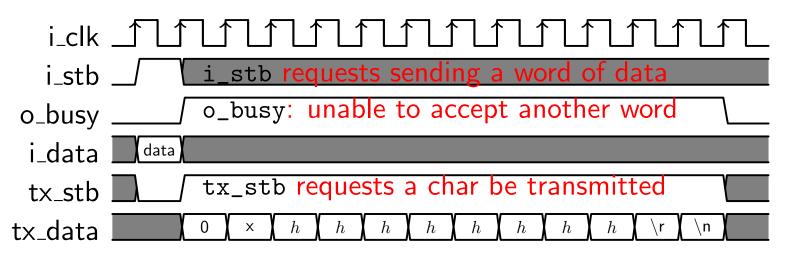
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Remember how we've built state machines before

- In this case, you have two triggers
  - One trigger, i\_stb, starts the process
  - A busy line from the serial port, tx\_busy (not shown),
     controls the movement from one character to the next
- This design will be the focus of this lesson



#### **Desired Structure**



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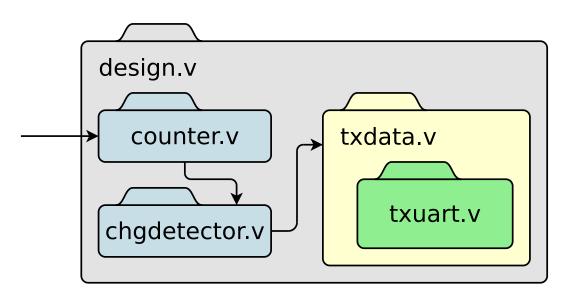
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Our overall design will look like this:



- Some event will trigger a counter
- A second module will detect that the counter has changed
- Finally we'll output the result
- We'll use txuart.v from the last exercise

Let's take a quick look at counter.v and chgdetector.v



# Creating a Counter



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You should already know how to make an event counter

Feel free to add a reset if you would like





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Conclusion

Detecting a change in the counter is also pretty easy

```
chgdetector(i_clk, i_data,
module
                 o_stb, o_data, i_busy);
        // ...
         initial \{ o_stb, o_data \} = 0;
        always @(posedge i_clk)
        if (!i_busy)
        begin
                 stb \ll 0:
                 if (o_data != i_data)
                 begin
                          stb \ll 1'b1:
                          o data <= i data:
                 end
        end
endmodule
```





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Conclusion

Detecting a change in the counter is also pretty easy

Nothing is allowed to change if i\_busy is true. That's the case where a request has been made, but it has yet to be accepted.

end

endmodule





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Detecting a change in the counter is also pretty easy

```
chgdetector(i_clk, i_data,
module
                  o_stb, o_data, i_busy);
         initial \{ o_stb, o_data \} = 0;
         always @(posedge i_clk)
         if (!i_busy)
         begin
                  stb \ll 0:
                  if (o_data != i_data)
                  begin
                           stb \ll 1'b1:
                           o_data <= i_data;
             Otherwise, anytime the data changes, we set up
         end
             a request to transmit the new data.
endmodule
```





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Conclusion

What formal properties might we use here?

Any output value should remain unchanged until accepted

```
// Remember this property?
always @(posedge i_clk)
if ((f_past_valid)
        &($past(o_stb))&&($past(i_busy)))
        assert((o_stb)&&($stable(o_data)));
```

Remember how this works? This says that . . .





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if ((f_past_valid)
    &($past(o_stb))&&($past(i_busy)))
    assert((o_stb)&&($stable(o_data)));
```

Remember how this works? This says that . . .

- If both o\_stb and i\_busy are true on the same clock cycle (i.e., the interface is stalled)
- Then request should remain outstanding on the next cycle
- ...and the data should be the same on that next cycle
- \$stable(o\_data) is shorthand for o\_data == \$past(o\_data)





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// Remember this property?
always @(posedge i_clk)
if ((f_past_valid)
        &($past(o_stb))&&($past(i_busy)))
        assert((o_stb)&&($stable(o_data)));
```

When o\_stb rises, o\_data should reflect the input

```
always @(posedge i_clk)
if ((f_past_valid)&&($rose(o_stb)))
    assert(o_data == $past(i_data));
```

**\$rose**(o\_stb) is shorthand for (o\_stb[0] && !**\$past**(o\_stb[0]))





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What formal properties might we use here?

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        assert(o_data == $past(i_data));
```

**\$rose**(o\_stb) is shorthand for (o\_stb[0] && !**\$past**(o\_stb[0]))

Can you think of any other properties we might need?





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This lesson will focus on txdata.v

- We've already built txuart.v
- You should have no problems designing counter.v or chgdetector.v
  - You are encouraged to do so on your own
  - If not, you can find counter.v and chgdetector.v in the course handouts

You should also have a good idea how to start on txdata.v.

- It's not all that different from txuart.v or helloworld.v
- The example in the course handouts is broken





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Here's the port list(s) we'll design to





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Here's the port list(s) we'll design to

- If i\_stb is true, we have a new value to send
- i\_data will then contain that 32-bit value
- o\_busy means we cannot accept data
- o\_uart\_tx is the 1-bit serial port output





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Conclusion

Here's the port list(s) we'll design to

- tx\_stb requests data be transmitted
- tx\_data is the 8-bit character to transmit
- tx\_busy means the serial port transmitter is busy and cannot accept data



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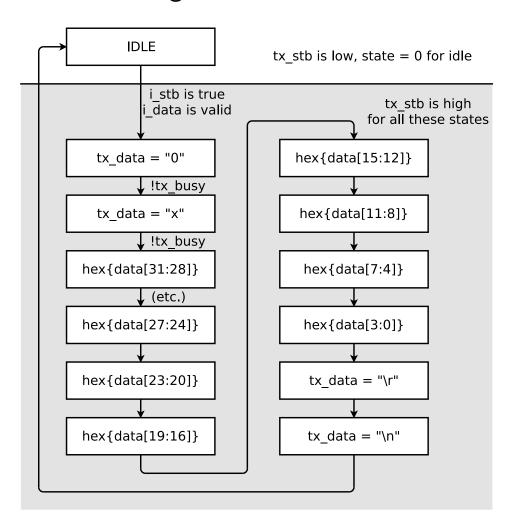
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Conclusion

We can create a state diagram for this state machine too





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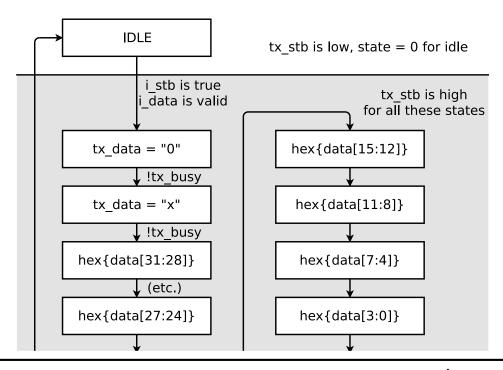
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We can create a state diagram for this state machine too



We'll start sending our message upon request (i\_stb is true), and advance to the next character any time the transmitter is not busy (tx\_busy is false)



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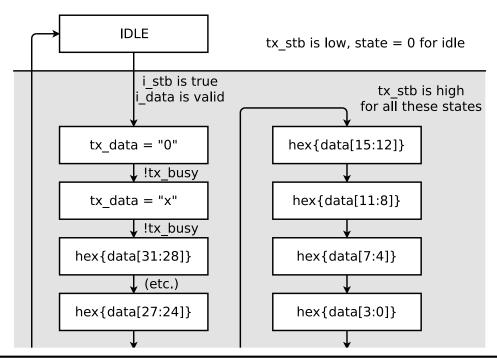
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We can create a state diagram for this state machine too



In this chart, data is the 32-bit word we are sending, and hex{} just references the fact that we need to convert the various nibbles to hexadecimal before outputting them



 $\mathcal{M}$ 

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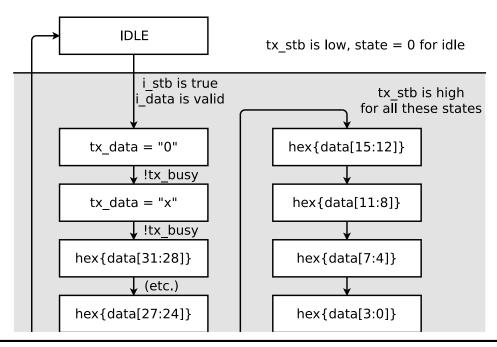
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We can create a state diagram for this state machine too



Remember, input data such as i\_data are only valid as long as the incoming request is valid (i\_stb is high). We'll need to make a copy of that data once the request is made, (i\_stb) && (!o\_busy), and then work off of that copy.





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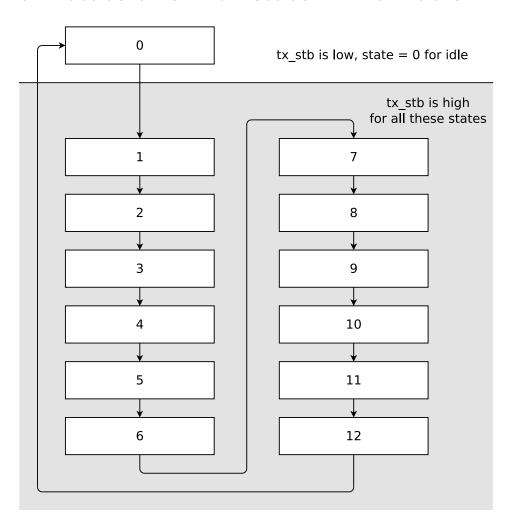
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We can even annotate this with state ID numbers







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The state machine should remind you of helloworld.v

```
always @(posedge i_clk)
if (!o_busy)
begin
        if (i_stb)
        begin
                 state <= 1:
                 tx_stb \ll 1:
        end // else state already == 0
end else if ((tx_stb)&&(!tx_busy))
begin
        state \le state + 1:
        if (state >= 4'hd)
        begin
                 tx_stb \ll 1'b0;
                 state \le 0:
```



# **Outgoing Data**



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The outgoing data is just a shift register

#### Question:

Why aren't we conditioning our load on i\_stb as well?



# **Outgoing Data**



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Converting to hex is very straight forward

```
always @(posedge i_clk)
case(sreg[31:28])
4'h0: hex <= "0":
4'h1: hex <= "1";
4'h2: hex <= "2";
4'h3: hex <= "3":
4'h9: hex <= "9";
4'ha: hex <= "a":
4'hb: hex <= "b";
4'hc: hex <= "c":
4'hd: hex <= "d";
4'he: hex <= "e":
4'hf: hex <= "f";
default: begin end
endcase
```



# **Outgoing Data**



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Put together, here's our code to transmit a byte

```
always @(posedge i_clk)
case(state)
if (!tx_busy)
        case(state)
        4'h1: tx_data \le "0"; // These are the
        4'h2: tx_data <= "x"; // values we'll
        4'h3: tx_data <= hex; // want to output
        4'h4: tx_data <= hex; // at each state
        // ...
        4'h9: tx_data \le hex;
        4'ha: tx_data <= hex;
        4'hb: tx_data \ll " \ r";
        4'hc: tx_data \ll " n";
        default: tx_data <= "Q"; // a bad value</pre>
        endcase
```



### **Simulation**



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Let's do simulation after formal verification

- It's easier to get a trace from formal
- Formal methods are often done faster
- etc.





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Conclusion

Our design is getting large

- We've already verified txuart.v
- It would be nice not to have to do it again

Let's simplify things instead!

- Let's replace txuart.v with something that ...
  - Might or might not act like txuart.v





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Let's simplify things instead!

- Let's replace txuart.v with something that ...
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  - ...at the solver's discretion





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Our design is getting large

- We've already verified txuart.v
- It would be nice not to have to do it again

Let's simplify things instead!

- Let's replace txuart.v with something that ...
  - Might or might not act like txuart.v
  - ...at the solver's discretion
  - Acting like txuart.v must remain a possibility

This is called abstraction





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- (\* anyseq \*) allows the solver to pick the values of serial\_busy and serial\_out
- (\* anyseq \*) values can change from one clock to the next
- They might match what txuart would've done





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Conclusion

- (\* anyseq \*) allows the solver to pick the values of serial\_busy and serial\_out
- (\* anyseq \*) values can change from one clock to the next
- They might match what txuart would've done, or they might not





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Conclusion

- (\* anyseq \*) allows the solver to pick the values of serial\_busy and serial\_out
- (\* anyseq \*) values can change from one clock to the next
- They might match what txuart would've done, or they might not
- If our design passes in spite of what this abstract txuart does





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```
"ifndef FORMAI.
    txuart #(UART_SETUP[23:0]) txuarti(i_clk,
        tx_stb, tx_data, o_uart_tx, tx_busy);
'else
        (* anyseq *) wire serial_busy, serial_out;
        assign o_uart_tx = serial_out;
        assign tx_busy = serial_busy;
```

- (\* anyseq \*) allows the solver to pick the values of serial\_busy and serial\_out
- (\* anyseq \*) values can change from one clock to the next
- They might match what txuart would've done, or they might not
- If our design passes in spite of what this abstract txuart does, then it will pass if txuart acts like it should



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We'll insist that our abstract UART is busy following any request

We can use f\_minbusy to force any transmit request to take at least four cycles before dropping the busy line



## **Formal Verification**

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Conclusion

We'll insist that our abstract UART is busy following any request

```
[1:0] f_minbusy;
reg
initial f_minbusy = 0;
always @(posedge i_clk)
if ((tx_stb)&&(!tx_busy))
         f_{minbusy} \ll 2'b01;
else if (f_minbusy != 2'b00)
         f_minbusy <= f_minbusy + 1'b1;</pre>
always \mathbb{Q}(*)
if (f_minbusy != 0)
        assume(tx_busy);
```

Since (\* <a href="mailto:anyseq">anyseq</a> \*) values act like inputs to our design, constraining them by an assumption is appropriate



### **Formal Verification**



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Conclusion

We'll also insist it doesn't become busy on its own

Now we can build a proof without re-verifying txuart.v!





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Let's see if this design works:

This would yield a trace with a reset

It works, but it's not very informative





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What if we except the reset?

A trace that yields all zeros is boring

It works, but it's not very informative





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What if we look for  $0x12345678\r\n$ ?

```
reg f_seen_data;
initial f_seen_data = 0;
always @(posedge i_clk)
if (i_reset)
        f seen data <= 1'b0:
else if ((i_stb)&&(!o_busy)
                \&\&(i_data = 32'h12345678))
        f seen data <= 1'b1:
always @(posedge i_clk)
if ((f_past_valid)&&(!$past(i_reset))
                \&\&(f_seen_data))
        cover($fell(o_busy));
```

Check out the trace.





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What if we look for  $0x12345678\r\n$ ?

```
reg f_seen_data;
initial f_seen_data = 0;
always @(posedge i_clk)
if (i_reset)
        f seen data <= 1'b0:
else if ((i_stb)&&(!o_busy)
                \&\&(i_data = 32'h12345678))
        f seen data <= 1'b1:
always @(posedge i_clk)
if ((f_past_valid)&&(!$past(i_reset))
                \&\&(f_seen_data))
        cover($fell(o_busy));
```

Check out the trace. Does your design work?





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What if we look for  $0x12345678\r\n$ ?

```
f seen data:
reg
initial f_seen_data = 0;
always @(posedge i_clk)
if (i_reset)
        f_seen_data <= 1'b0;
else if ((i_stb)&&(!o_busy)
                \&\&(i_data = 32'h12345678))
        f_seen_data <= 1'b1;
```

Caution: It's a snare to use something like f\_seen\_data outside of a cover context

- We aren't doing directed simulation
- The great power of formal is that it applies to all inputs
- We're just picking an interesting input for a trace



### **Assertions**



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Now, what assertions would be appropriate?

- We can assert state is legal
- $\Box$  That tx\_stb != (state == 0)
- Can we assert that the first data output is a "0"?
- That the second output is a "1"?

Your turn: what would make the most sense here?



## Sequence



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Yes, we can assert a sequence takes place!

```
[12:0] f_p1reg; // First property
reg
initial f_p1reg = 0;
always @(posedge i_clk)
if (i_reset)
        f_p1reg <= 0;
else if ((i_stb)&&(!o_busy))
begin
        f_p1reg \ll 1;
        assert(f_p1reg == 0);
end else if (!tx_busy)
        f_p1reg <= { f_p1reg[11:0], 1'b0 };
```

 $f_p1reg[x]$  will now be true for stage x of any output sequence



## Sequence

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Conclusion

Now we can make assertions about the different states in our sequence

```
always @(posedge i_clk)
if ((!tx\_busy)||(f\_minbusy == 0))
begin
        // If the serial port is ready for
        // the next character, or while we are
        // waiting for the next character, ...
        if (f_p1reg[0])
                assert((tx_data = "0")
                         &&(state == 1));
        if (f_p1reg[1])
                assert((tx_data == "x")
                         &&(state == 2));
        // etc.
end
```



### **Concurrent Assertions**



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Full System Verilog support would make this easier

This defines a sequence where

- (tx\_stb)&&... must be true
- while tx\_busy is true, and then
- until (and including) the clock where tx\_busy is false



## Sequence



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Conclusion

Full System Verilog support would make this easier

```
sequence SEND(A,B);
// ....
```

We could then string these sequences together in a **property** that could be asserted

```
assert property (@(posedge i_clk))
    disable iff (i_reset)
        (i_stb)&&(!o_busy)
        |=> SEND(1, "0") // First state
        ##1 SEND(2, "x") // Second, etc
        // ...
```

- $\blacksquare$  A  $\mid =>$  B means if A, then B is asserted true on the next clock
- ##1 here means one clock later



## Sequence



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Full System Verilog support would make this easier

```
sequence SEND(A,B);
// ....
```

We could then string these sequences together in a **property** that could be asserted

```
assert property (@(posedge i_clk))
    disable iff (i_reset)
        (i_stb)&&(!o_busy)
        |=> SEND(1, "0") // First state
        ##1 SEND(2, "x") // Second, etc
        // ...
```

SymbiYosys support for sequences requires a license



# Exercise #1



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Your turn!

Take a moment now to . . .

- Create your txdata.v, or
- Download my broken one, and then
- Formally verify it
  - Add such assertions as you deem fit
  - Make sure you get a trace showing it working

Does your design work?



## **Simulation**



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Let's move on to simulation

- Let's use the simulator to count key presses
- ncurses + Verilator offers a quick debugging environment
  - Every time a key is pressed, output a new count value
  - We'll use getch() to get key presses immediately

You may need to download and install neurses-dev

- We'll adjust uartsim() to print to the screen
- You can also examine internal register values with Verilator
  - While the design is running

Let's look at how we'd do these things





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Conclusion

ncurses is an old-fashioned text library

- It allows us easy access to key press information
- We can write to various locations of the screen
- □ etc.
- The original ZipCPU debugger was written with neurses

We'll only scratch the surface here





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Starting neurses requires some boilerplate

- This initializes the curses environment
- Turns off line handling and echo
- Decodes special keys (like escape) for us
- halfdelay (1) Doesn't wait for keypresses





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Our inner loop will start by checking for keypresses

```
do {
        done = false;
        tb -> m_core -> i_event = 0;
        // Ket a keypress
        chv = getch();
        if (chv == KEY_ESCAPE)
                 // Exit on escape
                 done = true;
        else if (chv != ERR)
                 // Key was pressed
                 tb -> m_core -> i_event = 1;
        tb->tick();
        (*uart)(tb->m_core->o_uart_tx);
 while(!done);
```





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We can speed this up too:

- $_{\square}$  getch() waits  $1/10^{th}$  of a second for a keypress
  - This is because we called **halfdelay** (1);
- This will run 1000 simulation ticks per getch() call





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Conclusion

We can also count given keypresses

```
do {
        // ...
        for(int k=0; k<1000; k++) {
                 tb->tick();
                 (*uart)(tb->m_core->o_uart_tx);
                 keypresses
                          += tb->m_core->i_event;
                 tb -> m_core -> i_event = 0;
} while(!done);
```

We'll print this number out before we are done





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Conclusion

We'll also need to replace the putchar() in uartsim.cpp

ncurses requires we use addch()

```
// if character received
if (m_rx_data != '\r')
        addch(m_rx_data);
```

- No flush is necessary, getch() handles that
- '\r' would clear our line, so we keep from printing it





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endwin() ends the ncurses environment

```
endwin();

printf("\n\nSimulation_complete\n");
printf("%4d_key_presses_sent\n", keypresses);
```

This is nice, but

wouldn't you also like a summary of keypresses the design counted?



### Verilator data



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Conclusion

Verilator maintains your entire design in a C++ object

- With a little work, we can find our variables
- A quick grep through Vthedesign.h reveals ...
- v\_DOT\_counterv contains our counter's value
  - Luse an older version of Verilator
  - Modern versions place this in **thedesign\_\_DOT\_\_counterv**
  - Supporting both requires a little work
- You can often find other values like this
  - Grep on your variables name
  - Be aware, Verilator will pick which of many names to give a value
  - Output wires may go by the name of their parent's value



### Verilator data

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Conclusion

This little adjustment will allow us to simplify the reference to our counter

```
#ifdef OLD_VERILATOR
#define VVAR(X) v__DOT_ ## A
#else
#define VVAR(X) thedesign__DOT_ ##A
#endif
#define counterv VVAR(_counterv)
```

- If OLD\_VERILATOR is defined (my old version)
  - counterv evaluates to v\_\_DOT\_\_counterv
- Otherwise counterv is replaced by
  - thedesign\_\_DOT\_\_counterv



### Verilator data



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Conclusion

We can now output our current counter



### Testbench build



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Conclusion

Two changes are required to our build script

- If you want to define NEW\_VERILATOR or OLD\_VERILATOR . . .
  - You'll need to do some processing on Verilator's version
  - The vversion.sh file does this, returning either
     DOLD VERILATOR or -DNEW VERILATOR
  - We can use this output in our g++ command line
  - Alternatively, you can just adjust the file for your version
- We need to reference -lncurses in our Makefile when building our executable



# Exercise #2



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#### Your turn!

Build and experiment with the simulation

- Using your txdata.v
- main() is found in thedesign\_tb.cpp in the handouts
- Experiment with . . .
  - Adjusting the number of tb->tick() calls between calls to getch()
  - Does this speed up or slow down your design?
  - Are all of your keypresses recognized?
  - What happens when you press the key while the design is busy?



# Exercise #3



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Conclusion

Only now is it time to test this in hardware

You'll need to test for button changes

- Does it work?
  - Does it count once per keypress?
  - Does the counter look reasonable?

My implementation experienced several anomalies.

We'll discuss those in the next lesson



### **Conclusion**



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What did we learn this lesson?

- How to formally verify a part of a design, and not just the leaf modules
- Creating interesting traces with cover
- Subtle timing differences can be annoying
- How to use Verilator with neurses
- Extracting an internal design value from within a Verilator simulation

We learned how to get information back out from within the hardware

We'll discuss the hazards of asynchronous inputs more in the next lesson