



- S takes 5 steps to reach node 5
- F takes $10 = 2k$ steps to reach node #...

$$K + K \text{ MOD } L = 5 + 2 = 7$$

- when S is at position k , F is $k \text{ mod } L$ steps ahead along the circle.
- when will they meet along the circle? when their positions along the circle are equal

"S POSITION"

"F POSITION"

$$0 + X \cdot S = k \text{ mod } L + 2 \cdot X \cdot S$$

$$X \cdot S = -(k \text{ mod } L) \quad \text{FIRST MEETING}$$

∴ The meeting point will be $k \text{ mod } L$ steps behind of the start of the loop

• I.E. the meeting point will be $L - k \text{ mod } L$ steps ahead the meeting point.

• S is moved to HEAD, and F remains at the first meeting point.

• S is now k steps away from the loop start while F is $k \text{ mod } L$ steps away. If $k < L$, then $k = k \text{ mod } L$ and the two pointers advance simultaneously, 1 step at a time to the loop start. If $k > L$, then F must loop around the L -step loop after traveling $k \text{ mod } L$ steps. For F and S to meet, F must do this loop an integral number of times. No fractional loops are possible.

SECOND MEETING:

$$K = k \text{ mod } L + L \cdot C \quad \begin{matrix} \text{CONST} \\ \text{INT} \end{matrix}$$

EXAMPLE 1:

$$5 = 5 \text{ mod } 3 + 3 \cdot C$$

$$5 = 2 + 3 \cdot C \Rightarrow C = 1 \Rightarrow$$

F loops around once after traveling $5 \text{ mod } 3$ steps.

EXAMPLE 2:

$$3 = 3 \text{ mod } 8 + 8 \cdot C$$

$$3 = 3 + 8 \cdot C$$

$$\Rightarrow C = 0$$

⇒ F does not need to loop around at all before meeting S