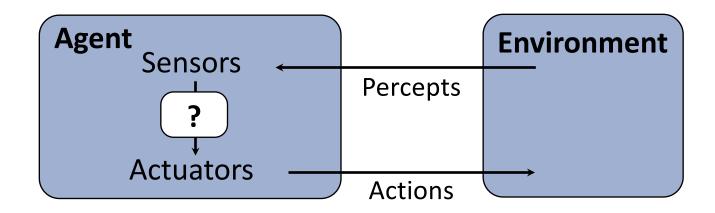
### Announcements

- HW1 is out, due on **Friday, September 29**, 11:59 PM
  - Topic: search algorithms to be taught today and next week

### Last Week: Agents



- An agent perceives its environment through sensors and acts upon it through actuators (or effectors, depending on whom you ask)
- The agent function maps percept sequences to actions
- It is generated by an agent program running on a machine

### Last Week: the Pacman Agent



### Last Week: The task environment - PEAS

#### Performance measure

-1 per step; + 10 food; +500 win; -500 die;+200 hit scared ghost

#### Environment

Pacman dynamics (incl ghost behavior)

#### Actuators

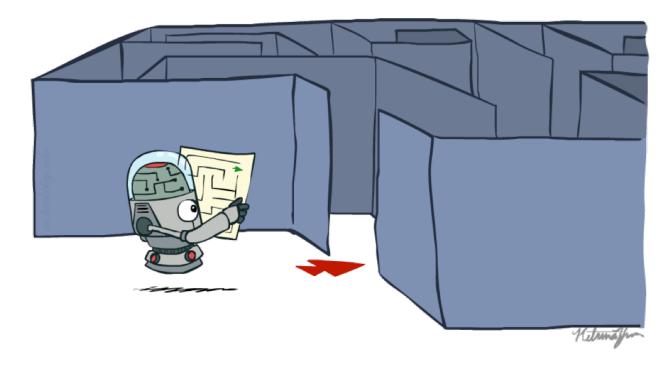
Left Right Up Down or NSEW

#### Sensors

Entire state is visible (except power pellet duration)



# CS 3317: Artificial Intelligence Search



Instructors: Cai Panpan

Shanghai Jiao Tong University

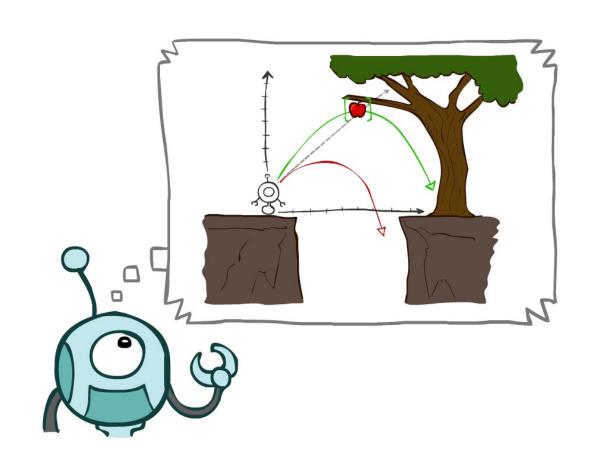
(slides adapted from UC Berkeley CS188)

### Today

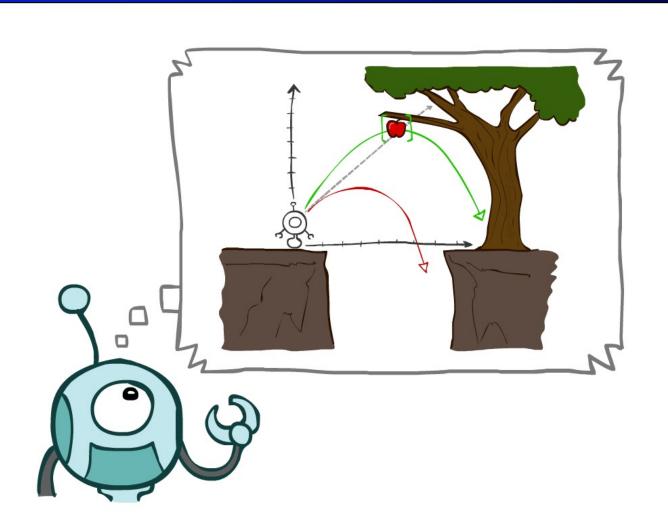
Agents that Plan Ahead

Search Problems

- Uninformed Search Methods
  - Depth-First Search
  - Breadth-First Search
  - Uniform-Cost Search



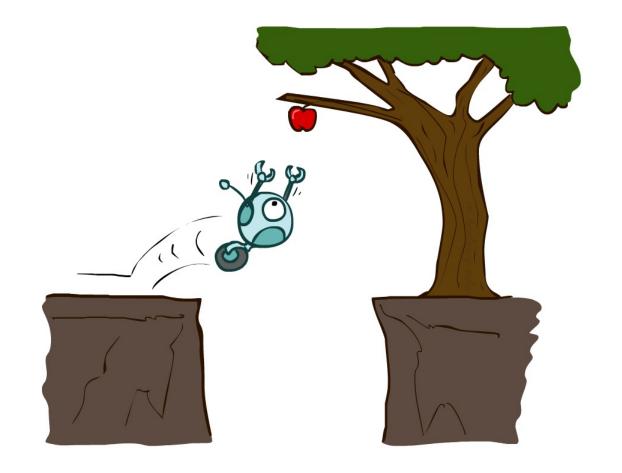
# Agents that Plan



### Reflex Agents

#### Reflex agents:

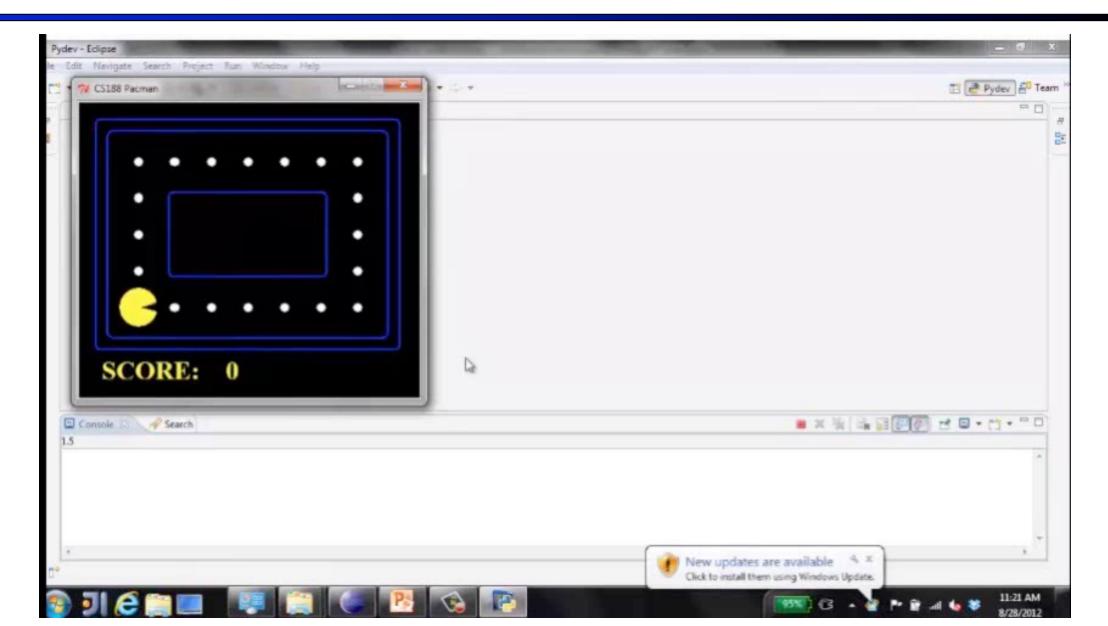
- Choose action based on current percept (and maybe memory)
- May have memory or a model of the world's current state
- Do not consider the future consequences of their actions
- Consider how the world IS
- Can a reflex agent be rational?



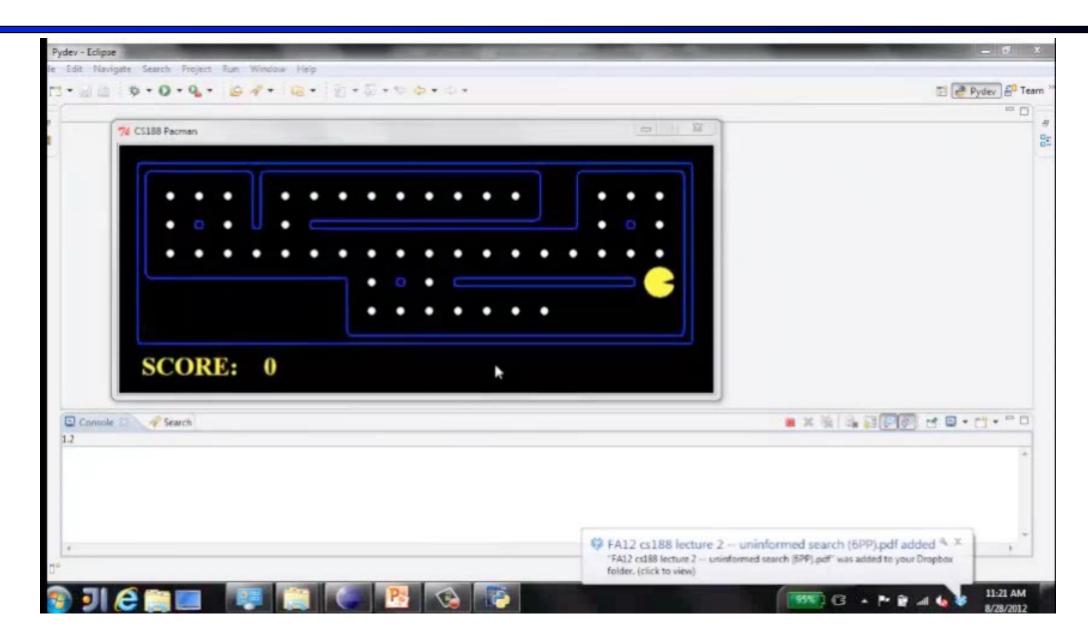
[Demo: reflex optimal (L2D1)]

[Demo: reflex optimal (L2D2)]

# Demo: Reflex Optimal

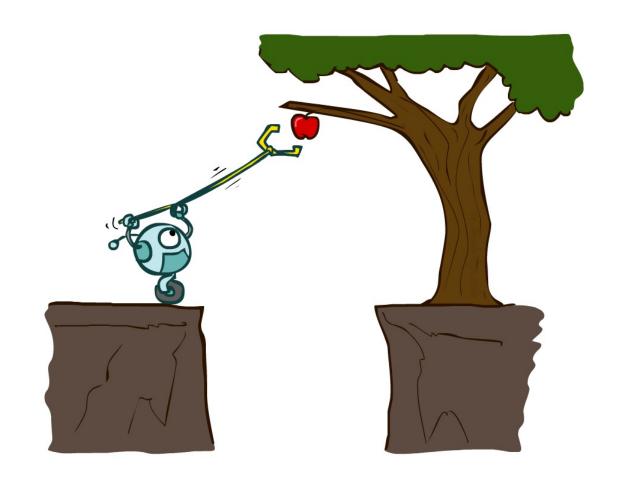


### Demo: Reflex Fails



### Planning Agents

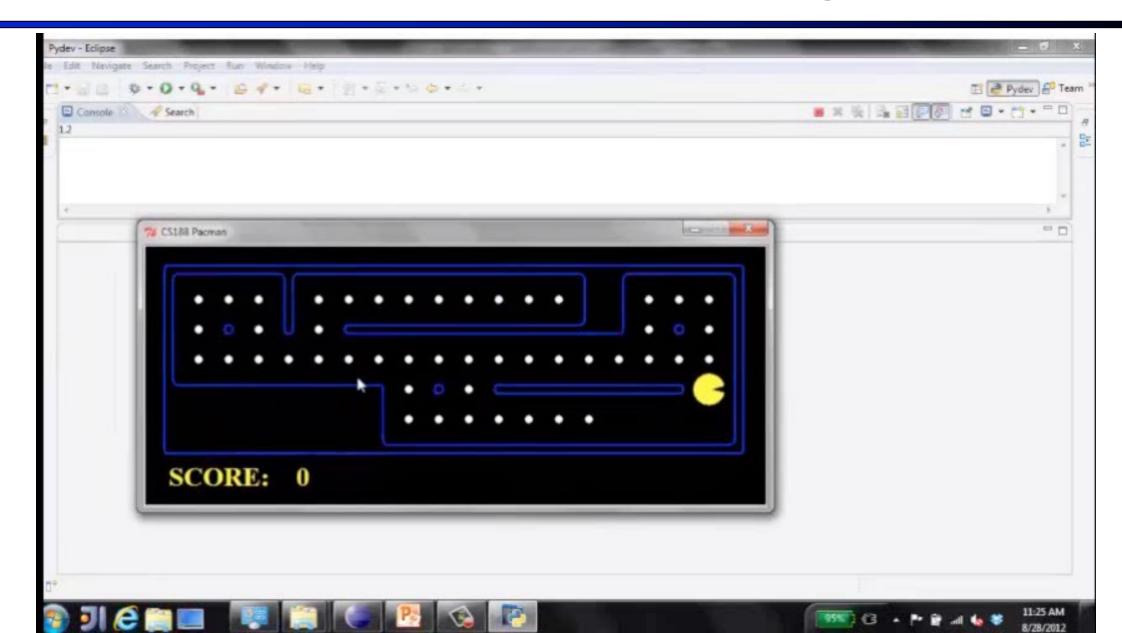
- Planning agents:
  - Ask "what if"
  - Decisions based on (hypothesized) consequences of actions
  - Must have a model of how the world evolves in response to actions
  - Must formulate a goal (test)
  - Consider how the world WOULD BE
- Optimal vs. complete planning
- Offline planning vs. online planning



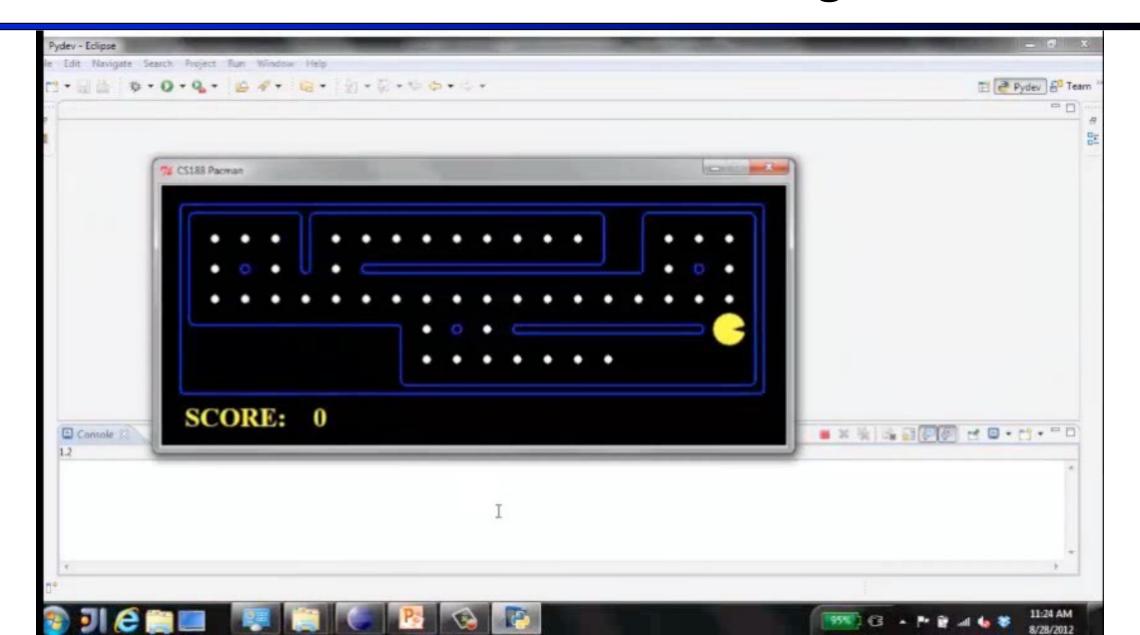
[Demo: re-planning (L2D3)]

[Demo: mastermind (L2D4)]

# Demo: Offline Planning



### Demo: Online Planning



### Search Problems

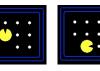


### Search Problems

- A search problem consists of:
  - A state space





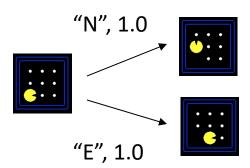








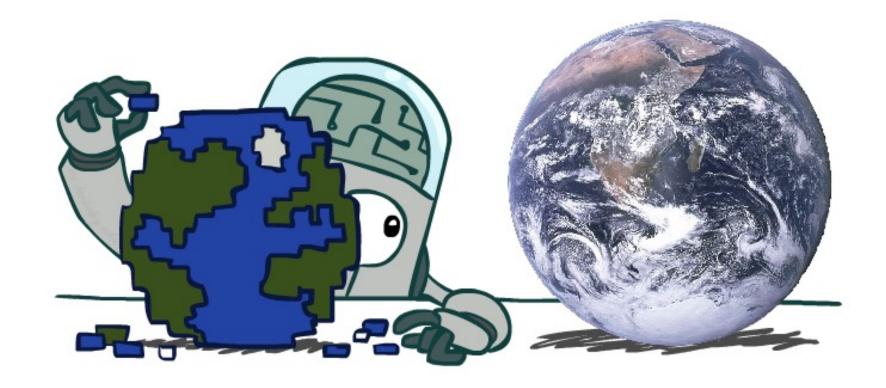
A successor function (with actions, costs)



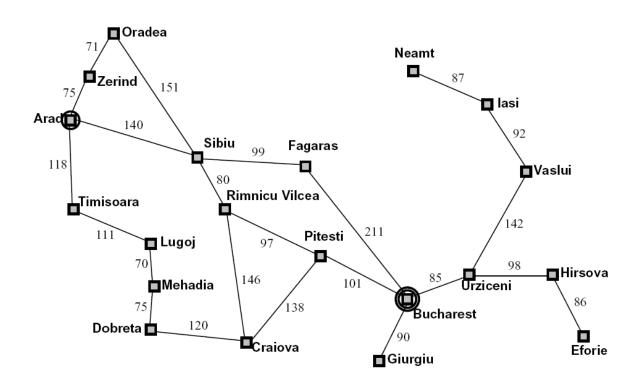
- A start state and a goal test
- A solution (a plan) is a sequence of actions which transforms the start state to a goal state

### Search Problems Are Models

"All models are wrong, some are useful." --George Box



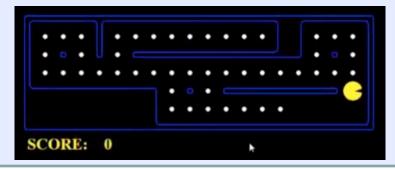
### Example: Traveling in Romania



- State space:
  - Cities
- Successor function:
  - Go to adjacent city
  - Cost = distance
- Start state:
  - Arad
- Goal test:
  - Is state == Bucharest?
- Solution?

### What's in a State Space?

The world state includes every last detail of the environment



A search state keeps only the details needed for planning (abstraction)

- Problem: Pathing
  - States: (x,y) location
  - Actions: NSEW
  - Successor: update location only
  - Goal test: is (x,y)=END

- Problem: Eat-All-Dots
  - States: {(x,y), dot booleans}
  - Actions: NSEW
  - Successor: update location and possibly a dot boolean
  - Goal test: dots all false

### State Space Sizes?

#### World state:

Agent positions: 120

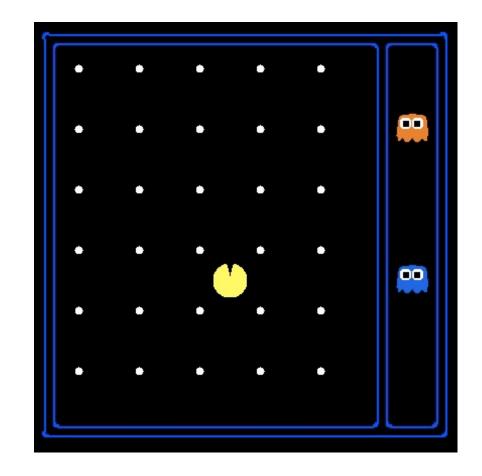
Food count: 30

Ghost positions: 12

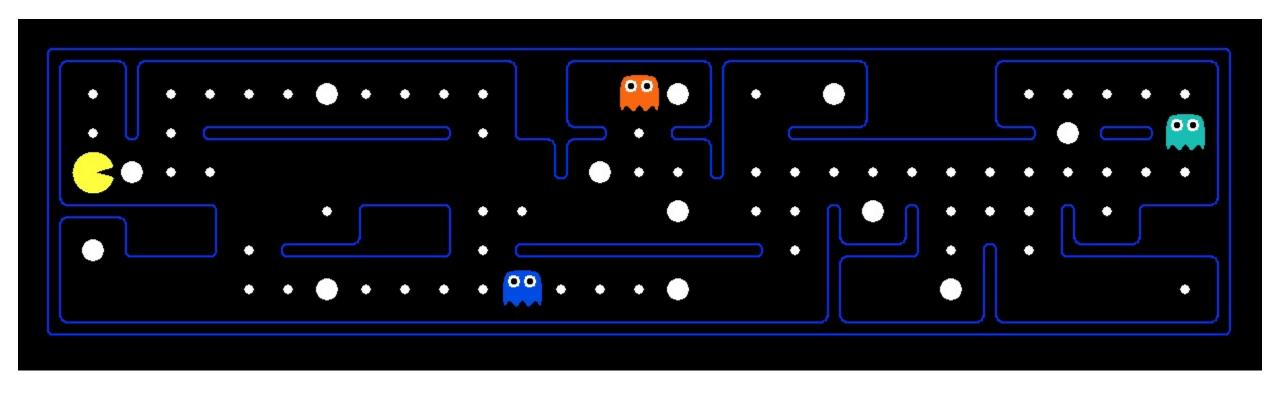
Agent facing: NSEW

#### How many

- World states?
   120x(2<sup>30</sup>)x(12<sup>2</sup>)x4
- States for pathing?120
- States for eat-all-dots?
   120x(2<sup>30</sup>)

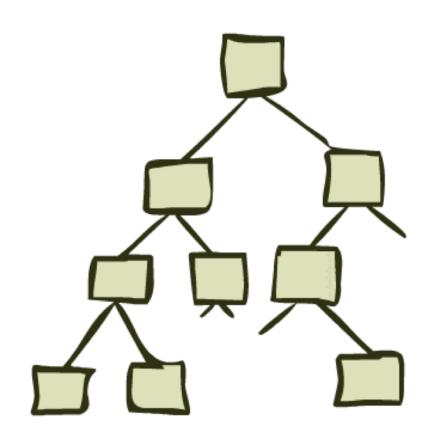


### Quiz: Safe Passage



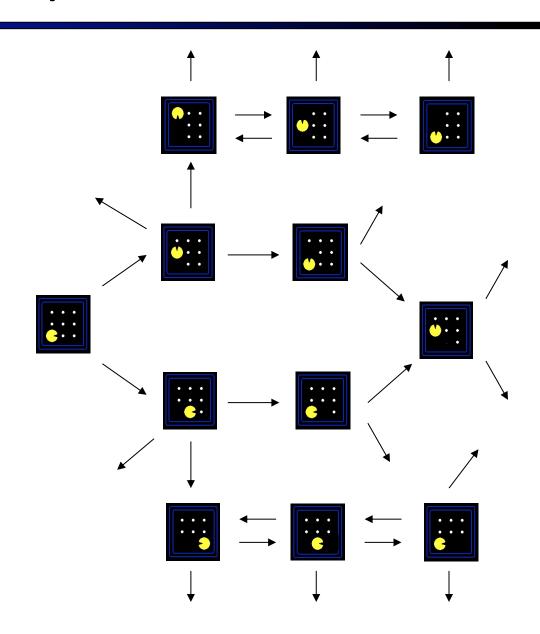
- Problem: eat all dots while keeping the ghosts always scared
- What does the state space have to specify?
  - (agent position, dot booleans, power pellet booleans, remaining scared time)

# State Space Graphs and Search Trees



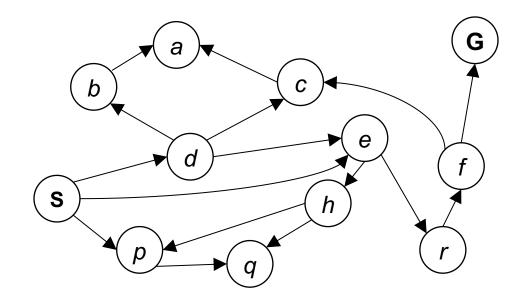
### State Space Graphs

- State space graph: A mathematical representation of a search problem
  - Nodes are (abstracted) world configurations
  - Arcs represent successors (action results)
  - The goal test is a set of goal nodes (maybe only one)
- In a state space graph, each state occurs only once!
- We can rarely build this full graph in memory (it's too big), but it's a useful idea



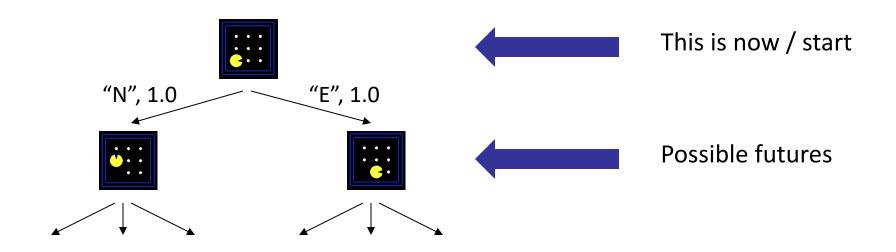
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Tiny state space graph for a tiny search problem

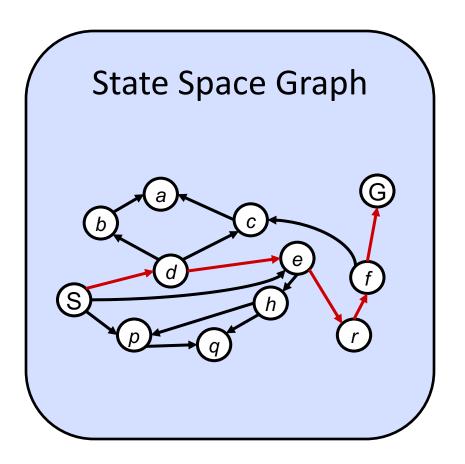
### Search Trees



#### A search tree:

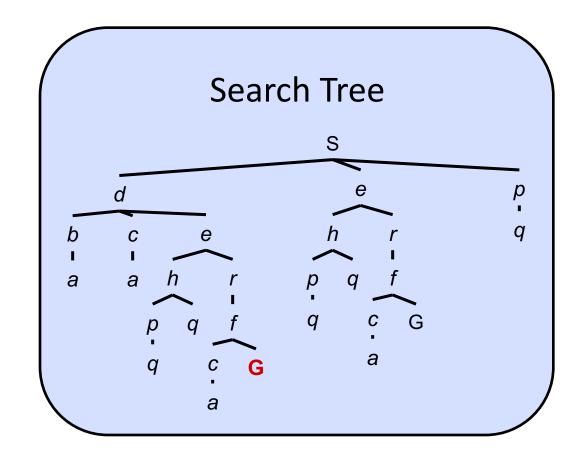
- A "what if" tree of plans and their outcomes
- The start state is the root node
- Children correspond to successors
- Nodes show states, but correspond to PLANS that achieve those states
- For most problems, we can never actually build the whole tree

### State Space Graphs vs. Search Trees



Each NODE in in the search tree is an entire PATH in the state space graph.

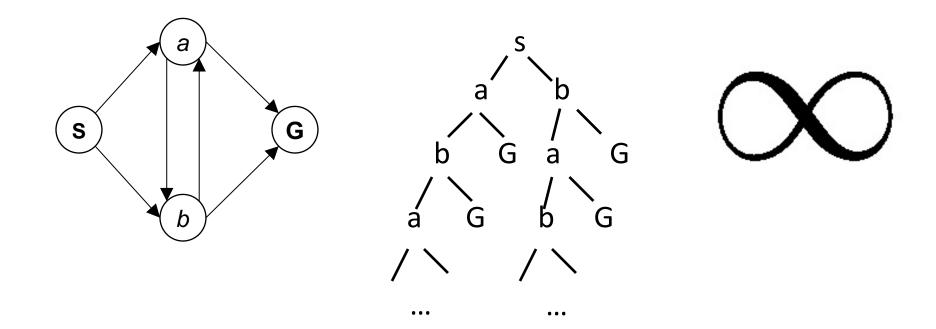
We construct both on demand – and we construct as little as possible.



### Quiz: State Space Graphs vs. Search Trees

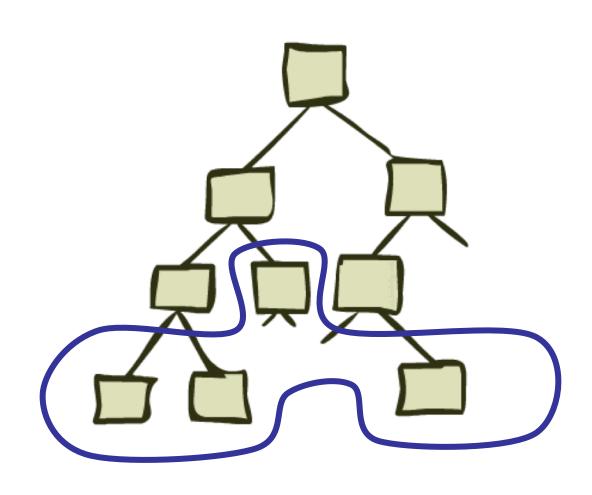
Consider this 4-state graph:

How big is its search tree (from S)?

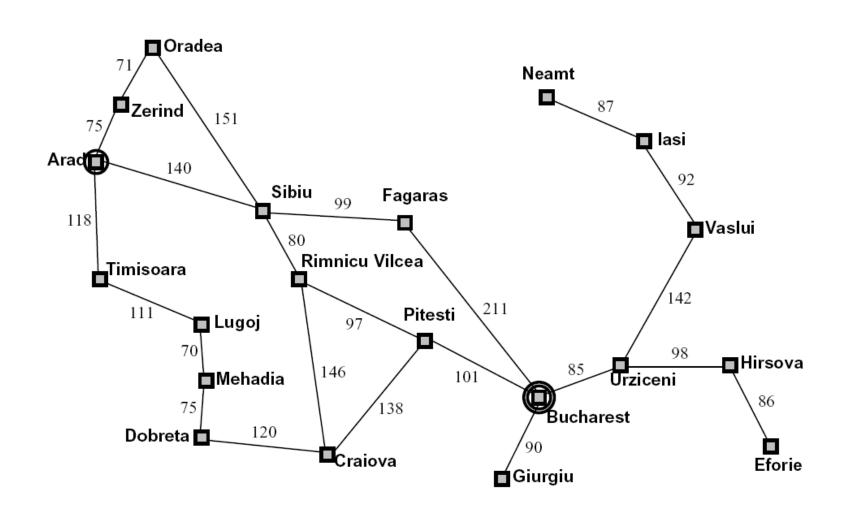


Caution: loops in the state-space graph can trap the search!

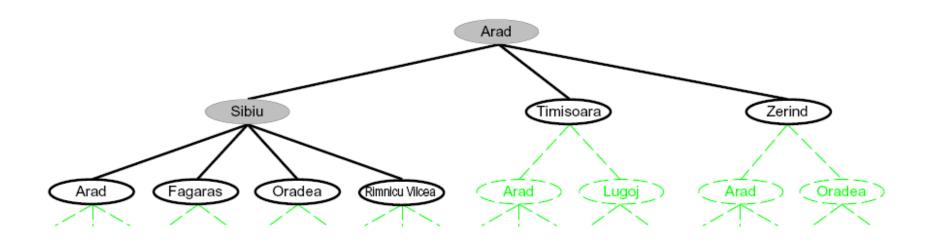
# Tree Search Algorithms



### Search Example: Romania



### Searching with a Search Tree



#### Search:

- Expand out potential plans (tree nodes)
- Maintain a fringe (list) of partial plans under consideration
- Try to expand as few tree nodes as possible

### **General Tree Search**

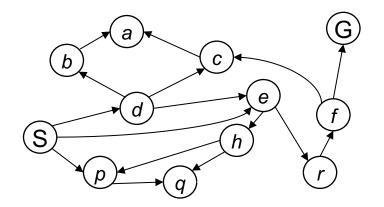
```
function TREE-SEARCH( problem, strategy) returns a solution, or failure initialize the search tree using the initial state of problem loop do

if there are no candidates for expansion then return failure choose a leaf node for expansion according to strategy

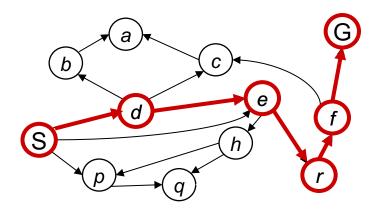
if the node contains a goal state then return the corresponding solution else expand the node and add the resulting nodes to the search tree end
```

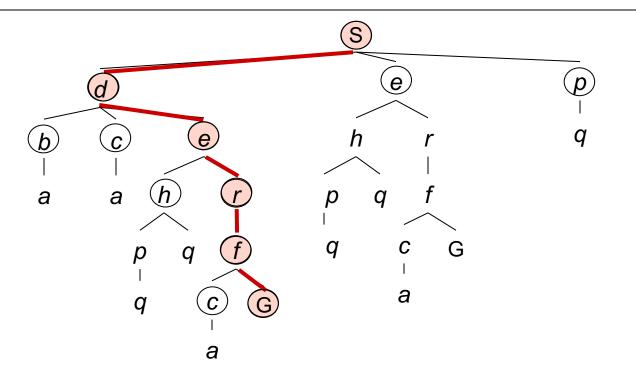
- Important ideas:
  - Fringe
  - Exploration strategy
  - Expansion
- Main question: which fringe nodes to explore?

# Example: Tree Search



### Example: Tree Search





```
s \rightarrow d

s \rightarrow e

s \rightarrow p

s \rightarrow d \rightarrow b

s \rightarrow d \rightarrow c

s \rightarrow d \rightarrow e

s \rightarrow d \rightarrow e \rightarrow h

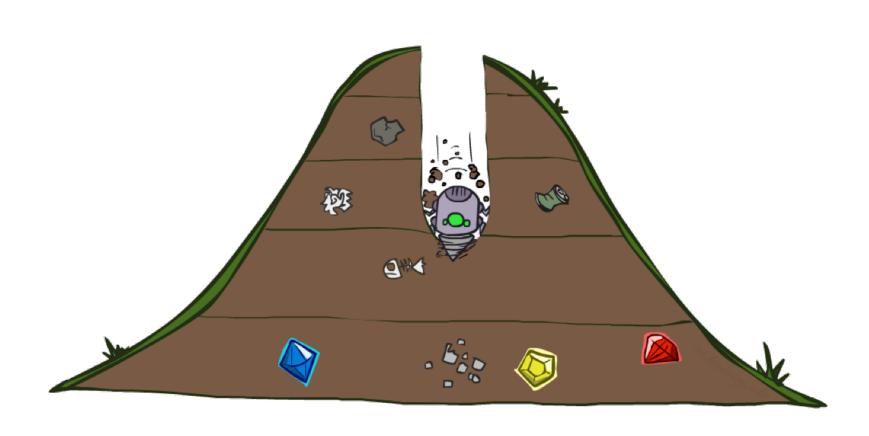
s \rightarrow d \rightarrow e \rightarrow r

s \rightarrow d \rightarrow e \rightarrow r \rightarrow f

s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow c

s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow c
```

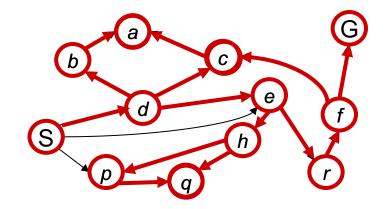
# Depth-First Search (DFS)

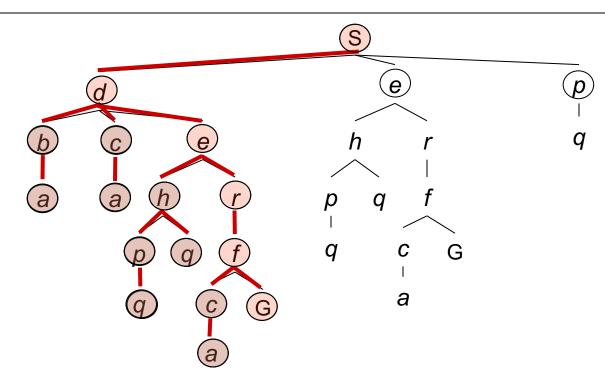


### Depth-First Search (DFS)

Strategy: expand a deepest node first

Implementation: Fringe is a LIFO stack



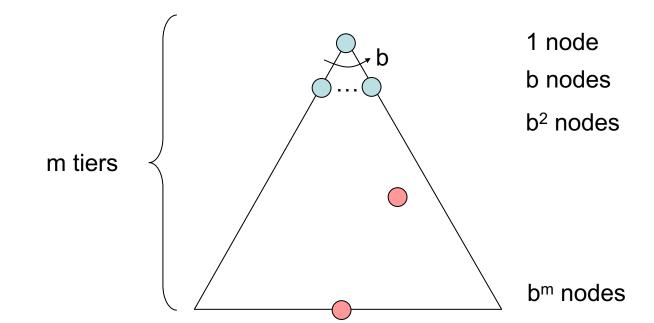


# Search Algorithm Properties



### Search Algorithm Properties

- Complete: Guaranteed to find a solution if one exists?
- Optimal: Guaranteed to find the least cost path?
- Time complexity?
- Space complexity?
- Cartoon of search tree:
  - b is the branching factor
  - m is the maximum depth
  - solutions at various depths



- Number of nodes in entire tree?
  - $1 + b + b^2 + .... b^m = O(b^m)$

# Depth-First Search (DFS) Properties

#### What nodes DFS expand?

- Some left prefix of the tree.
- Could process the whole tree!
- If m is finite, takes time O(b<sup>m</sup>)

#### • How much space does the fringe take?

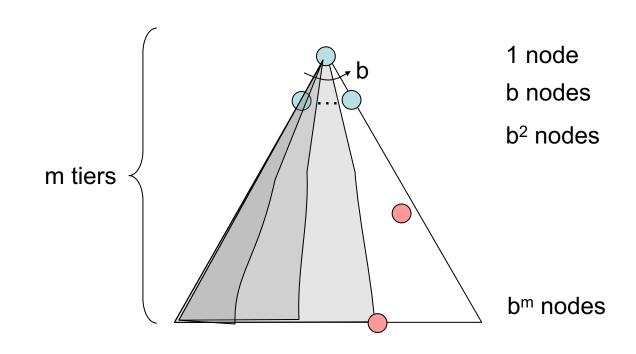
Only has siblings on path to root, so O(bm)

#### Is it complete?

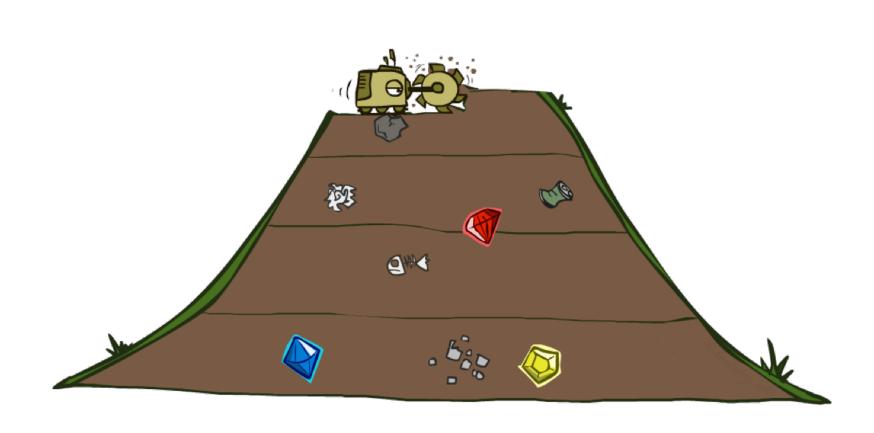
 m could be infinite, so only if we prevent cycles (more later)

#### Is it optimal?

 No, it finds the "leftmost" solution, regardless of depth or cost



# Breadth-First Search (BFS)

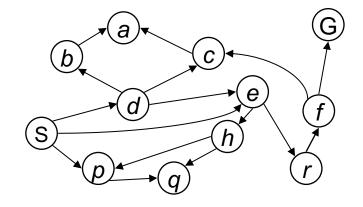


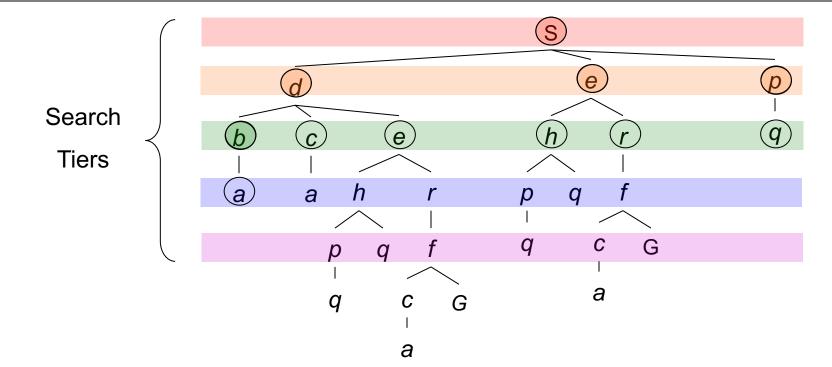
## **Breadth-First Search**

Strategy: expand a shallowest node first

*Implementation: Fringe* 

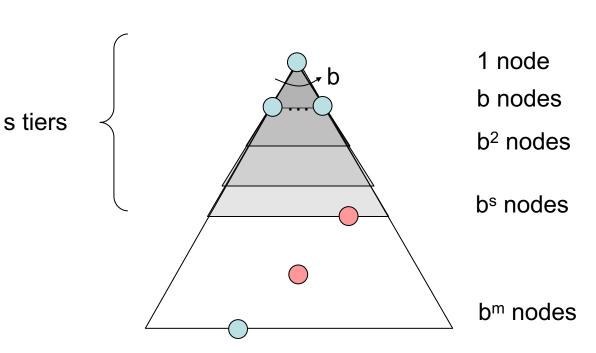
is a FIFO queue



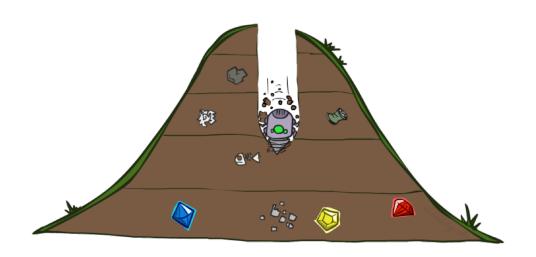


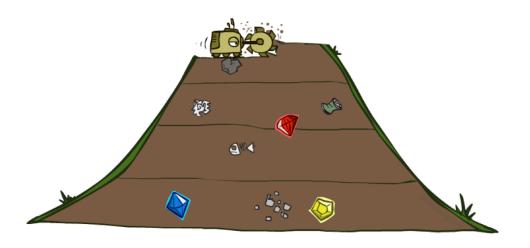
# Breadth-First Search (BFS) Properties

- What nodes does BFS expand?
  - Processes all nodes above shallowest solution
  - Let depth of shallowest solution be s
  - Search takes time O(b<sup>s</sup>)
- How much space does the fringe take?
  - Has roughly the last tier, so O(bs)
- Is it complete?
  - s must be finite if a solution exists, so yes!
- Is it optimal?
  - Only if costs are all 1 (more on costs later)



# Quiz: DFS vs BFS





## Quiz: DFS vs BFS

When will BFS outperform DFS?

When will DFS outperform BFS?

# Video of Demo Maze Water DFS/BFS (part 1)

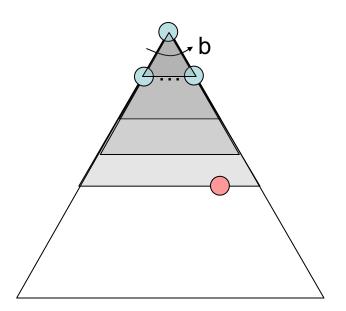


# Video of Demo Maze Water DFS/BFS (part 2)

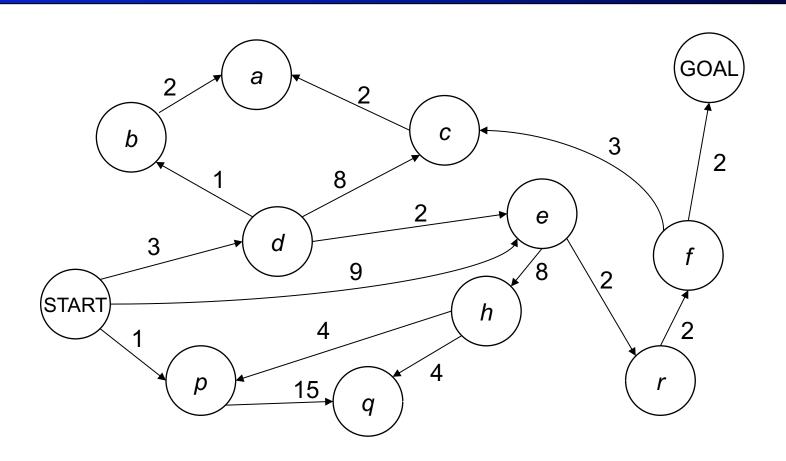


# **Iterative Deepening**

- Idea: get DFS's space advantage with BFS's time
   / shallow-solution advantages
  - Run a DFS with depth limit  $m_1 \ge 1$ . If no solution...
  - Run a DFS with depth limit  $m_2 > m_1$ . If no solution...
  - Run a DFS with depth limit  $m_3 > m_2$ . .....
- Isn't that wastefully redundant?
  - Generally most work happens in the lowest level searched, so not so bad!
  - You can also re-use the final fringe of the previous iteration as a warm start.

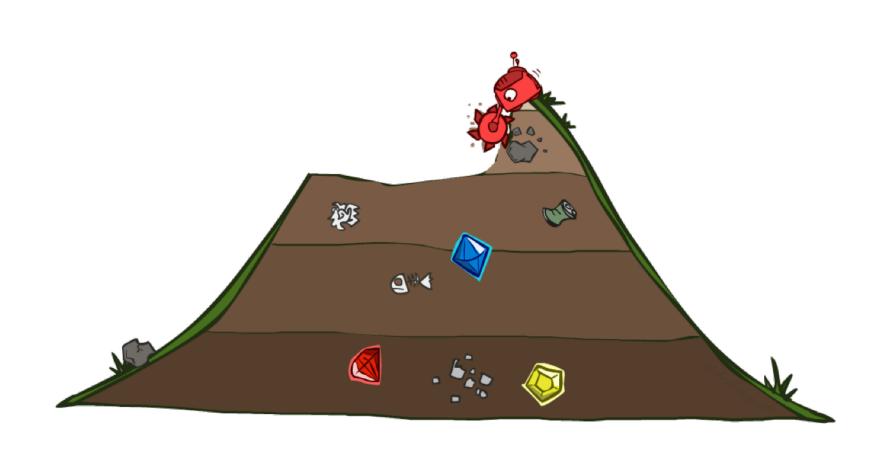


#### **Cost-Sensitive Search**



BFS finds the shortest path in terms of number of actions. It does not find the least-cost path. We will now cover a similar algorithm which does find the least-cost path.

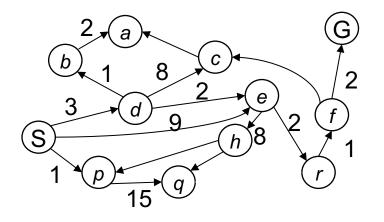
## **Uniform Cost Search**

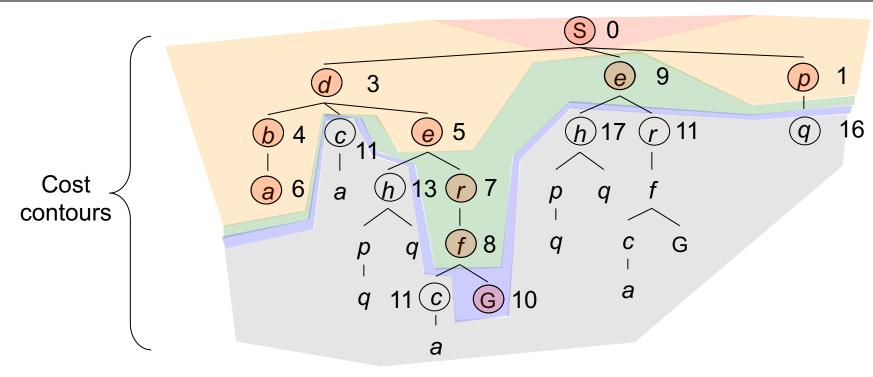


### **Uniform Cost Search**

Strategy: expand a cheapest node first:

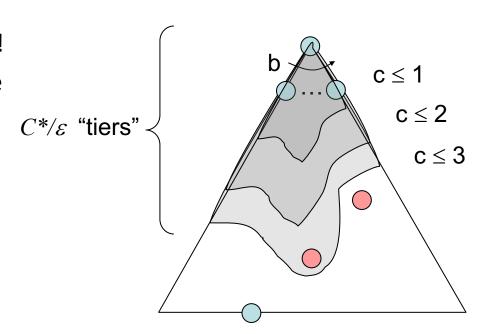
Fringe is a priority queue (priority: cumulative cost)





# Uniform Cost Search (UCS) Properties

- What nodes does UCS expand?
  - Processes all nodes with cost less than cheapest solution!
  - If that solution costs  $C^*$  and arcs cost at least  $\varepsilon$ , then the "effective depth" is roughly  $C^*/\varepsilon$
  - Takes time  $O(b^{C*/\varepsilon})$  (exponential in effective depth)
- How much space does the fringe take?
  - Has roughly the last tier, so  $O(b^{C*/\varepsilon})$
- Is it complete?
  - Assuming best solution has a finite cost and minimum arc cost is positive, yes!
- Is it optimal?
  - Yes! (Proof next lecture via A\*)



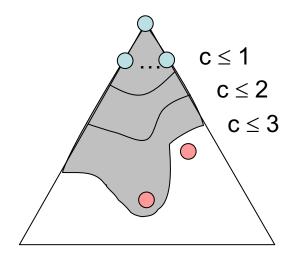
### **Uniform Cost Issues**

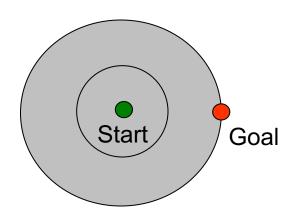
 Remember: UCS explores increasing cost contours

The good: UCS is complete and optimal!

- The bad:
  - Explores options in every "direction"
  - No information about goal location

We'll fix that soon!





[Demo: empty grid UCS (L2D5)] [Demo: maze with deep/shallow water DFS/BFS/UCS (L2D7)]

## Video of Demo Contours UCS Pacman Small Maze



# Video of Demo Empty UCS



Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 1)



Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 2)



Video of Demo Maze with Deep/Shallow Water --- DFS, BFS, or UCS? (part 3)



## One Queue for All

- All these search algorithms are the same except for fringe strategies
  - Conceptually, all fringes are priority queues (i.e. collections of nodes with attached priorities)
  - Practically, for DFS and BFS, you can avoid the log(n) overhead from an actual priority queue, by using stacks and queues
  - Can even code one implementation that takes a variable queuing object

