Bicolouring

The Problem

In 1976 the “Four Colour Map Theorem” was proven with the assistance of a computer. This theorem states that every map can be coloured using only four colours, in such a way that no region is coloured using the same colour as a neighbour region.

Here you are asked to solve a simpler similar problem. You have to decide whether a given arbitrary connected graph can be bicoloured. That is, if one can assign colours (from a palette of two) to the nodes in such a way that no two adjacent nodes have the same colour. To simplify the problem you can assume:

- no node will have an edge to itself.
- the graph is undirected. That is, if a node *a* is said to be connected to a node *b*, then you must assume that *b* is connected to *a*.
- the graph will be strongly connected. That is, there will be at least one path from any node to any other node.

The Input

The input consists of several test cases. Each test case starts with a line containing the number n ($1 < n < 200$) of different nodes. The second line contains the number of edges e. After this, e lines will follow, each containing two numbers that specify an edge between the two nodes that they represent. A node in the graph will be labeled using a number a ($0 <= a < n$). An input with $n = 0$ will mark the end of the input and is not to be processed.

The Output

You have to decide whether the input graph can be bicoloured or not, and print it as shown below.

Sample Input

```
3
3
0 1
1 2
2 0
9
8
0 1
0 2
0 3
0 4
0 5
0 6
0 7
0 8
0
```

Sample Output

```
NOT BICOLOURABLE.
BICOLOURABLE.
```