

# Dynamic Weighted Bit-Flipping Decoding Algorithms for LDPC Codes

Tofar C.-Y. Chang, *Student Member, IEEE*, and Yu T. Su, *Senior Member, IEEE*

**Abstract**—Bit-flipping (BF) decoding of low-density parity-check codes is of low complexity but gives inferior performance in general. To improve performance and provide new BF decoder options for complexity-performance tradeoffs, we propose new designs for the flipping function (FF), the flipped bit selection (FBS) rule and the checksum weight updating schedule. The new FF adjusts the checksum weights in every iteration while our FBS rules take more information into account. These two modifications represent efforts to track more closely the evolutions of both check and variable nodes' reliabilities. Two selective update schedules are proposed to offer more performance and complexity tradeoffs. The combinations of the new FBS rule and known FFs result in new BF decoders with improved performance and a modest complexity increase. On the other hand, combining the new FF and FBS rule gives a new decoder with performance comparable to that of the normalized min-sum algorithm while if we use a much simpler FBS rule instead, the decoder suffers little performance loss with reduced complexity. We also present a simple decision-theoretical argument to justify the new checksum weight formula and a time-expanded factor graph model to explain the proposed selective weight-updating schedules.

**Index Terms**—LDPC codes, belief propagation, bit-flipping decoding, flipped bit selection.

## I. INTRODUCTION

LOW-density parity-check (LDPC) codes have been shown to asymptotically give near-capacity performance when the sum-product algorithm (SPA) is used for decoding [1]. Gallager proposed two alternatives that use only hard-decision bits [2]. These so-called bit-flipping (BF) algorithms flip one or a group of bits based on the values of the flipping functions (FFs) computed in each iteration. The FF associated with a variable node (VN) is a *reliability metric* of the corresponding bit decision and depends on the binary-valued checksums of the VN's connected check nodes (CNs). Although BF algorithms are much simpler than the SPA, their performance is far from optimal. To reduce the performance gap, many variants of Gallager's BF algorithms have been proposed. Most of them

tried to improve the VN's reliability metric (the FF) and/or the method of selecting the flipped bits, achieving different degrees of bit error rate (BER) and convergence rate performance enhancements at the cost of higher complexity.

The class of weighted bit-flipping (WBF) algorithms [3]–[7] assign proper weights to the binary checksums. Each weight can be regarded as a reliability metric on the corresponding checksum and is a function of the associated soft received channel values. Another approach called gradient descent bit-flipping (GDBF) algorithm was proposed by Wadayama *et al.* [9]. Instead of using a weighted checksum based FF, the GDBF algorithm derives its FF by computing the gradient of a nonlinear objective function which is equivalent to the log-likelihood function of the bit decisions with checksum constraints. It was shown that the GDBF algorithm outperforms most known WBF algorithms when the VN degree is small. Sundararajan *et al.* [17] modified this FF by introducing a weighting on syndrome and a zero-mean Gaussian perturbation term. The resulting noisy GDBF (NGDBF) algorithm improves the performance of the GDBF algorithm which is further enhanced by adding a re-decoding process [18].

For the WBF algorithms, the weights are decided by the soft received channel values and remain unchanged throughout the decoding process. Since the weights reflect the decoder's belief on the checksums which, in turn, depend on those of the associated VNs' FF and bit decisions, the associated checksum weights should be updated accordingly. In [19], a reliability-based schedule is used in the initial decoding iteration to forward only reliable VN and CN's messages. Nguyen and Vasić [20] employed an extra bit to adaptively represent the reliabilities of VN and CN messages and developed a class of two-bit BF algorithms along with algorithm(s)-selection procedures. In this paper, we present dynamic weighted BF (DWBF) algorithms that assign dynamic checksum weights which are updated according to a nonlinear function of the associated VNs' FF values. As we shall show, the nonlinear function has the effects of quarantining unreliable checksums (which is similar to the method used in the first decoding iteration of [19]) while dynamically adjusts the more reliable checksums' weights. A simple decision theoretical interpretation is given to explain the effect of the nonlinear action and justify the threshold selection. We also suggest two selective weight-updating schedules which offer additional performance-complexity trade-offs. A time-expanded factor graph model is used to illustrate the weight-updating schedules.

The single-bit BF algorithms flip only the least reliable bit thus result in slow convergence rates. For this reason, many a multiple-flipped-bit selection rule was suggested [9]–[18].

Manuscript received December 25, 2014; revised May 20, 2015 and August 3, 2015; accepted August 12, 2015. Date of publication August 18, 2015; date of current version November 13, 2015. This work was supported in part by Taiwan's Ministry of Science and Technology under Grant NSC-100-2111-E-009-066-MY3. The associate editor coordinating the review of this paper and approving it for publication was D. Declercq.

T. C.-Y. Chang is with the Institute of Communications Engineering, National Chiao Tung University, Hsinchu 30010, Taiwan (e-mail: tofar.cm96g@nctu.edu.tw).

Y. T. Su is with the Department of Electrical and Computer Engineering, National Chiao Tung University, Hsinchu 30010, Taiwan (e-mail: ytsu@nctu.edu.tw).

Color versions of one or more of the figures in this paper are available online at <http://ieeexplore.ieee.org>.

Digital Object Identifier 10.1109/TCOMM.2015.2469780

By simultaneously flipping the selected bits, a BF decoder can offer rapid convergence but, sometimes, at the expense of performance loss. A bit selection rule may consist of simple threshold comparisons or include a number of steps involving different metrics. It is usually designed assuming a specific FF is used and may not be suitable when a different FF or metric is involved. Moreover, the FF value may not provide sufficient information for making a tentative bit decision, we propose a new flipped bit selection (FBS) rule that takes into account both the FF value and other information from related CNs.

FF, checksum weight computing, weight-updating schedules, and FBS rule are major constituent parts of a DWBF decoder. Our proposals on these parts offer a variety of new design options and tradeoffs. The efficiencies of *using the proposed schemes jointly or separately with existing designs* are evaluated by examining the corresponding numerical error rate and convergence behaviors. We show that our single-bit DWBF algorithm provides significant performance improvement over the existing single-bit GDBF and WBF algorithms. Our FBS rule works very well with different FFs and outperforms known FBS rules. Moreover, the selective weight-updating schedules suffer little performance degradation while offer significant complexity reduction when the CN and VN degrees are small.

Note that since the checksum weights are crucial parameters of a DWBF decoder's FF and their updates depend in turn on the FF values computed at the previous iteration, we henceforth mean both the FF and the associated weight computing formula whenever FF is mentioned.

The rest of this paper is organized as follows. In Section II, we define the basic system parameters and give a brief overview of various BF decoding algorithms, their FFs and FBS rules. In Section III, we introduce a new FF and its checksum weight-updating formula. A simple decision theoretical justification is given. We consider single-bit BF decoders and present two weight-updating schedules as well as their graphic models in Section IV. The performance of the our single-bit DWBF algorithm and some known single-bit BF algorithms are compared in the same section. We develop a new multi-bit FBS rule and present the error rate and convergence behaviors of various multi-bit BF decoder structures based on the new FBS rule in Section V. These decoders' complexities are analyzed in details for evaluating various performance and complexity tradeoffs. Finally, conclusion remarks are drawn in Section VI.

## II. BACKGROUND AND RELATED WORKS

### A. Notations and the Basic Algorithm

We denote by  $(N, K)(d_v, d_c)$  a regular binary LDPC code  $\mathcal{C}$  with VN degree  $d_v$  and CN degree  $d_c$ , i.e.,  $\mathcal{C}$  is the null space of an  $M \times N$  parity check matrix  $\mathbf{H} = [H_{mn}]$  which has  $d_v$  1's in each column and  $d_c$  1's in each row. Let  $\mathbf{u}$  be a codeword of  $\mathcal{C}$  and assume that the BPSK modulation is used so that a codeword  $\mathbf{u} = (u_0, u_1, \dots, u_{N-1})$ ,  $u_i \in \{0, 1\}$ , is mapped into a bipolar sequence  $\mathbf{x} = (x_0, x_1, \dots, x_{N-1}) = (1 - 2u_0, 1 - 2u_1, \dots, 1 - 2u_{N-1})$  for transmission. The equivalent base-band transmission channel is a binary-input Gaussian-output channel characterized by additive zero-mean white Gaussian

noise with two-sided power spectral density of  $N_0/2$  W/Hz. Let  $\mathbf{y} = (y_0, y_1, \dots, y_{N-1})$  be the sequence of soft channel values obtained at the receiver's coherent matched filter output. The sequence  $\mathbf{z} = (z_0, z_1, \dots, z_{N-1})$ , where  $z_i \in \{0, 1\}$ , is obtained by taking hard-decision on each components of  $\mathbf{y}$ . Let  $\hat{\mathbf{u}} = (\hat{u}_0, \hat{u}_1, \dots, \hat{u}_{N-1})$  be the tentative decoded binary sequence at the end of a BF decoding iteration. We compute the syndrome (checksum) vector  $\mathbf{s} = (s_0, s_1, \dots, s_{M-1})$  by  $\mathbf{s} = \hat{\mathbf{u}} \cdot \mathbf{H}^T \pmod{2}$ . We further denote the  $n$ th VN by  $v_n$ , the set of indices of its connecting CNs by  $\mathcal{M}(n)$ , and the set of indices of the VNs checked by the  $m$ th CN  $c_m$  by  $\mathcal{N}(m)$ . The indices of CNs in  $\mathcal{M}(n)$  are determined by the nonzero elements of the  $n$ th column of  $\mathbf{H}$  whereas those in  $\mathcal{N}(m)$  are by the  $m$ th row of  $\mathbf{H}$ .

A generic BF decoding algorithm can be described by **Algorithm 1** below which involves four important parameters:  $l$ , the iteration number;  $l_{\max}$ , the maximum iteration number;  $E_n$ , the FF; and  $\mathcal{B}$ , the index set of the flipped bits, or the flipped bit (FB) set for short. This algorithm performs two basic tasks: 1) computing  $E_n$ 's (**Step 2**) and 2) generating the FB set  $\mathcal{B}$  (**Step 3**). Most earlier works focused on improving either 1) or 2).

---

### Algorithm 1 A General BF Decoding Algorithm

---

**Initialization** Set  $l = 0$  and  $\hat{\mathbf{u}} = \mathbf{z}$ .

**Step 1**  $\forall m \in O_M \triangleq \{0, 1, \dots, M-1\}$ , compute

$$s_m = \sum_{n \in \mathcal{N}(m)} \hat{u}_n H_{mn} \pmod{2}. \quad (1)$$

If  $\mathbf{s} = \mathbf{0}$  or  $l = l_{\max}$ , stop decoding and output  $\hat{\mathbf{u}}$ ; otherwise,  $l \leftarrow l + 1$ .

**Step 2**  $\forall n \in O_N \triangleq \{0, 1, \dots, N-1\}$ , compute the FF  $E_n$ .

**Step 3** Use the FFs obtained in **Step 2** to update the flipped bit set  $\mathcal{B}$ .

**Step 4** Flip  $\hat{u}_n$  for all  $n \in \mathcal{B}$  and go to **Step 1**.

---

An FF, sometimes referred to as *cost function* or *inversion function* [9], is used as a reliability metric on a VN's tentative decision. Given the FF values and the FBS rule, we select a set of VNs and flip the corresponding tentative decisions (bits). Choosing the most unreliable bits or the bits whose FF values exceed a threshold are the two most popular rules. For the former rule, usually only one bit is flipped if the soft-valued channel information is employed in the FF, resulting in slow convergence. By contrast, the latter rule often gives faster convergence rate but possibly at the cost of performance loss. We briefly review the known FFs and FBS rules in the following paragraphs.

### B. Flipping Functions of BF Decoding Algorithms

Gallager proposed that a simple sum of binary checksums be used as the FF [2]

$$E_n = - \sum_{m \in \mathcal{M}(n)} (1 - 2s_m). \quad (2)$$

(2) implies that the FF value is inversely proportional to the bit decision reliability as it is an increasing function of the number of nonzero checksums (i.e., unsatisfied check nodes, UCNs).

By taking into account soft-valued channel information and assigning checksum weights, later modifications of Gallager's FF can be described by the following general formula

$$E_n = -\alpha_1 \cdot \phi(\hat{u}_n, y_n) - \sum_{m \in \mathcal{M}(n)} w_{mn}(1 - 2s_m), \quad (3)$$

where  $\alpha_1 > 0$ ,  $\phi(\hat{u}_n, y_n)$  is a reliability metric involving channel value and/or bit decision, and, to be consistent with (2),  $w_{mn} \geq 0$ .

For the WBF algorithm [4],  $\phi(\hat{u}_n, y_n) = 0$  and  $w_{mn}$  is

$$w_{mn} = \min_{n' \in \mathcal{N}(m)} |y_{n'}|, \quad (4)$$

The modified WBF (MWBF) algorithm [5] has  $\phi(\hat{u}_n, y_n) = |y_n|$  while the improved MWBF (IMWBF) algorithm [6] uses the same  $\phi(\hat{u}_n, y_n)$  but replaces the checksum weight by

$$w_{mn} = \min_{n' \in \mathcal{N}(m) \setminus n} |y_{n'}| \quad (5)$$

for the message passed from  $c_m$  to  $v_n$  should exclude that originated from  $v_n$ . For the reliability ratio based WBF (RRWBF) algorithm [7],  $\phi(\hat{u}_n, y_n) = 0$  and

$$w_{mn} = 1/w'_{mn} = \left( \beta \frac{|y_n|}{\max_{n' \in \mathcal{N}(m)} |y_{n'}|} \right)^{-1}, \quad (6)$$

where  $\beta$  is the normalizing factor to ensure that  $\sum_{n \in \mathcal{N}(m)} w'_{mn} = 1$ .

The GDBF algorithm of Wadayama *et al.* [9] applies the gradient descent method to minimize

$$E(\hat{\mathbf{u}}) = - \sum_{n=0}^{N-1} y_n(1 - 2\hat{u}_n) - \sum_{m=0}^{M-1} (1 - 2s_m) \quad (7)$$

with respect to  $(1 - 2\hat{u}_n)$  and obtains the FF

$$E_n = -y_n(1 - 2\hat{u}_n) - \sum_{m \in \mathcal{M}(n)} (1 - 2s_m), \quad (8)$$

which is equivalent to assigning  $\alpha_1 = 1$ ,  $\phi(\hat{u}_n, y_n) = y_n(1 - 2\hat{u}_n)$ , and  $w_{mn} = 1$  in (3). Recently, Sundararajan *et al.* [17] introduced the so-called noisy GDBF (NGDBF) algorithm based on

$$E_n = -y_n(1 - 2\hat{u}_n) - w \sum_{m \in \mathcal{M}(n)} (1 - 2s_m) + q_n, \quad (9)$$

where  $q_n$ 's in (9) are i.i.d. zero-mean Gaussian random perturbation with a signal-to-noise ratio (SNR) dependent variance and  $w$  is a constant syndrome weight.

For the above two FFs, (8) and (9),  $\phi(\hat{u}_n, y_n) = y_n(1 - 2\hat{u}_n)$  is equal to  $|y_n|$  when the bit decision  $\hat{u}_n$  is the same as the  $y_n$ -based hard decision,  $z_n$ ; otherwise, its value is the same as  $-|y_n|$ . This is consistent with the intuition that  $\hat{u}_n = z_n$  implies that  $\hat{u}_n$  is likely to be correct and since  $-E_n$  is a VN

reliability metric, a positive  $\phi(\hat{u}_n, y_n)$  helps increasing  $-E_n$  and preventing the  $\hat{u}_n$  from being flipped. In contrast,  $\phi(\hat{u}_n, y_n)$  is always positive ( $|y_n|$ ) in MWBF algorithms, which means that the MWBF algorithms tend to trust  $\hat{u}_n$  in spite of other evidence.

### C. Flipped Bit Selection Rules

For the algorithms mentioned in the Section II-B, only the bit(s) related to the VN having the largest FF value  $E_n$  is (are) flipped in each iteration, i.e., the FB set is

$$\mathcal{B} = \{n | n = \arg \max_i E_i\}. \quad (10)$$

As mentioned before,  $|\mathcal{B}| = 1$  and the corresponding convergence is often very slow if  $E_n$  has a soft-valued information term.

Flipping several bits in each iteration simultaneously can improve the convergence speed. The simplest FBS rule for multi-bit BF decoding uses the FB set

$$\mathcal{B} = \{n | E_n \geq \Delta\}, \quad (11)$$

where the threshold  $\Delta$  can be a constant or be adaptive. The optimal adaptive threshold was derived by Gallager [2], assuming that no cycle appears in the code graph. Since practical finite-length LDPC codes usually have cycles and the optimal thresholds can only be found through time-consuming simulations, two ad-hoc methods which automatically adjust  $\Delta$  were suggested in [12] and [13]. In the adaptive threshold BF (ATBF) algorithm [12], the initial  $\Delta$  is found by simulation and subsequent thresholds are a monotonically non-increasing function of the decoding iterations. The adaptive MWBF (AMWBF) algorithm [13] adjusts the threshold by

$$\Delta = E^* - |E^*| \left[ 1 - \frac{w_H(s)}{M} \right], \quad (12)$$

where  $E^* = \max_n E_n$  and  $w_H(s)$  is the Hamming weight of the syndrome vector  $s$ .

Sometimes, a tentative decoded vector  $\hat{\mathbf{u}}$  may reappear several times during the decoding process and form a decoding loop. This may be caused, for example, by the event that an even number of bits associated with a CN are flipped, leading to an unchanged checksum and then oscillating bit decisions. To eliminate the occurrence of loops, the AMWBF algorithm includes the loop detection scheme of [8] in its FBS rule so that if a loop is detected, the most reliable bit in  $\mathcal{B}$  is removed. The parallel weighted BF (PWBF) algorithm [10] tries to reduce the loop occurrence probability by having every UCN ( $s_m = 1$ ) send a constant *flip signal* (FS) to its least reliable linked VN (based on the FF of the IMWBF algorithm) and flips the bits in

$$\mathcal{B} = \{n | F_n \geq \Delta_{FS}\}, \quad (13)$$

where  $\Delta_{FS}$  is a constant optimized by simulations,

$$F_n = \sum_{m \in \mathcal{M}(n)} q_{mn} s_m, \quad (14)$$

and  $q_{mn}$  is given by

$$q_{mn} = \begin{cases} 1, & n = \arg \max_{n' \in \mathcal{N}(m)} E_{n'} \\ 0, & \text{otherwise.} \end{cases} \quad (15)$$

Since the above remedy can only eliminate loops with a certain probability, the improved PWBf (IPWBf) algorithm employs the loop detection scheme of [8] and when a loop is detected, it removes the bit(s) receiving the smallest  $F_n$  from  $\mathcal{B}$ . This algorithm also adds a bootstrapping step and a delay-handling procedure to further improve the bit selection accuracy but achieves limited improvement for the codes with high column degrees such as Euclidean geometry (EG) LDPC codes.

A hybrid GDBF (HGDBF) algorithm was proposed in [9]. In this algorithm, single- and multi-bit BF decoding is performed alternatively and an *escape process* is used for preventing the decoding process from being trapped in local minima/loops. Two extensions of the HGDBF algorithm were considered in [14] and [15] which require less complexity at the expense of inferior performance. A multi-bit GDBF algorithm with a probabilistic FBS rule was also suggested in [16] for hard-decision decoding. With the FF (9) and FBS rule of [15], the multi-bit NGDBF (M-NGDBF) algorithm [17] achieves the same BER performance as that of the HGDBF decoder with much less decoding iterations. However, [18] found that new trapping set conditions may exist in the M-NGDBF decoder but can be eliminated by re-decoding with a different perturbation sequence.

### III. CHECKSUM RELIABILITY AND DYNAMIC WEIGHTS

#### A. BF Decoding and Checksum Weights

In line with the belief propagation (BP) based SPA,  $E_n$  is similar to the total log-likelihood ratio (LLR) of  $v_n$  and  $-w_{mn}(1 - 2s_m)$  in (3) is analogous to the belief sent to  $v_n$  by  $c_m$ . Unlike SPA, however, for (4)–(6), the latter remains unchanged unless an  $\hat{u}_n, n \in \mathcal{N}(m)$  has been flipped, which leads to just a sign change of the belief. The general FF format, (3), includes two major terms that represent the decoder's confidence on a VN's tentative decision based respectively on its channel value (or the correlation of the channel value and tentative decision) and the reliabilities of the related checksums. Since the channel values remain fixed, the checksums should be given adjusted weights at least in the later iterations when the reliabilities of checksums change.

Although the flipping operation changes the reliability metric of  $\hat{u}_n$  and the related checksums, all the FFs used in known BF decoders use static  $w_{mn}$  thereby can neither reflect the dynamic of VNs' message passing nor offer self-adjustment capability in accurately updating bit reliability information. We present dynamic weight generation method in this section.

#### B. Flipping Function and Decision Reliability

The review on BF algorithms in Section II indicates clearly that the FF value is a proper explicit or implicit reliability metric of a VN's decision. As a checksum in turn is a function

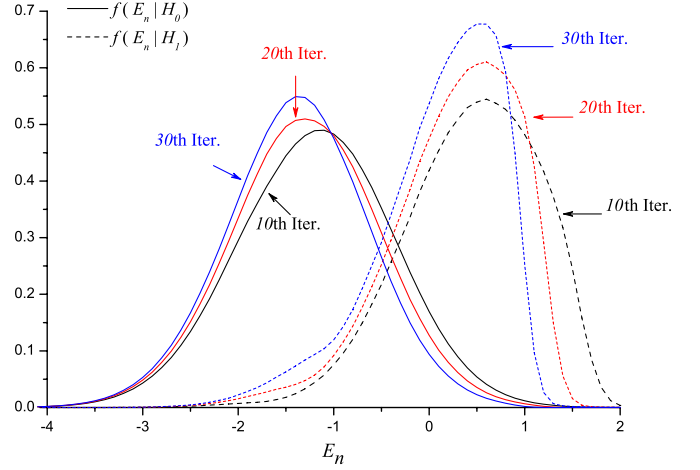


Fig. 1. Conditional FF distributions for the IMWBF algorithm in decoding MacKay (816, 272)(4, 6) LDPC code (816.1A4.845 [21]), SNR ( $E_b/N_0$ ) = 4 dB.

of the associated VNs' decisions, the corresponding checksum weights should be updated according to the current FF values. A reasonable candidate checksum weight is therefore given by

$$r_{mn}^{(l)} = \min_{n' \in \mathcal{N}(m) \setminus n} -E_{n'}^{(l)}, \quad (16)$$

where  $E_n^{(l)}$  is the FF value of  $v_n$  in the  $l$ th iteration. However, (16) may result in negative weights which is inconsistent with (2) and (3); both are increasing functions of the nonzero checksum number. To have proper positive checksum weights based on  $E_n$ , we consider the likelihood ratio

$$\Lambda(E_n) = \frac{f(E_n|H_0)}{f(E_n|H_1)}, \quad (17)$$

where  $H_0$  and  $H_1$  denote the hypotheses that  $\hat{u}_n = u_n$  and  $\hat{u}_n \neq u_n$ , respectively. The conditional probability density function (pdf),  $f(E_n|H_i)$ , is the pdf of those  $E_n$ 's associated with a correct or incorrect tentative bit decision at a given iteration. It is to be interpreted as a conditional pdf averaged over all VNs. The basic decision theory tells us that the optimal decision rule is given by

$$\Lambda(E_n) \underset{H_1}{\overset{H_0}{\gtrless}} \frac{\pi_1(C_{01} - C_{11})}{\pi_0(C_{10} - C_{00})} \quad (18)$$

where  $\pi_i = \Pr(H_i \text{ is true})$  and  $C_{ij}$  is the "cost" for accepting  $H_i$  while  $H_j$  is true. Unlike the conventional Bayesian minimizing error probability setting, both the costs and the a priori probabilities are difficult to assess. For a BF decoder, a tentative decision, except for the initial iteration, is determined by the previous decision and the flipping decision. On the other hand, both conditional pdfs in (17) depend on the definition of  $E_n$  and the FBS rule. They vary from one iteration to another. Although the dependence between  $\hat{u}_n$  and  $E_n$  is implicit as other parameters are also intertwined and the conditional pdfs are thus difficult if not impossible to derive analytically, they can be estimated numerically by simulations. Assuming that the all-zero codeword is transmitted in every frame, i.e.,  $\hat{u}_n = 0$  is the correct decision, we depict in Fig. 1 the evolutions of both

simulated conditional pdfs with all related factors considered and averaged—for the IMWBF algorithm with  $E_n$  being the updated FF values at, say the  $l$ th iteration, after the tentative decision at the  $(l-1)$ th iteration  $\hat{u}_n$  was made. This figure indicates that  $E_n$  does reflect the correctness of the corresponding tentative decision in that a smaller (negative) FF value is more likely to be associated with a correct decision. The probability of correct inference on  $\hat{u}_n$  based on  $E_n$  depends on the separation (distance) between two pdfs. The figure, however, demonstrates that the separation does not improve as the decoding iteration increases. We propose a new FF in the next subsection and show in Section IV that this new FF is capable of overcoming the shortcoming of the IMWBF algorithm.

### C. A New Flipping Function

Although  $\pi_i$  and  $C_{ij}$  are difficult to assess but the optimal Bayesian test, (18), can be simplified to a threshold test based on  $E_n$ . Given a suitable threshold  $\eta$ , Fig. 1 indicates that if  $-E_n > \eta$ , the decoder is more likely to have made a correct bit decision  $\hat{u}_n$  whose reliability is proportional to  $-E_n$ ; otherwise, the decision is probably incorrect.

The foregoing discussion and the aim to have a weight-updating rule reflecting a more accurate relation among bit decision, FF, and checksums, as elaborated in more details below, suggest that we modify (16) as

$$r_{mn}^{(l)} = \Omega \left( \min_{n' \in \mathcal{N}(m) \setminus n} -E_{n'}^{(l)} \right) = \min_{n' \in \mathcal{N}(m) \setminus n} \Omega \left( -E_{n'}^{(l)} \right), \quad (19)$$

where

$$\Omega(x) = \begin{cases} x - \eta, & x \geq \eta \\ 0, & x < \eta. \end{cases} \quad (20)$$

The clipping operator,  $\Omega(x)$ , besides ensuring only positive weights are used, can be interpreted as the decision for a CN to send no message to other linked VNs when the associated FF values fail to exceed the threshold, which bears the flavor of “stop-and-go” algorithms that pass a CN-to-VN message only if it is deemed reliable. Note that a checksum  $s_m$  is determined by  $d_c$  bit decisions, and if  $s_m = 0$  and there is only one unreliable decision  $\hat{u}_n$  ( $-E_n < \eta$ ) among VNs in  $\mathcal{N}(m)$ , the checksum is likely to be valid and the decision is in fact correct. Hence  $c_m$  should modify  $E_n$  to increase  $\Lambda(E_n)$  but not pass the message  $-r_{mn}^{(l)}(1 - 2s_m)$  to other connected (reliable) VNs ( $n' \in \mathcal{N}(m) \setminus n$ ). In doing so,  $E_n$  has a local (among  $\mathcal{N}(m)$ ) maximum FF decrease and the probability of reversing the bit decision is reduced. On the other hand, if  $s_m = 1$ ,  $\hat{u}_n$  is likely to be only local incorrect decision, the above rule will result in a local maximum FF increase and thus a higher probability of being flipped. When more than one  $-E_n$ ,  $n \in \mathcal{N}(m)$  are clipped, no message is sent from  $c_m$  as the checksum itself is unreliable. The temporary suspension of some message propagation induced by  $\Omega(x)$  also has the desired effect of containing the damage an incorrect message may have done and preventing the decoder from being trapped

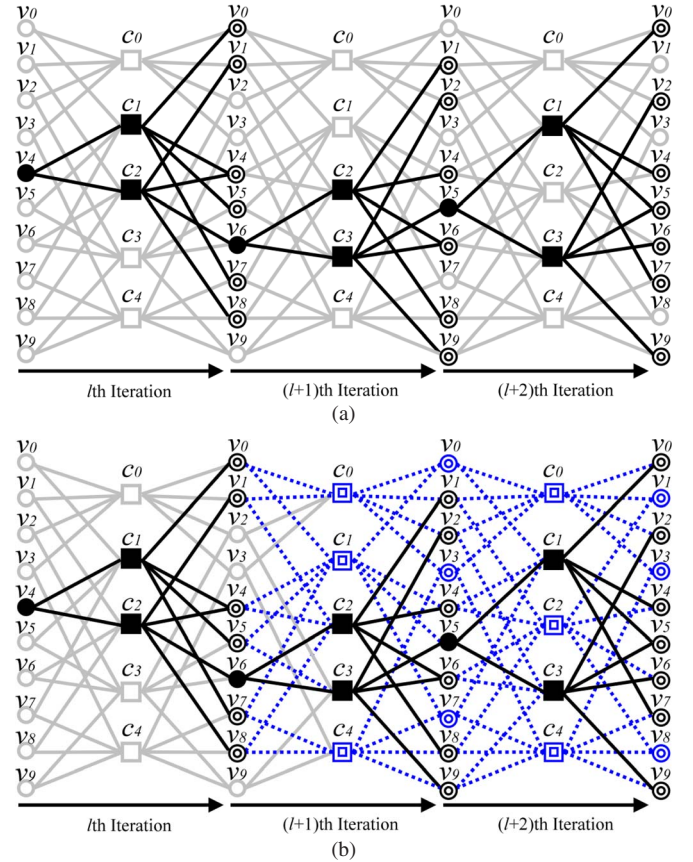


Fig. 2. The TEFGs for two different selective weight-updating schedules. (a) Selective weight-updating schedule A. (b) Selective weight-updating schedule B.

in a local minimum. The above discourse confirms that (20) does fulfill the goal that the weight updating should have FFs, checksums, and bit decisions join a cohesive effort in improving the performance of a BF decoder.

The FF defined by (3) using the recursive weight-updating rule (19) tends to make the *check reliability* part of the FF,  $-\sum_{m \in \mathcal{M}(n)} r_{mn}(1 - 2s_m)$ , starts to grow exponentially after most of the correctable bits were flipped and the number of UCNs decreases to just a few. It is conceivable that these reliability estimates should be given different weight with more remote estimates having less weights. This can be done by having the check reliabilities multiplied by a forgetting factor,  $0 < \alpha_2 < 1$ , as can be found in many recursive adaptive filters [22]. We are unable to determine the optimal clipping threshold  $\eta$  since a closed-form expressions for  $f(E_n|H_i)$  are practically unobtainable for reasons mentioned before. Some simulation efforts, however, indicate that  $\eta$  is close to 0 for several BF decoders, independent of SNR and the iteration of interest. With the above ideas in mind, we consider a new FF based on (19) and (20) using  $\eta = 0$ :

$$E_n^{(l)} = -y_n(1 - 2\hat{u}_n) - \alpha_2 \sum_{m \in \mathcal{M}(n)} r_{mn}^{(l-1)}(1 - 2s_m), \quad (21)$$

where  $0 < \alpha_2 < 1$  is a positive damping (forgetting) constant to be optimized by numerical experiments.



TABLE I  
COMPLEXITY OF TWO WEIGHT-UPDATING SCHEDULES

	SWUS-A	SWUS-B	FWUS
Number of Visited CNs	$d_v$	$\min\{d_v(d_v - 1)(d_c - 1) + d_v, M\}$	$M$
Number of Visited VNs	$\min\{d_v(d_c - 1) + 1, N\}$	$\min\{d_v(d_v - 1)(d_c - 1)^2 + [d_v(d_c - 1) + 1], N\}$	$N$

#### IV. NEW SINGLE-BIT BF DECODING ALGORITHMS

We define a (checksum) weight-updating schedule as a rule that selects a subset  $\mathcal{G}$  of  $O_M$ , which represents the set of CN indices, and updates only those  $r_{mn}^{(l)}, m \in \mathcal{G}$ . Such a rule determines the message-passing paths in the decoding process (see Fig. 2) hence is called a schedule. When the updated CN index set  $\mathcal{G} = O_M$ , we call the schedule as the *full weight-updating schedule* (FWUS). Alternate schedules with  $\mathcal{G} \neq O_M$  provide trade-offs between computational complexity and error-rate performance. In this section, we introduce a class of *single-bit dynamic weighted BF* decoding methods based on (19), (21), and different weight-updating schedules.

##### A. Single-Bit DWBF Decoding and Weight-Updating Schedules

Combining (19) and (21) with the consideration of the choice of a weight-updating schedule, we obtain the class of DWBF algorithms or, for simplicity, **Algorithm 2**.

---

##### Algorithm 2 A Class of DWBF Decoding Algorithms

---

**Initialization** Set  $l = 0$ ,  $\hat{\mathbf{u}} = \mathbf{z}$ , and  $\mathcal{G} = \mathcal{B} = \emptyset$ . Initialize  $r_{mn}^{(l)}$  by (5) for all  $n \in \mathcal{N}(m), m \in O_M$ . Set  $E_n^{(l)} = -y_n(1 - 2\hat{u}_n)$  for all  $n \in O_N$ .

**Step 1** Compute  $s_m$  for all  $m \in O_M$ . If  $s = \mathbf{0}$  or  $l = l_{\max}$ , stop decoding and output  $\hat{\mathbf{u}}$ ; otherwise,  $l \leftarrow l + 1$ .

**Step 2**  $\forall n \in O_N$ , compute  $E_n^{(l)}$  by (21).

**Step 3** Update  $\mathcal{B}$  and  $\forall n \in \mathcal{B}$ , flip  $\hat{u}_n$  and  $E_n^{(l)} \leftarrow -E_n^{(l)}$ .

**Step 4** Update  $\mathcal{G}$ . Then, update  $r_{mn}^{(l)}$  by (19)  $\forall n \in \mathcal{N}(m), m \in \mathcal{G}$  and set  $r_{mn}^{(l)} \leftarrow r_{mn}^{(l-1)} \forall n \in \mathcal{N}(m), m \in O_M \setminus \mathcal{G}$  and go to **Step 1**.

---

Algorithm 2 describes a general class of DWBF algorithms. For hard-decision decoding,  $-y_n(1 - 2\hat{u}_n)$  in (21) is replaced by  $-(1 - 2z_n)(1 - 2\hat{u}_n)$  and  $r_{mn}^{(l)}$  is initialized as 1. The FB set  $\mathcal{B}$  in **Step 3** is determined by the FBS rule used, it can be (10), **Algorithm 3** or 4 presented in Section V. When (10) is used, **Algorithm 2** is a single-bit DWBF algorithm and becomes a multi-bit DWBF algorithm if **Algorithm 3** or 4 is used as the FBS rule. In this section, we consider only the FBS rule (10), which implies that only one bit is flipped (i.e.,  $|\mathcal{B}| = 1$ ) at each iteration unless it is used for hard-decision decoding. Hence the resulting decoder is referred to as the single-bit dynamic weighted BF (S-DWBF) decoder.

As most FF values will change because of the recursive nature of (19) and (21) we may need to perform the FWUS in **Step 4**. This is one of the prices we have to pay when the dynamic weights instead of the conventional constant weights are

assigned to the checksums. We call the FWUS-based single-bit decoding algorithm as the S-DWBF-F algorithm for simplicity. To lessen the computing load of this algorithm, we reduce the size of  $\mathcal{G}$  by prioritizing the CNs and update only those with a higher priority.

We first notice that, to ensure that the newest updated information be broadcasted, the weights of the flipped bits' linked checksums should have the highest updating priority. Furthermore, for those VNs whose FF values change from one side of the clipping threshold  $\eta$  of (20) to the other side and undergo a *reliability inversion*, their related checksum weights should be renewed as well. With these considerations, the *selective weight-updating schedule A* (SWUS-A) updates only those checksums (CNs) whose indices lie in

$$\mathcal{G}_A^{(l)} \triangleq \{m | m \in \mathcal{M}(n), n \in \mathcal{B}\} \cup \left\{ m | m \in \mathcal{M}(n), \left( -E_n^{(l)} - \eta \right) \left( -E_n^{(l-1)} - \eta \right) < 0, n \in O_N \right\}. \quad (22)$$

The time-expanded factor graph (TEFG) shown in Fig. 2(a) is a simple example illustrating how SWUS-A behaves, assuming that the only VN which generates  $\mathcal{G}_A^{(l)} = \{1, 2\}$  is  $v_4$ . We denote this VN by  $\bullet$ , the CNs *visited* (selected) by the schedule by  $\blacksquare$ , and the VNs which receive new CN messages by  $\odot$ .

When SWUS-A is used in the S-DWBF algorithm, the resulting algorithm is called the S-DWBF-A algorithm which at the  $l$ th iteration updates the CN index set  $\mathcal{G}$  with  $\mathcal{G}_A^{(l)}$ .

Since only a few VNs received updated messages from the selected CNs, some  $E_n$ 's are likely to remain constant for many iterations (e.g.,  $v_3$ ) or even during the whole decoding process. To spread the updated messages to more VNs, we expand the updated CN set to include both  $\mathcal{G}_A^{(l)}$  and

$$\mathcal{G}_B^{(l)} \triangleq \left\{ m | m \in \mathcal{M}(n), n \in \mathcal{U}_A^{(l)} \right\} \quad (23)$$

where  $\mathcal{U}_A^{(l)} \triangleq \{n | n \in \mathcal{N}(m), m \in \mathcal{G}_A^{(l-1)}\}$  and call this updating schedule as the *selective weight-updating schedule B* (SWUS-B). In other words, the updated messages received by the VNs connected to the CNs in  $\mathcal{G}_A^{(l)}$  will also be forwarded to their connecting CNs in the following iteration, i.e.,  $\mathcal{G} \leftarrow \mathcal{G}_A^{(l)} \cup \mathcal{G}_B^{(l)}$ . Fig. 2(b) illustrates the expanded updating range by indicating the extra visited CNs (those whose indices belong to  $\mathcal{G}_B^{(l)}$ ) with the symbol  $\boxplus$ . Similarly, when this weight-updating schedule is used, we call the resulting algorithm as the S-DWBF-B algorithm.

Table I lists the average number of visited CNs (to update checksum weights) and VNs (to compute FF) per iteration per flipped bit or FF inversion for the different weight updating schedules. As expected, the SWUS-A/B need much less

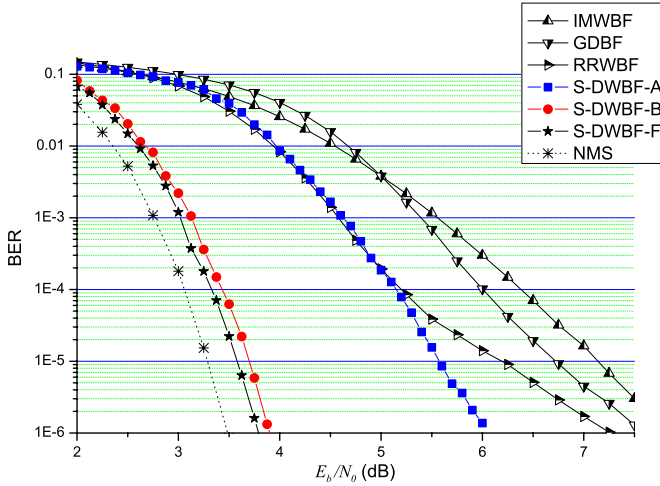


Fig. 3. BER performance of several single-bit BF (S-BF) decoders as a function of SNR ( $E_b/N_0$ ) for Code 1.

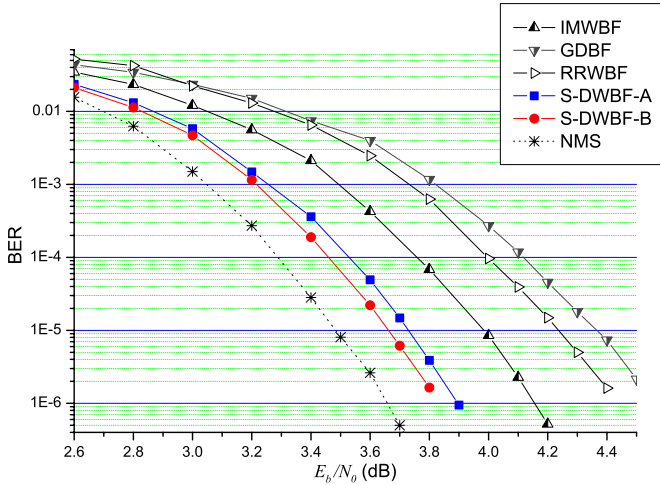


Fig. 4. BER performance of several S-BF decoder as a function of SNR for Code 2.

computational complexity and this reduction is more impressive when  $d_c$  and  $d_v$  are small. Note that for the expanded schedule SWUS-B, the VNs in  $\mathcal{U}_A^{(l)}$  may be linked to a common set of CNs and  $\mathcal{G}_A^{(l)} \cap \mathcal{G}_B^{(l)}$  can be nonempty if the code graph has short cycles. As a result, the practical average numbers of the visited CNs and VNs for SWUS-B is much less than those shown in Table I which assumes a cycle-free code. However, it can be shown that when the code girth is larger than 8 (10), the actual average visited CN (VN) number is equal to that of a cycle-free code.

### B. Performance of S-DWBF Decoding Algorithms

We apply the known single-bit BF algorithms and S-DWBF-A/B/F algorithms to decode two regular LDPC codes and present their performance in Figs. 3 and 4. The first code, MacKay's (816, 272)(4, 6) rate-0.333 LDPC code (816.1A4.845 [21]), is a typical low-rate, low-degree Gallager code with no special

structure. In contrast, the second code, the (1023, 781) (32, 32) rate-0.763 EG-LDPC code, is a high-rate, high-degree code whose performance has been evaluated in several WBF-related works [6], [8], [10], and [11]. For convenience, we refer to the above low- and high-rate codes as Code 1 and 2, respectively. It has been shown by simulations [3] that BF decoding is very effective in decoding (high-rate, high-degree) EG-LDPC codes. It would be interesting to examine the BF algorithms in decoding low-rate, low-degree codes. The performance of the normalized min-sum (NMS) algorithm [23] is also given for reference purpose.

Fig. 3 shows the BER performance of Code 1 with  $l_{\max} = 150$ . For the S-DWBF-A, S-DWBF-B, and S-DWBF-F decoders, the numerically-optimized  $\alpha_2$  values are 0.68, 0.44, and 0.35 whereas for the IMWBF decoder, we found  $\alpha_1 = 0.2$ . Our extensive simulation concluded that the optimal reliability threshold  $\eta$  in (20) is close to 0, whence  $\eta = 0$  is used for all DWBF algorithms. At  $\text{BER} = 10^{-5}$ , we observe that the S-DWBF-B and S-DWBF-F algorithms have 2.5 dB and 2.6 dB gains against the RRWBF algorithm; the simple S-DWBF-A algorithm achieves a much smaller 0.7 dB gain as it limit its weight update to a very limited range. The performance of Code 2 with  $l_{\max} = 50$  is shown in Fig. 4. Unlike Code 1, Code 2 has a much higher  $d_v$ , and the check reliability part of the GDBF algorithm's FF,  $-\sum_{m \in \mathcal{M}(n)} (1 - 2s_m)$ , thus dominates the FF value after a few iterations and its performance is similar to that of Gallager's BF algorithm, especially when SNR is high. To improve its performance we insert a damping factor  $\alpha_3$  so that (8) becomes

$$E_n = -y_n(1 - 2\hat{u}_n) - \alpha_3 \sum_{m \in \mathcal{M}(n)} (1 - 2s_m). \quad (24)$$

This modification multiplies the second summation of (8) by  $\alpha_3$ , which is analogous to the Lagrange multiplier in (the checksum) constrained optimization and when  $\alpha_3 = 1$ , (24) degenerates to (8). The optimal  $\alpha_3$  for (24) is close to 1/17 for Code 2. Referring to (3), the IMWBF algorithm uses  $\alpha_1 = 1.8$  and the S-DWBF-A (B) decoder uses  $\alpha_2 = 0.33$  (0.12) in (21). Due to the high VN/CN degrees of Code 2, almost all CNs are updated by S-DWBF-B algorithm after two or three iterations, yielding performance similar to that of the S-DWBF-F decoder. The same figure shows that the S-DWBF-A decoder provides about 0.25 dB performance gain with respect to the IMWBF decoder at  $\text{BER} = 10^{-5}$  and the S-DWBF-B algorithm offers additional 0.1 dB gain.

Table II presents the average number of CNs visited (in the associated ETFGs) by different schedules for Code 1 and Code 2 at selected iterations. Although the S-DWBF-F algorithm has the best BER performance among the single-bit algorithms, it requires higher computational complexity. By contrast, the S-DWBF-A/B algorithms provide trade-offs between complexity and performance. Furthermore, for both selective weight-updating schedules, the number of visited CNs decreases when one proceeds with more iterations as the numbers of flipped bits and reliability-inverted VNs decreases. We plot the FF value distributions for the S-DWBF-F algorithm in Fig. 5. In contrast to Fig. 1, where the separation between the two conditional

TABLE II  
AVERAGE NUMBER OF VISITED CANS IN S-DWBF ALGORITHMS

Code 1 ( $M = 544$ ); SNR = 4 dB		
Iteration	S-DWBF-A	S-DWBF-B
10	15.7	137.2
30	10.1	88.7
50	8.3	76.4
Code 2 ( $M = 1023$ ); SNR = 3.4 dB		
Iteration	S-DWBF-A	S-DWBF-B
5	105.3	1023
10	56.3	1023
20	38.6	1023

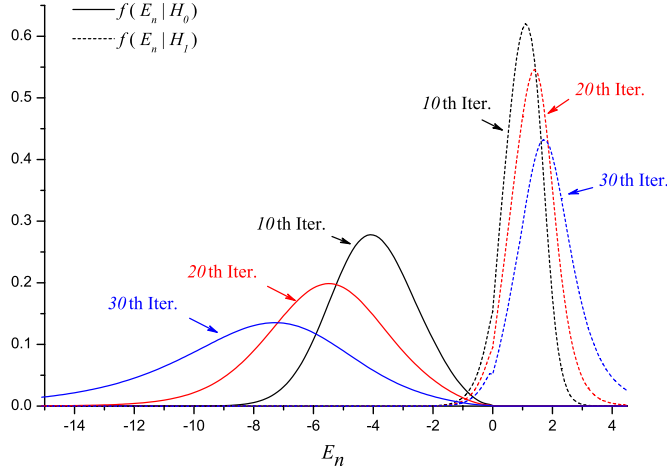


Fig. 5. Conditional FF distributions for the S-DWBF-F algorithm in decoding Code 1, SNR = 4 dB.

pdfs exhibits little variation, our DWBF algorithm is able to pull  $f(E_n|H_0)$  away from  $f(E_n|H_1)$  as the decoding process evolves. Since the reliability of a decoder's decisions based on  $E_n$  depends on the separation (distance) between the two pdfs, the improved separation is certainly welcome. As mentioned before, we use  $\eta = 0$  in (20) for all S-DWBF algorithms. Although the optimal clipping threshold is unknown, Fig. 5 does convince us that 0 is a valid and convenient choice and the FF with the proposed dynamic checksum weighting does give a much better VN reliability reference.

## V. NEW FLIPPED-BIT SELECTION METHODS AND MULTI-BIT BF DECODING ALGORITHMS

Multi-bit BF decoding algorithms were developed to accelerate the convergence performance. Most of these algorithms use the simple threshold comparison FBS rule (11) discussed in Section II [9], [12], [13], [15]–[18]. Even if the threshold  $\Delta$  has been optimized by numerical experiments, we find that it is necessary to add the option  $\mathcal{B} = \{n|n = \arg \max_i E_i\}$  in case  $\mathcal{B} = \emptyset$  whose occurrence probability during a decoding process is nonzero. In fact, simulation results indicate that, depending on SNR and the decoding/schedule algorithms used it can be higher than 0.2 for Code 1. When this option is included in FB decision, the resulting FBS rule is called **Algorithm 3** or the M1-FBS rule.

### Algorithm 3 Flipped Bit Selection Rule 1 (M1-FBS)

**Step 1** Find  $\mathcal{B} = \{n|E_n \geq \Delta\}$ . If  $\mathcal{B} \neq \emptyset$ , stop; otherwise, proceed to **Step 2**.  
**Step 2** Find  $\mathcal{B} = \{n|n = \arg \max_i E_i\}$

Recall that the PWB algorithm uses the FF of the IMWBF algorithm as the VN reliability metric with the FB set (13) determined by the FS count  $F_n$  of (14). That is, a CN sends a flip signal to its most unreliable linked VN only and the VNs which receive sufficient number of reliability warnings (flip signals) shall be flipped. It turns out that, with this extra filtering of VN-to-CN messages ( $E_n$ 's) and selective CN-to-VN flip signal passing, the PWB algorithm is able to outperform the IMWBF decoder in both convergence rate and error rate [10]. This performance gain motivates us to ponder if a more elaborated flipping decision strategy that uses more information can bring about further performance improvement for the BF decoders using either the proposed FF (21) or other FFs discussed in Section II.

#### A. Flipping Intensity

Let  $U_n \triangleq \sum_{m \in \mathcal{M}(n)} s_m$ ,  $\mu_m \triangleq \max_{n \in \mathcal{N}(m)} U_n$ ,  $\lambda_m \triangleq \arg \max_{n \in \mathcal{N}(m)} E_n$ , and  $\mathcal{M}'(n) = \{m|m \in \mathcal{M}(n), \lambda_m = n\}$ . With these notations, we define the *flipping intensity* (FI) of (received by)  $v_n$  as

$$\tilde{F}_n = \sum_{m \in \mathcal{M}'(n)} \theta_0 s_m \delta(U_{\lambda_m} - \mu_m) + \theta_1 s_m [1 - \delta(U_{\lambda_m} - \mu_m)], \quad (25)$$

where  $\theta_0 > \theta_1 \geq 0$  and  $\delta(x)$  is the Kronecker delta function. For simplicity, both  $\theta_0$  and  $\theta_1$  are confined to be integers so that FI is integer-valued. The above definition implicitly implies that  $\tilde{F}_n = 0$  if  $\mathcal{M}'(n) = \emptyset$  and only UCNs have a say in deciding FI. It also implies that a VN has a nonzero FI only if it has the largest FF value among  $\mathcal{N}(m)$  and if it is connected to the largest number of UCNs among its peers in  $\mathcal{N}(m)$ , the associated FI should be even higher ( $\theta_0 > \theta_1$ ). In both cases, a UCN  $c_m$  will send a non-negative message to the VN with the highest FF value in the set  $\mathcal{N}(m)$ . However, if  $c_m$  is a passed CN (PCN) ( $s_m = 0$ ) and  $d_c$  is small, it often implies that the tentative decisions of its linked VNs are all correct. Hence if the flipped bits are to be selected by checking whether the associated FI is greater than a threshold,  $v_{\lambda_m}$  should have a smaller probability of being chosen. This can be done by having the PCN send a drag message  $\theta_2(s_m - 1)$ . But if there is doubt that  $c_m$  is connected to even incorrect bit decisions, the PCN has better not sending such a message. We decide that this is likely to be the case if  $U_{\lambda_m} \neq \mu_m$  for this inequality means that at least one VN in  $\mathcal{N}(m)$  has more connected UCNs than  $v_{\lambda_m}$ . With UCNs and PCNs contributing opposite signals, we modify (25) for all  $n$ ,  $0 \leq n < N$  as

$$\tilde{F}_n = \sum_{m \in \mathcal{M}'(n)} [\theta_2(s_m - 1) + \theta_0 s_m] \delta(U_{\lambda_m} - \mu_m) + \theta_1 s_m [1 - \delta(U_{\lambda_m} - \mu_m)], \quad (26)$$



where  $\theta_2 \leq \theta_0$  is a nonnegative integer. On the other hand, when  $d_c$  is large, it is less likely that  $s_m = 0$  automatically implies correct decisions on all its linked bits and we thus stick to (25), having no PCN to contribute to FI. Although the thresholds  $\theta_i$ 's can be any nonnegative real numbers, to simplify implementation, we let  $\theta_i$ 's be nonnegative integers such that the FI is integer-valued.

### B. Flipped-Bit Selection Rule

A simple FI-based FBS rule is to flip the bits in the FB set  $\mathcal{B} = \{n | \tilde{F}_n \geq \Delta_{\text{FI}}\}$ . But the optimal threshold  $\Delta_{\text{FI}}$  is not easy to determine especially for a code with low VN degree. A smaller threshold may cause incorrect flipping decisions while a large threshold tends to slow down the convergence or even cause decoding failure as no VN meets the the flipping requirement. To overcome this dilemma, we select a relative high FI threshold and use the FB set  $\mathcal{B} = \{n | \tilde{F}_n \geq \Delta_{\text{FI}}\}$  if it is nonempty. Otherwise,  $\mathcal{B} = \{n | U_n = \max_{i \in \mathcal{T}} U_i\}$  where  $\mathcal{T} \triangleq \{n | \tilde{F}_n = \max_j \tilde{F}_j\}$ . We summarize below the new FBS rule as **Algorithm 4** which, for convenience of reference, is called the M2-FBS rule.

---

#### Algorithm 4 Flipped Bit Selection Rule 2 (M2-FBS)

---

- Step 1** For  $n = 0, 1, \dots, N - 1$ , compute  $\tilde{F}_n$  by (25) or (26).  
**Step 2** Find  $\mathcal{B} = \{n | \tilde{F}_n \geq \Delta_{\text{FI}}\}$ . If  $\mathcal{B} \neq \emptyset$ , stop; otherwise, proceed to **Step 3**.  
**Step 3** Update  $\mathcal{T}$  and find  $\mathcal{B} = \{n | U_n = \max_{i \in \mathcal{T}} U_i\}$ .
- 

Note this FBS rule is independent of the FF and can be used in conjunction with different FFs no matter whether the checksum weights are constant or not.

Loop-detection/breaking procedures can be included in our FBS algorithm if necessary. The loop detection scheme used [8] is an appropriate choice. When a loop is detected, we generate a disturbance on the tentative decoded sequence by switching to the FB set

$$\mathcal{B} = \{n | U_n = \max_i U_i\}. \quad (27)$$

### C. Numerical Results

Different combinations of the FBS rule, the FF, and the weight-updating schedule, used lead to different decoding algorithms. The error-rate performance and decoding speed of various decoders are presented in this subsection.

1) *Abbreviations*: For convenience of reference, we adopt a systematic labeling scheme similar to that used in Section IV-B to describe a decoding method. We denote a decoder by groups of capital letters separated by hyphens, specifying respectively the FBS rule, the FF and the weight-updating schedule used with the 3-field form, *FBS rule-FF-weight updating schedule*. That is, the first field is used to indicate if single (S) or multiple (M) bits are to be flipped in an iteration and, for the latter case, if the simple FF based (M1) or the more complicated FI-based

TABLE III  
SIMULATION PARAMETER SETTINGS

Algorithm	Code 1	Code 2
HGDBF	$\alpha_3 = 1$	$\alpha_3 = 1/17$
AMWBF	$\alpha_1 = 0.2$	$\alpha_1 = 1.8$
IPWBF	$\alpha_1 = 0.2, \Delta_{\text{FS}} = 1$	$\alpha_1 = 1.8, \Delta_{\text{FS}} = 10$
M2-IMWBF	$\alpha_1 = 0.2, \Delta_{\text{FI}} = 5$	$\alpha_1 = 3.2, \Delta_{\text{FI}} = 16$
M2-GDBF	$\alpha_3 = 1, \Delta_{\text{FI}} = 1$	$\alpha_3 = 1/17, \Delta_{\text{FI}} = 10$
M1-DWBF-A	$\alpha_2 = 0.7, \Delta = 0$	$\alpha_2 = 0.33, \Delta = 0$
M1-DWBF-B	$\alpha_2 = 0.35, \Delta = 0$	$\alpha_2 = 0.3, \Delta = 0$
M2-DWBF-A	$\alpha_2 = 0.58, \Delta_{\text{FI}} = 1$	$\alpha_2 = 0.33, \Delta_{\text{FI}} = 4$
M2-DWBF-B	$\alpha_2 = 0.35, \Delta_{\text{FI}} = 1$	$\alpha_2 = 0.3, \Delta_{\text{FI}} = 1$

(M2) FBS rule is adopted. The second field contains the abbreviation of the known or proposed algorithm such as IMWBF, GDBF or DWBF whose FF is used. The third field tells whether a selective (A or B) or the full (F) weight-updating schedule is used. Since only the DWBF algorithms need to update checksum weights, the third field is omitted for non-DWBF based decoders. Hence, M1-DWBF-A represents the decoder that uses the M1-FBS rule, the DWBF FF, and SWUS-A, and M2-IMWBF(-GDBF) denotes the decoder that uses the M2-FBS rule and the IMWBF (GDBF) algorithm's FF. For known constant weight algorithms without FBS modification and SWUS, we keep conventional abbreviations like AMWBF, IPWBF, and HGDBF only.

2) *Parameter Values Used*: For the decoders based on M2-FBS rule, we use the FI weights  $\theta_0 = 3, \theta_1 = 2$ , and  $\theta_2 = 1$ . Other major parameter values for different multi-bit BF algorithms are listed in Table III. The remaining parameters needed for the IPWBF algorithm follow those suggested in [11], and those associated with the escaping process in the HGDBF algorithm are also optimized. For simplicity, the FF clipping threshold  $\eta$  in (20) and the threshold  $\Delta$  used in (11) in M1-DWBF-A/B algorithms are set to zero.

Note that the parameters associated with a decoder are correlated, i.e., if the optimal value of a parameter is dependent on other parameters' values used, although the correlation may not be very high. Hence, we try to jointly optimize these parameters to minimize the converged error rate. Furthermore, simulation results indicate that the optimal parameter values are insensitive to SNR.

3) *BER and FER Performance*: Fig. 6 shows the BER performance of different multi-bit BF algorithms for Code 1 when  $l_{\text{max}} = 50$ . The effectiveness of the M2-FBS rule can also be verified by comparing the required  $E_b/N_0$  for  $\text{BER} = 10^{-5}$ : the M2-IMWBF decoder outperforms the IPWBF decoder by approximately 1.7 dB and the M2-GDBF algorithm has a 0.4 dB gain over the HGDBF algorithm. The DWBF algorithms yields better BER performance even with the simple M1 rule and, when the M2 rule is used, its performance becomes closer (0.4 dB) to that provided by the NMS algorithm. The convergence behaviors of these algorithms are shown in Figs. 7 and 8. The results show that the M2 rule gives better BER performance and, for both the DWBF and M2-GDBF algorithms, the convergence rate is improved as well.

Note that in Figs. 6–8, loop-detecting/breaking schemes are activated for all but the M2-DWBF-B algorithm. In general, loops are much less likely to occur in a DWBF decoder than in a

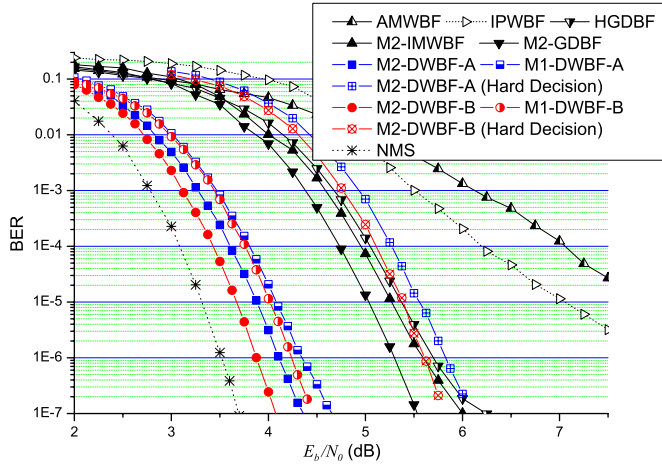


Fig. 6. BER performance of various multi-bit BF (M-BF) decoding algorithms as a function of SNR for Code 1.

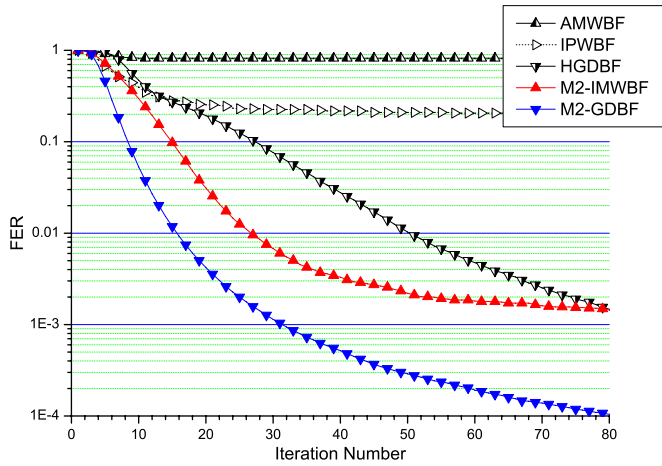


Fig. 7. Frame error rate (FER) convergence performance of various M-BF decoding algorithms using conventional or M2-FBS rule; Code 1, SNR = 5 dB.

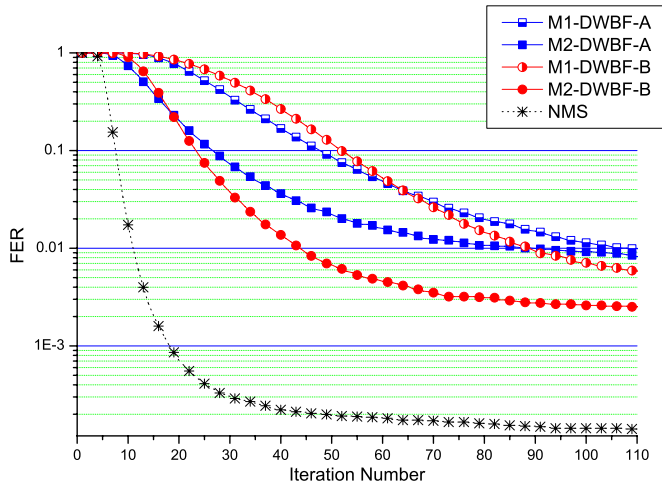


Fig. 8. FER convergence performance of multi-bit DWBF (M-DWBF) decoding algorithms; Code 1, SNR = 3.25 dB.

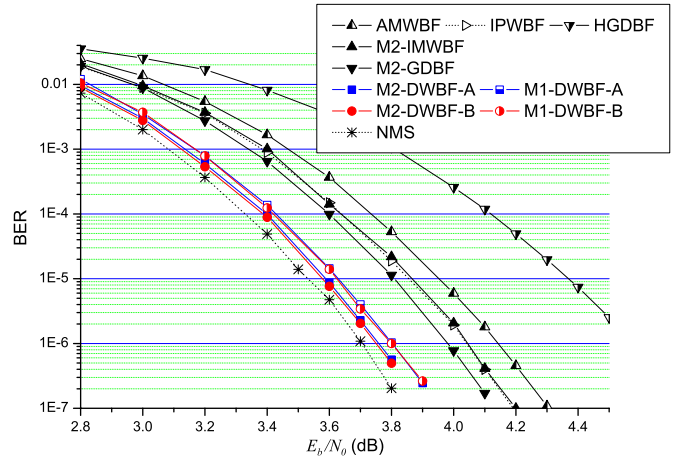


Fig. 9. BER performance of various M-BF decoding algorithms as a function of SNR for Code 2.

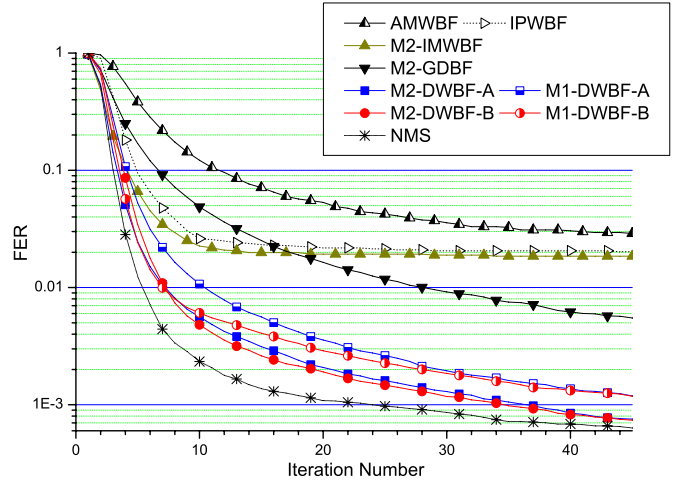


Fig. 10. FER convergence performance of several M-BF decoding algorithms; Code 2, SNR = 3.4 dB.

static CN weight decoder. When the FWUS or SWUS-B is used to decode Code 2, our simulations detect no loop for both codes whence there is no need for a loop breaker. This is because the time-varying checksum weights of the DWBF algorithm and wider message magnitude propagation ranges of the FWUS or SWUS-B schedule have made the BF decision related variables,  $E_n$ ,  $U_n$ ,  $\mu_m$ , and  $\tilde{F}_n$ , to have much larger dynamic ranges; see also Figs. 1 and 5.

The BER (with  $l_{\max} = 20$ ) and frame error rate (FER) convergence performance of various multi-bit BF decoders for Code 2 are respectively presented in Figs. 9 and 10. By comparing the two sets of BER curves, M2-GDBF versus HGDBF and M2-IMWBF versus AMWBF, we verify the effectiveness of the new FBS (M2) rule. Although the M2-IMWBF algorithm yields the same converged BER as that of the IPWBF decoder for this code, it gives better FER performance in the first few iterations. We further notice that the M2-DWBF-A (B) decoder is superior to the M1-DWBF-A (B) decoder in both BER performance and decoding speed. The former yields performance very close to

that of the NMS algorithm while the latter suffers only about 0.1 dB performance degradation against the NMS decoder at  $\text{BER} = 10^{-5}$ . Figs. 9 and 10 also show that the performance gap between the M1/M2-DWBF-A and M1/M2-DWBF-B decoders is smaller than that for Code 1. This is due to the high VN/CN degrees of Code 2: the high VN degree increases the probability that a CN is visited by the SWUS-A while the high CN degree helps spreading the updated weights to more VNs.

We want to remark that a) only the AMWBF and HGDBF decoders need a loop-breaker in decoding Code 2 since for the other decoders, loops are rarely detected and b) although the M1-FBS rule is simpler, our simulations indicate that the M2-FBS rule can significantly reduce the probability of decoding loops when it is used in conjunction with the DWBF, GDBF, or IMWBF algorithms. This is particular useful when using conventional FFs to decode low-degree codes.

#### D. Complexity Analysis

Besides the syndrome computing, which is the same for all algorithms, the computational complexity of a BF decoding algorithm consists mainly of three parts: i) FF update, ii) flipped bits selection, and iii) weight/message update. Once new CN messages,  $-w_{mn}(1 - 2s_m)$  in (3) or  $-r_{mn}(1 - 2s_m)$  in (21), are available, the FF update is just adding all returned CN messages and  $-\alpha_1|y_n|$  or  $-y_n(1 - 2\hat{u}_n)$ . There is little difference among the BF decoders in FF computing. The only exception is that used by GDBF algorithms, both (8) and (24) require integer additions only. iii) is needed for DWBF algorithms but not other BF algorithms which require only a sign change. Therefore, in the next three subsections, we consider ii) first, followed by the discussion of iii), the extra complexity requirement for DWBF algorithms, and finally compare the combined computational complexity of ii) and iii). The additional complexity such as that associated with a loop-breaking scheme, is addressed at the end of this section as well. Since most algorithms, except those using the M2-FBS rule, which need additional integer operations and memory for storing the FIs and UCN numbers, require approximately the same storage space, we discuss only the computational complexity.

1) *FBS Complexity*: The HGDBF decoder needs only  $N$  real comparisons in (11) in selecting the flipped bits. For the IPWBF decoder,  $d_c - 1$  real comparisons are required to find the most unreliable connected VN per UCN and a total of  $M_1^{(l)}$  integer additions and  $N$  integer comparisons are needed to compute the FS and decide the FB set (13), where  $M_1^{(l)}$  is the average UCN number in the  $l$ th iteration.

For the M1-FBS rule, the complexity of **Step 1** is ignored since we use  $\Delta = 0$  and a threshold comparison needs a sign-bit check only. The average complexity of **Step 2** is  $P_{S2}^{(l)}(N - 1)$  real comparisons, where  $P_{S2}^{(l)}$  is the probability that **Step 2** is activated in the  $l$ th iteration. For the M2-FBS rule (**Algorithm 4**), **Step 1** needs  $d_c - 1$  real and  $d_c - 1$  integer comparisons per CN in finding  $\lambda_m$  and checking if  $v_{\lambda_m}$  has the most connected UCNs. Each CN has to send an integer-valued message,  $\theta_0$ ,  $\theta_1$ , or  $-\theta_2$ , to one of its connected VNs, implying an integer addition in (26) or (25). Since the former involves both UCNs

and PCNs while the latter involves only UCNs, all  $M$  CNs have to perform all the above operations when decoding Code 1 in contrast to  $M_1^{(l)}$  CNs for Code 2. Moreover,  $N$  integer comparisons are required in **Step 2** for deciding the FB set. The average complexity of **Step 3** is approximately equal to  $P_{S3}^{(l)}N$  integer comparisons, where  $P_{S3}^{(l)}$  is the probability that **Step 3** is activated at the  $l$ th iteration.

2) *Weight Update Complexity*: Updating the weights associated with CN  $c_m$  in the M1-DWBF-A/B decoders require  $2d_c - 3$  real comparisons for finding the indices associated with the smallest and second smallest  $-E_n$ 's,  $n \in \mathcal{N}(m)$ . For M2-DWBF-A/B decoders, however, most of  $\lambda_m$ 's have been found in the FBS step, hence only the second smallest ones remain to be found for computing new weights. A more detailed analysis is given in the next two paragraphs.

We first consider a low CN degree code such as Code 1. After CNs compute  $\lambda_m$ 's and VNs compute their FIs via (26), the flipped bits are decided and flipped. The M2-DWBF-A/B algorithms then invert the associated FF values (**Step 3** of **Algorithm 2**) and update the CN weights (**Step 4** of **Algorithm 2**). As only a small portion of the visited CNs are connected to the flipped bits, most visited CNs require only  $d_c - 2$  real comparisons for finding the second smallest  $-E_n$ . For a visited CN that links to flipped bits, we only need to compare the connected flipped bits'  $E_n$ 's with the original smallest  $-E_n$  to find a new minimum  $-E_n$  since only the flipped bits'  $E_n$ 's are changed between the bit flipping and weight updating. As a result, for a visited CN linking to  $t$  flipped bits, only additional  $t$  real comparisons are required for updating the smallest  $-E_n$  before finding the second smallest  $-E_n$ . If we denote by  $M_{\text{WU}}^{(l)}$  and  $N_{\text{FB}}^{(l)}$  the average numbers of visited CNs and flipped bits at the  $l$ th iteration, respectively, we need, on the average, at most  $M_{\text{WU}}^{(l)}(d_c - 2) + N_{\text{FB}}^{(l)}d_v$  real comparisons for updating weights, where  $N_{\text{FB}}^{(l)}d_v$  accounts for the sum of all additional ( $t$ ) comparisons.

Decoding a high CN degree code requires  $2d_c - 3$  real comparisons for updating the weight of a visited PCN and  $d_c - 2 + t$  for an UCN, as the FI formula (25) involves only UCNs. Denote the average numbers of visited PCNs and UCNs at the  $l$ th iteration by  $M_{\text{WU},0}^{(l)}$  and  $M_{\text{WU},1}^{(l)}$ . We observe from simulations that almost all UCNs are visited (i.e.,  $M_{\text{WU},1}^{(l)} = M_1^{(l)}$ ). Hence, M2-DWBF-A/B decoders require an average of  $M_{\text{WU},0}^{(l)}(2d_c - 3)$  and at most  $M_1^{(l)}(d_c - 2) + N_{\text{FB}}^{(l)}d_v$  real comparisons per iteration for computing new CN weights of visited PCNs and UCNs, respectively. For both cases we ignore the complexity of the threshold comparison in (19) since  $\eta = 0$ .

3) *Complexity Summary*: Table IV presents the simulated average numbers of  $M_1^{(l)}$ ,  $M_{\text{WU}}^{(l)}$ ,  $M_{\text{WU},0}^{(l)}$ ,  $N_{\text{FB}}^{(l)}$ ,  $P_{S2}^{(l)}$ , and  $P_{S3}^{(l)}$  at selected iterations for the IPWBF, M2-IMWBF/GDBF, M1-DWBF-A/B, and M2-DWBF-A/B algorithms. Since the simulation results indicate that when decoding Code 2 with the M1-DWBF-A/B decoders, **Step 2** of the M1-FBS rule is *never* activated, we list the  $P_{S2}^{(l)}$  values for Code 1 only. Similarly, **Step 3** of the M2-FBS rule is needed *only if* the M2-IMWBF algorithm is used to decode Code 1, we thus specify the  $P_{S3}^{(l)}$  values for this case only.

TABLE IV  
AVERAGE UCN AND VISITED CN NUMBERS AT SOME SELECTED ITERATIONS

Code 1 ( $M = 544$ ); SNR = 3.25 dB														
Iteration ( $l$ )	IPWBF	M2-IMWBF	M2-GDBF	M1-DWBF-A		M1-DWBF-B		M2-DWBF-A			M2-DWBF-B			
	$M_1^{(l)}$	$M$	$P_{S3}^{(l)}$	$M$	$M_{WU}^{(l)}$	$P_{S2}^{(l)}$	$M_{WU}^{(l)}$	$P_{S2}^{(l)}$	$M$	$M_{WU}^{(l)}$	$N_{FB}^{(l)}$	$M$	$M_{WU}^{(l)}$	$N_{FB}^{(l)}$
5	91.7	544	0.36	544	115.2	$3.6 \times 10^{-5}$	386.7	$3.5 \times 10^{-6}$	544	135.2	34.1	544	506.6	17.2
10	78.3		0.52		110.6	0.0035	162.8	0.0014		80.1	19.1		414.2	8.8
15	71.0		0.67		107.7	0.0092	132.1	0.0019		69.0	16.5		338.1	6.3
20	68.6		0.65		96.5	0.0099	131.3	0.0012		66.9	15.8		328.7	5.9
Code 2 ( $M = 1023$ ); SNR = 3.4 dB														
Iteration ( $l$ )	IPWBF	M2-IMWBF	M2-GDBF	M1-DWBF-A		M1-DWBF-B		M2-DWBF-A			M2-DWBF-B			
	$M_1^{(l)}$	$M_1^{(l)}$	$M_1^{(l)}$	$M_{WU}^{(l)}$	$M_{WU}^{(l)}$	$M_{WU}^{(l)}$	$M^{(l)}$	$M_{WU,0}^{(l)}$	$N_{FB}^{(l)}$	$M_1^{(l)}$	$M_{WU,0}^{(l)}$	$N_{FB}^{(l)}$		
3	172.1	217.2	326.2	267.3	1023	281.6	382.4	18.3	332.4	690.6	24.6			
5	198.3	287.7	351.9	374.9	1023	318.5	386.9	26.1	322.7	700.3	27.0			
10	329.4	360.2	407.2	588.7	1023	425.8	415.4	46.4	436.5	586.5	50.7			
15	353.6	378.4	414.6	644.5	1023	438.1	421.0	48.2	445.0	577.0	59.5			

TABLE V  
COMPUTATIONAL COMPLEXITIES FOR VARIOUS DECODING ALGORITHMS (C1: CODE 1, C2: CODE 2)

Operation	HGDBF	IPWBF	M2-GDBF/IMWBF	M1-DWBF-A/B	M2-DWBF-A/B	NMS
Integer Additions	0	$M_1^{(l)}$	C1: $M$ C2: $M_1^{(l)}$	0	C1: $M$ C2: $M_1^{(l)}$	0
Real Number Additions	0	0	0	0	0	$Nd_v$
Integer Comparisons	0	$N$	C1 (M2-IMWBF): $M(d_c - 1) + N + P_{S3}^{(l)}N$ C1 (M2-GDBF): $M(d_c - 1) + N$ C2: $M_1^{(l)}(d_c - 1) + N$	0	C1: $M(d_c - 1) + N$ C2: $M_1^{(l)}(d_c - 1) + N$	0
Real Number Comparisons	$N$	$M_1^{(l)}(d_c - 1)$	C1: $M(d_c - 1)$ C2: $M_1^{(l)}(d_c - 1)$	C1: $M_{WU}^{(l)}(2d_c - 3) + P_{S2}^{(l)}(N - 1)$ C2: $M_{WU}^{(l)}(2d_c - 3)$	C1: $M(d_c - 1) + M_{WU}^{(l)}(d_c - 2) + N_{FB}^{(l)}d_v$ C2: $(M_1^{(l)} + M_{WU,0}^{(l)})(2d_c - 3) + N_{FB}^{(l)}d_v$	$M(2d_c - 3)$

TABLE VI  
AVERAGED OVERALL PER FRAME COMPLEXITIES ( $\times 10^3$ ) FOR ACHIEVING FER =  $10^{-3}$  (REAL: REAL COMPARISON OR ADDITION; INT: INTEGER COMPARISON OR ADDITION)

Code 1												
SNR	M1-DWBF-A		M1-DWBF-B		M2-DWBF-A			M2-DWBF-B			NMSA	
	$l_{\max}$	Real	$l_{\max}$	Real	$l_{\max}$	Real	Int.	$l_{\max}$	Real	Int.	$l_{\max}$	Real
3.5 dB	132	19.1	101	46.4	109	41.7	51.1	47	64.1	54.4	12	44.2
3.625 dB	89	15.3	76	41.7	66	36.2	49.6	35	57.9	49.5	10	41.6
3.75 dB	71	12.5	64	38.1	42	32.1	45.5	28	52.8	45.4	9	39.2
Code 2												
SNR	M1-DWBF-A		M1-DWBF-B		M2-DWBF-A			M2-DWBF-B			NMSA	
	$l_{\max}$	Real	$l_{\max}$	Real	$l_{\max}$	Real	Int.	$l_{\max}$	Real	Int.	$l_{\max}$	Real
3.4 dB	49	83.5	46	172.5	35	176.4	66.7	34	269.1	80.3	22	255.9
3.6 dB	11	59.6	8	135.2	8	130.7	50.4	7	210.7	62.3	6	218.2
3.7 dB	7	53.3	6	123.0	6	115.4	44.9	6	189.3	55.7	5	206.8

Considering both the FBS rules and weight updating, we summarize the computational complexity which includes real/integer additions and comparisons per iteration for various BF and the NMS algorithms in Table V. As computing the total LLRs in the NMS algorithm requires the same efforts as that of computing FF values in BF decoders, only the efforts needed for computing the CN-to-VN and VN-to-CN messages are listed in the table.

Table VI presents the simulated average complexity of various DWBF and the NMS algorithms for decoding a frame with a target FER of  $10^{-3}$  at different SNRs. As an integer (or real) comparison requires about the same computational complexity as that of an integer (or real) addition (hardware implementation of comparison can even be simpler than addition). Both are thus counted equally. We show in this table the average integer

and real operations and the maximum iteration number ( $l_{\max}$ ) needed.

Tables IV–VI and Figs. 6–10 provide useful information for studying tradeoffs between performance, complexity, and convergence rate when combining different FBS rules, FFs and weight-updating schedules. In particular, Table VI shows that when decoding Code 1, both M1-DWBF-A and M1-DWBF-B algorithms need less computational complexity to achieve FER =  $10^{-3}$  than that needed by the NMS algorithm in higher SNR (say, >3.625 dB) region. The M2-DWBF-A algorithm also needs less real operations in comparison with the NMS algorithm. For decoding Code 2, the M1-DWBF-A, M1-DWBF-B, and M2-DWBF-A algorithms need less total (real + integer) operations to achieve the FER requirement while the required iteration numbers are also comparable to that needed for the



NMS algorithm in higher SNR region. Furthermore, based on Tables IV, V, and Fig. 6, we conclude that the M2-IMWBF and M2-GDBF algorithms require far less complexity than that of the IPWBF algorithm in decoding Code 1.

Among the the decoding algorithms compared in Table V, the IPWBF algorithm uses a simpler FBS operation but it has to perform a delay-handling process in every iteration plus an initial bootstrapping step. These two extra operations need off-line computing effort in searching for the corresponding optimal parameter values. They also require additional storage and computational complexity. Although the HGDBF algorithm does not have to sort the FF values, three real thresholds, one for the multi-bit flipping mode and two for the escape (loop-breaking) process are required in its FBS rule, resulting extra off-line search and random variable generations. Other off-line efforts include the searches for  $\alpha_1$  (M2-IMWBF),  $\alpha_2$  (M1- and M2-DWBF-A/B), and  $\Delta_{FI}$  ( $\in [-d_v\theta_2, d_v\theta_0]$ ). As mentioned before, they must be jointly optimized. For the M2-GDBF algorithm, only the optimal  $\Delta_{FI}$  has to be found.

Our loop-breaking scheme (27) is simpler than those used by other decoding algorithm and more effective than the methods used by the IPWBF and AMWBF algorithms which remove the bit(s) having maximum  $F_n$  or  $E_n$  from  $\mathcal{B}$ ; when  $\mathcal{B} = \emptyset$ , the decoding process will be forced to terminate after the removal. Instead of reducing  $|\mathcal{B}|$ , the escape process (27) and that used by the HGDBF algorithm perturb the tentative decoded sequence to break a loop. The latter, however, has to generate Gaussian random variables.

## VI. CONCLUSION

We divide a typical BF LDPC code decoding algorithm into three major components, namely 1) VN decision reliability (FF) computing and the associated CN reliability (checksum weight) update formula; 2) the FBS rule; and 3) the checksum weight-updating schedule. These three components determine the performance and complexity of a BF decoder. We develop novel FF and FBS rules to improve the BF decoding performance. On the other hand, the checksum weight update operation is a complexity concern for the DWBF decoders, we propose selective weight-updating schedules to reduce the implementation complexity with little performance loss.

Different combinations of FF, checksum weight-updating method and schedule, FBS rule, result in different decoder structures. We simulate the error rate and convergence performance of various decoders, and the resulting numerical behaviors confirm the effectiveness of our new design proposals. We show that the combinations of the new multi-bit FBS rules with known BF algorithms achieve significant performance gain especially for a high-rate code. Detailed complexity analysis on various decoder structures is provided for complexity and performance tradeoff studies. We find that, compared with the NMS algorithm, the combination of the new FBS rules with our DWBF algorithms require less complexity in achieving a target FER if SNR is sufficiently high. We also find that the convergence rates are comparable when decoding a high-rate code.

## ACKNOWLEDGMENT

The authors would like to thank the editor, anonymous reviewers, and Jian-Jia Weng for their valuable comments and suggestions to improve the quality and readability of the paper.

## REFERENCES

- [1] D. J. C. MacKay, "Good error-correcting codes based on very sparse matrices," *IEEE Trans. Inf. Theory*, vol. 45, no. 2, pp. 399–431, Mar. 1999.
- [2] R. G. Gallager, *Low-Density Parity-Check Codes*. Cambridge, MA, USA: MIT Press, 1963.
- [3] W. E. Ryan and S. Lin, *Channel Codes: Classical and Modern*. Cambridge, U.K.: Cambridge Univ. Press, 2009, ch. 10, pp. 430–481.
- [4] Y. Kou, S. Lin, and M. Fossorier, "Low density parity check codes based on finite geometries: A rediscovery and more," *IEEE Trans. Inf. Theory*, vol. 47, no. 7, pp. 2711–2736, Nov. 2001.
- [5] J. Zhang and M. Fossorier, "A modified weighted bit-flipping decoding of low-density parity-check codes," *IEEE Commun. Lett.*, vol. 8, no. 3, pp. 165–167, Mar. 2004.
- [6] M. Jiang, C. Zhao, Z. Shi, and Y. Chen, "An improvement on the modified weighted bit flipping decoding algorithm for LDPC codes," *IEEE Commun. Lett.*, vol. 9, no. 9, pp. 814–816, Sep. 2005.
- [7] F. Guo and L. Hanzo, "Reliability ratio based weighted bit-flipping decoding for low-density parity-check codes," *Electron. Lett.*, vol. 40, no. 21, pp. 1356–1358, Oct. 2004.
- [8] T. Ngatched, F. Takawira, and M. Bossert, "A modified bit-flipping decoding algorithm for low-density parity-check codes," in *Proc. IEEE Int. Conf. Commun.*, Glasgow, U.K., 2007, pp. 653–658.
- [9] T. Wadayama *et al.*, "Gradient descent bit flipping algorithms for decoding LDPC codes," *IEEE Trans. Commun.*, vol. 58, no. 6, pp. 1610–1614, Jun. 2010.
- [10] X. Wu, C. Zhao, and X. You, "Parallel weighted bit-flipping decoding," *IEEE Commun. Lett.*, vol. 11, no. 8, pp. 671–673, Aug. 2007.
- [11] G. Li and G. Feng, "Improved parallel weighted bit-flipping decoding algorithm for LDPC codes," *IET Commun.*, vol. 3, no. 1, pp. 91–99, Jun. 2009.
- [12] J. Cho and W. Sung, "Adaptive threshold technique for bit-flipping decoding of low-density parity-check codes," *IEEE Commun. Lett.*, vol. 14, no. 9, pp. 857–859, Sep. 2010.
- [13] T. Chen, "Adaptive-weighted multibit-flipping decoding of low-density parity-check codes based on ordered statistics," *IET Commun.*, vol. 7, no. 14, pp. 1517–1521, Jul. 2013.
- [14] R. Haga and S. Usami, "Multi-bit flip type gradient descent bit flipping decoding using no thresholds," in *Proc. Int. Symp. Inf. Theory Appl.*, Honolulu, HI, USA, 2012, pp. 6–10.
- [15] M. Ismail *et al.*, "Low latency low power bit flipping algorithms for LDPC decoding," in *Proc. IEEE 21st Int. Symp. Pers., Indoor, Mobile Radio Commun.*, Istanbul, Turkey, 2010, pp. 278–282.
- [16] O. A. Rasheed *et al.*, "Fault-tolerant probabilistic gradient-descent bit flipping decoder," *IEEE Commun. Lett.*, vol. 18, no. 9, pp. 1487–1490, Sep. 2014.
- [17] G. Sundararajan, C. Winstead, and E. Boutillon, "Noisy gradient descent bit-flip decoding for LDPC codes," *IEEE Trans. Commun.*, vol. 62, no. 10, pp. 3385–3400, Oct. 2014.
- [18] T. Tithi, C. Winstead, and G. Sundararajan, "Decoding LDPC codes via noisy gradient descent bit-flipping with re-decoding," 2015, to be published. [Online]. Available: <http://arxiv.org/abs/1503.08913>
- [19] A. Nouh and A. H. Banihashemi, "Reliability-based schedule for bit-flipping decoding of low-density parity-check codes," *IEEE Trans. Commun.*, vol. 52, no. 12, pp. 2038–2040, Oct. 2004.
- [20] D. V. Nguyen and B. Vasić, "Two-bit bit flipping algorithms for LDPC codes and collective error correction" *IEEE Trans. Commun.*, vol. 62, no. 4, pp. 1153–1163, Apr. 2014.
- [21] D. J. C. MacKay, Encyclopedia of Sparse Graph Codes. [Online]. Available: <http://www.inference.phy.cam.ac.uk/mackay/codes/data.html>
- [22] S. Haykin, "Recursive least-square adaptive filters," in *Adaptive Filter Theory*, 4th ed. Upper Saddle River, NJ, USA: Prentice-Hall, 2002, ch. 9, pp. 436–465.
- [23] J. Chen and M. Fossorier, "Near optimum universal belief propagation based decoding of low-density parity check codes," *IEEE Trans. Commun.*, vol. 50, no. 3, pp. 406–414, Mar. 2002.



**Tofar C.-Y. Chang** (S'09) received the B.S. degree in communications engineering from the National Chiao Tung University, Hsinchu, Taiwan, in 2007. Currently, he is pursuing the Ph.D. degree at the National Chiao Tung University. His current research interests include channel coding and wireless communications.



**Yu T. Su** (S'81–M'83–SM'06) received the B.S. degree in electrical engineering from Tatung Institute of Technology, Taipei, Taiwan, in 1974, and the M.S. and Ph.D. degrees from the University of Southern California, Los Angeles, USA, in 1983. From 1983 to 1989, he was with LinCom Corporation, Los Angeles, CA, USA, where his last position was a Corporate Scientist and was involved in various satellite communication systems design and verification. Since late 1989, he has been with the National Chiao Tung University (NCTU), Hsinchu, Taiwan, where he is currently a Professor in the Department of Electrical and Computer Engineering. He was an Associative Dean of the College of Electrical and Computer Engineering from 2004 to 2007, and Head of the Communications Engineering Department from 2001 to 2003. He is also affiliated with the Microelectronic and Information Systems Research Center of NCTU and served as a Deputy Director during 1997–2000. From 2005 to 2008, he was the Area Coordinator of the National Science Council's Telecommunications Programme. His main research interests include communication theory and statistical signal processing.