# You could have invented monads

Roeland Matthijssens

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# **Functional Programming**

#### Why?

See Jeffrey's chalk talk

#### Concepts

- No state change
- No side effects
- Function composition is crucial
- What about IO?





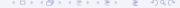
# Demystification

- A monad is not a specific feature.
- They are simple data structures that happen to implement the monadic interface
- Being a monad is simply adhering to a certain(monadic) interface (bind, unit)

# Monads

#### Disclaimer

No guarantees.



# Slightly contrived example

#### Something you may want to do

- Add debugging statements to function calls
- Remember: no destructive changes (adding to a global log list)

# Intuitive implementation

```
f x \rightarrow (x + 1, 'f was called')
g x \rightarrow (x + 1, 'g was called')
```

f and g both take a float and return a tuple (float, string)

## How to use this in composition $f \cdot g$

```
y, logString_g = g(x)
z, logString_f = f(y)
log = logString_g + logString_f
return (z, log)
```

It can be done, but it is ugly and long winded

# if f and g simply took a *float* and returned a *float* we could write $f \cdot g$

```
f x \rightarrow x+1

g x \rightarrow x+2

f \cdot g 3 \rightarrow 6
```

# But they take a *float* and output a (*float*, *string*). Simply composing them would not work

```
f x -> (x+1, "f called.")
g x -> (x+2, "g called.")
f . q 3 ->
```

Error: f expects a float but received a tuple

### Lets redefine f and g

```
f' x log -> (x+1, log + "f called.")
g' x log -> (x+2, log + "g called.")
f' . g' 3 "" -> (6, "g called.f called.")
```

So now this works. Jeej Unfortunatly we don't have an f' but an f which is inconvenient

# Small Tangent

## Simple type definition

f is a function that takes a float and returns a
(float, string)

```
f :: float -> (float, string)
```

#### More complex type definition

f is a function that takes a function g and a float and returns a string the function g that is an argument of f will take a float and

return a (float, string)
f :: q -> float -> string

```
f :: (float->(float, string))->float->string
```

#### AKA currying

# Bind

The main problem is that *f* does not take (*float*, *string*) but just *float*. Lets create *bind*, a higher order function that can help us

### Typedef of bind

- Inputs: bind takes a function that takes a float and returns a (float, string) and a (float, string)
- Output: a new function that will take a (float, string) and returns a (float, string)

```
bind ::
    (float -> (float, string)) ->
    (float, string) ->
    ((float, string) -> (float, string))
```

#### Define bind

```
bind f (float, string) =
  let (new_float, new_string) = f(float)
  in (new_float, string + new_string)
```

## Cool, but what good does it do?

- Given a pair of debuggable functions f and g of type  $float \rightarrow (float, string)$
- We can easily compose them using (bind f)  $\cdot g$
- Or simply bind f · q
- Lets call this composition function >>=
- f >>= q >>= h 3

- We now have half of our formal definition.
- Bear with me a bit longer
- Lets ask ourselves if we can create an identity (aka unit) function for this bind method we just defined
- So that (unit >>= f) = (f >>= unit) = f
- Of course there is
- We can just define unit as *unit*  $x \rightarrow (x, "")$

- The fun thing about this *unit* function is that we can use it to *lift* a normal function to a debuggable function
- $lift f = unit \cdot f$
- This means that I can now compose a normal function with a debuggable functions.
- debugable >>= lift normalfunction
- Of cource the debug output of our lifted normal function is just the empty string (thats how we defined unit

## Summary of example

- We started with simple functions
- We wanted to make them debuggable by having them output an additional string
- We need to be able to compose small functions into larger ones
- We did not want to write a lot of boiler plate
- We accomplished this by implementing bind, unit, lift and >>=

#### Second small example

Before we give the formal definition lets do another quick example

# Multiple return values

```
double :: Int -> [Int]
double x -> [x, x]
primeFactors :: Int -> [Int]
primeFactors 30 -> [2, 3, 5]

double . primeFactors = ???
```

#### How to do double · primeFactors

- First we get the primeFactors
- For each result in the returned list we apply double
- This returns a new list where each element is a list of length 2
- To form a single list of all results we simply flatten the list

## Sounds familiar

```
bind: f x = concat(map f x)
unit: x = [x]
```

- We saw debuggable functions
   Debuggable: a → (a, string)
- We saw multivalued functions  $Multivalued : a \rightarrow [a]$
- Make Debuggable or Multivalued abstract and call it m
- In both cases you get a function  $a \rightarrow m a$
- For that specific m we define bind and unit (the others are in function of these two)

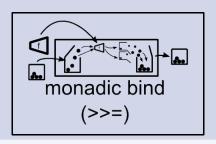
# Formal definition

A monad is a triplet (*m*, *bind*, *unit*). *bind* and *unit* must satisfy 3 mathematical laws

- left-unity
- right-unity
- associativity

# Informal definition

- A monad is a box
- You define how you put a new item in this box (unit)
- You define how you can open a box to apply a function to the element inside of it (bind)
- If those two methods follow some basic rules (commutativity, ...) you can chain them nicely





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#### What are they used for

The monad ('box') abstracts away a lot of computations that would be boiler plate. In our example passing the logging messages to each function is refactored in the monad's bind method

#### Other useful monads

- Maybe Monad: Abstracts away null checks
- State Monad: Abstracts away creating new changed states and passing it arround your functions (e.g. SeededRandom)
- Continuation Monad: Abstracts away call stack manipulations
- IO monad: Abstract away screen state manipulation
- ...

## Caveat of box analogy

All 'boxes' can be seen as monads. But not all monads can be seen as boxes.

#### Questions?

- Need additional examples?
- Something not clear?
- Want a different explanation?
- You know where to find me !!

And when you understand monads fully, you are ready to learn about Arrows :)