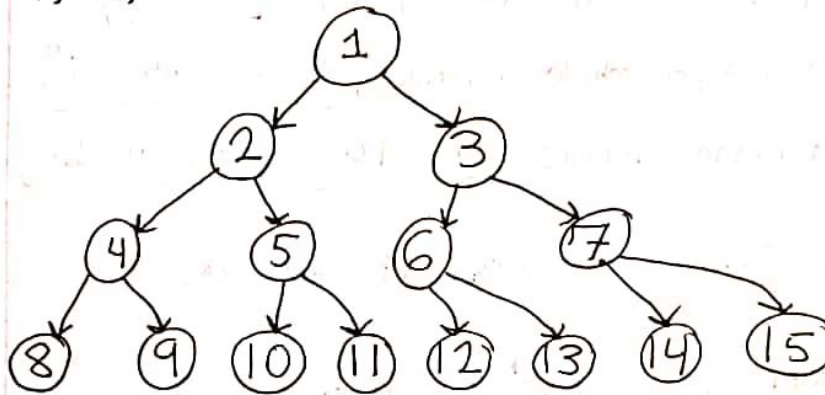


# Project #1 Solutions

1) a)



b) Breadth-first search:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 10 \rightarrow 11$

Depth-limited Search (3):  $1 \rightarrow 2 \rightarrow 4 \rightarrow 8 \rightarrow 9 \rightarrow 5 \rightarrow 10 \rightarrow 11$

Iter Deep:  $1 \rightarrow 2 \rightarrow 3$ ;  $1 \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 3 \rightarrow 6 \rightarrow 7$ ;

$1 \rightarrow 2 \rightarrow 4 \rightarrow 8 \rightarrow 9 \rightarrow 5 \rightarrow 10 \rightarrow 11$

c) Bidirectional search works well. Forward  $b=2$ , backward  $b=1$

d) Yes. Backwards Search from goal use floor operator as transition.

e) Binary Representation of Goal state, ignore Most Significant Bit  
0 = Left, 1 = Right

ex:  $8 \rightarrow 1000 \rightarrow 000 \rightarrow \text{Left, Left, Left}$

$11 \rightarrow 1011 \rightarrow 011 \rightarrow \text{Left, Right, Right}$

$14 \rightarrow 1110 \rightarrow 110 \rightarrow \text{Right, Right, Left}$

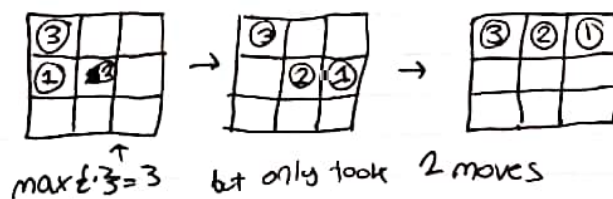
2) a)  $n$  vehicles,  $n^2$  spaces,  $(n^2)^n = n^{2n}$  states

b) 5 actions per vehicle,  $n$  vehicles, branching factor =  $5^n$

c) Manhattan distance - 
$$h_i = |(n-i+1) - x_i| + |n - y_i|$$

d) i) & ii)  $\sum_{i=1}^n h_i \geq \max\{h_1, \dots, h_n\}$  (if ii) isn't admissible, i) isn't either)

Counter Example



iii) • Define work,  $W$ , as the total distance moved by all vehicles, for a given solution.

• At each step the most work that can be done is  $n$

• So, the min # of steps to do work  $W$  is  $\frac{W}{n}$

• Note that  $W \geq \sum_{i=1}^n h_i \geq n \min\{h_1, \dots, h_n\}$

• So,  $(\# \text{ of steps}) \geq \frac{W}{n} \geq \frac{n \min\{h_1, \dots, h_n\}}{n} = \min\{h_1, \dots, h_n\}$

$\min\{h_1, \dots, h_n\} \leq (\# \text{ of steps})$

So, iii) is the only admissible heuristic

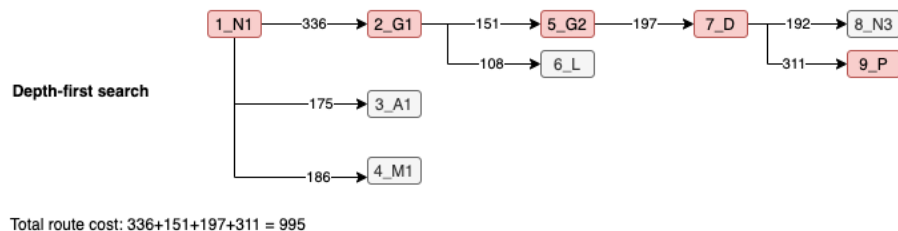
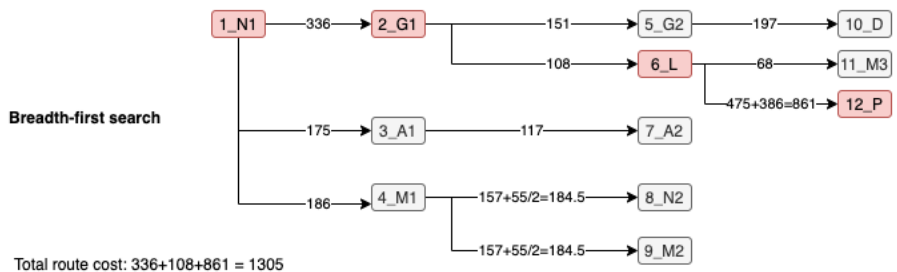
**Question 3 (Student submission with one possible correct solution; formatted nicely... It is also OK to assume cities \*are\* added in the search tree again)**

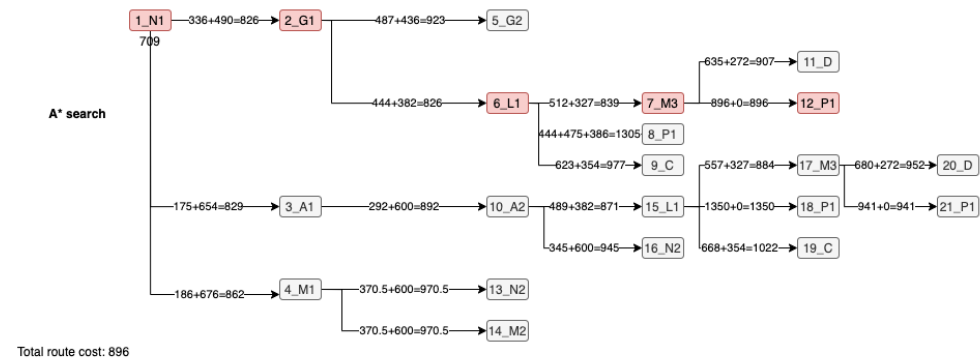
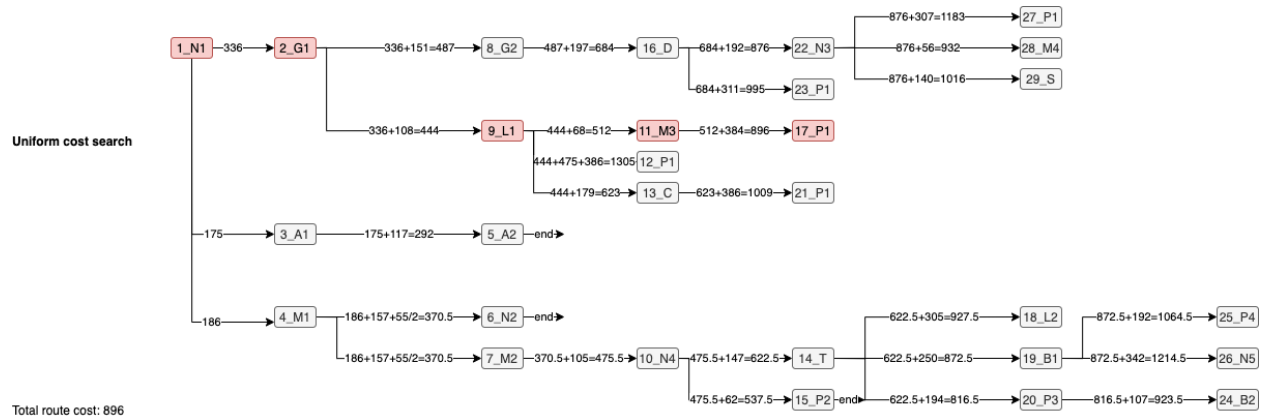
Note: for breadth-first search, depth-first search, and uniform-cost search, assume nodes (cities) that are already in the search tree won't be added again (as suggested by piazza post). Breadth-first and depth-first search are not necessarily optimal, because the search stops at the first solution. Uniform-cost search is optimal in theory, but in this case, since the nodes that are already in the search tree won't be added again, it might not necessarily reach to optimal solution due to the order when nodes are visited (luckily it is still optimal in this example). A\* is optimal.

Heuristic used in this example:

City to Paris	
Nice	709
Toulouse	600
Grenoble	490
Aix	654
Marseille	676
Genève	436
Lyon	382
Macon	327
Clermont-Ferrand	354
Avignon	600
Dijon	272
Nimes	600
Montpellier	600

N1 = Nice, N2 = Nimes, N3 = Nancy, N4 = Narbonne,  
G1 = Grenobl, G2 = Geneve, M1 = Marseille, M2 = Montpellier, M3 = Macon, A1 = Aix, A2 = Avignon





A1 = Aix, A2 = Avignon,  
 B1 = Boreaus, B2 = Bayunne  
 G1 = Grenobl, G2 = Geneve,  
 L1 = Lyon, L2 = Limoges  
 M1 = Marseille, M2 = Montpellier, M3 = Macon, M4 = Metz  
 N1 = Nice, N2 = Nimes, N3 = Nancy, N4 = Narbonne, N5 = Nantes  
 P1 = Paris, P2 = Perpignan, P3 = Pau, P4 = Pailiers

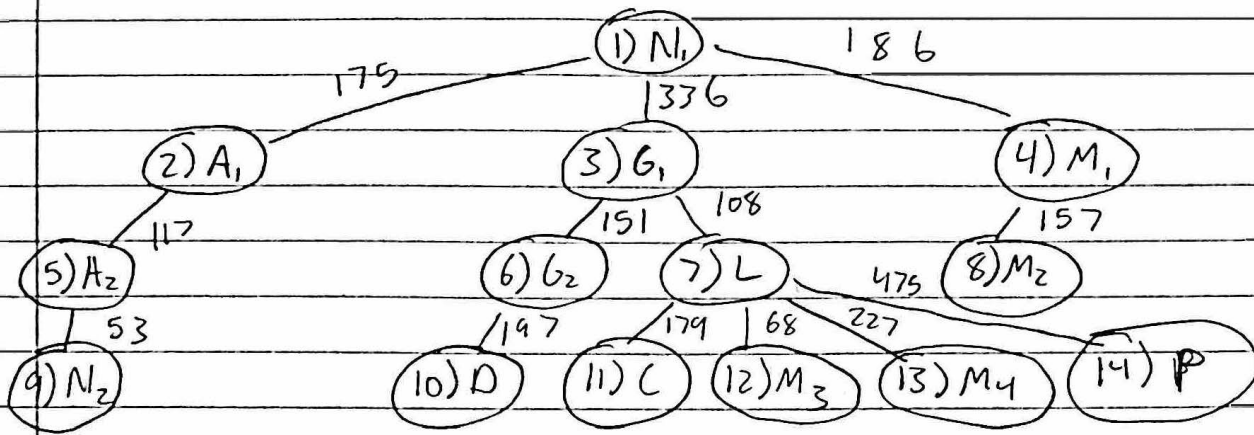
$A_1 = Aix$        $G_1 = Grenoble$        $M_1 = Marseille$        $N_1 = Nice$   
 $A_2 = Avignon$        $G_2 = Geneva$        $M_2 = Montpellier$        $N_2 = Nîmes$   
 $M_3 = Macon$   
 $M_4 = Montlucon$

Q3

Solution:

Version 2

3 a) Breadth first search



Path: Nice  $\rightarrow$  Grenoble  $\rightarrow$  Lyon  $\rightarrow$  Paris

cost =  $336 + 108 + 475 = 919\text{km}$

Solution is not optimal

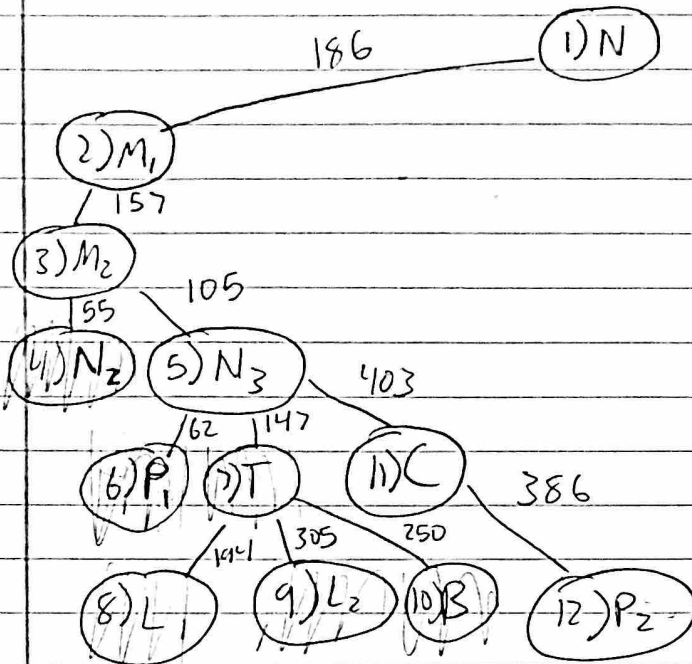
$M_1 = \text{Marseille}$   
 $M_2 = \text{Montpellier}$

$N_1 = \text{Nice}$   
 $N_2 = \text{Nîmes}$   
 $N_3 = \text{Narbonne}$

$P_1 = \text{Perpignan}$   
 $P_2 = \text{Paris}$

$L_1 = \text{Limoges}$   
 $L_2 = \text{Lourdes}$

3 b) Depth first search: Depth limit = 5



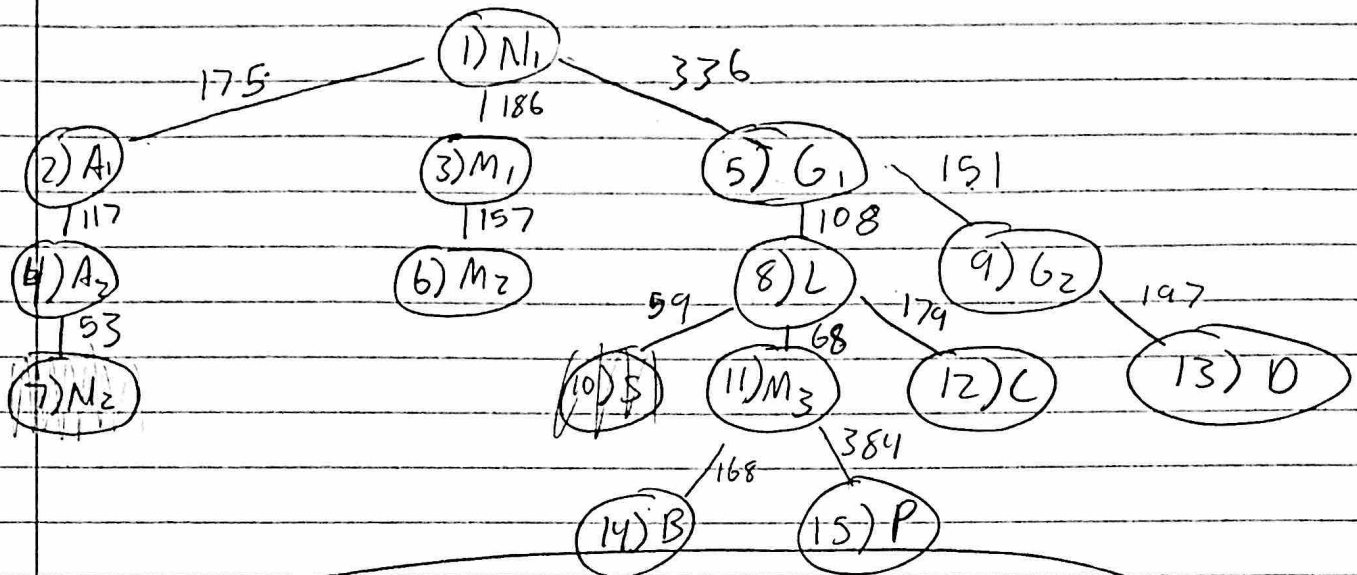
Path: Nice  $\rightarrow$  Marseille  $\rightarrow$  Montpellier  $\rightarrow$  Narbonne  
 $\rightarrow$  Clermont-Ferrand  $\rightarrow$  Paris

Cost:  $186 + 157 + 105 + 403 + 386 = 1237 \text{ km}$

Solution is not optimal

$M_1 = \text{Marseille}$      $A_1 = \text{Aix}$      $N_1 = \text{Nice}$      $G_1 = \text{Grenoble}$   
 $M_2 = \text{Montpellier}$      $A_2 = \text{Avignon}$      $N_2 = \text{Nîmes}$      $G_2 = \text{Geneva}$   
 $M_3 = \text{Macon}$

### 3 c) Uniform cost search

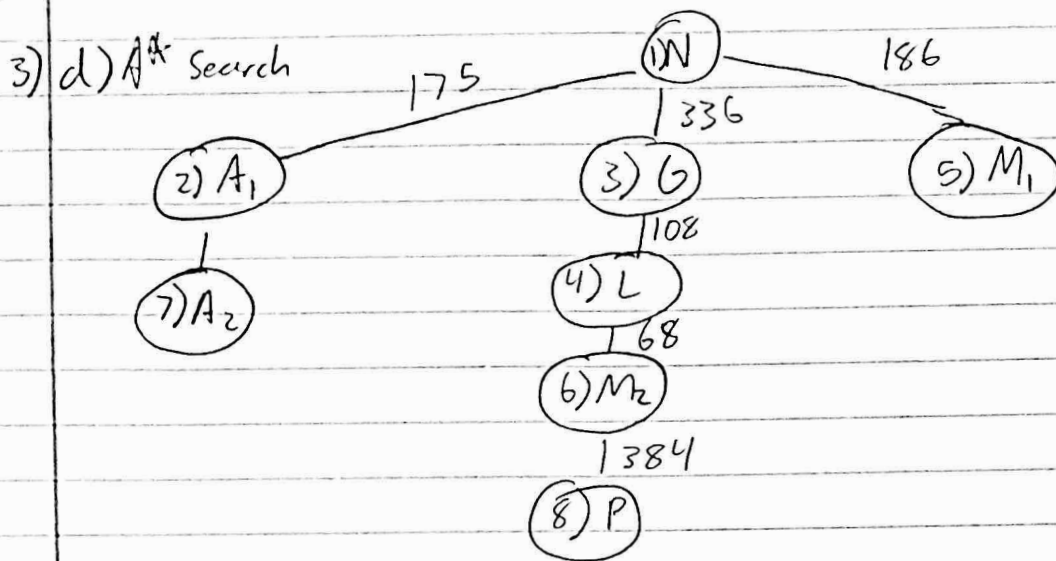


Path: Nice  $\rightarrow$  Grenoble  $\rightarrow$  Lyon  $\rightarrow$  Macon  $\rightarrow$  Paris

$$\text{Cost} = 336 + 108 + 68 + 384 = 896$$

solution is optimal

$A_1 = \text{Aix}$        $M_1 = \text{Marseille}$   
 $A_2 = \text{Avignon}$      $M_2 = \text{Macon}$



Path: Nice  $\rightarrow$  Grenoble  $\rightarrow$  Lyon  $\rightarrow$  Macon  $\rightarrow$  Paris

$$\text{Cost} = 336 + 108 + 68 + 384 = 896 \text{ km}$$

This is the optimal path

Straight line  
 Distance to Paris  


---

 Grenoble : 481 km  
 Aix : 638 km  
 Marseille : 660 km  
 Avignon : 577 km  
 Lyon : 391 km  
 Geneva : 897 km  
 Macon : 339 km  
 St-Etienne : 409 km