

$$L_1 = \{a^n b^n : n \geq 0\} \Rightarrow \text{not regular}$$

$$L_2 = \{w : w \text{ has an equal \# of } a\text{'s} + b\text{'s}\} \Rightarrow \text{not regular}$$

$$L_3 = \{w : w \text{ has equal occurrences of } ab \text{ and } ba \text{ as substrings}\} \Rightarrow \text{regular}$$

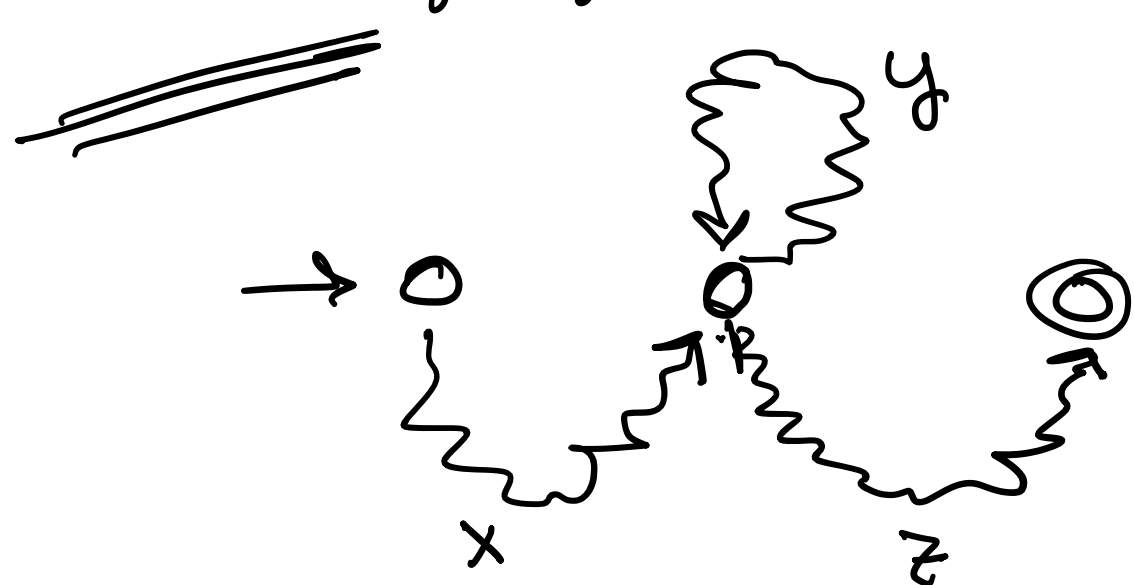
Q: prove a  $L$  is regular  $\Rightarrow$

1. create a dfa
2. create a nfa
3. create a r.e.
4. create a regular grammar.
- [5.] finite language.

Q: prove a  $L$  is not regular  $\Rightarrow$

1. using the pumping lemma for regular lang.
- create a proof by contradiction.

In general: pumping lemma for regular languages is a set of characteristics of infinite regular languages that are guaranteed.

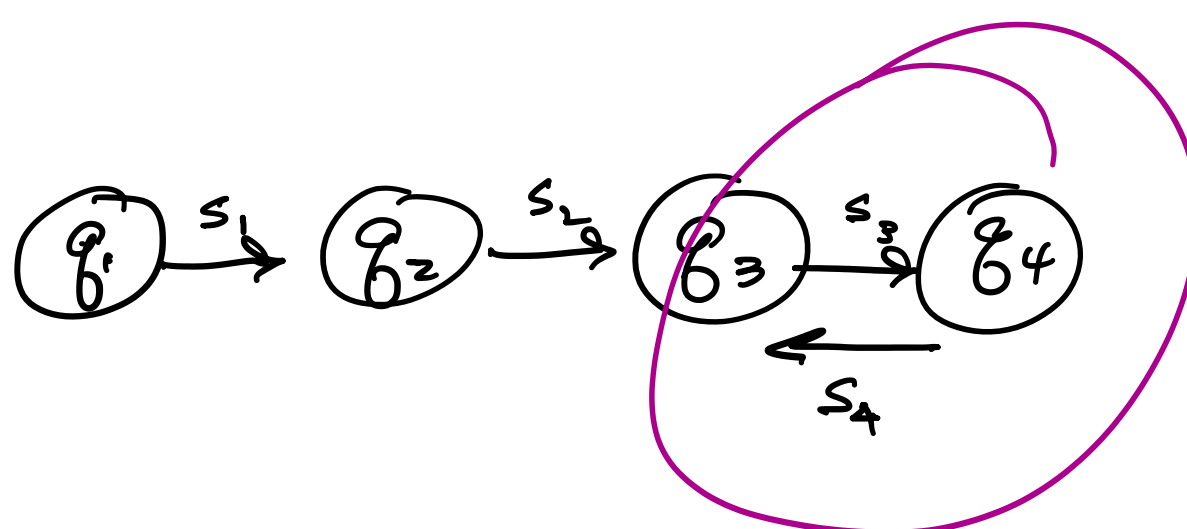


$$\begin{array}{l} \boxed{s_1 = xyz \in L} \\ s_2 = xy^2z \in L \\ s_3 = xy^3z = xy^3z \in L \\ s_0 = xz = xy^0z \in L \end{array} \quad \begin{array}{l} \forall i \geq 2, s_i \in L \\ i = 0, s_0 \in L \end{array} \quad \left. \begin{array}{l} \text{pumping lemma} \\ \uparrow \\ \text{pumping down} \end{array} \right\} \forall i \geq 0, s_i \in L$$

$p$  = number of states  
 $|s| \geq p$

$$p = 4$$

$$|s| \geq 4$$



pumping lemma for regular languages:

$\forall s \in L$ , where  $L$  is infinite regular language,

$|s| \geq p$ , where  $p$  is positive integer,

$s = xyz \Rightarrow |xy| \leq p, |y| \geq 1 \Rightarrow$

$$s_i = xy^i z \in L \quad \forall i \geq 0.$$

proof using pumping lemma: [proof by contradiction]

1. Assume  $L$  is regular  $\Rightarrow$  pumping lemma for r.e. holds i.e.,  $\forall s \in L, |s| \geq p$  (pos. int.)  $\Rightarrow$

$$s = xyz \Rightarrow |xy| \leq p, |y| \geq 1 \Rightarrow \underline{s_i = xy^i z \in L \quad \forall i \geq 0.}$$

2. pick  $s \in L, |s| \geq p$ .

$$\text{ex. } L = \{a^n b^n : n \geq 0\}$$

$$\text{pick } s \in L, |s| \geq p$$

$$\text{let } s = a^p b^p$$

$$[|s| = p + p = 2p \geq p, s \in L]$$

~~$$\text{bad ex. of } L: s = a^{p-1} b^{p-1}$$~~

3. show  $xy$  equal to.

$$\text{show: } |xy| \leq p, |y| \geq 1$$

$$\text{ex. } y = a^k \quad 1 \leq k \leq p$$

$$\begin{array}{l} x = \lambda \\ y = a^p \\ z = b^p \end{array}$$

$$\begin{array}{l} x = a^{p-2} \\ y = a \\ z = ab^p \end{array}$$

$$\begin{array}{l} x = \lambda \\ y = a \\ z = a^{p-1} b^p \end{array}$$

~~$$\begin{array}{l} x = a^{p-1} \\ y = ab \\ z = b^{p-1} \end{array}$$~~

$$\Rightarrow y = a^k \quad 1 \leq k \leq p$$

4. pick  $i$  and pump string

$$\text{ex. let } i = 2$$

$$s_2 = xy^2z \quad \text{if } y = a^k$$

$$s = \frac{a^k}{y} \frac{a^{p-k} b^p}{xz [!y]}$$

$$s_i = xy^i z$$

$$s_i = a^{ki} a^{p-k} b^p$$

$$s_2 = \frac{a^{2k}}{y} \frac{a^{p-k} b^p}{!y} = a^{p+k} b^p$$

$$\text{but } k \geq 1$$

$$\Rightarrow n_a(s_2) > n_b(s_2)$$

$$\therefore s_2 \notin L$$

$$\therefore L \text{ is not regular}$$