#### CPE201 Digital Design

By Benjamin Haas

Class 25: Counter Design

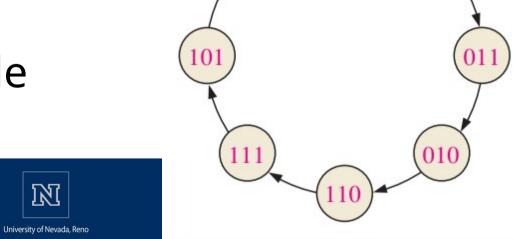


#### Outline

- Designing Counters
  - Examples
- System Design Example



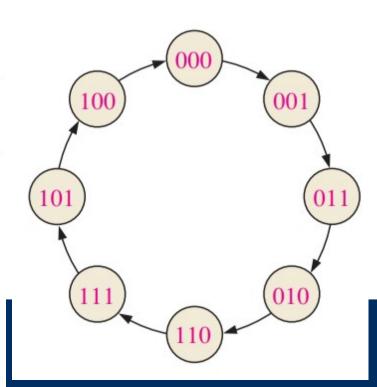
- Step 1 Create state machine
  - For counters, the state is the output
- Ex: 3-bit gray code counter



#### Step 2 – Next state table

Next-state table for 3-bit Gray code counter.

<b>Present State</b>				<b>Next State</b>	
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
0	0	0	0	0	1
0	0	1	0	1	1
0	1	1	0	1	0
0	1	0	1	1	0
1	1	0	1	1	1
1	1	1	1	0	1
1	0	1	1	0	0
1	0	0	0	0	0



# Step 3 – Map state transitions to circuit inputs

Next-state table for 3-bit Gray code counter.

	Present St	ate		Next State	
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
0	0	0	0	0	1
0	0	1	0	1	1
0	1	1	0	1	0
0	1	0	1	1	0
1	1	0	1	1	1
1	1	1	1	0	1
1	0	1	1	0	0
1	0	0	0	0	0

his case

Transition table for a J-K flip-flop.

	Output Tran	<b>Flip-Flop Inputs</b>			
$Q_N$		$Q_{N+1}$	J	K	
0	$\longrightarrow$	0	0	X	
0	$\longrightarrow$	1	1	X	
1	$\longrightarrow$	0	X	1	
1	$\longrightarrow$	1	X	0	

 $Q_N$ : present state

 $Q_{N+1}$ : next state X: "don't care"

Step 3b – Table of inputs to create next

Jo	K <sub>o</sub>	Q <sub>2</sub>	$\mathbf{Q}_1$	$Q_0$	Q <sub>o</sub> next
1	Χ	0	0	0	1
Χ	0	0	0	1	1
Χ	1	0	1	1	0
0	Χ	0	1	0	0
1	Χ	1	1	0	1
X	0	1	1	1	1
V	1	1	0	1	0

Transition table for a J-K flip-flop.

	Output Tran	Flip-Flop Inputs			
$Q_N$		$Q_{N+1}$	J	K	
0	$\longrightarrow$	0	0	X	
0	$\longrightarrow$	1	1	X	
1	$\longrightarrow$	0	X	1	
1	$\longrightarrow$	1	X	0	



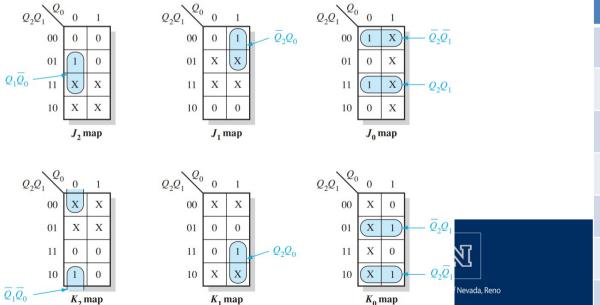
 $Q_N$ : present state

 $Q_{N+1}$ : next state

X: "don't care"

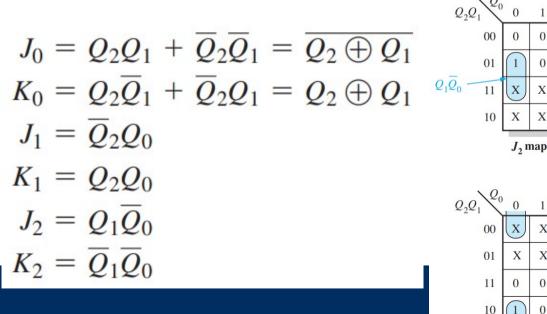
Step 4 – Karnaugh Maps

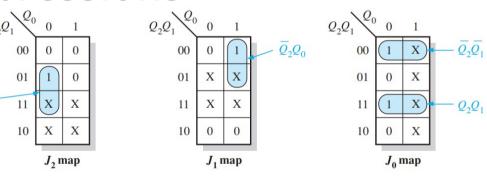
– JK FFs are ACTIVE HIGH, so !



Jo	K <sub>o</sub>	Q <sub>2</sub>	$\mathbf{Q}_1$	$\mathbf{Q}_{0}$	Q <sub>0</sub> next
1	X	0	0	0	1
Χ	0	0	0	1	1
Χ	1	0	1	1	0
0	Χ	0	1	0	0
1	Χ	1	1	0	1
Χ	0	1	1	1	1
Χ	1	1	0	1	0

Step 5 – Logical Expressions





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 $K_1$  map

K, map

 $\overline{Q}_2Q_1$ 

 $K_0$  map

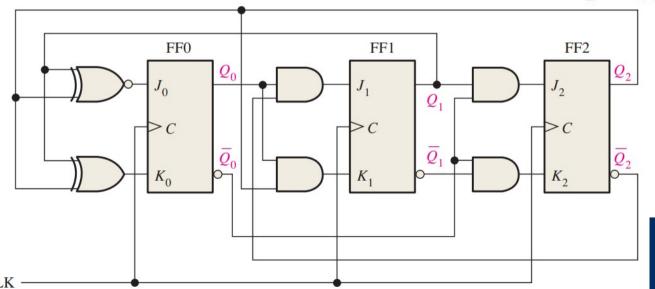
Design 
$$J_0 = Q_2Q_1 + \overline{Q}_2\overline{Q}_1 = \overline{Q}_2 \oplus \overline{Q}_1$$
  
 $K_0 = Q_2\overline{Q}_1 + \overline{Q}_2Q_1 = Q_2 \oplus Q_1$   
 $J_1 = \overline{Q}_2Q_0$ 

Step 6 - Implementation

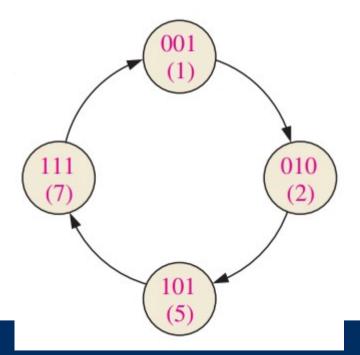
$$K_1 = Q_2 Q_0$$

$$J_2 = Q_1 \overline{Q}_0$$

$$K_2 = \overline{Q}_1 \overline{Q}_0$$



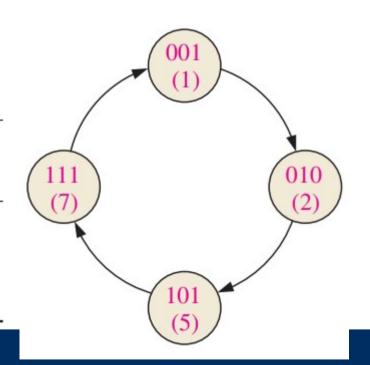
- 3-bit irregular counter
  - State machine given (step
  - Implement with D FFs



#### Step 2 – next state table

Next-state table.

<b>Present State</b>				Next State	e
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
0	0	1	0	1	0
0	1	0	1	0	1
1	0	1	1	1	1
1	1	1	0	0	1





Step 3b - Table of inputs to create next

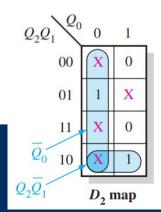
Transition table for a D flip-flop.

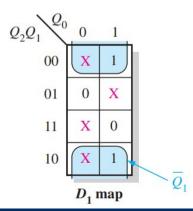
Ou	tput Trans	Flip-Flop Input	
$Q_N$		$Q_{N+1}$	D
0	$\longrightarrow$	0	0
0	$\longrightarrow$	1	1
1	$\longrightarrow$	0	0
1	$\longrightarrow$	1	1

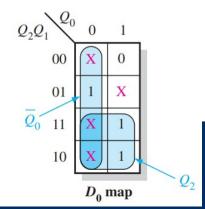
D <sub>0</sub>	Q <sub>2</sub>	Q <sub>1</sub>	$Q_0$	Q₀ next
0	0	0	1	0
1	0	1	0	1
1	1	1	1	1
1	0	0	1	1

- Step 4 Karnaugh Maps
  - Pink X's are Don't Care because the state is not in the state machine
  - D FFs are ACTIVE HIGH, so SOP

D <sub>0</sub>	Q <sub>2</sub>	$\mathbf{Q}_1$	$Q_0$	Q₀ next
0	0	0	1	0
1	0	1	0	1
1	1	1	1	1



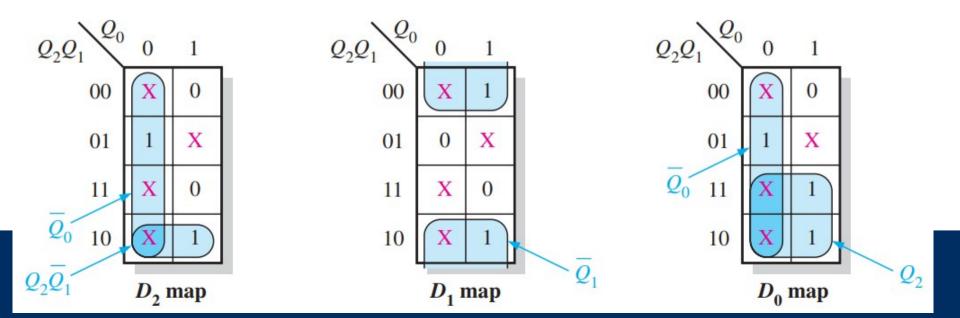




# Example $D_0 = \overline{Q}_0 + Q_2$

$$D_0 = \overline{Q}_0 + Q_2$$
$$D_1 = \overline{Q}_1$$

• Step 5 – Logical Expressions  $D_2 = \overline{Q}_0 + Q_2 \overline{Q}_1$ 

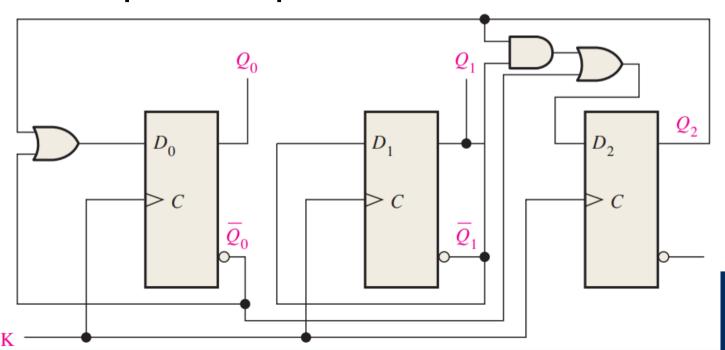


$$D_0 = \overline{Q}_0 + Q_2$$

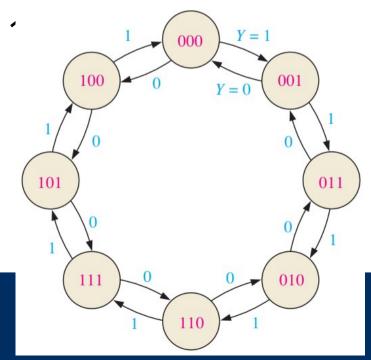
$$D_1 = \overline{Q}_1$$

Step 6 - Implementation

$$D_2 = \overline{\overline{Q}}_0 + Q_2 \overline{Q}_1$$

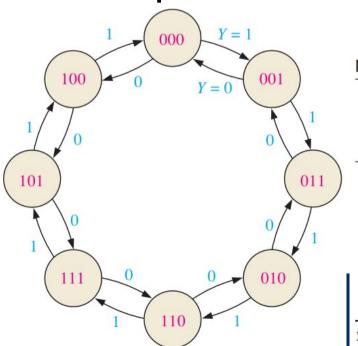


- 3-bit Gray Code up/down counter
  - State machine given (step '
  - Implement with JK FFs





#### Step 2 – Next state table



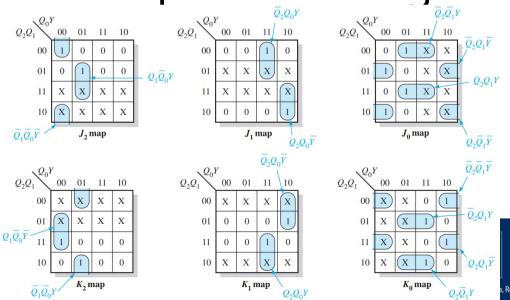
Next-state table for 3-bit up/down Gray code counter.

			Next State						
Pı	<b>Present State</b>			Y = 0  (DOWN)			Y = 1  (UP)		
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$	
0	0	0	1	0	0	0	0	1	
0	0	1	0	0	0	0	1	1	
0	1	1	0	0	1	0	1	0	
0	1	0	0	1	1	1	1	0	
1	1	0	0	1	0	1	1	1	
1	1	1	1	1	0	1	0	1	
1	0	1	1	1	1	1	0	0	
1	0	0	1	0	1	0	0	0	

 $Y = UP/\overline{DOWN}$  control input.

					1	TOME DELLE	
	Y	$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
	0	0	0	0	1	0	0
Y = 0  (DOWN)	0	0	0	1	0	0	0
	0	0	1	1	0	0	1
Y = 1 (UP)	0	0	1	0	0	1	1
	0	1	1	0	0	1	0
	0	1	1	1	1	1	0
	0	1	0	1	1	1	1
	0	1	0	0	1	0	1
	1	0	0	0	0	0	1
	1	0	0	1	0	1	1
	1	0	1	1	0	1	0
	1	0	1	0	1	1	0
	1	1	1	0	1	1	1
	1	1	1	1	1	0	1
	1	1	0	1	1	0	0
	1	1	0	0	0	0	0

- Step 3b Same process as before
- Step 4 Karnaugh Maps for 4 inputs



Step 5 and Step 6 are the same process as before

- Lights at a busy road perpendicular to a side street
- Requirements:
  - Green light on main for at least 25s before a change
  - Green light on side street as long as there are vehicles, up to 25s max
  - Yellow for 4s on both streets

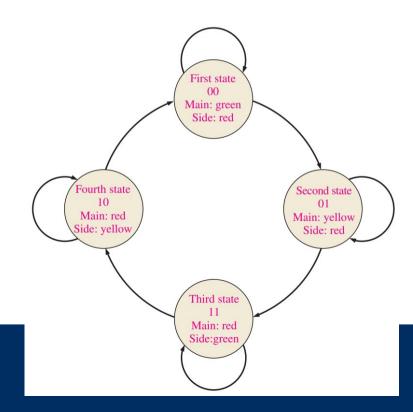


#### Inputs:

- Vehicle sensor input (1 = vehicle present)
- System clock (1Hz)
- Outputs:
  - Red, yellow, green signals for both streets(1 = on)



- Basic state machine
  - You know how traffic lights work
  - Also added gray code to represent each state as a number



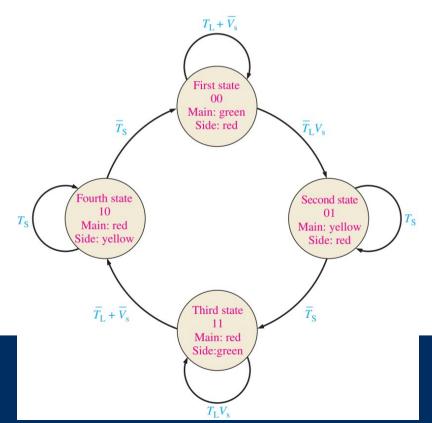


- Let's describe the transitions
  - First state is 'normal', takes a vehicle on side street and 25s to have passed to change to next state
  - Second state, wait for 4s on yellow
  - Third state, stay here as long as a vehicle is present and it is less than 25s in this state
  - Fourth state, same as second state



- Now make some inputs/variables for the state machine
  - Make from state transitions, make vars binary
  - $-V_s$  = vehicle present on side street
  - $-T_L = long timer (25s) is on$
  - $-T_s = \text{short timer } (4s)$  on

 Note that each transition condition is literally the NOT of the looping condition

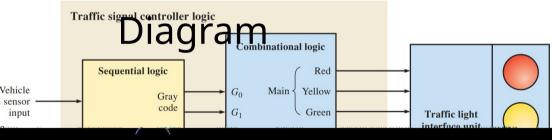




- The controller can be broken into several parts
  - Sequential logic takes in all inputs on state machine and outputs state gray code
  - Combinational Logic decodes state, triggers timers, and creates light outputs
  - Timers input triggers and clock, outputs for if each timer is running (creates T<sub>L</sub> and T<sub>S</sub>)

Break each block down until it does one thing

- Sequential is co
- Timers
- Combination st does too much

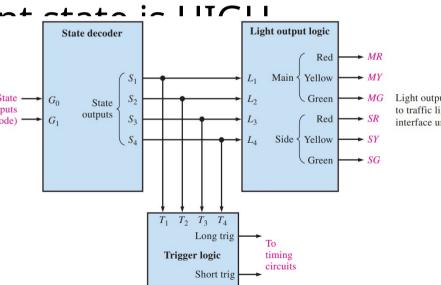


Do your thinking in blocks to simplify problems

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State decoder – cur

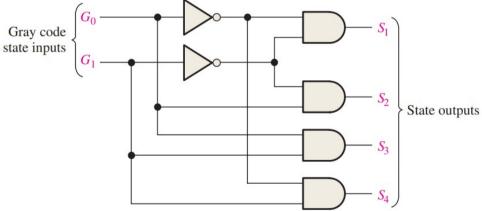
State directly ties
 to light colors and
 starts a timer



State decoder from truth table

$G_1$ outputs $S_1$ $S_3$ $S_4$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
	$S_1$

Truth table for the state decoder.						
State Inputs (Gray Code)		State Outputs				
$G_1$	$G_0$	$S_1$	$S_2$	$S_3$	$S_4$	
0	0	1	0	0	0	
0	1	0	1	0	0	
1	1	0	0	1	0	
1	0	0	0	0	1	

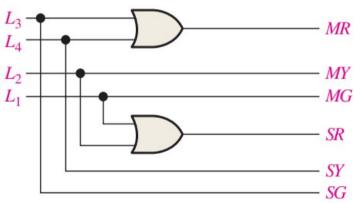


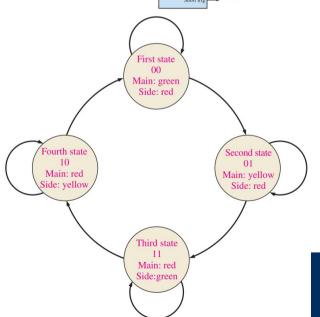
# Traffic Signal Designal Designation Designation

Inputs/Outputs are not usually named the same

• 
$$MR = L3 + L4$$

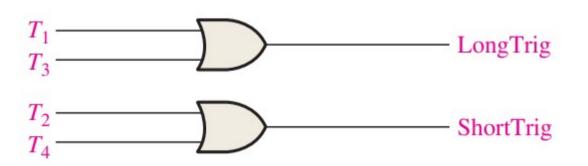
• 
$$SR = L1 + L2$$

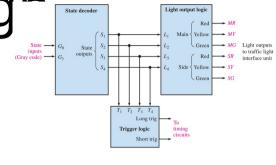


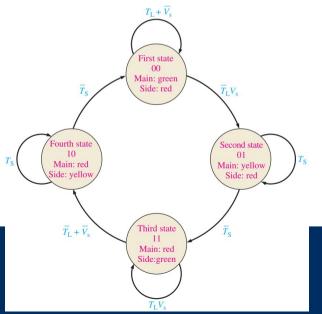


 $(Gray code) \longrightarrow G_1$ 

- S1 and S3 start long timer
- S2 and S4 start short timer

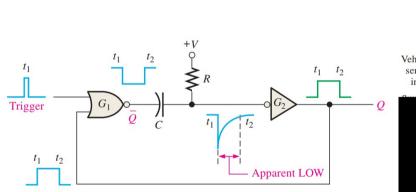


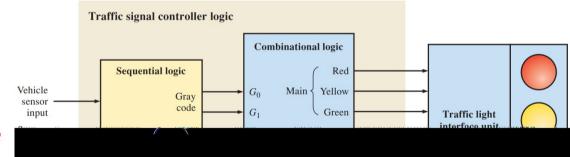






 TS and TL are high when timer is running, just like a one shot

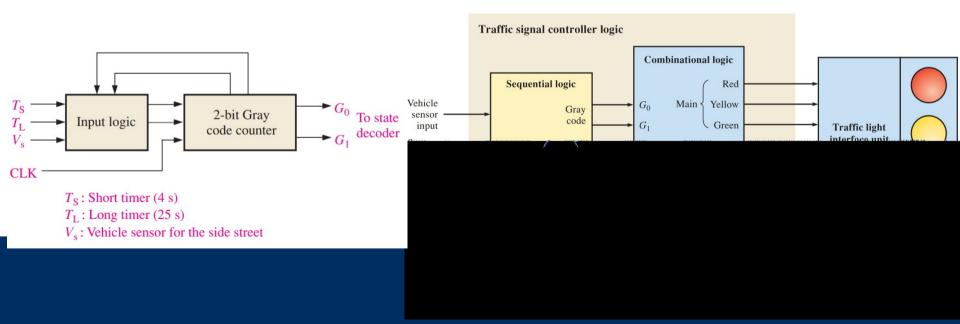




- $t_w = 0.7RC$
- $4s = 0.7R_{4s}C_{4s}$
- If  $C_{4s} = 1000 \mu F$ , then  $R_{4s} = 5,714 \Omega$
- $25s = 0.7R_{25s}C_{25s}$
- If  $C_{25s} = 1000 \mu F$ , then  $R_{25s} = 35,714 \Omega$

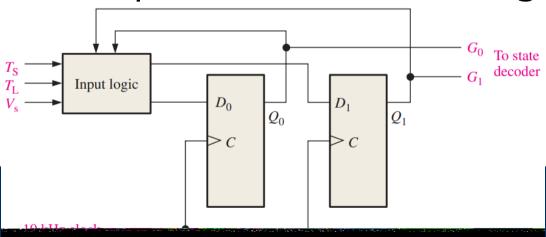


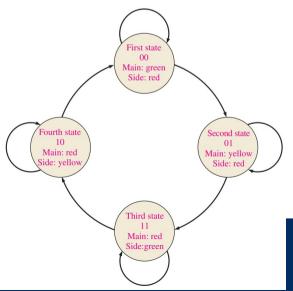
One chunk left, it's a 2-bit counter



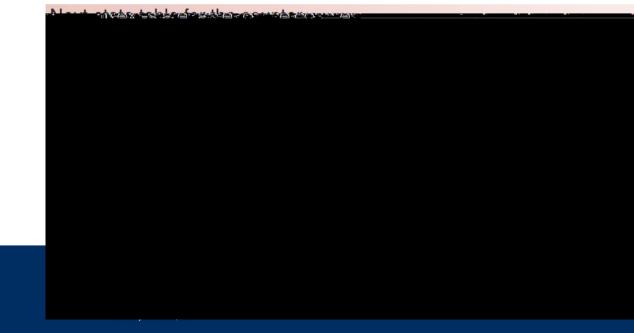
- We'll use D FFs
  - Arbitrarily pick system clock at 10kHz

Step 1 - State machine is giver





Step 2- Next state table

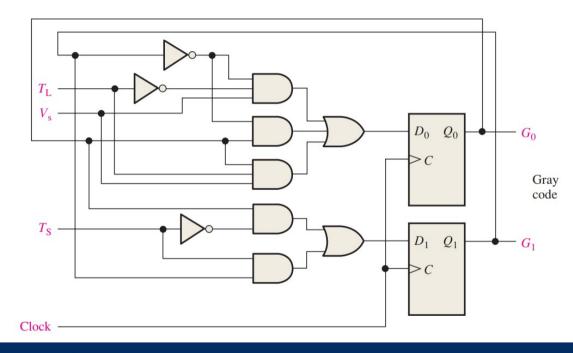


- Karnaugh maps would be very complicated with 5 variables (G<sub>1</sub>, G<sub>0</sub>, V<sub>s</sub>, T<sub>L</sub>, T<sub>s</sub>)
  - Also a sparse table, so you can opt to not simplify the circuit to be quick
  - Write out the terms that make  $D_0 = 1$



- $D_0 = G_1G_0'T_L'V_S + G_1G_0'T_S + G_1G_0T_S' + G_1G_0T_LV_S$

That's all the pieces!



#### Reading

- This lecture
  - 9.5, Ch6 and Ch7 Applied Logic
- Next lecture
  - Sections 12.1-12.3