

proof by contradiction:

1. Assume L is regular
2. \Rightarrow state the pumping lemma
- * 3. pick string $s \in L$, $|s| \geq p$
4. show all possible decompositions of s ($s = xyz$) based $|xy| \leq p$, $|y| \geq 1$
- * 5. pick i - how many times pumping where $i \geq 0$
6. pump up or down based on your i .
7. show pumped string $\notin L$
8. $\therefore L$ must not be regular.

ex show $L = \{a^n b^n : n \geq 0\}$ is not regular.

Assume L is regular $\Rightarrow \forall s \in L$, $|s| \geq p$
 (p = pos. int. called pumping length),
 $s = xyz$, $|xy| \leq p$, $|y| \geq 1$, and

$\rightarrow \underline{s_i = xy^i z \in L \quad \forall i \geq 0.}$

let $s = a^p b^p$

$s \in L$, $|s| = p + p = 2p \geq p$.

$$s = a^{2p} b^{2p}$$

$$s_i = a^{p+2i} b^{p+2i}$$

since $|xy| \leq p$, $|y| \geq 1$

$\Rightarrow y = a^k \quad 1 \leq k \leq p$

$x = a^m \quad 0 \leq m \leq p-1$

$$\Rightarrow s = \underbrace{a^m}_x \underbrace{a^k}_y \underbrace{a^{p-m-k} b^p}_z$$

$$s_i = a^m a^{ki} a^{p-m-k} b^p$$

let $i = 0$

$$s_0 = \underbrace{a^m}_x a^{p-m-k} b^p = \underline{a^{p-k} b^p}$$

\Rightarrow since $k \geq 1 \Rightarrow n_a(s_0) < n_b(s_0)$

$\Rightarrow s_0 \notin L$ because $\#a \neq \#b$

$\therefore L$ must not be regular!

or

$$s_i = a^m a^{ki} a^{p-m-k} b^p$$

let $i = 2$

$$s_2 = a^m a^{2k} a^{p-m-k} b^p = \underline{a^{p+k} b^p}$$

since $k \geq 1 \Rightarrow n_a(s_2) > n_b(s_2)$

$\therefore s_2 \notin L \Rightarrow L$ is not regular.

$$s = \underbrace{a^k}_y \underbrace{a^{p-k} b^p}_z$$

$$s_i = a^{ki} a^{p-k} b^p$$

ex show $L = \{w \in \{a,b\}^* : n_a(w) \neq n_b(w)\}$

let $s = a^p b^{p+1}$

$s \in L$, $|s| = p + p + 1 = 2p + 1 \geq p$

$y = a^k \quad 1 \leq k \leq p$

aside:

$$\begin{array}{c} | \text{a a a a a b b b b b} \\ \hline \xrightarrow{\quad} \xrightarrow{\quad} \xrightarrow{\quad} \\ \begin{array}{ccc} x & y & z \\ | & | & | \\ \leq p & \geq 1 & \end{array} \end{array}$$