

prove $dfa \equiv nfa \Rightarrow dfa \rightarrow nfa \wedge nfa \rightarrow dfa$

aside: $\delta_{nfa}: \varnothing \times (\Sigma \cup \{\lambda\}) \rightarrow \boxed{2^P}$

$\Sigma = \{a, b\}$ $\Sigma \cup \{\lambda\} = \{a, b, \lambda\}$
 $\varnothing = \{\emptyset, \emptyset\}$
 $(q_0, a) = \dots$ $2^P = \{\emptyset, \{\emptyset\}, \{\emptyset, \emptyset\}, \{\emptyset, \emptyset, \emptyset\}\}$
 $(q_0, b) = \dots$ $\uparrow \quad \uparrow \quad \uparrow \quad \uparrow$
 $(q_0, \lambda) = \{\emptyset, \emptyset\}$
 $(q_1, a) = \emptyset$
 $(q_1, b) = \{\emptyset\}$
 $(q_1, \lambda) = \{\emptyset\}$

① $nfa \rightarrow dfa \Rightarrow L(nfa) = L(dfa)$

nfa \rightarrow dfa conversion rules:

- 1. create a graph G_0 w/ vertex $\{q_0\}$ = initial vertex

$\rightarrow \{q_0\}$

2. repeat for all edges:

- a. take any vertex, $\{q_i, \dots, q_k\}$, of G_0 that has outgoing edge for some $a \in \Sigma$.

$\rightarrow \{q_0\} \xrightarrow{a}$

- b. compute $\delta_n^*(q_i, a), \delta_n^*(q_j, a), \dots, \delta_n^*(q_k, a)$

nfa $\rightarrow \{q_0\} \xrightarrow{a} \{q_1\}$

$\delta_n^*(q_0, a) = \{q_0, q_1\}$

- c. form union of all $\delta_n^* \xrightarrow{a} \{q_2, \dots, q_n\}$

- d. create a vertex for G_0 labeled $\{q_2, \dots, q_n\}$ if vertex does not already exist.

- e. add to G_0 an edge from $\{q_i, \dots, q_k\} \xrightarrow{a} \{q_2, \dots, q_n\}$

nfa $\rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \xrightarrow{b} \{q_2\}$

$(q_0, a) = \underline{\{q_0, q_1\}}$

$G_0: \rightarrow \{q_0\} \xrightarrow{a} \{q_0, q_1\} \xrightarrow{b} \{q_1, q_2\}$

$(q_0, b) = \{q_2\}$

(2b) $(\{q_0, q_1\}, b) = \delta_n^*(q_0, b), \delta_n^*(q_1, b)$

(2c)

$\{q_2\} \cup \{q_1\} = \{q_1, q_2\}$

- 3. every state of G_0 w/ label $q_f \in F_N$ is final vertex

nfa: $\rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \dots$

$G_0: \rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \xrightarrow{b} \{q_1, q_2\} \dots$

- 4. if M_n accepts λ , vertex $\{q_0\}$ in G_0 is a final vertex.

nfa $\rightarrow \{q_0\} \xrightarrow{a, \lambda} \{q_1\} \xrightarrow{a} \{q_1\}$
 $G_0: \rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \xrightarrow{b} \{q_1, q_2\} \dots$

\therefore for \forall nfa using nfa \rightarrow dfa conversion rules $\Rightarrow \exists$ dfa, $\ni L(nfa) = L(dfa)$
 $\Rightarrow nfa \rightarrow dfa$

proven ② nfa \rightarrow dfa

proved ① dfa \rightarrow nfa and ② nfa \rightarrow dfa
 $\therefore nfa \equiv dfa$.

remember:

$\Sigma = \{a, b\}$

nfa $\rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \xrightarrow{a} \{q_1\}$

$\rightarrow \{q_0\} \xrightarrow{b} \{\emptyset\} \xrightarrow{a, b} \{\emptyset\}$

$(\emptyset, a) = \emptyset$
 $(\emptyset, b) = \emptyset$

nfa $\rightarrow \{q_0\} \xrightarrow{a} \{q_1\} \xrightarrow{a, b} \{q_2\} \xrightarrow{b} \{q_2\}$

dfa $\rightarrow \{q_0\} \xrightarrow{a} \{q_0, q_1\} \xrightarrow{b} \{q_0, q_1, q_2\}$

$\delta_n^*(q_0, a) = \{q_0, q_1\}$

$\delta_n^*(\{q_0, q_1\}, a)$
 $= \delta_n^*(q_0, a) \cup \delta_n^*(q_1, a)$
 $= \{q_0, q_1\} \cup \{q_2\}$
 $= \{q_0, q_1, q_2\}$

$\delta_n^*(q_0, b) = \emptyset$

$\delta_n^*(\emptyset, a) = \emptyset$

$\delta_n^*(\emptyset, b) = \emptyset$