

5.1 - The Backtracking Technique

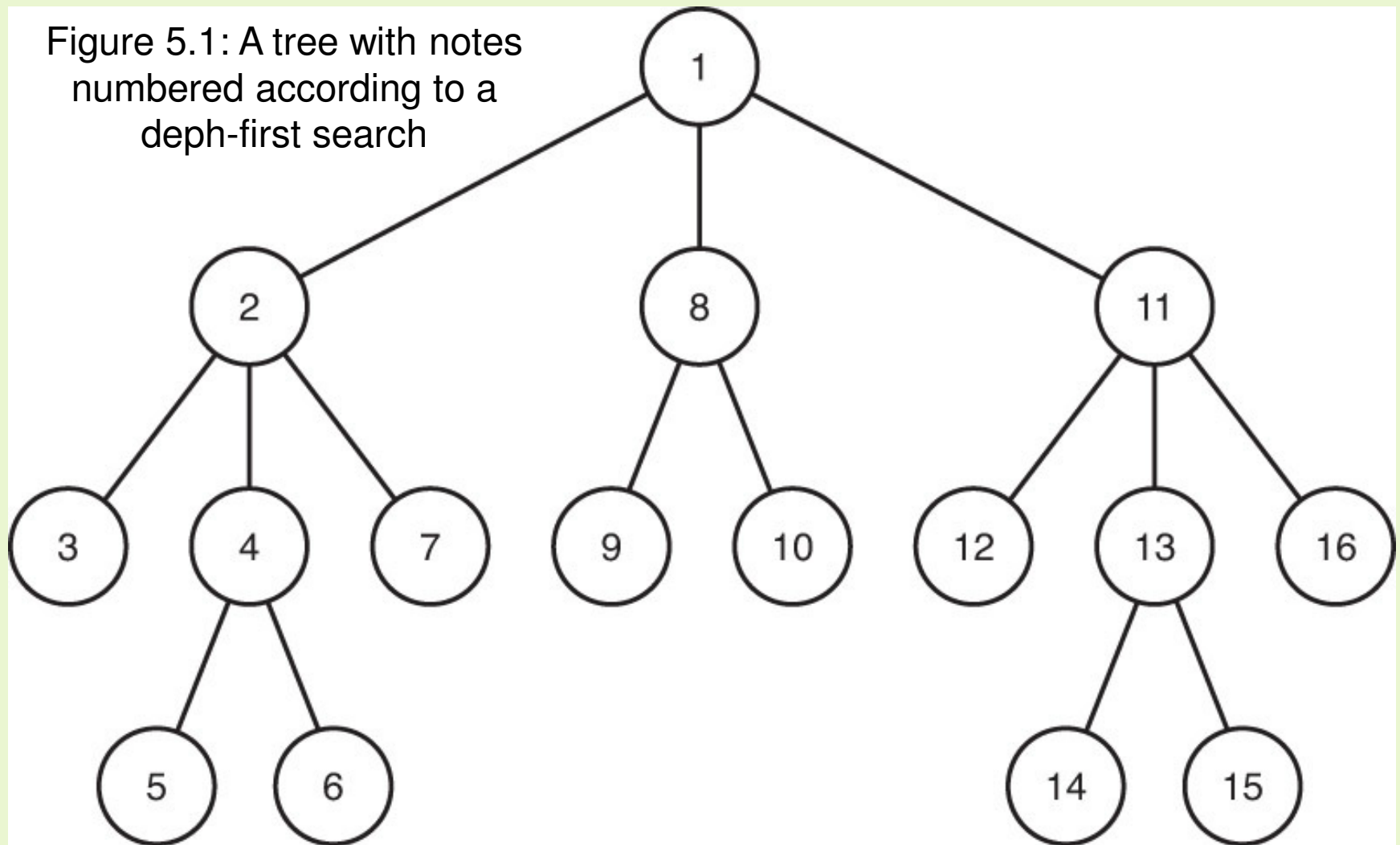
- ▶ Used to solve problems in which a *sequence* of objects is chosen from a set, so that the sequence satisfies a criterion
- ▶ A sequence of choices are made at various decision points while trying to find a solution
- ▶ If a dead end is reached, the algorithm has to **backtrack** to a previous decision point and try a different path
- ▶ In the worst case, performance is exponential, but better results often achieved in practice

Depth-First Search

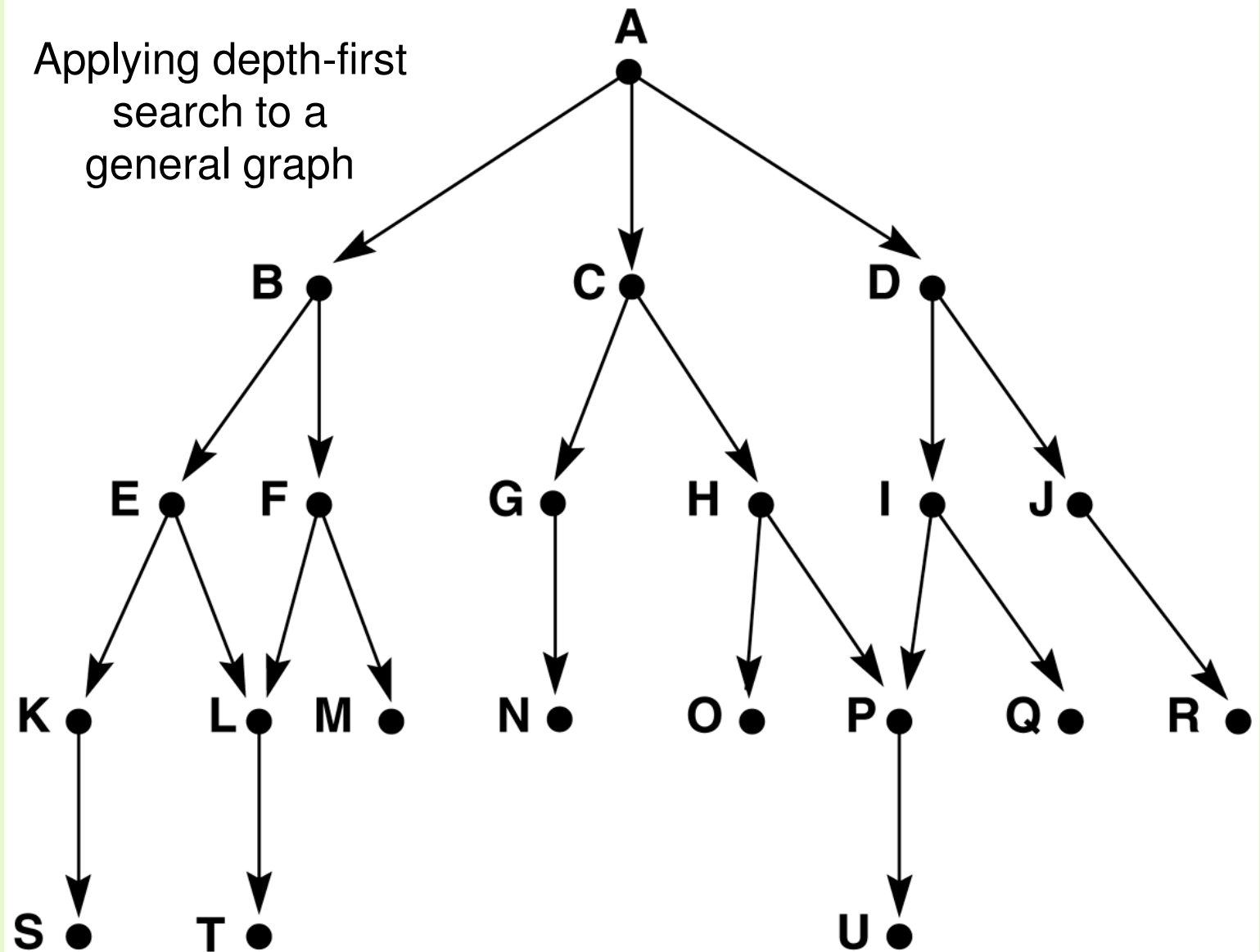
- ▶ A backtracking algorithm is essentially a modified depth-first search
- ▶ DFS is a generalized preorder traversal, in which a node is visited first, followed by a recursive traversal of its children.

```
void depth_first_tree_search (node v)
{
    node u;
    visit v;
    for (each child u of v)
        depth_first_tree_search(u);
}
```

Figure 5.1: A tree with notes numbered according to a depth-first search

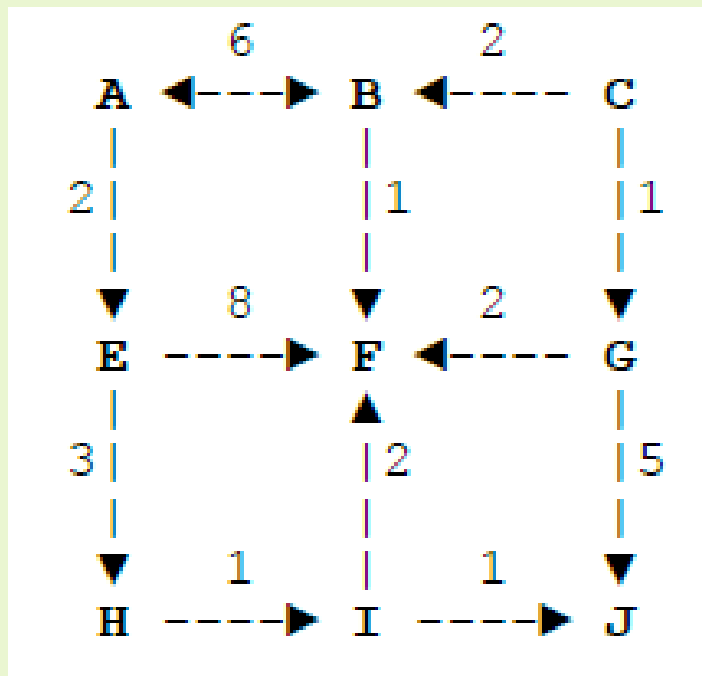


Applying depth-first
search to a
general graph



Applying depth-first search to a general graph

(Assume that adjacent vertices are visited in alphabetical order)



The n-Queens Problem

- The n-queens problem asks how n queens can be placed on an $n \times n$ chess board so that no two queens attack each other
- Instead of looking at all combinations, we notice two queens cannot be in the same row.

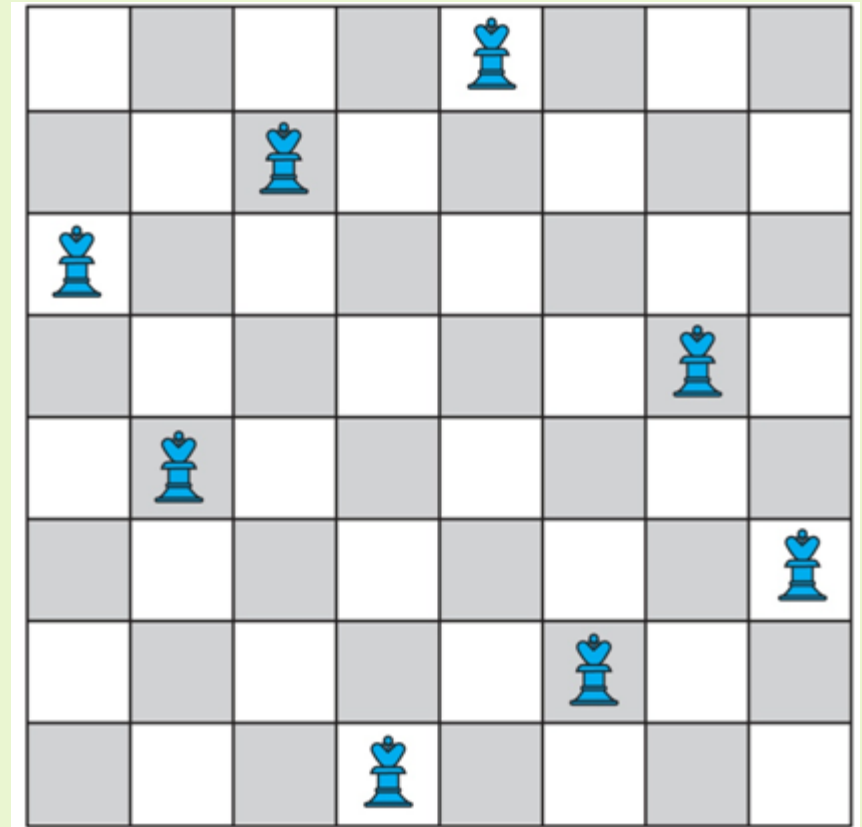
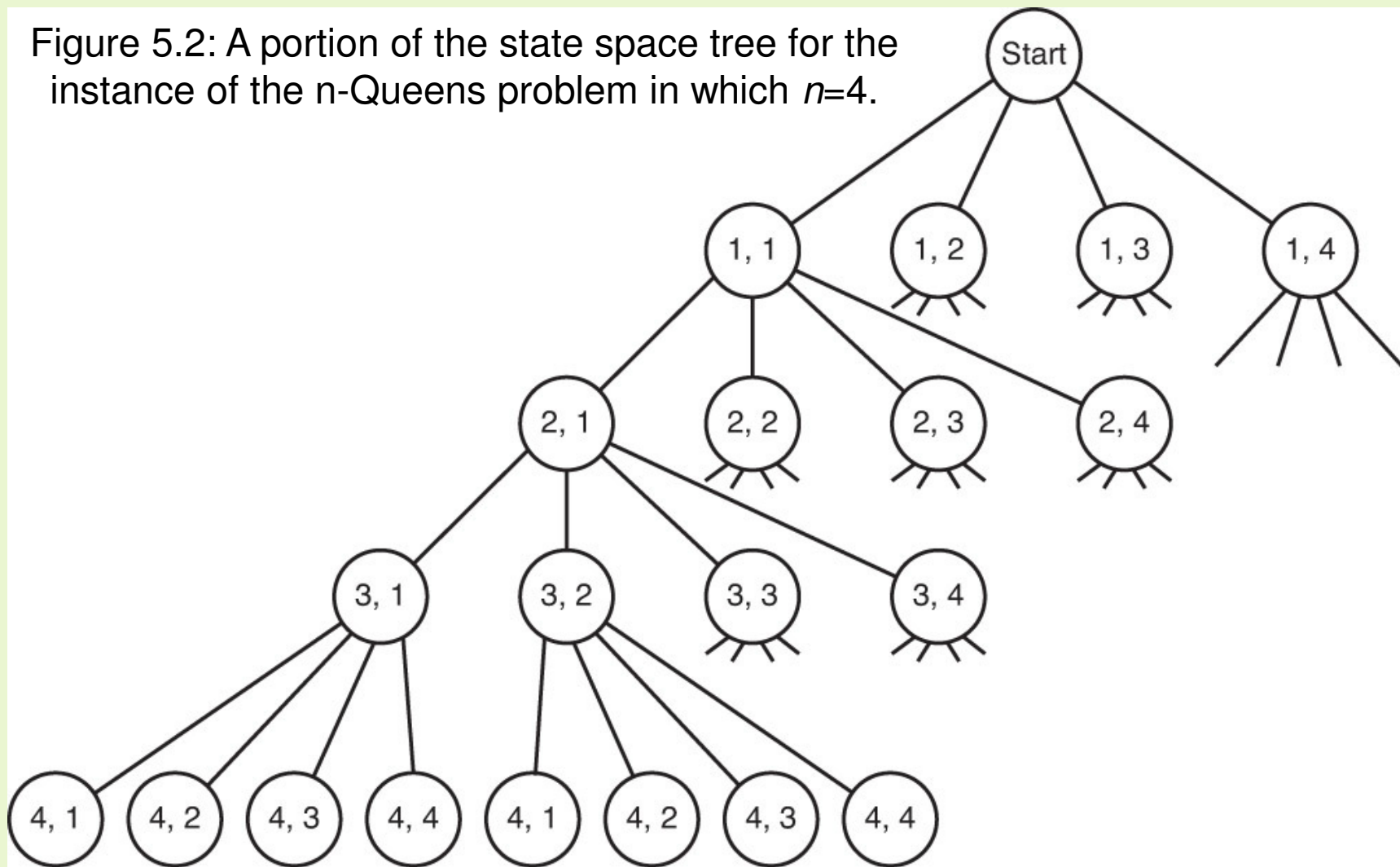


Figure 5.2: A portion of the state space tree for the instance of the n -Queens problem in which $n=4$.



State Space Trees

- ▶ Figure 5.2 is an example of a state space tree, in which each node represents a possible partial solution –or state- to a problem.
- ▶ A simple depth-first search would generate and check every path from the root to a leaf until a solution is found ($4^4 = 256$ candidate solutions!)

```
[< 1, 1 >, < 2, 1 >, < 3, 1 >, < 4, 1 >]  
[< 1, 1 >, < 2, 1 >, < 3, 1 >, < 4, 2 >]  
[< 1, 1 >, < 2, 1 >, < 3, 1 >, < 4, 3 >]  
[< 1, 1 >, < 2, 1 >, < 3, 1 >, < 4, 4 >]  
[< 1, 1 >, < 2, 1 >, < 3, 2 >, < 4, 1 >]
```


The Backtracking Advantage

- In backtracking, we take advantage of the fact that some paths cannot possibly lead to a solution:

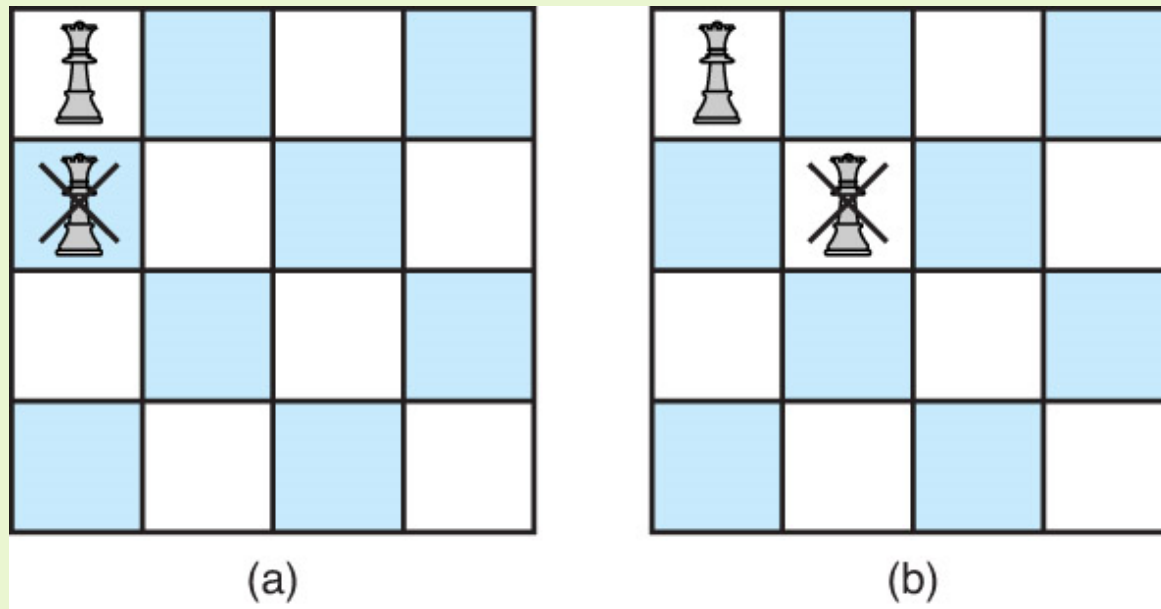


Figure 5.3: Diagram showing that if the first queen is placed in column 1, the second queen cannot be placed in column 1 (a) or column 2 (b).

The Backtracking Algorithm

In backtracking, each node visited via DFS is examined and found to be either

- ▶ Nonpromising: cannot lead to a solution, OR
- ▶ Promising: needs to be explored further.

Paths involving a **nonpromising** node are **pruned** from further consideration, resulting in a **pruned** state space tree.

```
void checknode (node v)
{
    node u;

    if (promising(v))
        if (there is a solution at v)
            write the solution;
        else
            for (each child u of v)
                checknode(u);
}
```

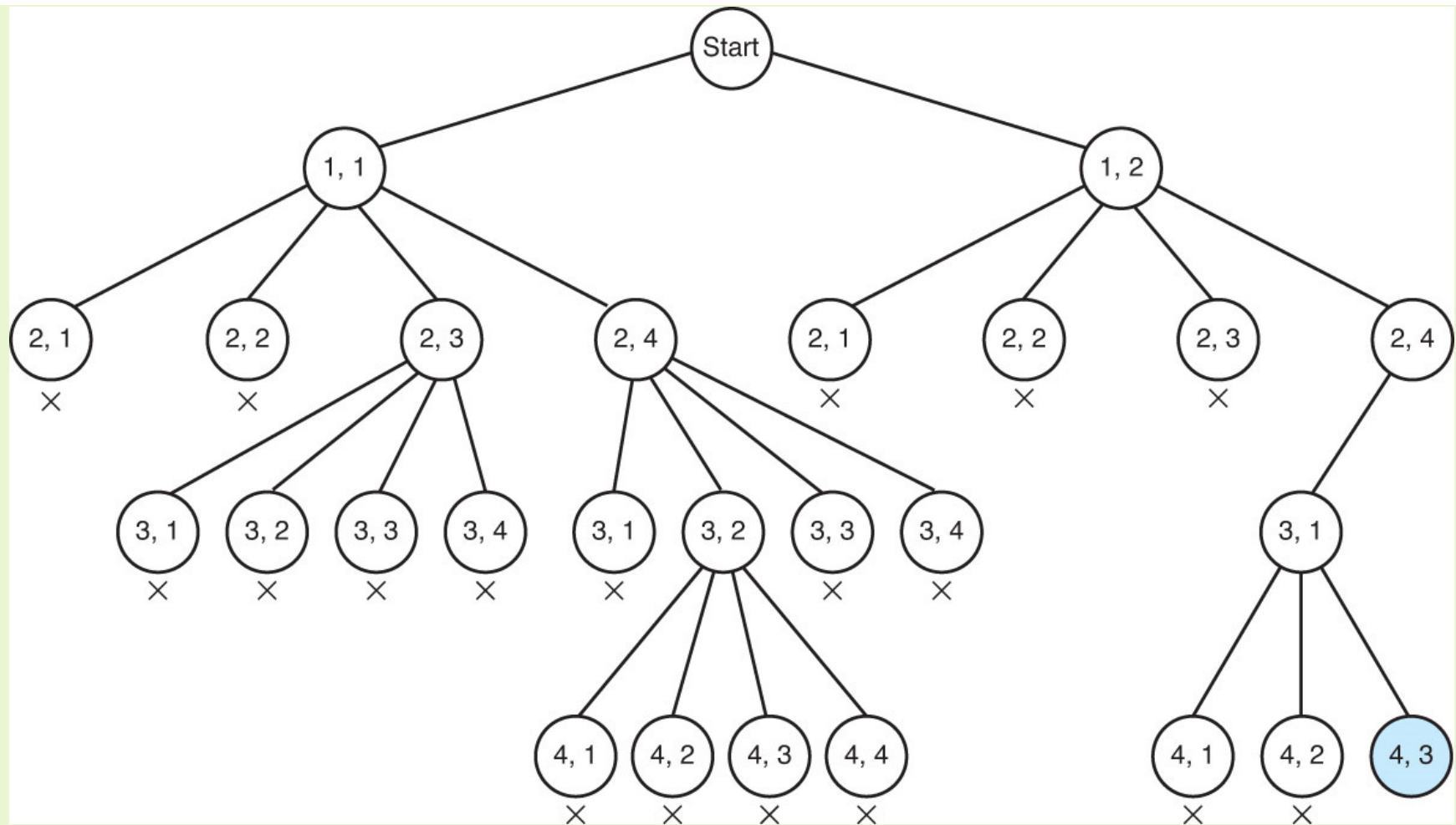


Figure 5.4: A portion of the pruned state space tree produced when backtracking is used to solve the instance of the n -Queens problems in which $n=4$.

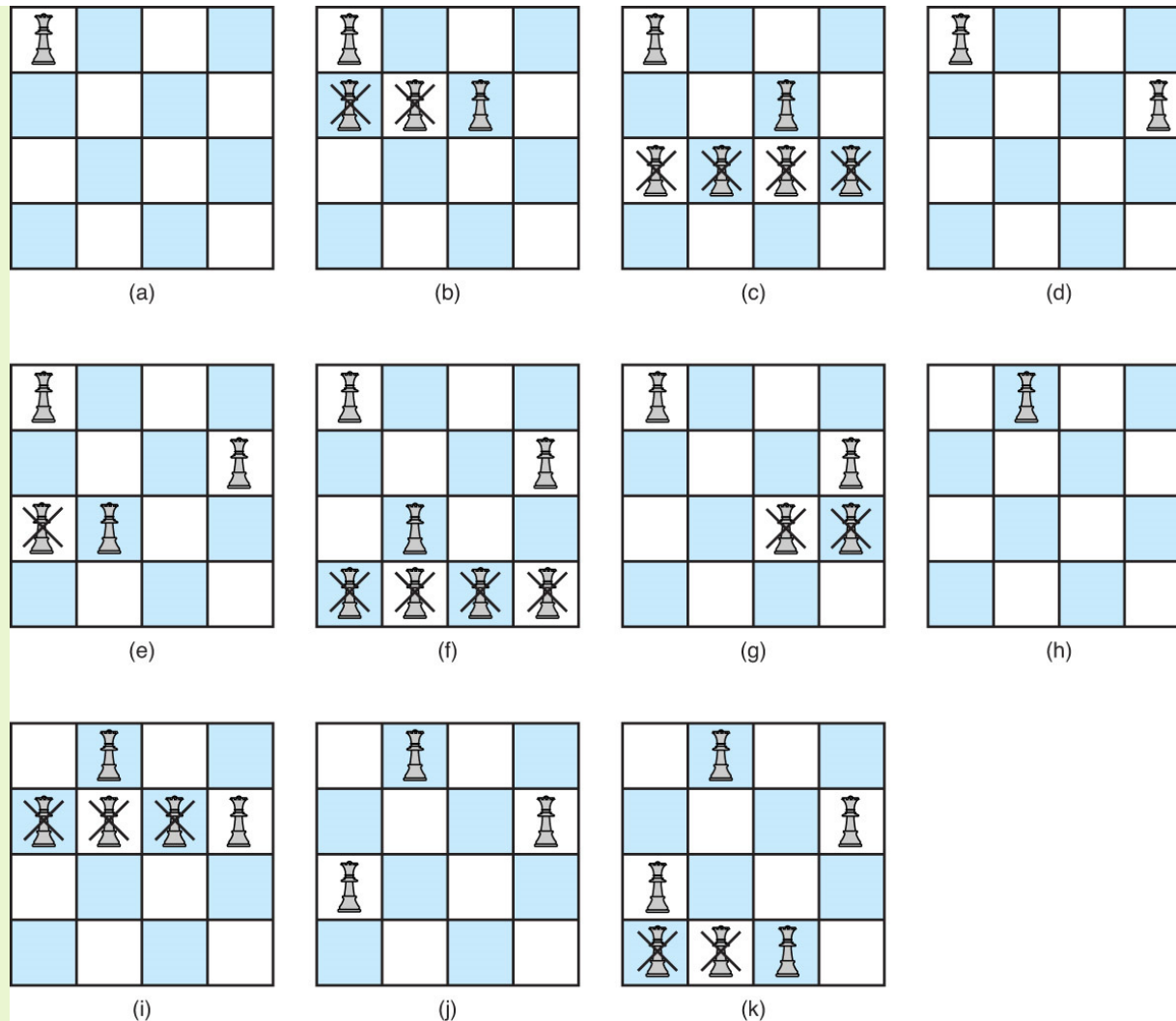


Figure 5.5: The actual chessboard positions that are tried when backtracking is used to solve the instance of the n -Queens problem in which $n=4$. Each non promising position is marked with a cross.