## **CS425 MP3: Simple Distributed File System Report**

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## General Algorithm:

There are 3 types of machines: mDNS, mServer, and mClient. mDNS is a fixed machine (assumed to not fail) that stores the current leader's IP and Port. mClient will resolve the master machine through mDNS. Each mServer instance is associated with an mServerType (either MASTER or SERVANT), and will execute accordingly. mClient can only communicate with mDNS and the master mServer.

Upon a PUT, the master server will split the file and distribute evenly within the SDFS system with a replication factor of 2. For example, if mClient uploads a 1GB file and there are 4 servers, each server will have 1 original and 1 replica each of size 250MB. Thus, our system employs an **even-distribution load-balancing scheme**. The original and replica chunks stored on each machine are not the same chunk, however (the primary and secondary replicas are never stored on same machine). For example, if there are 4 machines, a file will be split into 4 chunks and for each mServer instance *i*, it will have chunk *i* of type original and chunk *i+1* of type replica. The distribution scheme is explained in more details within the code.

Upon a GET, the master server will send any relevant chunks it may have, and then forward the request to the replication servers who will then accordingly connect to mClient and **push** the download.

The failure detection protocol employed is the **Ring Heartbeat Protocol**; when an mServer instance detects predecessor failure, it TCP multicasts the failure. The leader election protocol is **custom**: It is guaranteed in our system that the membership list for each mServer instance is consistent; therefore, it is logically sound for each mServer instance to simply elect the first non-faulty member in their membership list. The type of replication strategy employed is **active replication**; as long as there at least 2 servers, it is guaranteed there will be two replicas.

Using MP1 would theoretically be helpful in debugging MP3; however, since we created a custom mprintf (re-implementation of printf) that allowed us to write both to screen and log file, its use was extremely limited.

## Measurements:

- i. Re-replication time and bandwidth upon a failure (1GB file):
  - Re-replication time = number of replicas (2) \* time it takes to transfer and write chunk from mServer to mServer  $\cong$  9.2 seconds
  - Bandwidth upon failure is equivalent to size of chunks contained on the failed mServer  $\rightarrow$  250 \* 2 = 500 MB
- ii. Time Between Master Failure and Leader Elected
  - Let variable ntwrk\_latency represent the time it takes for a simple message to travel from one mServer instance to another. As discussed earlier in the General Algorithm section, our custom leader election allows for an upper time bound of (failure\_detection\_time + ntwrk\_latency).
- iii. Read/Write 1MB, 10MB, 100MB, 1GB → Tables and Graphs below
- iv. Wikipedia Corpus PUT: 122.7557702 Seconds

Table 1. Write (aka PUT)

Testfile	Trial	mClient PUT	Master mServer	Servant mServer	Total/System PUT	AVERAGE	STDEV
			PUT	PUT			
testfile1MB	Trial 1	0.002812	0.00476	0.0080427	0.0156147	0.01424454	0.001340859
	Trial 2	0.003018	0.003902	0.006736	0.013656		
	Trial 3	0.001388	0.004063	0.010315	0.015766		
	Trial 4	0.001774	0.003667	0.007674	0.013115		
	Trial 5	0.001765	0.003773	0.007533	0.013071		
testfile10MB	Trial 1	0.023063	0.035106	0.066476	0.124645	0.1015658	0.014261339
	Trial 2	0.015198	0.028689	0.052268	0.096155		
	Trial 3	0.015238	0.028556	0.043069	0.086863		
	Trial 4	0.016042	0.02746	0.060521	0.104023		
	Trial 5	0.016171	0.027527	0.052445	0.096143		

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testfile100MB	Trial 1	0.189124	0.371491	0.56943	1.130045	1.0735724	0.040349392
	Trial 2	0.136523	0.340994	0.616084	1.093601		
	Trial 3	0.129555	0.330522	0.562921	1.022998		
	Trial 4	0.130769	0.329924	0.603666	1.064359		
	Trial 5	0.130418	0.323096	0.603345	1.056859		
testfile1GB	Trial 1	2.83686	5.572365	9.24126	17.650485	17.5365386	0.557910235
	Trial 2	2.54789	5.987083	9.16573	17.700703		
	Trial 3	2.89768	5.763524	8.75632	17.417524		
	Trial 4	2.53265	5.556989	8.59979	16.689429		
	Trial 5	2.98009	5.878982	9.36548	18.224552		

Table 2: Read (aka GET)

Testfile	Trial	System GET	AVERAGE	STDEV
testfile1MB	Trial 1	0.013151	0.0133384	0.000414915
	Trial 2	0.013655		
	Trial 3	0.012938		
	Trial 4	0.013052		
	Trial 5	0.013896		
testfile10MB	Trial 1	0.106148	0.1182562	0.010801131
	Trial 2	0.123598		
	Trial 3	0.124365		
	Trial 4	0.129854		
	Trial 5	0.107316		
testfile100MB	Trial 1	2.766047	2.7928132	0.035879466
	Trial 2	2.837556		
	Trial 3	2.776873		
	Trial 4	2.824771		
	Trial 5	2.758819		
21.105				
testfile1GB	Trial 1	33.378858	33.7042124	0.644378996
	Trial 2	34.547876		
	Trial 3	34.236113		
	Trial 4	33.112456		
	Trial 5	33.245759		

Note: Overall, the graphs for these data points are useless due to the huge variance between (1MB, 1MB, 100MB) and 1GB. Therefore, just for the sake of showing trends (the standard deviation bars are indeed there, just too small to see):

