

# Computer Networks

## COL 334/672

Inter-domain Routing

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*Slides adapted from KR*

Sem 1, 2024-25

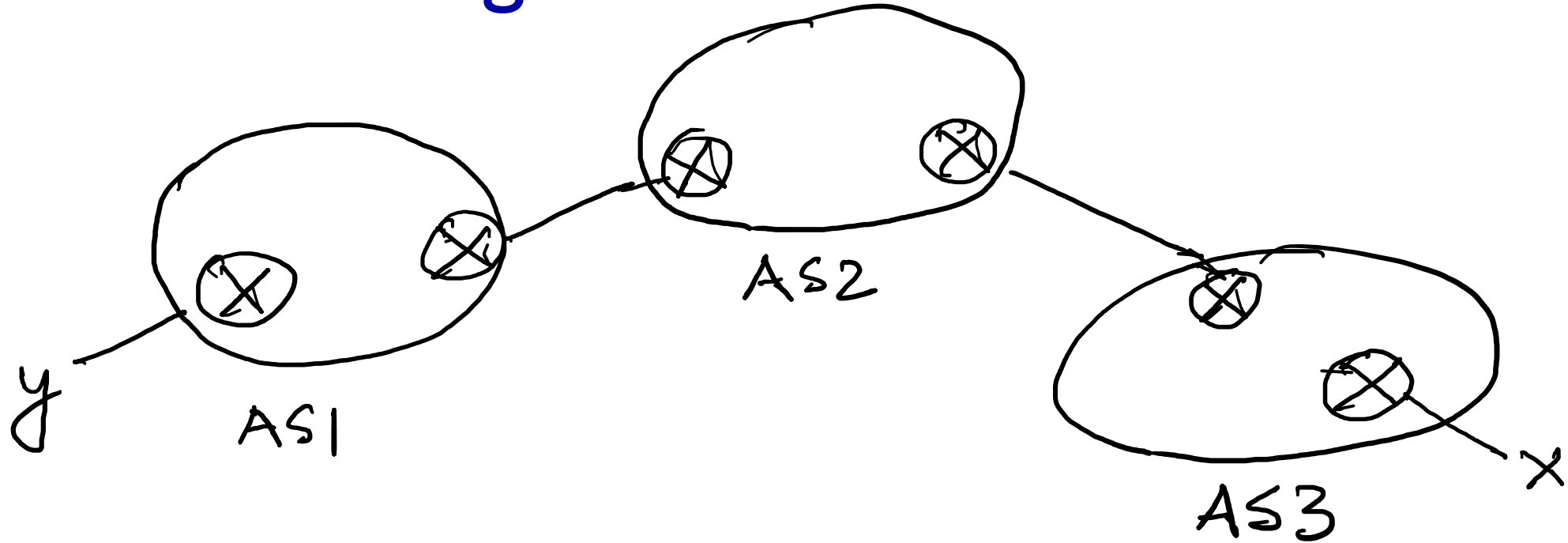
## Moodle Quiz: BananaFritters

# Recap: Intra-domain routing

Most common protocols for routing within an Autonomous System (AS)

- **RIP: Routing Information Protocol** [RFC 1723]
  - classic DV: DVs exchanged every 30 secs
  - no longer widely used
- **EIGRP: Enhanced Interior Gateway Routing Protocol**
  - DV based
  - formerly Cisco-proprietary for decades (became open in 2013 [RFC 7868])
- **OSPF: Open Shortest Path First** [RFC 2328]
  - link-state routing
  - IS-IS protocol (ISO standard, not RFC standard) essentially same as OSPF

# Inter-AS routing



*How does Internet route packet from x to y?*





What are the requirements for inter-AS routing?

- Common addressing scheme
- Common routing protocol

# IP addressing: introduction

- **IP address:** 32-bit identifier associated with each host or router *interface*
- dotted-decimal IP address notation:

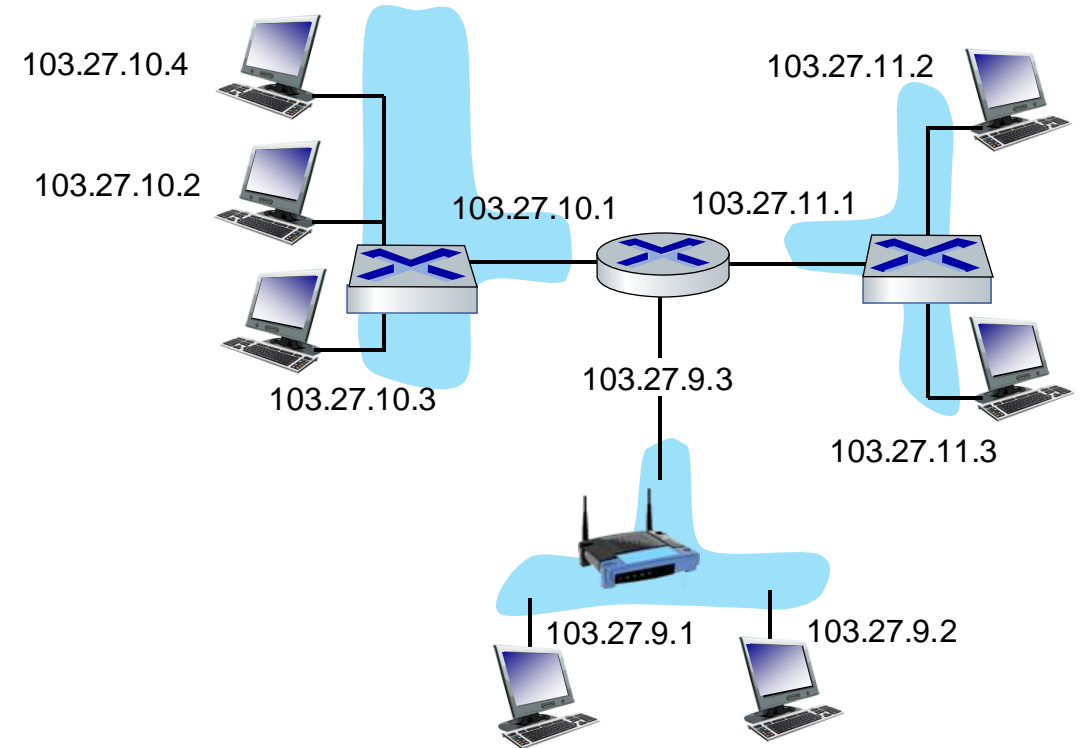
103.27.10.1 = 01100111 00011011 00001010 00000001

			
103	27	10	1

- How does an AS obtain IP address?
  - Need a coordinating agency
  - Internet Assigned Numbers Authority (IANA) and Internet Corporation for Assigned Names and Numbers (ICANN)

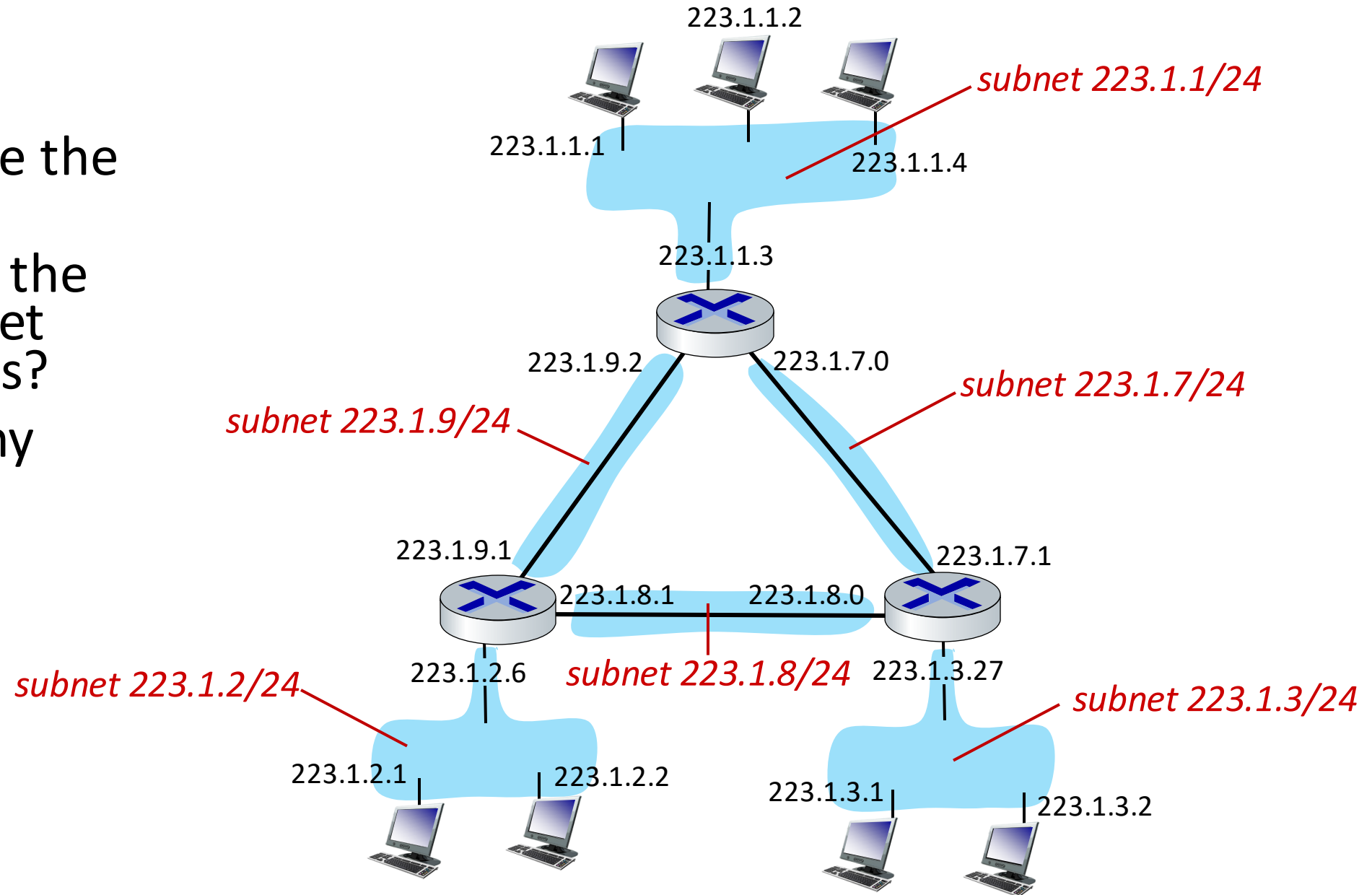
# IP addressing: Subnets

- *What's a subnet ?*
  - device interfaces that can physically reach each other **without passing through an intervening router**
- IP addresses have structure:
  - **subnet part**: devices in same subnet have common high order bits
  - **host part**: **remaining** low order bits



# Subnets

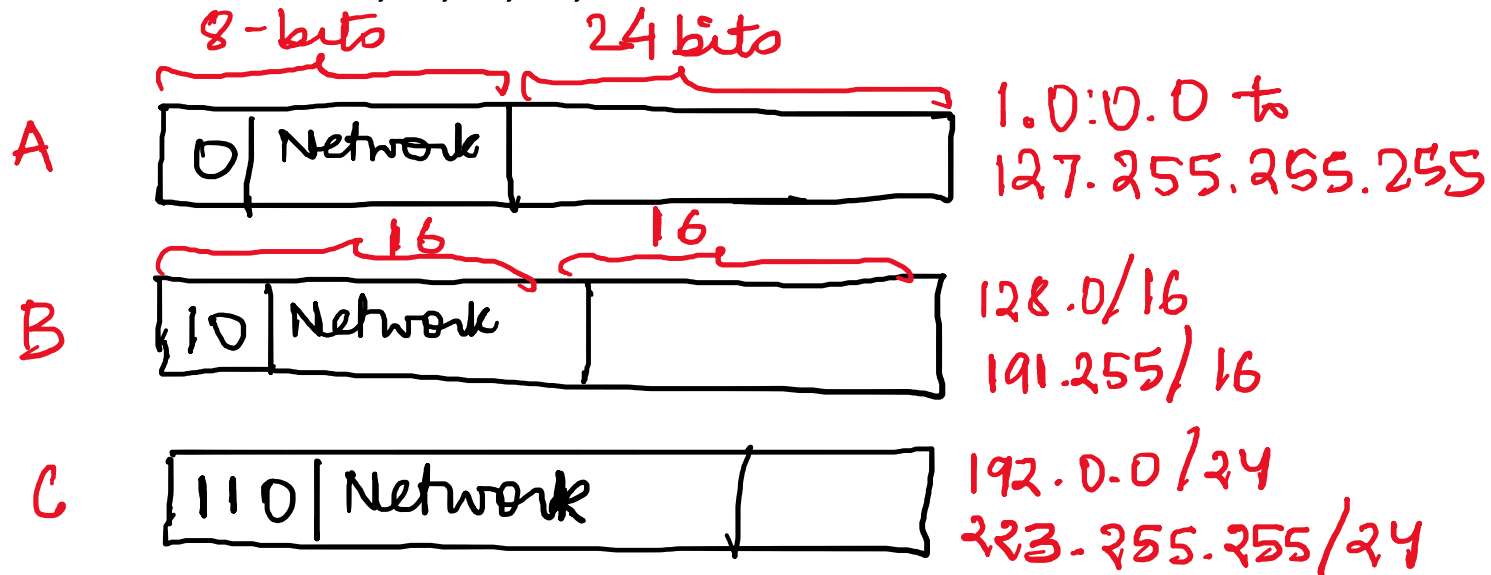
- where are the subnets?
- what are the /24 subnet addresses?
- how many subnets?



# Classful Addressing

- Why do we split the IP address into 4 parts?
- In the beginning, IP addresses were divided into 5 categories, called **classful addressing**

- Classes: A, B, C, D, and E



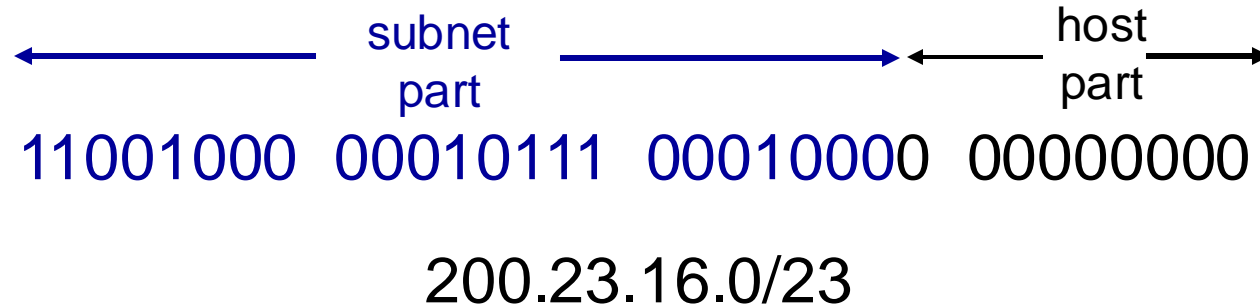
**Advantages:** Routing is easier  
**Limitations:** Wasteful



# IP addressing: CIDR


**CIDR: C**lassless **I**nter**D**omain **R**outing (pronounced “cider”)

- subnet portion of address of arbitrary length
- address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



# Example: IP Addresses of two AS

## AS3 – Massachusetts Institute of Technology

Country  [United States](#) ⓘ

Website [mit.edu](https://mit.edu)

Hosted domains 718

Number of IPv4 1,836,288

Number of IPv6  $6.34 \times 10^{29}$


ASN type Education

Registry ARIN

Allocated 55 years ago on Jan 01, 1970

Updated 14 years ago on Sep 27, 2010

## AS132780 – Indian Institute of Technology Delhi

Country  [India](#) ⓘ

Website [iitd.ac.in](https://iitd.ac.in)

Hosted domains 13

Number of IPv4 1,024

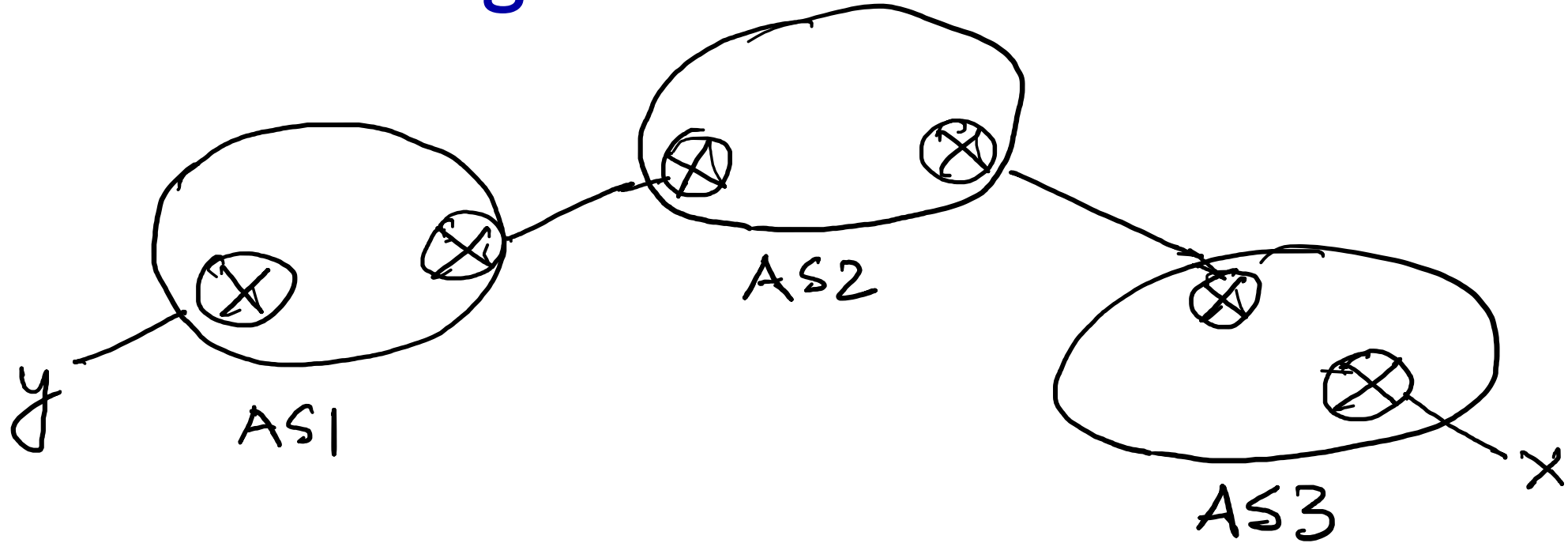
Number of IPv6  $1.21 \times 10^{24}$

ASN type Education

Registry APNIC

Updated 2 years ago on Aug 24, 2022

# Inter-AS routing



*How does Internet route packet from x to y?*

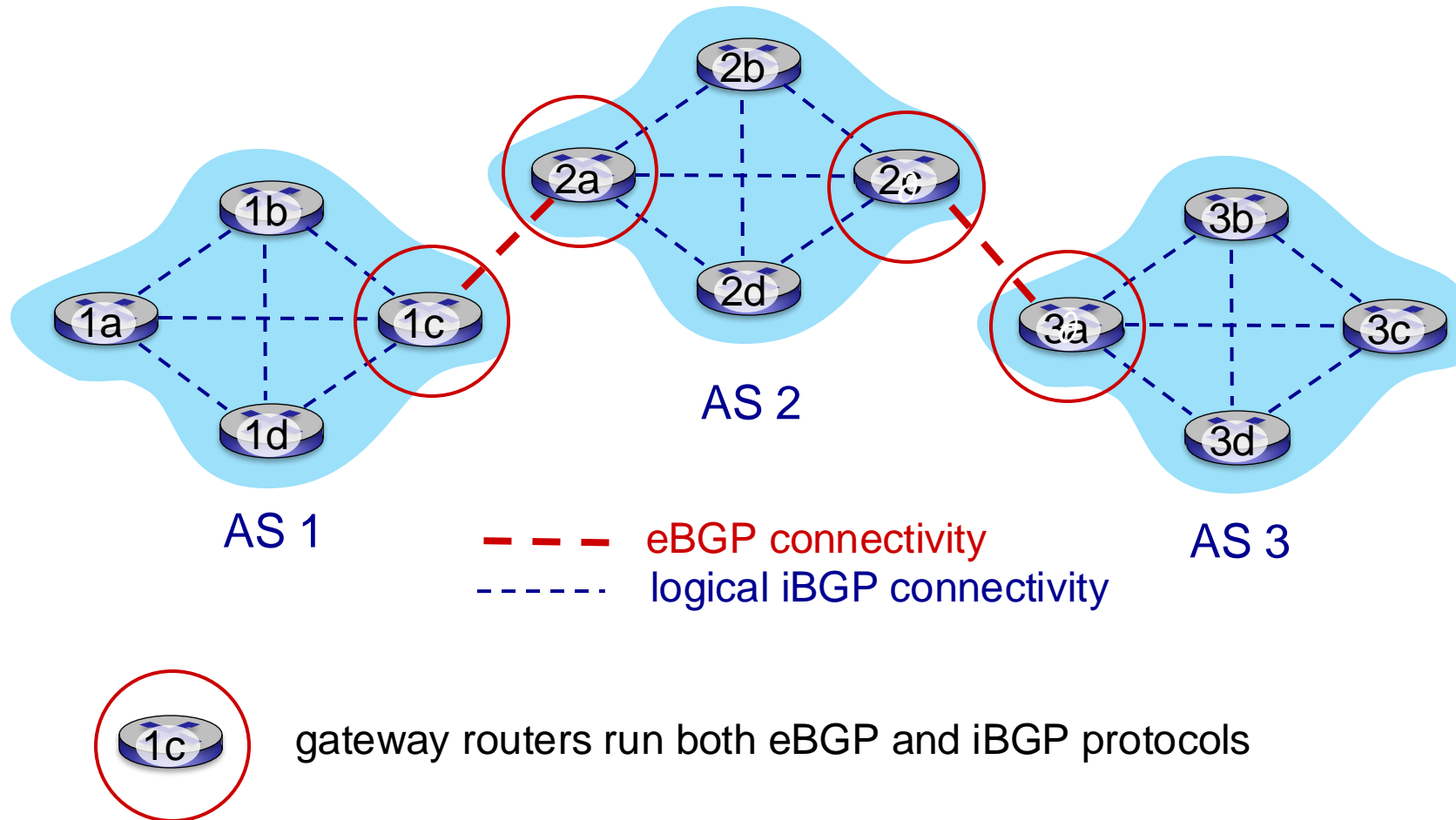
What are the requirements for inter-AS routing?

- Common addressing scheme
- **Common routing protocol**

# Internet inter-AS routing: BGP

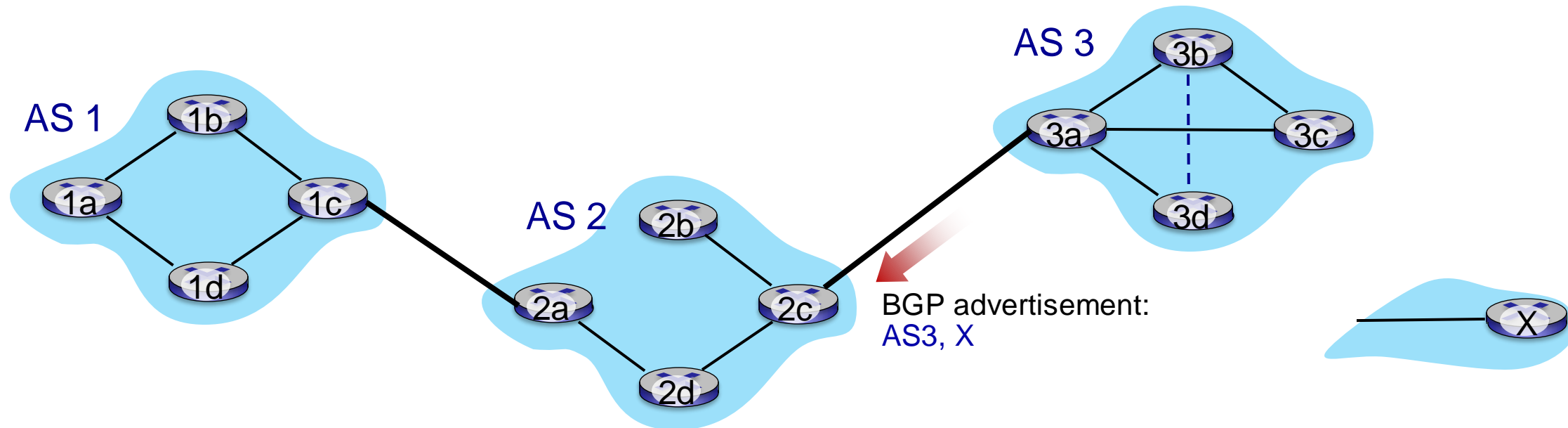
- **BGP (Border Gateway Protocol)**: *the* de facto inter-domain routing protocol
  - “glue that holds the Internet together”
- allows subnet to advertise its existence, and the destinations it can reach, to rest of Internet: *“I am here, here is who I can reach, and how”*
- BGP provides each AS a means to:
  - obtain destination network reachability info from neighboring ASes (**eBGP**)
  - determine routes to other networks based on reachability information and *policy*
  - propagate reachability information to all AS-internal routers (**iBGP**)
  - **advertise** (to neighboring networks) destination reachability info

# eBGP, iBGP connections



# BGP basics

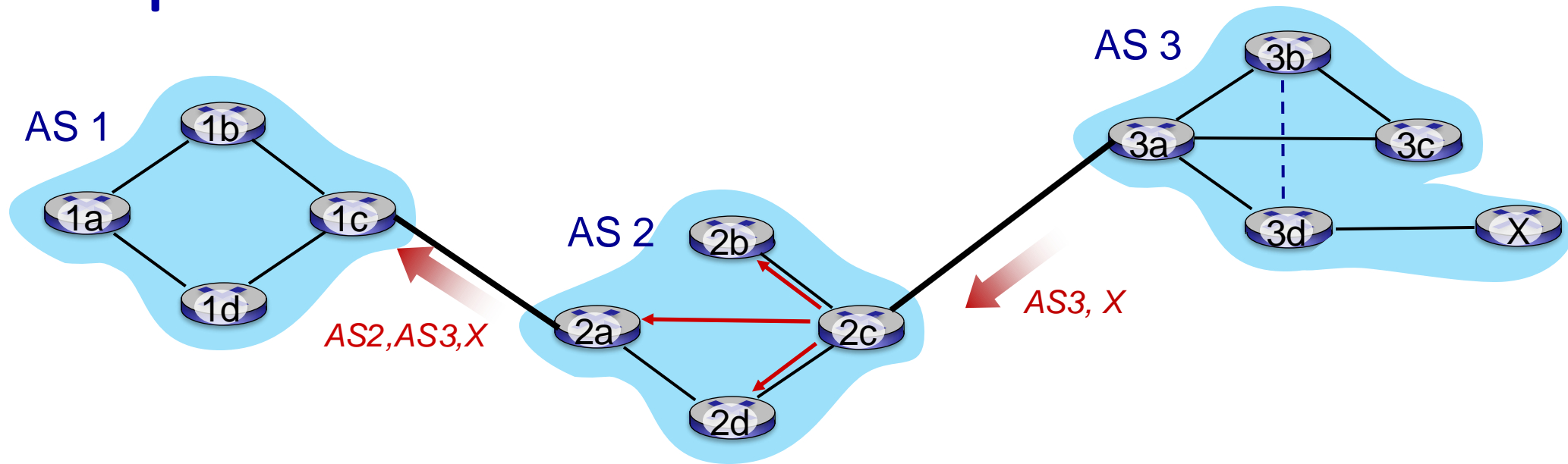
- **BGP session:** two BGP routers (“peers”) exchange BGP messages over semi-permanent TCP connection:
  - advertising *paths* to different destination network prefixes (BGP is a “path vector” protocol)
- when AS3 gateway 3a advertises *path AS3,X* to AS2 gateway 2c:
  - AS3 *promises* to AS2 it will forward datagrams towards X



# Path attributes and BGP routes

- BGP advertised route: prefix + attributes
  - prefix: destination being advertised
  - two important attributes:
    - **AS-PATH**: list of ASes through which prefix advertisement has passed
    - **NEXT-HOP**: indicates specific internal-AS router to next-hop AS
- **policy-based routing**:
  - gateway receiving route advertisement uses *import policy* to accept/decline path (e.g., never route through AS Y).
  - AS policy also determines whether to *advertise* path to other neighboring ASes

# BGP path advertisement



- AS2 router 2c receives path advertisement **AS3,X** (via eBGP) from AS3 router 3a
- based on AS2 policy, AS2 router 2c accepts path AS3,X, propagates (via iBGP) to all AS2 routers
- based on AS2 policy, AS2 router 2a advertises (via eBGP) path **AS2, AS3, X** to AS1 router 1c