

# CS335 Introduction to AI



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# First Order Logic

Can intelligent systems deduce stuff?

One Sherlock Holmes

Another Sherlock

BBC One's Sherlock (2010-) with Benedict Cumberbatch and Martin Freeman.

- ▶ Whorf (1956) suggest that communities determine lang. categories
- ▶ Wanner (1975) subject remember the **content** of what they read better than the actual words.
- ▶ Mitchell et al. (2008) could predict –with above chance accuracy– the areas of the brain that would activate with certain words(fMRI)

- ▶ Objects (cat, dog, house, John, etc.)
- ▶ Relations (has color, bigger than, comes between, etc.)
- ▶ Facts: (One value for a given input: has father, has head, can swim)

Facts have a truth value. *true or false*

# Formal Languages

## Ontological and Epistemological Commitments

Language	Ontological Commitment <sup>1</sup>	Epistemological Commitment <sup>2</sup>
Propositional Logic	facts	true/false/unknown
First-Order Logic	facts, objects, relations	true/false/unknown
Temporal Logic	facts, objects, relations, time	true/false/unknown
Probability Theory	facts	degree of belief $\in [0, 1]$
Fuzzy Logic	facts with degree of truth	known interval value

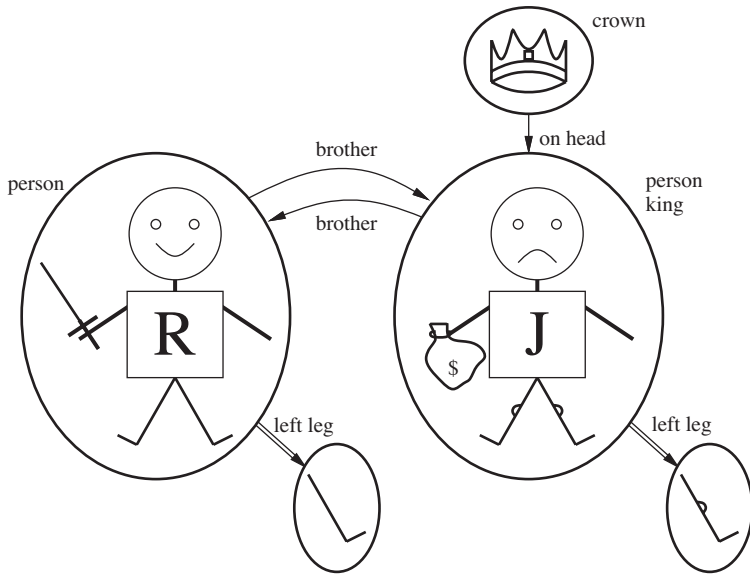
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<sup>1</sup>What exists in the world

<sup>2</sup>Agent's beliefs about facts

# Relationships

## Models for first order logic



# First-Order-Logic

## Syntax

<i>Sentence</i>	→	<i>AtomicSentence</i>   <i>ComplexSentence</i>
<i>AtomicSentence</i>	→	<i>Predicate</i>   <i>Predicate</i> ( <i>Term</i> , ... )  <i>Term</i> = <i>Term</i>
<i>ComplexSentence</i>	→	( <i>Sentence</i> ) [ <i>Sentence</i> ]
		$\neg$ <i>Sentence</i>
		<i>Sentence</i> $\wedge$ <i>Sentence</i>
		<i>Sentence</i> $\vee$ <i>Sentence</i>
		<i>Sentence</i> $\Rightarrow$ <i>Sentence</i>
		<i>Sentence</i> $\Leftrightarrow$ <i>Sentence</i>
		<i>Quantifier</i> <i>Variable</i> , ... <i>Sentence</i>
<i>Term</i>	→	<i>Function</i> ( <i>Term</i> , ... )
		<i>Constant</i>
		<i>Variable</i>
<i>Quantifier</i>	→	$\forall$   $\exists$
<i>Constant</i>	→	<i>A</i>   <i>X<sub>1</sub></i>   <i>John</i>   ...
<i>Variable</i>	→	<i>a</i>   <i>x</i>   <i>s</i>   ...
<i>Predicate</i>	→	<i>True</i>   <i>False</i>   <i>After</i>   <i>Loves</i>   <i>Raining</i>   ...
<i>Function</i>	→	<i>Mother</i>   <i>Left leg</i>   ...
Operator Precedence <sup>3</sup>	:	$\neg$ , $\wedge$ , $\vee$ , $\Rightarrow$ , $\Leftrightarrow$

<sup>3</sup>Otherwise the grammar is ambiguous

- ▶ Three kinds of symbols
  - ▶ Constant: objects
  - ▶ Predicate: relations
  - ▶ Function: functions (i.e. can return values other than truth vals.)
- ▶ Predicate and Function have **arity**.
- ▶ Symbols have an interpretation
- ▶ Terms: *LeftLeg(John)*
- ▶ Atomic Sentences state facts: *Brother(Richard, John)*
- ▶ Complex Sentence: *Brother(R, J)  $\wedge$  Brother(J, R)* or  *$\neg$ King(Richard)  $\Rightarrow$  King(John)*
- ▶ Universal Quantifiers:  $\forall King(x) \Rightarrow Person(x)$
- ▶ Existential Quantifiers:  $\exists Crown(x) \wedge OnHead(x, John)$



What is the interpretation for:

- ▶  $King(Richard) \vee King(John)$
- ▶  $\neg Brother(LeftLeg(Richard), John)$
- ▶  $\forall x \forall y Brother(x, y) \Rightarrow Sibling(x, y)$
- ▶  $In(Paris, France) \wedge In(Marseilles, France)$
- ▶  $\forall c Country(c) \wedge Border(c, Ecuador) \Rightarrow In(c, SouthAmerica)$
- ▶  $\exists Country(c) \wedge Border(c, Spain) \wedge Border(c, Italy)$

## First Order Logic

### More Facts

- Richard has only two brothers, John and Geoffrey:

$Brother(John, Richard) \wedge Brother(Geoffrey, Richard) \wedge John \neq Geoffrey \wedge$   
 $\forall x Brother(x, Richard) \Rightarrow (x = John \vee x = Geoffrey)$

- No Region in South America borders any region in Europe

$\forall c, d In(c, SouthAmerica) \wedge In(d, Europe) \Rightarrow \neg Border(c, d)$

- No two adjacent countries have the same map color

$\forall x, y Country(x) \wedge Country(y) \wedge Border(x, y) \Rightarrow$   
 $\neg (Color(x) = Color(y)) \wedge \neg (x = y)$

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## Assertions and Queries in FOL

### ASK and TELL

- ▶  $TELL(KB, King(John))$
- ▶  $TELL(KB, \forall x King(x) \Rightarrow Person(x))$
- ▶  $ASK(KB, King(John))$  return True
- ▶  $ASK(KB, \exists x Person(x))$  return True
- ▶  $ASKVARS(KB, Person(x))$  yields  $\{x/John, x/Richard\}$ , a binding list



“The son of my father is my brother”; “One’s grandmother is the mother of one’s parent”; etc.

- ▶ Domain: People.
- ▶ Unary predicates: *Male, Female*
- ▶ Relations:  
*Parent, Sibling, Brother, Sister, Child, Daughter, Son, Spouse, Wife, Husband, Grandparent, Grandchild, Cousin, Aunt, Uncle*
- ▶ Functions: *Mother, Father*

“One’s mother is one’s female parent”

$\forall m, c \text{ Mother}(c) = m \leftrightarrow \text{Female}(m) \wedge \text{Parent}(m, c)$

“A sibling is another child of one’s parents”

$\forall x, y \text{ Sibling}(x, y) \leftrightarrow x \neq y \wedge \exists p \text{ Parent}(p, x) \wedge \text{Parent}(p, y)$

“Wendy is female”

$\text{Female}(\text{wendy})$

These are **axioms**

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These are **axioms**

Can be represented more concisely

- ▶ at time step  
3: *Percept*([*Stench*, *Breeze*, *Glitter*, *None*, *None*], 3)
- ▶ at time step  
6: *Percept*([*None*, *Breeze*, *None*, *None*, *Scream*], 6)
- ▶ Actions can be:  
*Turn*(*Right*), *Turn*(*Left*), *Forward*, *Shoot*, *Grab*, *Climb*

And we can ASK the best action at time step 5

*ASKVARS*( $\exists a$  *BestAction*(*a*, 5))

Can encode:

- ▶ Raw percepts:

$$\forall t, s, g, m, c \text{ Percept}([s, b, \textit{Glitter}, m, c], t) \Rightarrow \textit{Glitter}(t)$$

- ▶ Reflex actions:  $\forall t \text{ Glitter}(t) \Rightarrow \textit{BestAction}(\textit{Grab}, t)$

Instead of encoding stuff like:

*Adjacent*(*Square*<sub>1,2</sub>, *Square*<sub>1,1</sub>)

*Adjacent*(*Square*<sub>3,4</sub>, *Square*<sub>4,4</sub>)

Encode:

$$\forall x, y, a, b \text{ Adjacent}([x, y], [a, b]) \Leftrightarrow$$

$$(x = a \wedge (y = b - 1 \vee y = b + 1)) \vee (y = b \wedge (x = a - 1 \vee x = a + 1))$$



- ▶ Define pieces:  $Long(p)$ ,  $Short(p)$ , etc.
- ▶ and restrictions:  $\forall p \text{ } Long(p) \leftrightarrow \neg(Long(p) \wedge Short(p))$
- ▶ Define rules:  $\forall x, y, z \text{ } CanConnectWithOverlap(x, y, z) \leftrightarrow x \neq y \wedge Piece(x) \wedge Piece(y) \wedge Number(z) \wedge Value(z) \leq Overlap(x, y) \dots$
- ▶  $\forall x, y \text{ } Short(x) \wedge Short(y) \wedge Overlap(x, y) < 1 \Rightarrow WeakLink(x, y)$

Alain Colmerauer (1972)

Download SWI Prolog

And a quick tutorial

## Creating a Knowledge Base

- ▶ Identify the Task
- ▶ Assemble the relevant knowledge
- ▶ Decide on a vocabulary of predicates, functions and constants
- ▶ Encode general knowledge about the domain (rules)
- ▶ Encode a description of the problem
- ▶ Pose queries to the inference procedure and get answers
- ▶ Debug the KB

Given  $\forall x \text{ King}(x) \wedge \text{Greedy}(x) \Rightarrow \text{Evil}(x)$

One can infer

- ▶  $\text{King}(\text{John}) \wedge \text{Greedy}(\text{John}) \Rightarrow \text{Evil}(\text{John})$
- ▶  $\text{King}(\text{Richard}) \wedge \text{Greedy}(\text{Richard}) \Rightarrow \text{Evil}(\text{Richard})$
- ▶  $\text{King}(\text{Father}(\text{John})) \wedge \text{Greedy}(\text{Father}(\text{John})) \Rightarrow \text{Evil}(\text{Father}(\text{John}))$

## Inference in First Order Logic

- ▶ Universal Instantiation (in a  $\forall$  rule, substitute all symbols)
- ▶ Existential Instantiation (in a  $\exists$  rule, substitute one symbol, use the rule and discard)
- ▶ Skolem Constants is a new variable that represents a new inference.

Suppose KB:

- ▶  $\forall x \text{ King}(x) \wedge \text{Greedy}(x) \Rightarrow \text{Evil}(x)$
- ▶  $\text{King}(\text{John})$
- ▶  $\text{Greedy}(\text{John})$
- ▶  $\text{Brother}(\text{Richard}, \text{John})$

Apply UI using  $\{x/\text{John}\}$  and  $\{x/\text{Richard}\}$

- ▶  $\text{King}(\text{John}) \wedge \text{Greedy}(\text{John}) \Rightarrow \text{Evil}(\text{John})$
- ▶  $\text{King}(\text{Richard}) \wedge \text{Greedy}(\text{Richard}) \Rightarrow \text{Evil}(\text{Richard})$

And discard the Universally quantified sentence. We can get the KB to be propositions.

## Inference in First Order Logic

Suppose KB:

- ▶  $\forall x \text{ King}(x) \wedge \text{Greedy}(x) \Rightarrow \text{Evil}(x)$
- ▶  $\text{King}(\text{John})$
- ▶  $\forall y \text{ Greedy}(y)$

Apply UI using  $\{x/\text{John}\}$  and  $\{y/\text{John}\}$

## Inference

### Generalized Modus Ponens

for atomic sentences  $p_i, p'_i$  and  $q$ , where there is a substitution  $\theta$  such that  $SUBST(\theta, p'_i) = SUBST(\theta, p_i)$ , for all  $i$

$$\frac{p'_1, p'_2, \dots, p'_n, (p_1 \wedge p_2 \wedge \dots \wedge p_n \Rightarrow q)}{SUBST(\theta, q)}$$

$$p'_1 = King(John)$$

$$p_1 = King(x)$$

$$p'_2 = Greedy(y)$$

$$p_2 = Greedy(x)$$

$$\theta = \{x/John, y/John\}$$

$$q = Evil(x)$$

$$SUBST(\theta, q)$$

.



$UNIFY(p, q) = \theta$  Where  $SUBST(\theta, p) = SUBST(\theta, q)$

For example:

- ▶ We ask  $ASKVARS(Knows(John, x))$  (Whom does John know?)
- ▶  $UNIFY(Knows(John, x), Knows(John, Jane)) = \{x/Jane\}$
- ▶  $UNIFY(Knows(John, x), Knows(y, Bill)) = \{y/John, x/Bill\}$
- ▶  $UNIFY(Knows(John, x), Knows(y, Mother(y))) = \{y/John, x/Mother(John)\}$

Algorithm in the book (goes variable by variable recursively unifying)

## Inference

Putting it all together

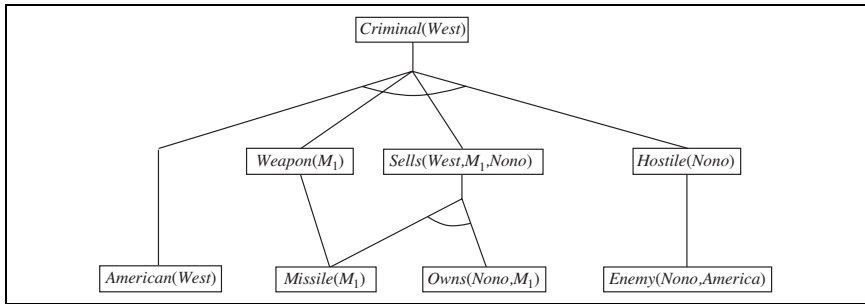
“The Law says that it is a crime for an American to sell weapons to hostile nations. The country Nono, and enemy of America, has some missiles , and all of its missiles were sold to it by Colonel West, who is American”

Prove that Colonel Wes is a Criminal

“The Law says that it is a crime for an American to sell weapons to hostile nations. The country Nono, and enemy of America, has some missiles , and all of its missiles were sold to it by Colonel West, who is American”

- ▶ R1:  
 $American(x) \wedge Weapon(y) \wedge Sells(x, y, z) \wedge Hostile(z) \Rightarrow Criminal(x)$
- ▶ R2:  $Owns(Nono, M_1)$  Nono has some missiles
- ▶ R3:  $Missile(M_1)$
- ▶ R4:  $Missile(x) \Rightarrow Weapon(x)$  A missile is a weapon
- ▶ R5:  $Missile(x) \wedge Owns(Nono, x) \Rightarrow Sells(West, x, Nono)$  All missiles sold by west
- ▶ R6:  $Enemy(x, America) \Rightarrow Hostile(x)$  Enemies of America are hostile
- ▶ R7:  $American(West)$  West is american
- ▶ R8:  $Enemy(Nono, America)$

# Inference Graph



## Forward Chaining ASK Iterations

### Iteration 1:

- ▶ R5 satisfied with  $\{x/M_1\}$  and R9: *Sells(West,  $M_1$ , Nono)* is added
- ▶ R4 satisfied with  $\{x/M_1\}$  and R10: *Weapon( $M_1$ )* is added
- ▶ R6 satisfied with  $\{x/Nono\}$  and R11: *Hostile(Nono)* is added

### Iteration 2:

- ▶ R1 is satisfied with  $\{x/West, y/M_1, z/Nono\}$  and *Criminal(West)* is added.

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### Iteration 1:

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## Inference in First Order Logic

### Discussion

- ▶ Once we have facts that evaluate to T or F
- ▶ We can apply Forward Chaining, Backwards Chaining and Resolution
- ▶ The key is to understand Unification
- ▶ Very similar to Logical agents.