

RSA Cryptosystem

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RSA cryptosystem

First published:

- Scientific American, Aug. 1977 (Patent up to Sept 21, 2000)

Currently the “work horse” of Internet security:

- Most Public Key Infrastructure (PKI) products.
- SSL/TLS: Certificates and key-exchange.
- Secure e-mail: PGP, Outlook, ...

Transport Layer Security
Secure Sockets Layer

RSA trapdoor 1-to-1 function

Parameters: $N=pq$. $N \approx 1024$ bits. $p, q \approx 512$ bits.
 e – encryption exponent. $\gcd(e, \phi(N)) = 1$.

1-to-1 function: $\text{RSA}(M) = M^e \pmod{N}$ where $M \in \mathbb{Z}_N^*$

Trapdoor: d – decryption exponent.

Where $e \cdot d = 1 \pmod{\phi(N)}$

Inversion: $\text{RSA}(M)^d = M^{ed} = M^{k\phi(N)+1} = M \pmod{N}$

(n, e, t, ϵ) -RSA Assumption: For any t -time algorithm A :

$$\Pr \left[A(N, e, x) = x^{1/e} \pmod{N} : \begin{array}{l} p, q \stackrel{R}{\leftarrow} \text{n-bit primes,} \\ N \leftarrow pq, \quad x \stackrel{R}{\leftarrow} \mathbb{Z}_N^* \end{array} \right] < \epsilon$$

$$\begin{aligned} e \cdot d &= 1 \pmod{\phi(N)} \\ ed &= k\phi(N) + 1 \\ a &= 1 \pmod{m} \end{aligned}$$

Example 1 - Key Setup

$$\begin{aligned} e^{-1} \bmod \phi(N) \\ 7^{-1} \bmod 160 \\ = \end{aligned}$$

1. Select primes: $p=17$ & $q=11$
2. Calculate $N = pq = 17 \times 11 = 187$
3. Calculate $\phi(N) = (p-1)(q-1) = 16 \times 10 = 160$
4. Select e : $\gcd(e, 160) = 1$; choose $e=7$
5. ? Determine d : $de = 1 \bmod 160$ and $d < 160$
6. Value is $d=23$ since $23 \times 7 = 161 = 10 \times 160 + 1$
7. Publish public key $PU = \{7, 187\} = \{e, N\}$
8. Keep secret private key $PR = \{23, 187\} = \{d, N\}$

Example - RSA En/Decryption

1. Publish public key
PU = { 7, 187 }
2. Keep secret private key
PR = { 23, 187 }

➤ RSA encryption/decryption is:

➤ Given message M = 88

➤ Encryption:

$$C = 88^7 \bmod 187 = 11$$

➤ Decryption:

$$M = 11^{23} \bmod 187 = \underline{\underline{88}}$$

Handwritten calculations in red:

$$\begin{aligned} C &= m^e \bmod N \\ m &= C^d \bmod N \\ (88^7)^{23} &= (88) \\ &= 88^{7 \times 23} \bmod 187 \\ &= 88^{161} \bmod 187 \\ &= 88 \bmod 187 \end{aligned}$$

Example 2

$$p = 11, q = 7, N = 77, \Phi(N) = 60$$

$$e = 37 \quad (\text{ed} = 481; \text{ed mod } 60 = 1)$$

What is d ?

$$d = 13$$

Let $M = 15$.

$$\text{Then } C \equiv M^e \pmod{N}$$

- $C \equiv 15^{37} \pmod{77} = 71$

$$M \equiv C^d \pmod{n}$$

- $M \equiv 71^{13} \pmod{77} = 15$

Handwritten notes:

$$\begin{aligned} C &= M^e \pmod{N} \\ &= 15^{37} \pmod{77} \\ &= 71 \end{aligned}$$

encryptions.

Example 3

$$(e, \phi(N)) = 2.$$

Parameters:

- $p = 3, q = 5, N = pq = 15$
- $\Phi(N) = ?$ 8

Let $e = 3$, what is d ?

Given $M=2$, what is C ?

How to decrypt?

$$\phi(N) = (p-1)(q-1)$$

$$= 2 \times 4 = 8$$

$$(d) = 3, \quad ed = 3 \times 3 \pmod{8} = 1.$$

$$C = M^e = 2^3 = 8 \pmod{15} = 8$$

$$M = C^d = 8^3 = 2 \pmod{15}$$

Example

Alice \rightarrow Bob, as
 $\{n, e\} = \{187, 7\}$
public.

Suppose Alice wishes to send a plaintext message M to Bob using the RSA algorithm.

Bob's public-key is $(n, e) = (187; 7)$. Note that $187 = 17 * 11$.

Alice uses an alphabet set of only 10 letters and encode them as

$A = 0; C = 1; D = 2; E = 3; I = 4; N = 5; O = 6; R = 7; T = 8; U = 9$.

Alice transmits the message in blocks. Each block corresponding to two letters which are encoded into their numerical equivalent, e.g., NO encodes as [56] and then it is encrypted using RSA.

If Alice wants to send the text "NO", what ciphertext will be received by Bob ?

$$78 = RT$$

Q: Suppose Bob receives [11], then what was the message transmitted by Alice?

Ans: [88]

→ $\{n, e\}$ - public

$\{n, d\}$ - private

$$e, (e, \phi(n)) \rightarrow \begin{cases} n = p \times q \\ \phi(n) = (p-1)(q-1) \end{cases}$$

$de \equiv 1 \pmod{\phi(n)}$

$$C = m^e \pmod{n}$$

$$m = C^d \pmod{n}$$

Q: $\phi(n)$ - secret | ~~public~~ ??

We can find $d = e^{-1} \pmod{\phi(n)}$??

$\Phi(N)$ implies factorization

Knowing both n and $\Phi(N)$, one knows

$$\underline{N = pq}$$

$$\Phi(N) = (p-1)(q-1) = \underline{pq - p - q + 1}$$

$$= N - p - N/p + 1$$

$$p\Phi(N) = Np - p^2 - N + p$$

$$p^2 - Np + \Phi(N)p - p + N = 0$$

$$p^2 - (N - \Phi(N) + 1)p + N = 0$$

There are two solutions of p in the above equation.

Both p and q are solutions.

Thank You