

Substitution

→ Substitution

→ Permutation.

Substitution cipher

* Suppose S_n is a set of 'n' symbols.

$\pi: S_n \rightarrow S_n$ is called permutation if it is a bijective mapping.

$$\pi = (\pi(1), \pi(2), \dots, \pi(n)).$$

(or)

$$\pi = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ \pi(1) & \pi(2) & \pi(3) & \dots & \pi(n) \end{pmatrix}$$

$$f = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 1 & 3 & 4 \end{pmatrix} \quad g = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 1 & 3 & 2 & 4 \end{pmatrix}$$

Note:-

$S_n \rightarrow$ set of all permutation with 'n' symbols

$$S = \{1, 2, 3, \dots, n\}$$

$$\Rightarrow f, g \in S_n$$

$$\Rightarrow (f \circ g)(x) = f(g(x))$$

$$f \circ g = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 3 & 1 & 4 \end{pmatrix} \rightarrow \text{Ans}$$

$$\begin{aligned} (f \circ g)(1) &= f(g(1)) \\ &= f(1) \\ &= 2 \end{aligned}$$

$$\begin{aligned} \therefore f \circ g(2) &= f(g(2)) \\ &= f(3) \\ &= 3 \end{aligned}$$

$$f \circ g(3) = f(g(3))$$

$$= f(2) = 2$$

$$f \circ g(4) = f(g(4))$$

$$= f(4) = 4$$

Note :-

$f \circ g \neq g \circ f \Rightarrow$ composition is not commutative.

* Transposition cipher:

$$P = C = (\mathbb{Z}_{26})^b \quad b\text{-size of message block.}$$

(Key) K - set of all permutations over $(1, 2, \dots, b)$.

$$\pi = (\pi(1), \pi(2), \dots, \pi(b))!$$

$$\pi^{-1} \circ \pi \Rightarrow \pi^{-1}(\pi(i)) = i, \forall i$$

(inverse)

$$\Rightarrow \text{block } (x_1, x_2, \dots, x_b) \in P.$$

then encryption $E_{\pi}(x_1, x_2, \dots, x_b)$,

$$= (x_{\pi(1)}, x_{\pi(2)}, \dots, x_{\pi(b)}).$$

decryption. $d_{\pi}(y_1, y_2, \dots, y_b)$.

$$= (y_{\pi^{-1}(1)}, y_{\pi^{-1}(2)}, \dots, y_{\pi^{-1}(b)}).$$

$$\Rightarrow y = f(x)$$

$$x = \bar{f}(y).$$

Eg:-

$$f = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 3 & 1 & 4 \end{pmatrix}$$

$$\bar{f} = \begin{pmatrix} 2 & 3 & 1 & 4 \\ 1 & 2 & 3 & 4 \end{pmatrix}$$

$$f \circ \bar{f} = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 3 & 1 & 2 & 4 \end{pmatrix} \xrightarrow{\text{for } 1 \ 2 \ 3 \ 4 - 1/p} \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 4 & 1 & 3 \end{pmatrix}$$

$$f \circ \bar{f} = (1 \ 2 \ 3 \ 4)$$

written input
in this order

$$\text{Let } K = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 1 & 3 & 4 & 2 \end{pmatrix} \xrightarrow{\text{Key order.}}$$

$$x = \begin{matrix} acd \\ x_1 x_2 x_3 x_4 \end{matrix}$$

$$y = \underline{\underline{22}}$$

ge

Sol

~~if~~ π is

$$x = \begin{matrix} acd \\ 13042 \end{matrix} \xrightarrow{\text{if } \pi \text{ is one-one}} y = \underline{\underline{22}}$$

$$\text{affinity to } \text{enc}_K(acdb) = x_{K(1)} \cdot x_{K(2)} \cdot x_{K(3)} \cdot x_{K(4)}$$

$$\begin{aligned} x_{\pi(1)} &= x_1 x_3 x_4 x_2 \\ &= \text{acd b.c.} \\ x_{\pi(2)} &= y_1 y_2 y_3 y_4 \\ \text{and } x_{\pi(3)} &= \text{abcd} \\ &= x_{\pi(4)}. \end{aligned}$$

This mapping provides an \rightarrow one-one correspondence.

$$\begin{aligned} \text{enc}_K(acdb) &= y_{K(1)} \cdot y_{K(2)} \cdot y_{K(3)} \cdot y_{K(4)} \\ &= y_1 y_4 y_2 y_3 \\ &= \text{abcd.} \end{aligned}$$

$$\text{Eq: } \Rightarrow x = \text{acdb } \text{rs. } \text{cd.}$$

Note :- → Transposition is not one-one substitution.
 → Here, 'c' will substitute with different letters.

* Binary operation :-

→ A '*' is called binary operation defined over non-empty 'S' if $a * b \in S$, $\forall a, b \in S$. we denote it with $(S, *)$.

Ex:- $(N, -)$ \Rightarrow non-binary operation over \mathbb{N} .

Eg:-

$$3 - 4 \notin N$$

$(\mathbb{Z}, -)$ \Rightarrow yes, binary operation (not identity) $3 - 0 = 0 - 3$

Ex:- (S_n, \circ) \Rightarrow Yes, binary operation.

(Composition of two permutation is always a permutation).

* Group :-

A non-empty set 'G' is binary operation with

* is called a group. if it satisfy.

(i) closure.

(ii) associative.

(iii) Identity.

(iv) Inverse.

Eg:- (S_n, \circ) \Rightarrow group.

* Lagrange's theorem :-

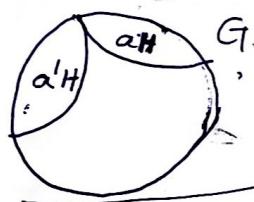
If 'G' is a group and 'H' is a subgroup of 'G', then $\frac{o(H)}{o(G)}$ divides.

Pf :-

(i). For any $a \neq a'$, if $a \notin a'H$, then

$$aH \cap a'H = \emptyset.$$

(ii). $o(aH) = o(H)$.



Consider,

$$a = ae = a(c\bar{c}^1).$$

$$= (ac)\cdot \bar{c}^1 \rightarrow ①$$

Suppose $\exists b \in aH \cap a'H$

$$\text{So, } \exists c, c' \in H \text{ such that } ac = b = a'c'$$

① Continue : $a = (ac)\cdot \bar{c}^1$

$$= b \cdot \bar{c}^1$$

$$= (a'c') \cdot \bar{c}^1$$

$$\text{Hence } a = a' \cdot (c' \cdot \bar{c}^1).$$

$$\therefore a \in a'H. \text{ (Contradiction)}$$

ii) We have to prove $o(aH) = o(H)$.

Suppose $b, c \in aH$ with $b \neq c$, and

$$\leftarrow ab = ac.$$

Since, \bar{a}^1 exists in G ,

$$\bar{a}^1(ab) = \bar{a}^1(ac).$$

$$(\bar{a}^1a) \cdot b = (\bar{a}^1a) \cdot c.$$

$$\Rightarrow eb = ec \Rightarrow \boxed{b=c} \Rightarrow \text{(Contradiction)}$$

$$\therefore O(aH) = O(H).$$

$$\Rightarrow O(H) \mid O(G)$$

divides

Ex :- $(\mathbb{Z}_4, +)$.

$$\mathbb{Z}_4 = \{0, 1, 2, 3\}$$

$$H = \{0, 2\}.$$

$$\Rightarrow \mathbb{Z}_{12} = \{0, 1, 2, \dots, 11\}$$

$$H_1 = \{0, 2, 4, 6, 8, 10\}$$

$$H_2 = \{0, 4, 8\}.$$

* Quotient group :-

Let 'G' be a group (abelian), and $H \subseteq G$.

The quotient group of 'G' modulo 'H' denoted by.

$G \setminus H$, is the set of all cosets 'aH' with 'a' ranging over 'G', with the group operation $*$, defined by $(aH) * (bH) = a^b H$. and with the identity element being 'eH'.

* Cyclic groups :-

Let $(G, +)$, be a group. we say that $a \in G$ is a generator of 'G' if

$$G = \{a^n \mid n > 0, n \in \mathbb{Z}\}. \Rightarrow \text{denoted by.}$$

Ex :-

$$\mathbb{Z}_6 = \{0, 1, 2, 3, 4, 5\}, (\mathbb{Z}_6, +)$$

$$G = \langle a \rangle.$$

Algebraic properties

$$\Rightarrow O(0) = 1.$$

$$O(1) = 6.$$

$$1^1 = 1$$

$$1^2 = 1+1 = 2$$

$$1^3 = 1+1+1 = 3$$

$$1^4 = 4$$

$$1^5 = 5$$

$$1^6 = 6 \bmod 6 = 0$$

\Rightarrow cyclic.

Example: $\langle 0, 1, 2, 3, 4, 5 \rangle \Rightarrow$ generator.

Ex:

$$(z_6^*, \times) \Rightarrow z_6^* = \{1, 5\}$$

generator

$$5^1 = 5$$

$$5^2 = 5 \times 5 = 25 \bmod 6 = 1$$

\Rightarrow cyclic.

Ex: $S_n =$ set of all permutations

over n -symbols $\{1, 2, 3, \dots, n\}$.

$$S_3 = \{f_0, f_1, f_2, f_3, f_4, f_5\}$$

Not cyclic $\Rightarrow (0)(1, 2)$ is a generator

- ④ Any group is commutative then it is not a cyclic group.

Consider (G, \times) is a group with identity element e and multiplication operation \times .

Let $a, b \in G$ such that (a, b) is not a generator of G .

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* Discrete logarithm problem:- (DLP)

Let 'G' be a group. and $g \in G$. Finding smallest +ve integer 'k', for given $h \in \langle g \rangle \Rightarrow h = g^k$

$$K = \log_g(h) \Rightarrow (\text{hard problem for large base values})$$

invertible

* $\mathbb{Z}_n^* = \{a / a \in \mathbb{Z}_n \text{ and } \gcd(a, n) = 1\} \Rightarrow \text{all non-zero elements}$

$(\mathbb{Z}_n^*, \times_n)$ is group / abelian grp.

Suppose, $n = pq$, where 'p' & 'q' are primes.

\Downarrow
(integer factorisation \Rightarrow hard problem).

* Any element in \mathbb{Z}_n^* , where $n = pq$, has an order dividing $\text{LCM}(p-1, q-1)$.

* What is order of \mathbb{Z}_n^* ?

$$\phi(\mathbb{Z}_n^*) = \phi(n) \rightarrow \text{Euler quotient.}$$

$$n = pq \Rightarrow \phi(n) = (p-1)(q-1).$$

* Ring:-

→ Let 'R' be a non-empty set and binary operations are addition (+), multiplication (x).

Then. $(R, +, x)$. is called, a 'Ring' if

① $(R, +)$, is abelian grp.

② (R, x) is closure, & associative.

$$\left. \begin{array}{l} a \in \mathbb{Z}_n^* \\ p \rightarrow a^{p-1} \bmod p = 1 \\ q \rightarrow a^{q-1} \bmod q = 1 \end{array} \right\} \lambda = \text{lcm}(q-1, p-1).$$

$$\left. \begin{array}{l} p-1 \rightarrow \lambda \\ q-1 \rightarrow \lambda \end{array} \right\} \lambda$$

$$\left. \begin{array}{l} a \rightarrow a^{p-1} \bmod p = 1 \\ a \rightarrow a^{q-1} \bmod q = 1 \end{array} \right\}$$

③ Distributive laws. : $a \times (b+c) = a \times b + a \times c$.
 $(a+b) \times c = a \times c + b \times c$.

* Eg:-
 $S = \{0, \pm 2, \pm 4, \dots\}$.
 $(S, +, \times)$.

- ① $(S, +) \rightarrow$ it is a abelian grp.
 - ② $(S, \times) \rightarrow$ closure & associative.
 - ③ Distributive laws:
- commutative ?? (yes)
identity exists \Rightarrow w.r.t \times exists ?? (No)
 \downarrow
 $(S, +, \times)$ does not contain 1)
- Ring
with
commutativ
e.
(no identity)

Eg:-
Set of n -square matrices.

- i) Matrix addition. \Rightarrow abelian.
 - ii) Matrix multiplication \Rightarrow closure, associative.
- \Downarrow
- Ring with identity. (no commutative)

* Integral domain:-

A ring with the following properties:
A commutative ring. $(R, +, \times)$, with the following properties:

- i) Multiplicative identity: if $a, b \in R$, and $ab = 0$.
- ii) No zero divisors: if $a, b \in R$, then $a=0$, (or) $b=0$.

Eg:- $Z_6 = \{0, 1, 2, 3, 4, 5\}$, $+_6, \times_6$.
 $a=3, b=4$.
 $ab = 3 \times 4 = 12 \bmod 6 = 0 \Rightarrow 3, 4$ zero divisors.

Note:-

→ whenever an element is a zero divisor then inverse does not exist.

Eg:-

→ contd...

Here, '3' and '4' are not co-prime to '6'. (non-invertible elements).

• $\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$: is an integral domain.

here, all non-zero elements co-prime to '5'. So, inverse exists.

Eg:-

$(\mathbb{Z}, +, \times)$. → integral domain ??

★ Field:

We say $(F, +, \times)$ is a field if it is an integral domain with all non-zero elements are invertible.

Eg:- $(\mathbb{R}, +, \times)$ → field.

$(\mathbb{Z}_p, +_p, \times_p)$ where 'p' is prime. → field.

$(\mathbb{C}, +, \times)$ 'p' is not prime → ring.

Complex

Generator:

$(G, *)$ - group.

$g \in G, - G = \langle g \rangle$.

* Polynomial over \mathbb{R} :

$$f(x) = \sum_{i=0}^n a_i x^i ; a_i \in \mathbb{R}$$

$$g(x) = \sum_{i=0}^m b_i x^i ; b_i \in \mathbb{R}$$

Eg:- $f(x) = x^2 + 1$.

$$g(x) = x^3 + x + 2$$

$$f(x) + g(x) = x^3 + x^2 + x + 3$$

$$f(x) \cdot g(x) = ??$$

* Irreducible polynomial:

$f(x)$ over a field ' F ' is called irreducible, if there exists no root in ' F ' (or) we cannot write it as a product of lower degree polynomials.

' α ' is a root of $f(x)$.

$(x-\alpha)$ is a factor.

$$f(x) = (x-\alpha) \cdot g(x)$$

Eg:- $f(x) = x^2 + 1$ over $\mathbb{Z}_2 = \{0, 1\}$.

$$f(1) = 1^2 + 1 = 2 \bmod 2 \\ = 0$$

$$f(x) = (x+1) \cdot (x+1)$$

* G.F.:- (Galois field).

Set of all polynomials of degree $(n-1)$. (or)

less than or equal to $n-1$, over the field \mathbb{Z}_p , forms a field.

modulo p an irreducible polynomial $m(x)$ of

degree n .

Note:-

- Order of $G.F(P^n)$ is P^n .
- $G.F(P)$ is called prime field.
 - $G.F(2^n)$ is called binary field.

Eg:- $G.F(2^3)$.

$$O(G.F(2^3)) = 2^3 = 8$$

$$\mathbb{Z}_2 = \{0, 1\}$$

$$f(x) = a_2x^2 + a_1x + a_0$$

$$a_2, a_1, a_0$$

$$0 - 000$$

$$1 - 001$$

$$2 - 010$$

$$3 - 011$$

$$4 - 100$$

$$5 - 101$$

$$6 - 110$$

$$7 - 111$$

Eg:-

$$f = 2$$

$$f+g = ??$$

$$g = 5$$

$$0 \text{ } 10$$

$$\begin{array}{r} 101 \\ + 011 \\ \hline 111 \end{array} \Rightarrow 7$$

$$\begin{array}{r} 110 \\ + 011 \\ \hline 011 \end{array} \Rightarrow 2$$

$$f = 4$$

$$g = 6$$

$$f \cdot g = ??$$

$$4 \times 6 = 24 \bmod 8$$

$$= 0$$

Eg:- $f=2 \Rightarrow 010$
 $g=5 \Rightarrow 101$



~~$f = x$~~

$$g = x^2 + 1$$

$$f \cdot g = -x^3 + x \text{, mod } m(x) = ??$$

$$\Rightarrow 110 = \text{G}(x)$$

Eg:- $f=4 = 100 = x^2$ $\xrightarrow{\text{GF}(2^3) \text{ over } }$
 $g=6 = 110 = x^2 + x$. $m(x) = x^3 + x + 1$.

$$f \cdot g = x^4 + x^3 \text{, mod } m(x).$$

$$= (x^4 + x^3) \text{ mod } (x^3 + x + 1)$$

$$\begin{array}{r} x^3 + x + 1) \overline{) x^4 + x^3} \\ \cancel{x^4 + x^2 + x} \\ \hline x^3 + x^2 + x \\ \cancel{x^3 + x + 1} \\ \hline (x^2 + 1) \end{array}$$

$$\Rightarrow f \cdot g = x^2 + 1$$

$$= 5$$

* If $\gcd(f(x), g(x)) = 1$, then find inverse $(f(x))^{-1} = ?$

Eg:- AES used:
 $\text{GF}(2^8)$. over $m(x) = x^8 + x^4 + x^3 + x + 1$.

* Let F_{p^n} is a field, (there exists GF, F,)

→ $(F_{p^n})^*$ is a cyclic group, over multiplication.

→ A generator of $(F_{p^n})^*$ is called primitive root.