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- Residue Classes
- Finding Inverse Modulo m

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- General Caesar Cipher

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If m does not divide b-a, we say a and b are not congruent mod m, and write  $a \not\equiv b \pmod{m}$ 

•  $2 \not\equiv 7 \pmod{3}$ , because 3 does not divide 7 - 2 = 5

- $a = q_1 m + r_1$  and  $b = q_2 m + r_2$ , where  $0 \le r_1 \le m - 1$  and  $0 \le r_2 \le m - 1$ .
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Set of residue class  $\{0, 1, 2, \dots, m-1\}$  modulo m is denoted by  $\mathbb{Z}_m$ , that is,

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- Addition and Multiplication works exactly like real addition and multiplication, except reduce modulo m.
- $11 \times 13 = 143$  in  $\mathbb{Z}_{16}$ , and reduce it to modulo 16:  $143 = 8 \times 16 + 15$ , so  $143 \pmod{16} = 15 \in \mathbb{Z}_{16}$

## Definition (Inverse of an element)

Suppose  $a \in \mathbb{Z}_m$ . The multiplicative inverse of a is an element  $a^{-1} \in \mathbb{Z}_m$  such that  $aa^{-1} = a^{-1}a = 1 \pmod{m}$ 

# Finding Inverse Modulo m

Theorem (Multiplicative Inverse Modulo *m*)

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## Theorem (Multiplicative Inverse Modulo *m*)

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#### Proof.

Suppose a and m are relatively prime

Then, GCD(a, m) = 1

There exists x and y such that 1 = ax + my

Now apply modulo m, we get  $1 = (ax + 0) \pmod{m}$ 

That is,  $1 = ax \pmod{m}$ 

Means, there exists x such that  $ax = 1 \pmod{m}$ 

Therefore, x is inverse of a modulo m

Find multiplicative inverse of a = 8 modulo m = 11

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# Finding 8<sup>-1</sup> (mod 11) using Euclidean Algorithm

$$m = qa + r$$
$$a = q_1r + r_1$$

- $11 = (1) \times 8 + 3$
- $8 = (2) \times 3 + 2$
- $3 = (1) \times 2 + 1$
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- $3 = 11 (1) \times 8$
- $2 = 8 (2) \times 3$
- $1 = 3 (1) \times 2$

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Reverse the process:

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$$11 = (1) \times 8 + 3$$

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$$8 = (2) \times 3 + 2$$

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$$3 = (1) \times 2 + 1$$

• 
$$2 = (2) \times 1 + 0$$

• 
$$3 = 11 - (1) \times 8$$

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$$2 = 8 - (2) \times 3$$

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$$1 = 3 - (1) \times 2$$

Find multiplicative inverse of a = 8 modulo m = 11

# Finding $8^{-1} \pmod{11}$ using Euclidean Algorithm

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Reverse the process: find 
$$1 = 8x + 11y$$
 form

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$$2 = (2) \times 1 + 0$$

#### Rewrite

• 
$$3 = 11 - (1) \times 8$$

• 
$$2 = 8 - (2) \times 3$$

• 
$$1 = 3 - (1) \times 2$$

# Reverse the process: find 1 = 8x + 11y form

$$1 = 3 - (1) \times 2$$

$$= 3 - (1) \times [8 - (2) \times 3]$$

$$= (-1) \times 8 + (3) \times 3$$

$$= (1) \times 0 + (0) \times 0$$

$$= (-1) \times 8 + (3) \times [11 - (1) \times 8]$$
$$= (-4) \times 8 + (3) \times 11$$

$$x = -4 \pmod{11} = 7 = 8^{-1} \pmod{11}$$

#### Find Inverse of 7 modulo 26

#### Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

$$5 = (2) \times 2 + 1$$

$$2 = (2) \times 1 + 0$$

$$5 = 26 - (3) \times 7$$

$$2 = 7 - (1) \times 5$$

$$1 = 5 - (2) \times 2$$

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#### Rewrite

$$5 = 26 - (3) \times 7$$

$$2 = 7 - (1) \times 5$$

$$1 = 5 - (2) \times 2$$

## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

$$1 = (-2) \times 7 + (3) \times 5$$

$$1 = (-2) \times 7 + (3) \times [26 - (3) \times 7]$$

$$1 = (-11) \times 7 + (3) \times 26$$

$$x = -11 \pmod{26} = 15 = 7^{-1} \pmod{26}$$

#### Find Inverse of 7 modulo 26

#### Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

$$5 = (2) \times 2 + 1$$

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#### Rewrite

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$$2 = 7 - (1) \times 5$$

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## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

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## Finding Inverse

### Finding 5<sup>-1</sup> (mod 26) using Euclidean Algorithm

• 
$$26 = 5 \times 5 + 1$$

• 
$$5 = 5 \times 1 + 0$$

### • 1 = 5x + 26ywhere x = -5 and y = 1

#### Rewrite

• 
$$1 = 26 - 5 \times 5$$

## Finding Inverse

### Finding 5<sup>-1</sup> (mod 26) using Euclidean Algorithm

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$$26 = 5 \times 5 + 1$$

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$$1 = 26 - 5 \times 5$$

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$$1 = 5x + 26y$$
  
where  $x = -5$  and  $y = 1$ 

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$$1 = 5x \pmod{26}$$
, that is,  
 $x = 5^{-1} = -5 \pmod{26} = 21$ 

## Finding Inverse

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$$\bullet \ \mathcal{P} = \mathcal{C} = \mathcal{K} = Z_{26} = \{0, 1, 2, \dots, 25\}$$

- $\mathcal{P}=\mathcal{C}=\mathcal{K}=Z_{26}=\{0,1,2,\ldots,25\}$   $Z_{26}$  set of remainders when divide by 26
- Encryption function  $E_k : \mathcal{P} \to \mathcal{C}$  and decryption function  $D_k : \mathcal{C} \to \mathcal{P}$ , where  $k \in \mathcal{K}$ , defined as follows:

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$$C = E_k(m) = (m+k) \pmod{26}$$
  
 $m = D_k(C) = (C-k) \pmod{26}$ 

 Assign numerical value from 0 - 25 to each letter of plaintext alphabet a - z, respectively.

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Note that, if key k = 3, it is simply a Caesar cipher.



### Example

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- Assume that  $m = m_1 m_2 \dots m_n$ the plaintext message with n letters  $m_1$  to  $m_n$
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- Then  $m_1 = c$ ,  $m_2 = r$ ,  $m_3 = y$ ,  $m_4 = p$ ,  $m_5 = t$ ,  $m_6 = o$
- Suppose the corresponding ciphertext letters are C<sub>1</sub> to C<sub>n</sub>

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	I	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	1	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	6	17	18	1	9	20	21	22	23	24	25

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Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1:	9	20	21	22	23	24	25

$$m = \text{"crypto"}$$
 and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

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$$C_1 = E_k(m_1) = (2+3) \pmod{26}$$

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Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
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$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26}$$

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
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Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
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$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = F$$
  
 $C_2 = E_k(m_2) = (17+3) \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	I	m	n
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$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	I	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	C	1	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	6	17	18	1	9	20	21	22	23	24	25

$$m = "crypto" and C_i = E_k(m_i) = (m_i + k) \pmod{26}$$

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	1	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	I	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1	9	20	21	22	23	24	25

$$m =$$
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 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	1	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	16	17	18	1:	9	20	21	22	23	24	25

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 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
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 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	1	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
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$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
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 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26} = 22 = W$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26} = 22 = W$ 
 $C_6 = E_k(m_6) = (14+3) \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 1	16	17	18	1:	9	20	21	22	23	24	25

$$m =$$
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 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26} = 22 = W$ 
 $C_6 = E_k(m_6) = (14+3) \pmod{26} = 17 \pmod{26}$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	I	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1	9	20	21	22	23	24	25

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26} = 22 = W$ 
 $C_6 = E_k(m_6) = (14+3) \pmod{26} = 17 \pmod{26} = 17 = R$ 

plaintext (m)	а	b	С	d	е	f	g	h	i	j	k	Ι	m	n
Assigned No.	0	1	2	3	4	5	6	7	8	9	10	11	12	13
plaintext (m)	0	р	(	ŗ	r	S	t		u	٧	W	Х	У	Z
Assigned No.	14	15	5 -	16	17	18	1:	9	20	21	22	23	24	25

### Encryption algorithm works as follows:

$$m =$$
"crypto" and  $C_i = E_k(m_i) = (m_i + k) \pmod{26}$ 

$$C_1 = E_k(m_1) = (2+3) \pmod{26} = 5 \pmod{26} = 5 = F$$
 $C_2 = E_k(m_2) = (17+3) \pmod{26} = 20 \pmod{26} = 20 = U$ 
 $C_3 = E_k(m_3) = (24+3) \pmod{26} = 27 \pmod{26} = 1 = B$ 
 $C_4 = E_k(m_4) = (15+3) \pmod{26} = 18 \pmod{26} = 18 = S$ 
 $C_5 = E_k(m_5) = (19+3) \pmod{26} = 22 \pmod{26} = 22 = W$ 
 $C_6 = E_k(m_6) = (14+3) \pmod{26} = 17 \pmod{26} = 17 = R$ 

The ciphertext C is "FUBSWR", that is,  $E_3(crypto) = FUBSWR$ 

### **Affine Cipher**

where  $m, C \in \mathbb{Z}_{26}$ 

Let 
$$\mathcal{P} = \mathcal{C} = \mathbb{Z}_{26}$$
, and  $\mathcal{K} = \{(a,b) \in \mathbb{Z}_{26} \times \mathbb{Z}_{26} : GCD(a,26) = 1\}$  
$$C = E_k(m) = (am+b) \pmod{26}$$
 
$$m = D_k(C) = a^{-1}(C-b) \pmod{26}$$

# Affine Cipher

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## Correctness proof:

$$D_k(E_k(m)) = D_k(am+b) \pmod{26}$$

# Affine Cipher

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, and  $\mathcal{K} = \{(a,b) \in \mathbb{Z}_{26} \times \mathbb{Z}_{26} : GCD(a,26) = 1\}$  
$$C = E_k(m) = (am+b) \pmod{26}$$
  $m = D_k(C) = a^{-1}(C-b) \pmod{26}$ 

where  $m, C \in \mathbb{Z}_{26}$ 

## Correctness proof:

$$D_k(E_k(m)) = D_k(am+b) \pmod{26}$$
  
=  $a^{-1}((am+b)-b) \pmod{26}$ 

# Affine Cipher

Let 
$$\mathcal{P} = \mathcal{C} = \mathbb{Z}_{26}$$
, and  $\mathcal{K} = \{(a,b) \in \mathbb{Z}_{26} \times \mathbb{Z}_{26} : GCD(a,26) = 1\}$ 

$$C = E_k(m) = (am+b) \pmod{26}$$

$$m = D_k(C) = a^{-1}(C-b) \pmod{26}$$

where  $m, C \in \mathbb{Z}_{26}$ 

## Correctness proof:

$$D_{k}(E_{k}(m)) = D_{k}(am+b) \pmod{26}$$

$$= a^{-1}((am+b)-b) \pmod{26}$$

$$= a^{-1}(am) \pmod{26}$$

$$= (a^{-1}a)m \pmod{26}$$

$$= m \pmod{26}$$

$$= m$$

Suppose k = (7,3), then

• 
$$C = E_k(m) = 7m + 3$$
 (mod 26)

Suppose k = (7,3), then

•  $C = E_k(m) = 7m + 3$  (mod 26)

## Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

$$5 = (2) \times 2 + 1$$

$$2 = (2) \times 1 + 0$$

Suppose k = (7,3), then

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## Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

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## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

$$1 = (-2) \times 7 + (3) \times 5$$

$$1 = (-2) \times 7 + (3) \times [26 - (3) \times 7]$$

$$1 = (-11) \times 7 + (3) \times 26$$

$$x = -11 \pmod{26} = 15$$

Suppose k = (7,3), then

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$$26 = (3) \times 7 + 5$$

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$$1 = (-11) \times 7 + (3) \times 26$$

$$x = -11 \pmod{26} = 15$$

$$D_k(C) = 15(C-3) \pmod{26}$$

Suppose k = (7,3), then

• 
$$C = E_k(m) = 7m + 3$$
 (mod 26)

## Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

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## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

$$1 = (-2) \times 7 + (3) \times 5$$

$$1 = (-2) \times 7 + (3) \times [26 - (3) \times 7]$$

$$1 = (-11) \times 7 + (3) \times 26$$

$$x = -11 \pmod{26} = 15$$

$$D_k(C) = 15(C-3) \pmod{26}$$
  
=  $15([7m+3]-3) \pmod{26}$ 

Suppose k = (7,3), then

• 
$$C = E_k(m) = 7m + 3$$
 (mod 26)

## Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

$$5 = (2) \times 2 + 1$$

$$2 = (2) \times 1 + 0$$

## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

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$$1 = (-11) \times 7 + (3) \times 26$$

$$x = -11 \pmod{26} = 15$$

$$D_k(C) = 15(C-3) \pmod{26}$$
  
=  $15([7m+3]-3) \pmod{26}$   
=  $105m \pmod{26}$ 

Suppose k = (7,3), then

• 
$$C = E_k(m) = 7m + 3$$
 (mod 26)

## Remainder Form

$$26 = (3) \times 7 + 5$$

$$7 = (1) \times 5 + 2$$

$$5 = (2) \times 2 + 1$$

$$2 = (2) \times 1 + 0$$

## **Reverse Process**

$$1 = 5 - (2) \times 2$$

$$1 = 5 - (2) \times [7 - (1) \times 5]$$

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$$D_k(C) = 15(C-3) \pmod{26}$$
  
=  $15([7m+3]-3) \pmod{26}$   
=  $105m \pmod{26}$   
=  $m$ 

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

# Encryption | c r y p t o

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

Encryption							
plaintext <i>m</i>	с 2	r	у	р	t	0	

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

# Encryption plaintext c r y p t o m 2 17

- the plaintext message *m*: "*crypto*"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

Encryption							
plaintext m	c 2	r 17	у 24	р	t	0	

- the plaintext message *m*: "*crypto*"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

Encryption							
plaintext m	с 2	r 17	у 24	р 15	t	0	

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

Encryption						
plaintext	С	r	у	р	t	0
m	2	17	24	15	10	

- the plaintext message *m*: "*crypto*"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

Encryption							
plaintext	С	r	у	р	t	0	
m	2	17	24	15	19	14	

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

```
plaintext | c r y p t o m | 2 17 24 15 19 14 5 m + 2
```

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	у	р	t	0
m	2	17	24	15	19	14
5m + 2	12					

## Find the Affine cipher for given

- the plaintext message *m*: "*crypto*"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	у	р	t	0
m	2	17	24	15	19	14
5m + 2	12	87				

#### Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	У	р	t	0
m	2	17	24	15	19	14
5m + 2	12	87	122			

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	У	р	t	0
m	2	17	24	15	19	14
			122			

#### Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	У	р	t	0
m	2	17	24 122	15	19	14
5 <i>m</i> + 2	12	87	122	77	97	

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

plaintext	С	r	У	р	t	0
m	2	17	24	15	19	14
5m + 2	12	87	122	77	97	72

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

```
plaintext | c r y p t o m | 2 17 24 15 19 14 5m+2 | 12 87 122 77 97 72 (5m+2) \pmod{26}
```

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

```
plaintext c r y p t o m 2 17 24 15 19 14 5m+2 (5m+2) (mod 26) 12
```

## Find the Affine cipher for given

- the plaintext message m: "crypto"
- key k = (a, b) = (5, 2), then  $C = E_k(m) = 5m + 2 \pmod{26}$

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plaintext c r y p t o m 2 17 24 15 19 14 5m+2 (5m+2) (mod 26) 12 9
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5 <i>m</i> + 2			122			
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That is,  $E_K(crypto) = MJSZTU$ .

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Finding 5<sup>-1</sup> (mod 26) using Euclidean Algorithm

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•  $1 = 26 - 5 \times 5$ 

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$$26 = 5 \times 5 + 1$$

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cip	nertext	
C		

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ciphertext
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21(C-2)

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 21( $C$  – 2)

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ciphertext
C
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М	J	S	Z	Т	U
12	9	18	25	19	20
210	147	336	483	357	378

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cipnertext	IVI	J	S	Z	ı	U
C	12	9	18	25	19	20
21( <i>C</i> – 2)	210	147	336	483	357	378
$21(C-2) \pmod{26}$	2					

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cipnertext	IVI	J	5	_	ı	U
C	12	9	18	25	19	20
21( <i>C</i> – 2)	210	147	336	483	357	378
$21(C-2) \pmod{26}$	2	17				

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21( <i>C</i> – 2)	210	147	336	483	357	378
$21(C-2) \pmod{26}$	2	17	24	15	19	14
plaintext	С					

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Decryption	1
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ciphertext	M	J	S	Z	Т	U
С	12	9	18	25	19	20
21( <i>C</i> – 2)	210	147	336	483	357	378
$21(C-2) \pmod{26}$	2	17	24	15	19	14
plaintext	С	r	у	р	t	0

**Remark:** If a = 1, the Affine cipher becomes simply a Caesar cipher, that is,  $C = E_K(m) = x + b \pmod{26}$ .

# Thank You