

Data Structure



Design ...

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Seat Reservation Manager...

Companies :-

∞ Meta

1845. Seat Reservation Manager

Hint

Medium

704

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Design a system that manages the reservation state of n seats that are numbered from 1 to n .

Implement the `SeatManager` class:

- `SeatManager(int n)` Initializes a `SeatManager` object that will manage n seats numbered from 1 to n . All seats are initially available.
- `int reserve()` Fetches the **smallest-numbered** **unreserved** seat, reserves it, and **returns its number**.
- `void unreserve(int seatNumber)` Unreserves the seat with the given `seatNumber`.

Example 1:

Input

-1	-1	-1	
----	----	----	--

`["SeatManager", "reserve", "reserve", "unreserve", "reserve", "reserve", "reserve", "unreserve"]`

`[[5], [], [], [2], [], [], [], [], [5]]`

Output

`[null, 1, 2, null, 2, 3, 4, 5, null]`

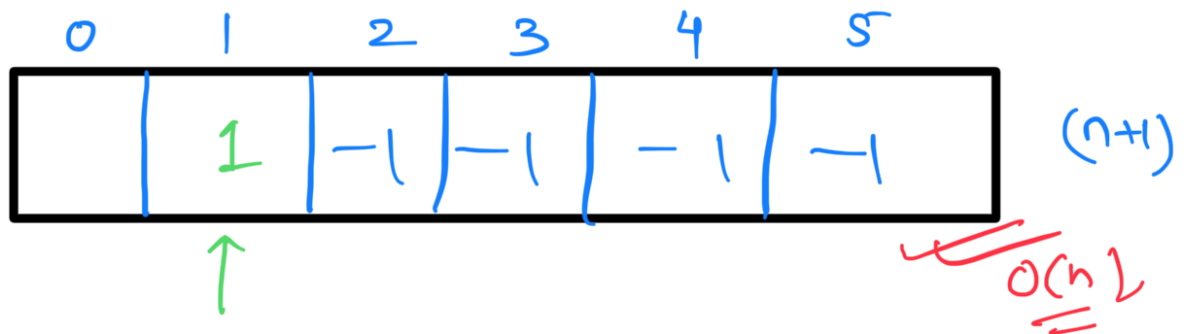
Brute Force :-

☆ We just have to keep track of seats. \rightarrow Unreserved = -1 \times

→ Reserved = 1 ✓

☆ Take an array to denote seats."

$n=5$ (1 2 ... n)



m times

```
reserve() {  
    // array traversal //  $O(n)$  ←  
    // pick smallest seat with -1 .  
}
```

$O(m \times n)$


TLE

```
unreserve(seat) {  
    arr[seat] = -1 ;  $O(1)$   
}
```

Optimal Approach

Intuition → Why → where was the problem?

\swarrow m times
 reserve() { $O(m \times n)$
 // array traversal $\leftarrow O(n)$
 // pick smallest seat with -1 .
 }


 "We need a Data Structure which can find me the smallest seat in better time complexity".
 ⇒ "Min heap"

⇒ ① Pq → min-heap
 $m \log n$ \rightarrow for (i = 1; i ≤ n; i++) \leftarrow
 Pq.Push(i) $\leftarrow \log(n)$
 $O(n \log n)$

② \downarrow m
 reserve()

1
2
3
4
5

Pq.



$O(m \times \log(n))$

reserve \rightarrow smallest
 $seat = \underline{Pq.top()}; // 1 \rightarrow \underline{O(1)}$
 $\underline{Pq.pop()}; \leftarrow O(\log(n))$
 $next\ seat; // 1$

(3) $\downarrow m$
unreserve (seat)

$Pq.push(seat); \underline{O(\log(n))}$

Improvement :-

Without Pre-filling the Pq.

Example 1:

Input

["SeatManager", "reserve", "reserve", "unreserve", "reserve",
"reserve", "reserve", "reserve", "unreserve"]

[5], [], [], [2], [], [], [], [], [5]]

Output

[null, 1, 2, null, 2, 3, 4, 5, null]

1 2 3 4 5

Seat-marker = 1 2 3

// SeatManager (int n) {

Seat-marker = 1; // O(1)

}

// int reserve () {

{ if (!pq.empty()) {
Seat = pq.top(); pq.pop(); $\leftarrow \underline{\underline{\log(n)}}$
return Seat;

}

Seat = seat-marker; // 3
seat-marker++; // 4 $\leftarrow \int \underline{\underline{O(1)}}$
return Seat;

void unreserve (int seatNumber) {

pq.push (seatNumber); // $\log(n)$.

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