





Median Of Two Sorted Arrays (PART-1)

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4. Median of Two Sorted Arrays

Given two sorted arrays nums1 and nums2 of size m and n respectively, return the median of the two sorted arrays.

The overall run time complexity should be $O(\log (m+n))$.

(Pent-2)

Example 1:

Input: nums1 = [1,3], nums2 = [2]

Output: 2.00000

Explanation: merged array = [1,2,3] and median is 2.

Example 2:

Input: nums1 = [1,2], nums2 = [3,4]

Output: 2.50000

Explanation: merged array = [1,2,3,4] and median is (2 + 3) / 2 = 2.5.

(Y)edian :- (middle number)

$$\frac{1}{3}, \frac{3}{5}, \frac{4}{5}, \frac{8}{5}$$

$$\frac{3+5}{2} = \frac{9}{5}$$

$$\frac{3+5}{2} = \frac{3}{5}$$

$$\frac{5}{5}, \frac{4}{5}, \frac{8}{5}$$

$$\frac{7}{5}, \frac{7}{5}, \frac{7}{5}, \frac{7}{5}$$

$$\frac{7}{5}, \frac{7}{5}, \frac{7}{5}, \frac{8}{5}$$

$$\frac{7}{5}, \frac{7}{5}, \frac{7$$

Brufe Force :-

nums =
$$\begin{bmatrix} 1 & 3 \end{bmatrix}^{i} \rightarrow m$$

nums = $\begin{bmatrix} 2 \end{bmatrix}^{i} \rightarrow m$

nums = $\begin{bmatrix} 2 \end{bmatrix}^{i} \rightarrow m$

nums = $\begin{bmatrix} 2 \end{bmatrix}^{i} \rightarrow m$

nums = $\begin{bmatrix} 1 & 2 \end{bmatrix}^{i}$

nums = $\begin{bmatrix} 1 & 3 \end{bmatrix}^{i}$

nums = $\begin{bmatrix} 1 & 3 \end{bmatrix}^{i}$

nums = $\begin{bmatrix} 1 & 3 \end{bmatrix}^{i}$

Nums 2 =
$$\begin{bmatrix} 8, 9, 10, 11 \end{bmatrix}$$

tem $p = \begin{bmatrix} 3 & 7 & 8 & 9 & 10 & 11 \end{bmatrix}$
Size = $m + m = 7$
 $\frac{1}{2} \left(6ize \frac{1}{2} - \frac{1}{2} \right) = 1$
Hohn tem $p \left(5ize/2 \right)$;

Better Approach:

$$n_{\text{UML}} = \begin{cases} 1, & 3, & 7 \\ 8, & 9, & 10, & 11,12 \end{cases}, n = 3$$

$$n_{\text{UML}} = \begin{cases} 8, & 9, & 10, & 11,12 \end{cases}, n = 5$$

$$\frac{|a_{\text{UML}}|}{|a_{\text{UML}}|} = \frac{3}{3} \frac{|s_{\text{SIZ}}|}{|s_{\text{SIZ}}|}$$

$$|c_{\text{UML}}| = \frac{3}{3} \frac{|s_{\text{SIZ}}|}{|s_{\text{SI$$

```
(TPProach-1
 i] (nums,[i] < num2[j]) }
                                  i] (numsi[i] < nums2 [j]) {
                                        i) (K = = idx1)
     temp(K++) = noms([i++]);
                                           elevet = nums [i];
 ) حلام لا
                                        \sqrt{(K==id_{N2})}
     temp[K++] = nums 2 [j++];
                                           elun2 = numsi(i);
  4
                                  Jelse of
                                      i](K = = idx1)
                                        elmefl
                                               = nome(i);
while (i<m) }
                                     ief(K==idre)
                                          elu2 = mor []/
    ten(K++)=m1(i++);
                                 K++;
while (j(n)
   tr(c~) = we(j+4);
                               while ( i < m) of
                                      if (K == idx1)
                                         Ch! = noman(i)
                                      il (K=idny
                                         duz = nuli)
                                      K++;
                               while (j<n)
```

Numis(j);