

CPSC-354 Report

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Abstract

This report contains assignments throughout the fall 2024 semester and is intended for the purpose of documenting my work and showing my progress in CPSC 354 - Programming Languages, taught by Jonathan Weinberg.

Contents

1 Introduction

This report, prepared for CPSC 354 - Programming Languages at Chapman University, is a comprehensive account of my academic voyage over the semester. It includes a detailed compilation of my notes, homework solutions, and critical reflections on the coursework. This report serves as a bridge between the theoretical knowledge imparted in lectures and the practical skills essential for future pursuits in both graduate studies and the software industry.

2 Week by Week

2.1 Week 1

Notes

In week 1 we learnt about Lean as a programming language and its correlation to discrete math. We also learnt about other proof assistants. We then shifted our focus to the NNG tutorial world as you can see below.

Homework

Tutorial world

Level 5 / 8 : Adding zero

Active Goal

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$a + (b + 0) + (c + 0) = a + b + c$

rw [add_zero]

Active Goal

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$a + b + (c + 0) = a + b + c$

rw [add_zero]

Active Goal

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$a + b + c = a + b + c$

rfl

level completed! 🏆

Level 5:

Discrete math's lemmas tell us that anything added to 0 will give the result of that number itself. So; $A+0=A$. Using this we can bring the left hand side down to $a+b+c$. From here we can use the property of reflexivity to show that both sides are equal, hence solving the puzzle.

Level 6 / 8 : Precision rewriting

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$$a + (b + 0) + (c + 0) = a + b + c$$

```
rw [add_zero c]
```

Active Goal

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$$a + (b + 0) + c = a + b + c$$

```
rw [add_zero b]
```

Active Goal

Objects:

$a\ b\ c : \mathbb{N}$

Goal:

$$a + b + c = a + b + c$$

```
rfl
```

level completed! 🏆

Level 6:

Level 7 / 8 : add_succ

Theorem `succ_eq_add_one`: For all natural numbers a , we have $\text{succ}(a) = a + 1$.

Active Goal

Objects:
 $n : \mathbb{N}$

Goal:
 $\text{succ } n = n + 1$

`rw [one_eq_succ_zero]`

Active Goal

Objects:
 $n : \mathbb{N}$

Goal:
 $\text{succ } n = n + \text{succ } 0$

`rw [add_succ]`

Active Goal

Objects:
 $n : \mathbb{N}$

Goal:
 $\text{succ } n = \text{succ } (n + 0)$

`rw [add_zero]`

Level 7:

Level 8 / 8 : 2+2=4

$2 + 2 = 4.$
example : $(2 : \mathbb{N}) + 2 = 4 := \text{by}$

```

2 rw[three_eq_succ_two]
3 rw[two_eq_succ_one]
4 rw[one_eq_succ_zero]
5 rw[succ_eq_add_one]
6 rw[one_eq_succ_zero]
7 rw[add_succ]
8 rw[add_zero]
9 rw[add_succ]
10 rw[add_succ]
11 rw[add_zero]
12 rfl

```

Level completed! 🎉

No Goals

Level 8:

2.2 Week 2

Notes

In week 2 we learnt about recursion and its application in other problems such as the Towers of Hanoi game we played. We also learnt about its various benefits such as breaking down complexity of problems and being more concise.

Homework

Addition world

Level 1 / 5 : zero_add

Theorem zero_add: For all natural numbers n , we have $0 + n = n$.

```

theorem zero_add (n : ℕ) : 0 + n = n := by
  1 induction n with d hd
  2 rw[add_zero]
  3 rfl
  4 rw[add_succ]
  5 rw[hd]
  6 rfl
  7 |

```

Level 1:

Level completed! 🎉

Level 2 / 5 : succ_add

Theorem succ_add: For all natural numbers a, b , we have $\text{succ}(a) + b = \text{succ}(a + b)$.

```

theorem succ_add (a b : ℕ) : succ a + b = succ (a + b) := by
  1 induction b with d hd
  2 rw[add_zero]
  3 rw[add_zero]
  4 rfl
  5 rw[add_succ, add_succ]
  6 rw[hd]
  7 rfl
  8 |

```

Level 2:

Level completed! 🎉

Level 3 / 5 : add_comm (level boss)

Theorem add_comm: On the set of natural numbers, addition is commutative. In other words, if a and b are arbitrary natural numbers, then $a + b = b + a$.

```

theorem add_comm (a b : ℕ) : a + b = b + a := by
  1 induction b with d hd
  2 rw[add_zero, zero_add]
  3 rfl
  4 rw[add_succ, succ_add, hd]
  5 rfl
  6 |

```

Level 3:

Level completed! 🎉

Level 4 / 5 : add_assoc (associativity of addition)

Theorem add_assoc: On the set of natural numbers, addition is associative. In other words, if a, b and c are arbitrary natural numbers, we have $(a + b) + c = a + (b + c)$.

```

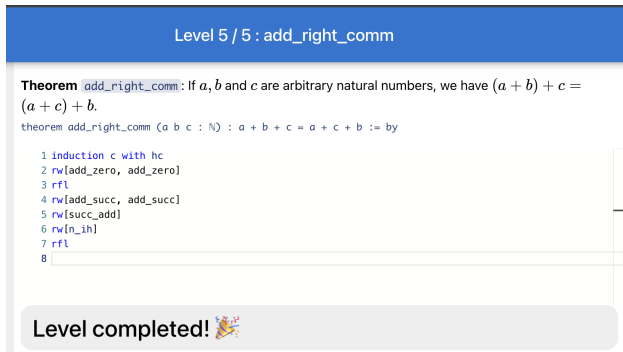
theorem add_assoc (a b c : ℕ) : a + b + c = a + (b + c) := by
  1 induction c with c hc
  2 rw[add_zero, add_zero]
  3 rfl
  4 rw[add_succ, add_succ]
  5 rw[add_succ]
  6 rw[hc]
  7 rfl
  8 |

```

Level 4:

Level completed! 🎉

Using induction on c we can initially create an easier medium to use reflexivity to solve for $a+b$. Then solving the other side we just use the mathematical definition of a successor function, until we can use the induction again to get the equation to the point where we can use reflexivity to prove it. This is a clear example of mathematical proofs by induction.



Level 5:

Once again this is proof by mathematical induction similar to the previous one. This time we add the zeroes and then use reflexivity for the first part of the proof. Then to prove the second part we use the successor function until bringing back the induction we used like the previous question. Then to complete the proof we use reflexivity again.

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Discord Question: Since recursion has so many benefits and also breaks down the complexity of problems, why aren't we taught to use it as our primary method? In other words, why isn't it the first method of problem solving we're taught?

2.3 Week 3

Notes

In week 3 we spoke about recursion further, and focused on our calculators in python. Eventually connecting the dots for our first Assignment that was due at the same time as Homework 4, where I used recursion in my python calculator in order to get it to function as efficiently as possible. We also discussed parsing, and derivation trees which is what we practiced in Homework 4.

Homework

For homework 4 we did some practice on derivation trees for the strings you can see in the handwritten work below, along with the respective answers;

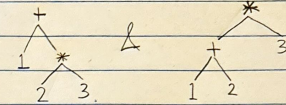
CPSC 354 - Programming Languages
Homework-4

Q. Write out the derivation trees for the following strings:

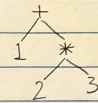
a) $2+1$



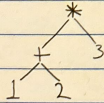
b) $1+2\cdot3$



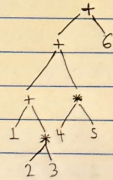
c) $1+(2\cdot3)$



d) $(1+2)\cdot3$



e) $1+2\cdot3+4\cdot5+6$



Discord Question: Why do programming languages need different types of parsers, and how does this choice impact the way a computer understands the code?

2.4 Week 4

Notes

In week 4 we spoke about parsing and trees further which is what we practiced in last weeks Homework. But now we delved into the more notation heavy side of it. This weeks homework contained the lean logic game that helped us put some of that into practice.

Homework

For homework 5 we completed the Lean logic game tutorial world. I have provided the answers to the same below.

1. Exhibit evidence that you're planning a party.

Level 1:

```
example (P : Prop) (todo_list : P) : P := by
```

Solution:

```
exact todo_list
```

Level 2:

```
example (P S : Prop) (p : P) (s : S) : P S := by
```

Solution:

```
exact and.intro p s
```

Level 3:

```
example (A I O U : Prop) (a : A) (i : I) (o : O) (u : U) : (A I) O U := by
```

Solution:

```
have ou := and_intro o u
have ai := and_intro a i
exact and_intro ai ou
```

Level 4:

```
example (P S : Prop) (vm: P S) : P := by
```

Solution:

```
exact vm.left
```

Level 5:

```
example (P Q : Prop) (h: P Q) : Q := by
```

Solution:

```
exact h.right
```

Level 6:

```
example (A I O U : Prop) (h1 : A I) (h2 : O U) : A U := by
```

Solution:

```
exact and_intro h1.left h2.right
```

Level 7:

```
example (C L : Prop) (h: (L ((L C) L) L L L)) (L L) L : C := by
```

Solution:


```

have a:= h.left
have b:= a.right
have c:= b.left
have d:= c.left
exact d.right

```

Level 8:

```

example (A C I O P S U : Prop)(h: ((P S) A) ¬I (C ¬O) ¬U) : A C P S := by

```

Solution:

```

have a:=h.left
have fin:=a.left
have a:=a.right
have r:=h.right
have t:=r.right
have k:= t.left
have c:=k.left
have e:= and_intro c fin
exact and_intro a e

```

Explanation for Solution 8:

We ultimately need $A \ C \ P \ S$, which means that we need to retrieve A , C , P , and S from the original expression.

- (1) Seeing that A , P and S are on the left side of $h \rightarrow$ extract $h.left$ to get $(P \ S) \ A$.
- (2) Seeing that $(P \ S)$ is already clearly made on the left side we can store $(P \ S)$ together as 'f'.
- (3) Now from the same expression we can extract the right side to retrieve and store A as 'a'.
- (4) Now that we have A , P and S stored and easily accessible, we need C from the right side of h .
- (5) After storing $h.right$ in another expression, 'r', we take the right side of r to further narrow down our search.
- (6) Now we take the left side of that expression for the same purpose, store it as k .
- (7) Finally we can retrieve C by storing the $k.left$ as 'c'.
- (8) Now to finally present our solution we can start by using `and_intro` on c and fin , storing the result as e .
- (9) To complete our solution we use `exact and_intro` on a and e .

Discord Question: How does the choice of a parsing technique (e.g., top-down vs. bottom-up) impact the efficiency and clarity of the resulting derivation tree? In other words, when should I choose between the two approaches so that my chosen approach is significantly more beneficial than the other?

3 Lessons from the Assignments

4 Conclusion

References

[BLA] Author, [Title](#), Publisher, Year.