lecture 9 :- Resolution principle

literal :- A variable or its negation p q
¬p ¬q.

clause :- A disjunction of literals

¬p ∨ ¬q
p ∨ q ∨ r

¬p ∨ q

Step 1 :- For each premise find out corresponding clauses.

P1 :- p ∧ q c1 :- p
c2 :- q

P2 :- p → q c1 :- ¬p ∨ q

Step 2 :- Find the negation of conclusion. and determine clause.

Step 3 :- Repeatedly Apply PR.

p ∨ q
¬q ∨ r
∴ p ∨ r.

P1 p
P2 p → Q
C ∴ Q.

c1 : p ✓
c2 :- ¬P ∨ Q. ✓
c3 :- ¬Q. ✓
c4 :- Q ✓ from c1 and c2.
c5 :- □ from c3 and c4.

Ex 11 P65

P1 T → (M ∨ E)
P2 S → ¬E
P3 T ∧ S
C. ∴ M

c1 :- ¬T ∨ M ∨ E ✓
c2 :- ¬S ∨ ¬E ✓

C1:- ¬T ∨ M ∨ E ✓
C2:- ¬S ∨ ¬E ✓
C3:- T ✓
C4:- S ✓
C5:- ¬M ✓
C6:- ¬T ∨ M ∨ ¬S ✓ from C1 and C2.
C7:- M ∨ ¬S ✓ " C3 " C6.
C8:- M ✓ " C4 " C7.
C9:- □ " C5 " C8.

Argument is valid.

Ex7:- P62:-

P1 $p \to q$
P2 $\neg p \to r$
P3 $r \to s$
C $\therefore \neg q \to s$

C1 $\neg p \vee q$ ✓
C2 $p \vee r$ ✓
C3 $\neg r \vee s$ ✓
C4 $\neg q$ ✓
C5 $\neg s$. ✓
C6 $q \vee r$ ✓ from C1, C2.
C7. $q \vee s$. ✓ from C3, C6.
C8 s ✓ " C4, C7.
C9 □ " C5, C8.

$\neg(\neg q \to s)$
$\equiv \neg(q \vee s)$
$= \neg q \wedge \neg s$.

Argument valid.

HW.
P70 – 72.
Exercise 1–30

Quiz #5

P1 $L \to A$
P2 $E \to \neg I$
P3 $A \to E$
C. $\therefore L \to \neg I$.

Prove or disproof.

RELATIONS

→ Set :- A collection of distinct objects.

→ SET : A Collection of distinct objects.

(ظاہری شکل) Syntax. { }. جامع - 1
مانع - 2

Semantics Repeation Not allowed.

$A \times B$ $A = \{1,2,3\}$. $B = \{a,b\}$.

$A \times B = \{(1,a), (1,b), (2,a), (2,b), (3,a), (3,b)\}$.

$B \times A = \{(a,1), (a,2), (a,3), (b,1), (b,2), (b,3)\}$.

Cardinality of a Set $|A| = 3$
" " a " $|B| = 2$

$|B \times A| = |A \times B| = |A| \times |B| = 3 \times 2 = 6$.

Subset $A \subseteq B$.

$\{1,2,3,4\} = \{2,3,4,1\}$.

Power Set (A) = All Subsets of A.

$A = \{1,2,3\}$.

$PS(A) = \{\phi, \{1\}, \{2\}, \{3\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\}\}$.

$|PS(A)| = 2^{|A|} = 2^3 = 8$.

$PS(A \times B) = 2^{|A \times B|} = 2^{3 \times 2} = 2^6 = 64$.

$= \{\phi, \{(1,a)\}, \{(1,b)\} \cdots\}$.