Fast & Sharing-aware Live Migration of Virtual Machines

Roja Eswaran, PhD

Agenda

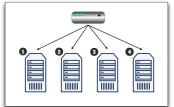
- Motivation
- Problem Statement
- Background
- Inter-host
 - Template-aware Live Migration of Virtual Machines (TLM)
 - Sharing-aware Live Migration of Virtual Machines (SLM)
- Intra-host TLM
- Future Research Directions
- Conclusion



Motivation

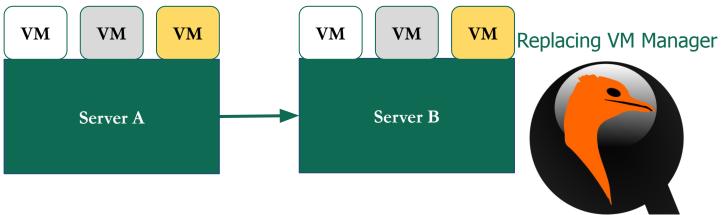
Hardware Failure



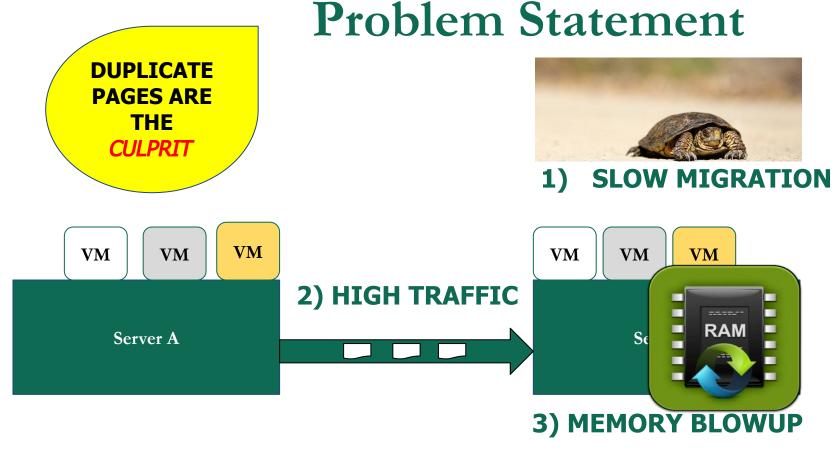








Live Migration is crucial technique for moving VMs around the datacenter



em Statement **Current Live Migration is** unaware of the pre-existing deduplication optimization (Pre-copy) (Pre-copy) 2500 Source(Post-copy) Destination(Post-copy) Usage(MB) 2000 1500 Memory 1000 500 Increase in migration time Numb and network traffic Fig: Memory Usage of V

SOLUTION Cognizant of Host/Hypervisor Page Sharing Optimization! **Fast and Reliable** VM VM**VM** VM**VM** $\mathbf{V}\mathbf{M}$ 2) High **Low Traffic** Server A Server B

3) Efficient memory utilization (BONUS!)

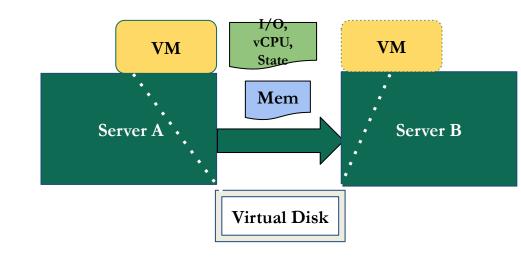
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Inter-host Live Migration

Moving VMs between two different hosts

 Useful for load-balancing, maintenance, and hardware failures



Pre-copy Live Migration

- Pre-copy
 - Transfer *memory first* and then transfer states

Downtime

- Transfer remaining memory and state
- Resume VM at destination and
 stop VM at source

Clark, C., Fraser, K., Hand, S., Hansen, J.G., Jul, E., Limpach, C., Pratt, I. and Warfield, A., 2005, May. Live migration of virtual machines.

In Proceedings of conference on Symposium on Networked Systems Design & Implementation.

Mem VM I/O, vCPU, State Source Destination Virtual Disk **First** Round: Second Round: nth Round: **Transfers Transfer Pages All Dirty** Transfer dirtied in the Pages dirtied **Pages** first round in the n-1th round

Background(2/6)

Post-copy Live Migration

Post-copy

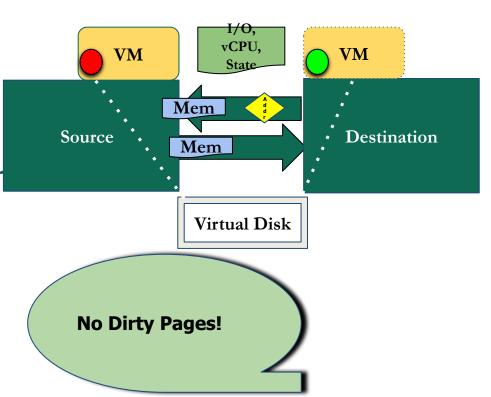
Transfer states first
 resume at destination
 and then transfer
 memory

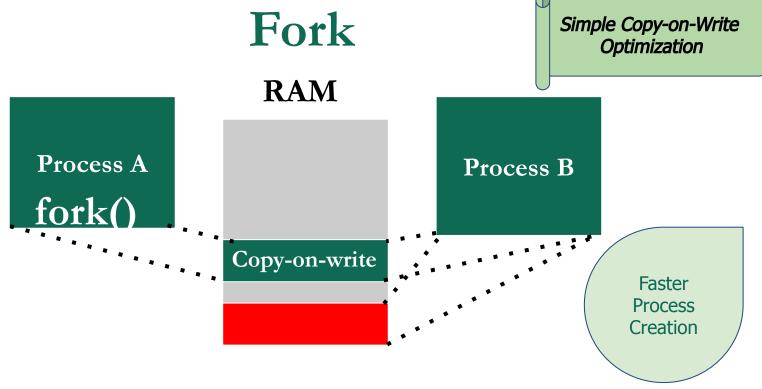
Two mechanism for memory transfer

- Demand Paging
- Active Pushing

Hines, M.R., Deshpande, U. and Gopalan, K., 2009. *Post-copy live migration of virtual machines.* ACM SIGOPS operating systems review.

Background(3/6)



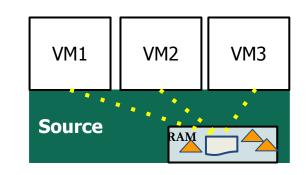


Fork() creates a child process with *copy-on-write* sharing with the parent memory pages

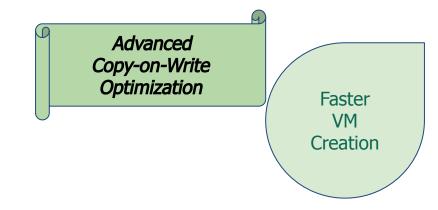
Background(4/6)

VM Templating

 VM templating is a technique to quickly instantiate multiple lightweight VMs from a shared copy-on-write (COW) template



 The additional pages dirtied by the VMs are referred to as *delta*



Each virtual page has its physical page

Kernel Samepage Merging (KSM)

Fig:1 Without KSM

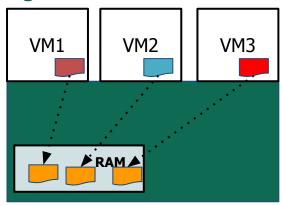
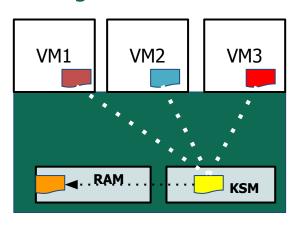


Fig:2 With KSM



KSM merges duplicate pages and create a single physical page with

copy-on-write mapping

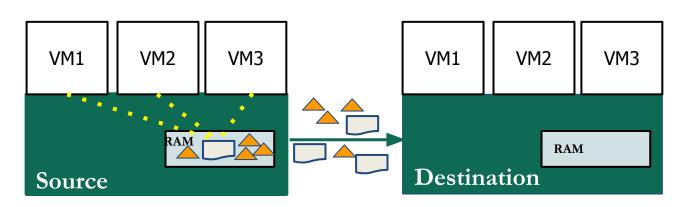
Arcangeli, A., Eidus, I. and Wright, C., 2009, July. *Increasing memory density by using KSM.*In Proceedings of the linux symposium.

Advanced Copy-on-Write Optimization

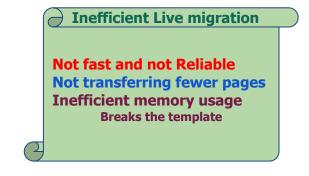
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 - Template-aware Live Migration of Virtual Machines (TLM)
 [Edgecomm: SEC'23]

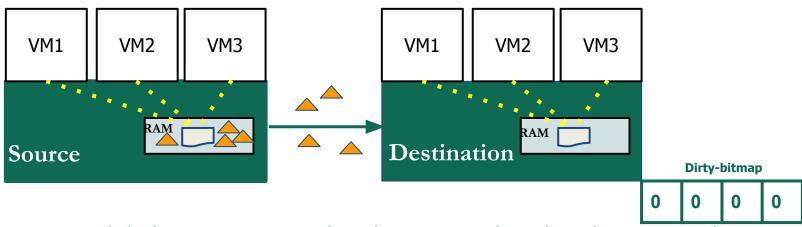
Key Insight: TLM



- Current live migration lacks the awareness of the VM Template
 - Transfers delta and the base template
 - Breaks the template
 - Increases network-traffic
 - Increases Total Migration Time



Design: TLM



- 1. At source, while booting up templated VMs, initialize dirty-bitmap tracking
- 2. Copy the template file to destination before migration
- 3. Transfer Delta by looking up the dirty-bitmap

Dirty-bitmap

Eswaran, R., Yan, M. and Gopalan, K., 2023, December. Template-Aware Live Migration of Virtual Machines. In 2023 IEEE/ACM Symposium on Edge Computing (SEC) (pp. 336-340). IEEE Computer Society.

Evaluation Setup

CPU: Intel Xeon E5-2620 v2 processors.

Mem: 128GB DRAM

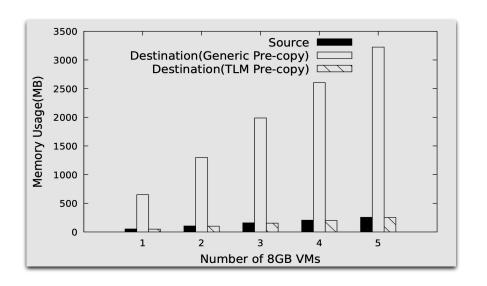
Hypervisor: QEMU/KVM

Host and Guest Kernel: linux-4.10.1

Live Migration technique: Pre-copy

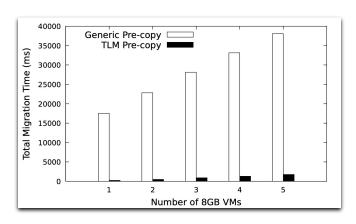
Generic - Traditional Pre-copy

Evaluation: Memory Footprint

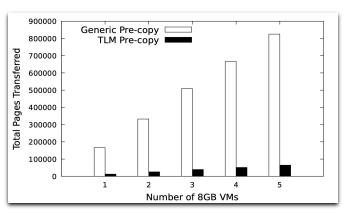


The source and destination have the same memory usage retaining the templating benefits with the help of TLM

Evaluation: Total Migration Time & Total Pages Transferred



96% reduction in total migration time



93% reduction in total pages transferred

Summary

- Besides reducing the memory footprint, TLM
 - Improves the total migration time and network-traffic

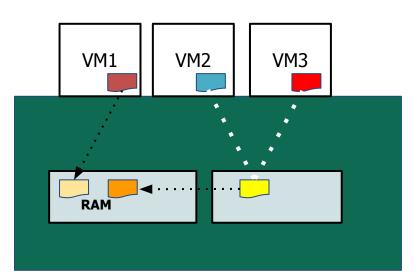
- Limitations:
 - Works only for Templated VMs and unaware of KSM or Fork

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 [Edgecomm: SEC'23]
 - Sharing-aware Live Migration of Virtual Machines (SLM)
 [CCGRID'24]



Key Insight: SLM



Page	Virtual Page Number (VPN)	Physical Frame Number (PFN)	
	0x11	0xAA	
	0x22	0xAA	
	0x33	0xBB	

Merged pages have the *same physical frame numbers*

Eswaran, R., Yan, M. and Gopalan, K., Tackling memory footprint expansion during live migration of virtual machines," in Proc. of International Symposium on Cluster, Cloud and Grid Computing (CCGrid), 2024

Existing Solution: Content-based Hashing

 Compute the hash of the page content to find the duplicate pages thereby avoiding their retransmission

Deshpande, U., Wang, X. and Gopalan, K., 2011, June.

Live gang migration of virtual machines.

In Proceedings of international symposium on High performance distributed computing.

. Compute-intensive

. Unaware of existing COW Pages

Zhang, X., Huo, Z., Ma, J. and Meng, D., 2010, September.

Exploiting data deduplication to accelerate live virtual machine migration.

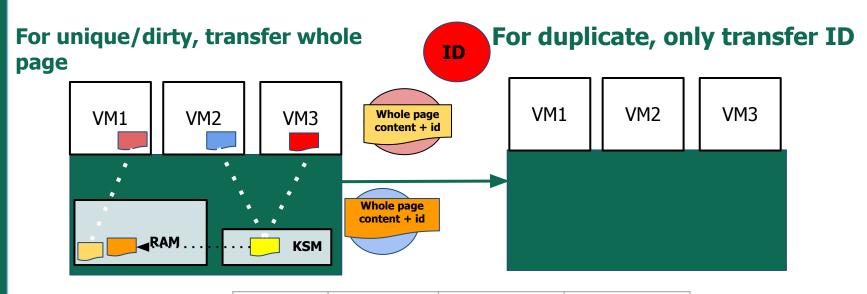
In IEEE international conference on cluster computing.

Chiang, J.H., Li, H.L. and Chiueh, T.C., 2013.

Introspection-based memory de-duplication and migration.

ACM SIGPLAN Notices

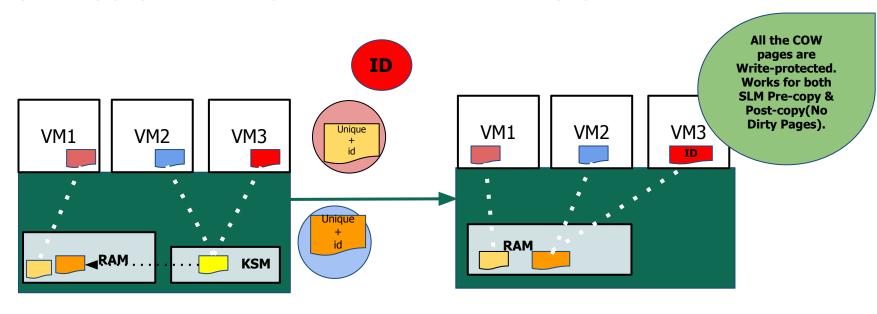
SLM Source



Page	Virtual Page Number (VPN)	Physical Frame Number (PFN)	PageType
-	0x11	0xAA	Unique
	0x22	0xBB	Unique
	0x33	0xBB	Duplicate

SLM Destination

Unique/Dirty pages have unique RAM offsets to write the page content at the destination



Duplicate pages COW map into the unique page content using the identifier

Evaluation Setup

CPU: Intel Xeon E5-2620 v2 processors

Mem: 128GB DRAM

Hypervisor: QEMU/KVM

Host and Guest Kernel: linux-4.10.1

Live Migration technique: Pre-copy, Post-copy

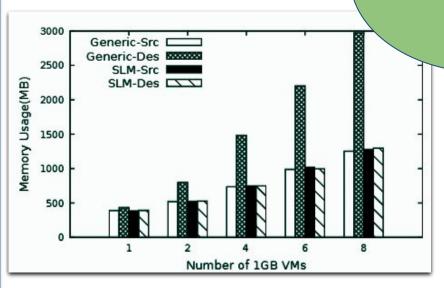
Generic - Traditional Pre/Post copy

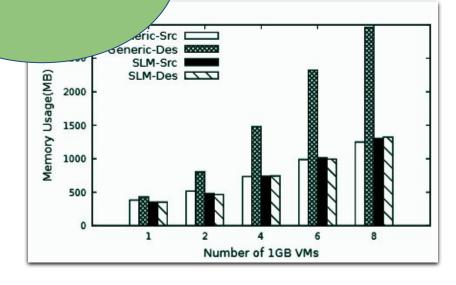
Evaluation: Memory Footprint



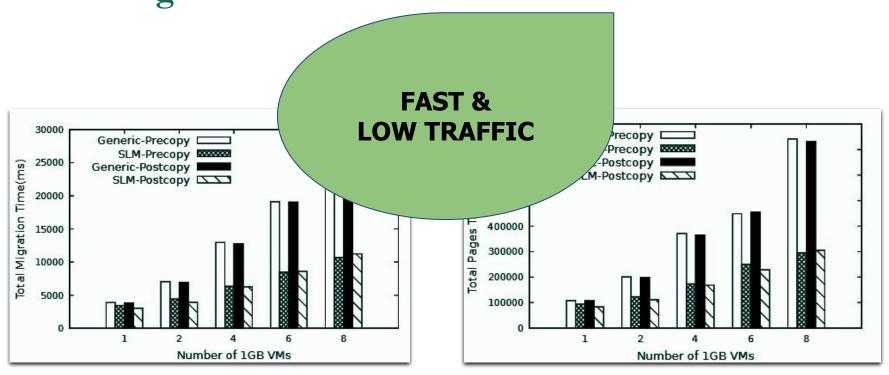
NO MEMORY BLOW-UP

: Post-copy





Evaluation: Total Migration Time & Total Pages Transferred



59% reduction in total migration time

62% reduction in total pages transferred

Summary

- SLM retains the *COW shared pages irrespective of mechanisms*
 - Fork, KSM and others
 - Works for both pre-copy and post-copy live migrations

- SLM works for both regular non-templated and templated VMs
 - besides reducing the total migration time and total pages transferred

- Both Iperf and Redis shows that SLM doesn't introduce any additional performance overhead
 - Infact, compared to the baseline it finishes the migration faster due to less page transfer.

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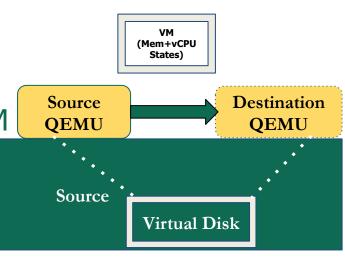
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- Intra-host TLM [TPDS'24 under-review]

Intra-host Live Migration

 Moving VMs from one QEMU to another within the same host

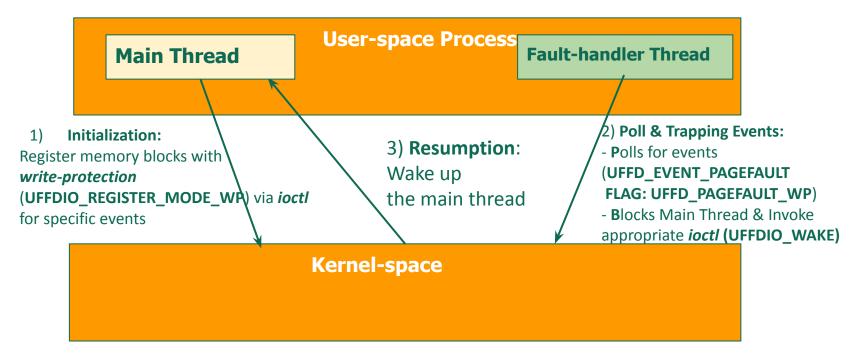
 Useful for live replacement of VM Manager process (QEMU)

- New Features and Functionalities
- Compatibility updates
- Performance improvements



Userfaultfd

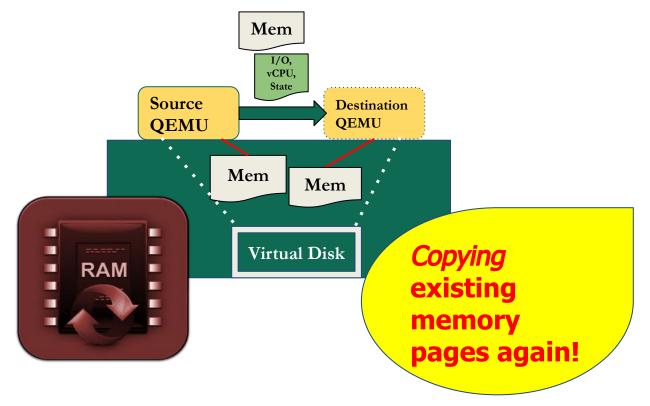
• Userfaultfd mechanism lets the user-space process (**via fault-handler thread**) handles memory-related events especially pagefault without **kernel intervention**



BACKGROUND (2/2)

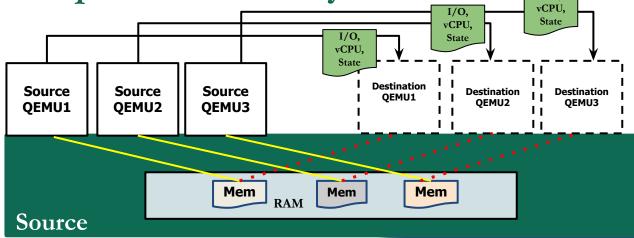


Problem 1: Copying Existing Memory Pages



B

Problem 2: Duplicated Memory Backend-files

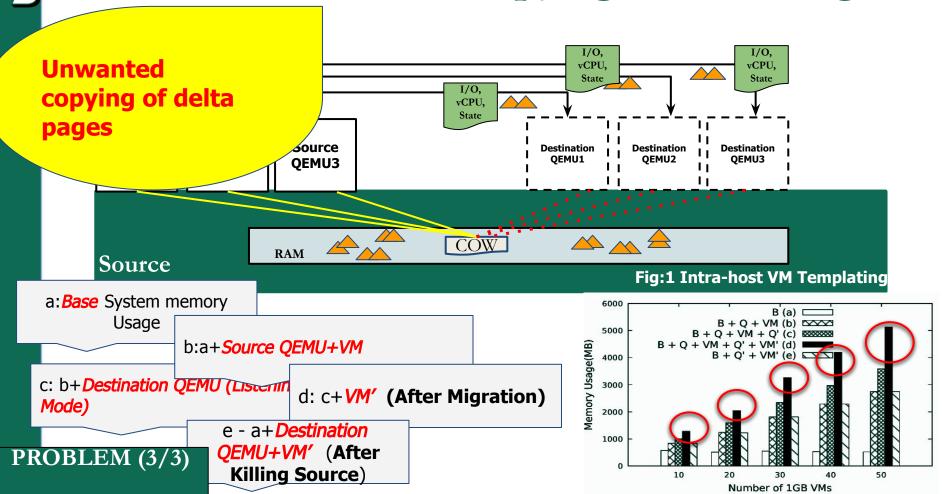


Using "by-pass" flag, the transfer of memory can be eliminated

- High *Memory Usage* due to duplicated back-end file
- Essential to integrate VM
 Templating for efficient
 memory usage

PROBLEM (2/3)

Problem 3: Unwanted Copying of Delta Pages



Intra-host TLM Design: Phase 1 - Initialization

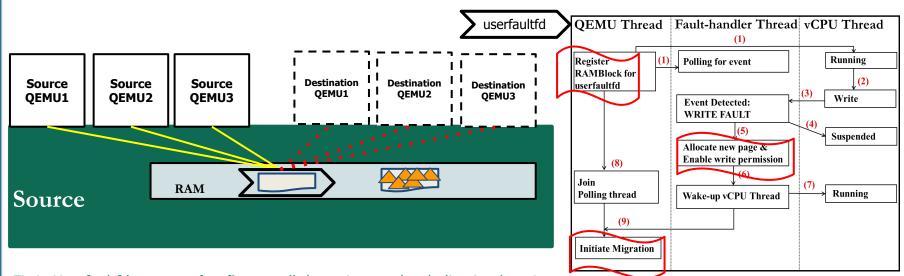


Fig1: Userfaultfd traps and redirects all the writes to the dedication location

Fig2: Workflow of Initialization Phase

Intra-host TLM (1/6)

Intra-host TLM Design: Phase 2 - Migration

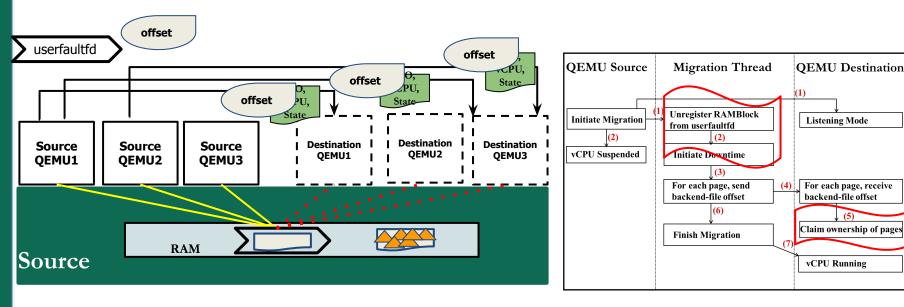


Fig1: Intra-host TLM *transfers the ownership of pages* instead of copying them again by only sending the *vCPU states and offsets*

Fig2: Workflow of Migration

Intra-host TLM (2/6)

Performance Evaluation: Intra-host TLM

Fig 1a: Memory Usage - Generic TLM

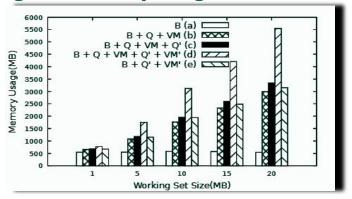
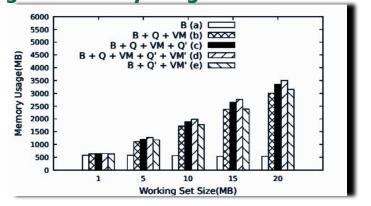


Fig 1b: Memory Usage - Intra-host TLM



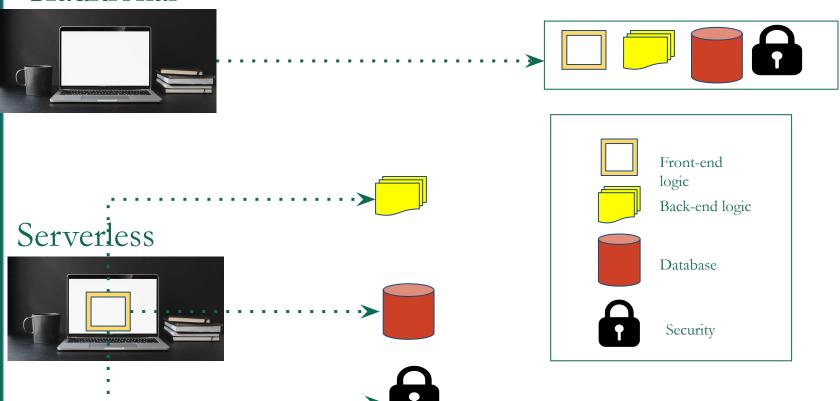
Intra-host TLM **eliminates the intermediate memory bump** thereby reducing the memory usage

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- Future Research Directions

Traditional vs Serverless

Traditional



Serverless Functions and Limitations

- Serverless functions (Stateful/Stateless), due to their scalability and lightweight nature, are deployed on edge nodes.
 - Functions can be wrapped within a Webassembly(WASM)/container/VMs
- . Existing stateful serverless function has following limitations
 - Doesn't support migration
 - Cold-start penalty
 - Thin isolation
 - Lacking awareness of hardware extensions

Future Work

 Using AI/ML Infrastructure tools to decide the optimal isolation strategy for serverless functions

 To investigate the live migration implications for stateful serverless functions wrapped within VMs/containers

 Inter-communication latency for co-located serverless functions wrapped within wasm/containers/VMs

Conclusion

 We identified the limitation of current state-of-art inter/intra host live migration techniques

We introduced *Inter-host TLM* that retains *COW* benefits from VM templating across the machines
 during live migration

we identified the limitation of *Inter-host TLM*, and implemented a more generic technique called *SLM*, that retains COW shared pages from *any COW memory optimization techniques* and works for *all types of VMs*

Conclusion

Finally, we also introduced a more generic technique,
 Intra-host TLM for efficient live migration of all VM types within the same machines avoiding unnecessary copy of pages and concluded with future research directions

Thank you for the attention!

Questions?

MISC SLIDES

Evaluation: Iperf

Fig 1: Pre-copy

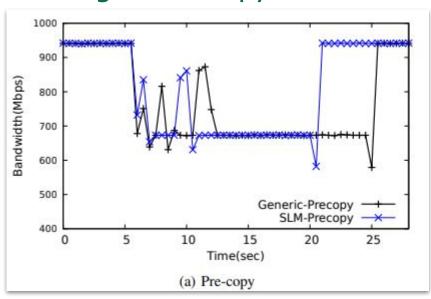
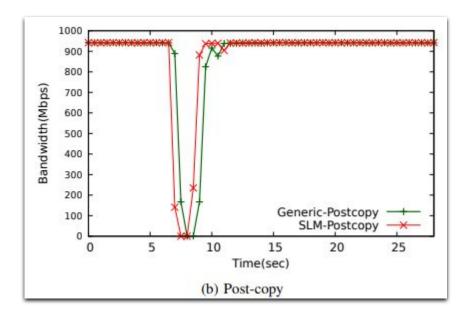


Fig 2: Post-copy



Evaluation: Total Migration Time & Total Pages Transferred



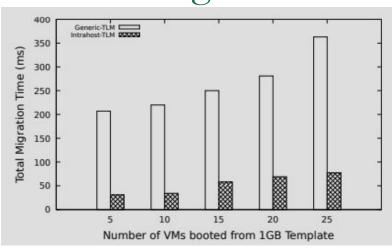
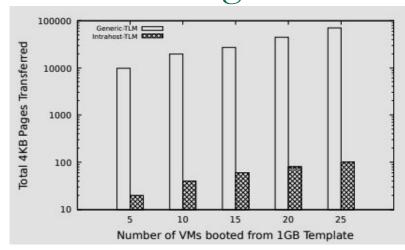
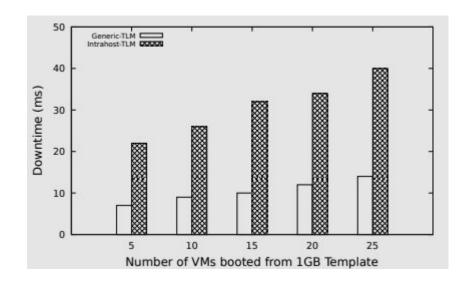


Fig:2



Evaluation: Downtime

 Downtime Overhead is due to transfer of backend-file offsets.



Evaluation: Memory Footprint

Fig:1

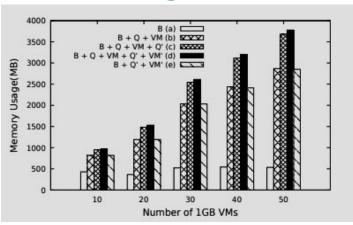
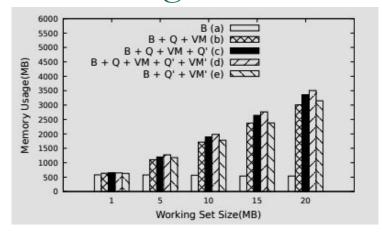
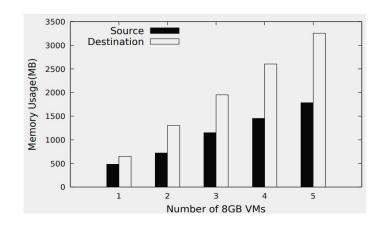


Fig:2



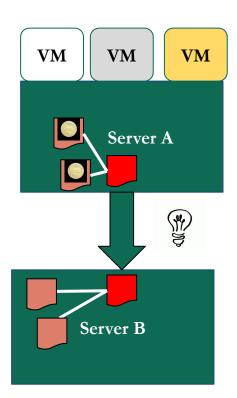
VM Templating & Live Migration

- Current live migration is not aware of VM templates
 - Transfer shared base template pages increasing
 - Total Migration Time
 - Network Traffic

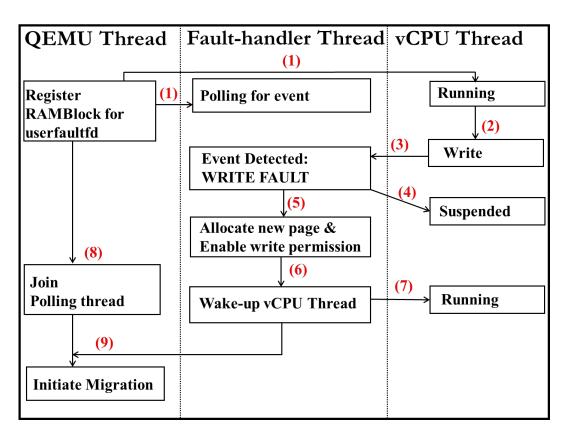


Key Insight: TLM

- Improving live migration awareness of duplicated pages using existing host/hypervisor Copy-on-Write (COW) optimization
 - Reducing Total Migration Time
 - Reducing Network Traffic

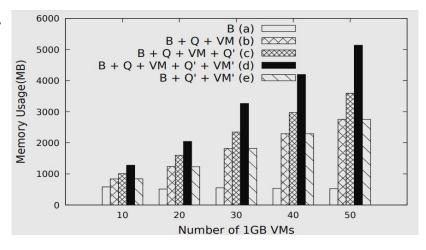


Design: Source



Key Insight: Intrahost TLM

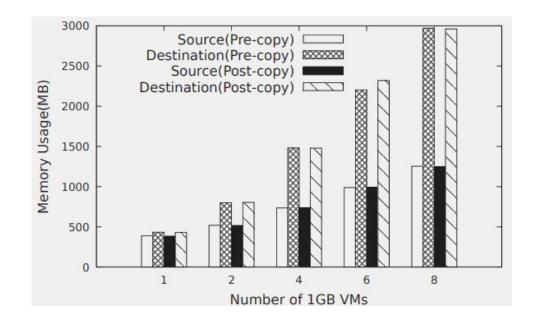
- Current Pre-copy is inefficient for migration within the same host because it copies the pages twice
 - Increasing memory footprint
 - Total Migration Time



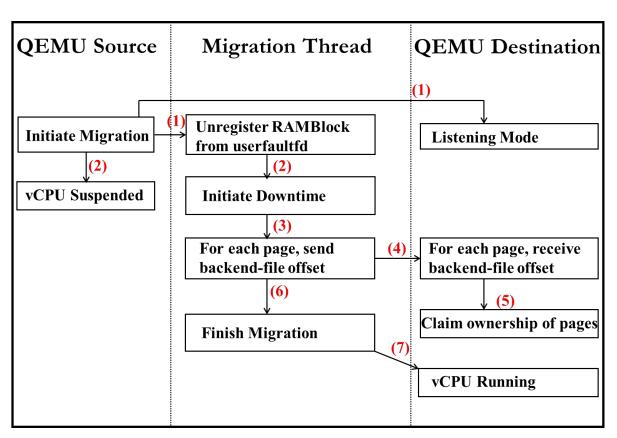
Key Insight: SLM

Current live migration is not aware of COW shared pages irrespective of the underlying memory optimization, increasing

- Total migration time
- Network Traffic



Design: Destination



Future Work

- . Serverless functions, due to their scalability and lightweight nature, is deployed on edge nodes.
- For strong isolation, they needed to be encapsulate inside a VM.

- Our Future work focuses on addressing:
 - The invocation latency of serverless functions within the VM.
 - Inter-VM communication latency for the co-located serverless functions within the same node.