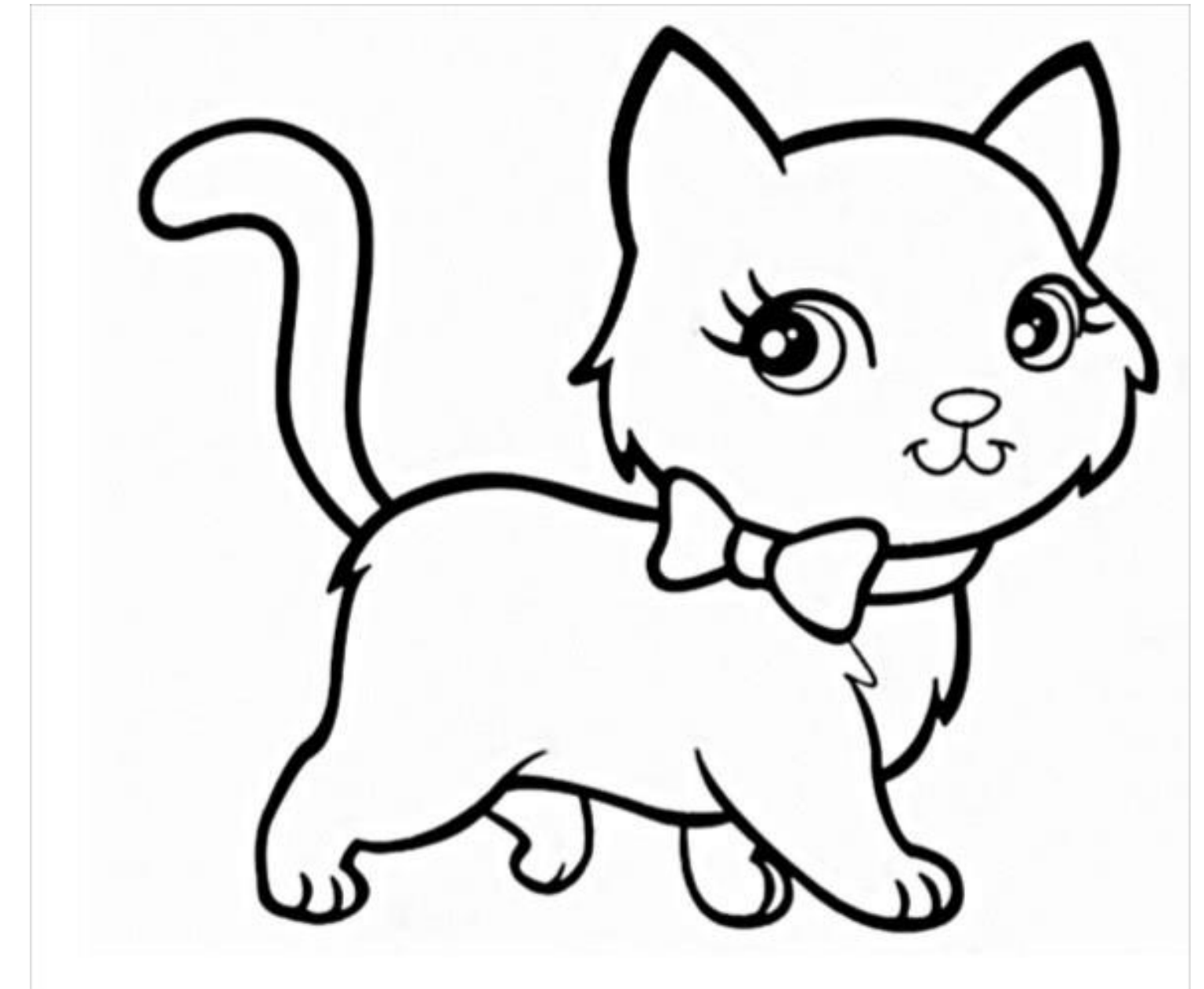
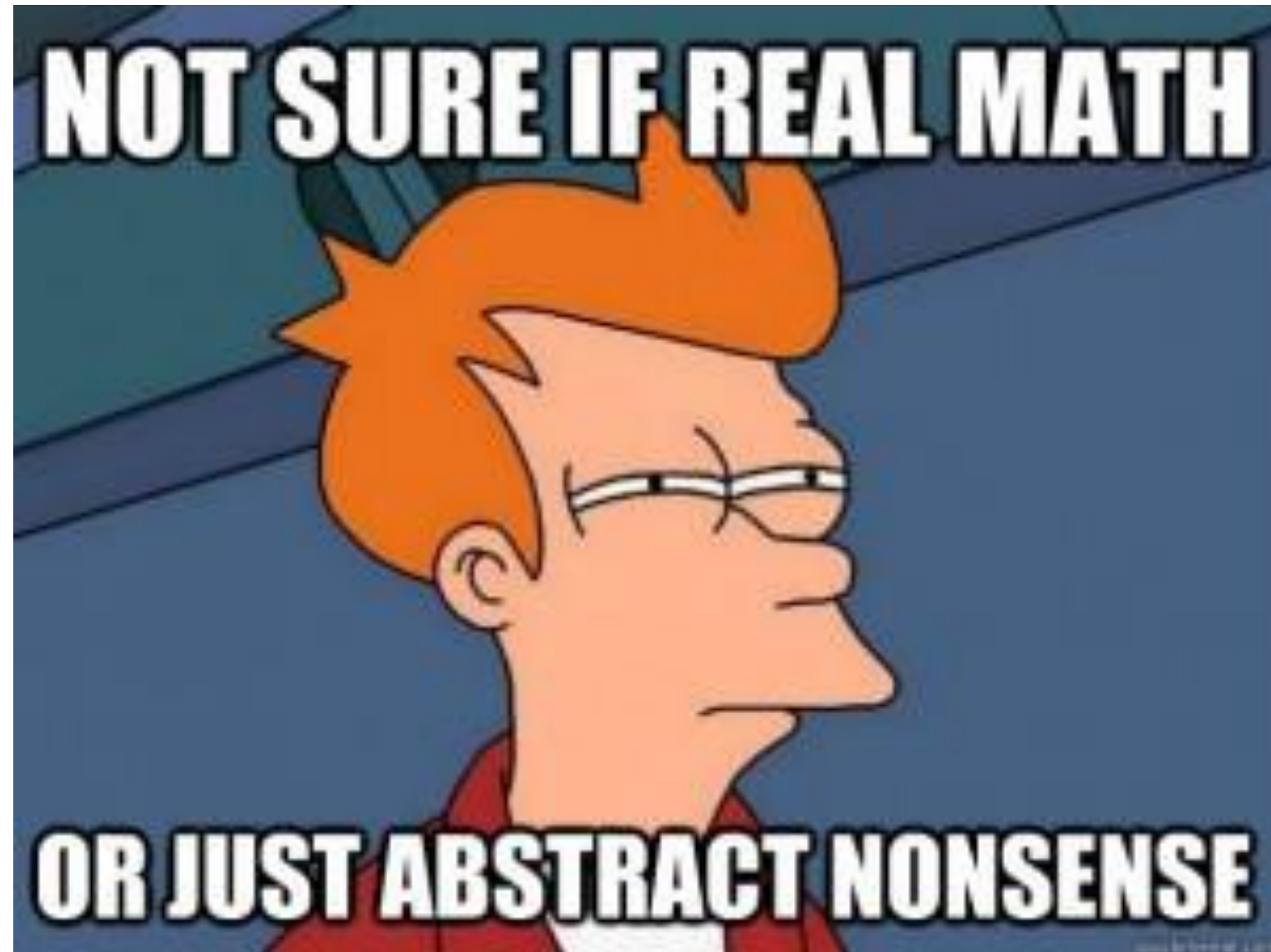


Category Theory: You already know it

Romain Berthon @romaintrm
Emilien Pecoul @ouarzy



Disclaimer



- We're not mathematicians
- Just curious programmers



Question?

www.slido.com

#223494

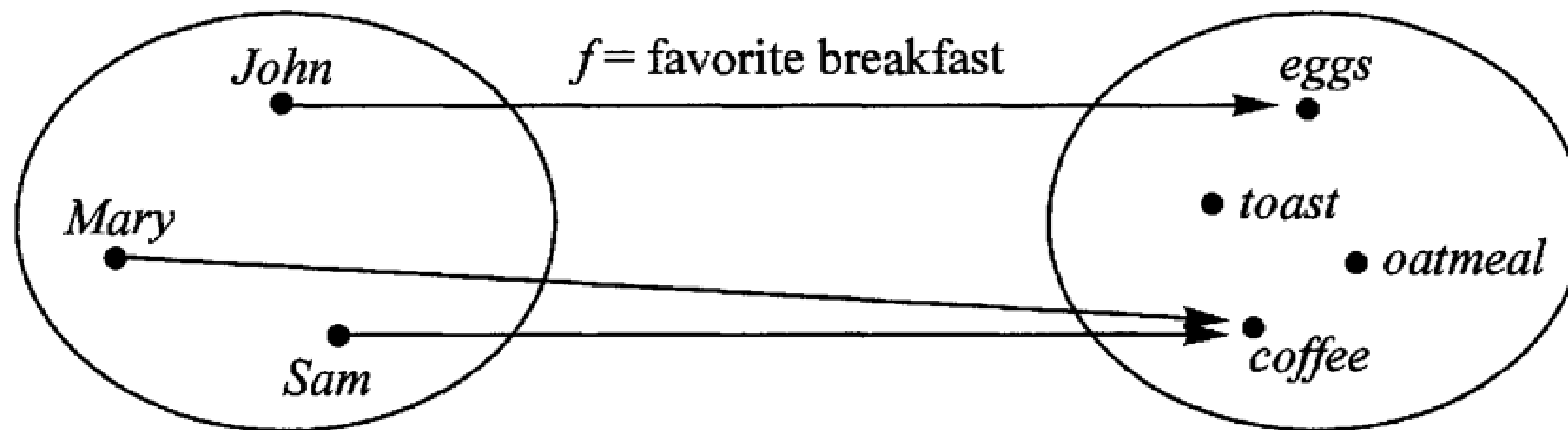


Why Category Theory ?

Common abstract language

Behaviors analysis instead of object analysis

It's all about structure



Code review

Romain BERTHON

```
public Amount GetTotalAmountOfSuspiciousOperations(IReadOnlyList<AccountLine> lines)
{
    var suspiciousOperations : IReadOnlyList<AccountLine> = GetSuspiciousOperations(lines);
    return GetTotalAmount(suspiciousOperations);           // # 1 Composition
}
```

1 usage Romain BERTHON

```
public IReadOnlyList<AccountLine> GetSuspiciousOperations(IReadOnlyList<AccountLine> lines) =>
    lines // IReadOnlyList<AccountLine>
        .Select(line => (line, isSuspicious: IsSuspiciousAmount(line))) // # 3 Functor & map
                                                // Product
        .SelectMany(x:(line,isSuspicious) => x.isSuspicious // Where
            ? new List<AccountLine> { x.line } // Monad & Bind
            : new List<AccountLine>()) // IEnumerable<AccountLine>
        .ToList(); // # 1 Composition
```

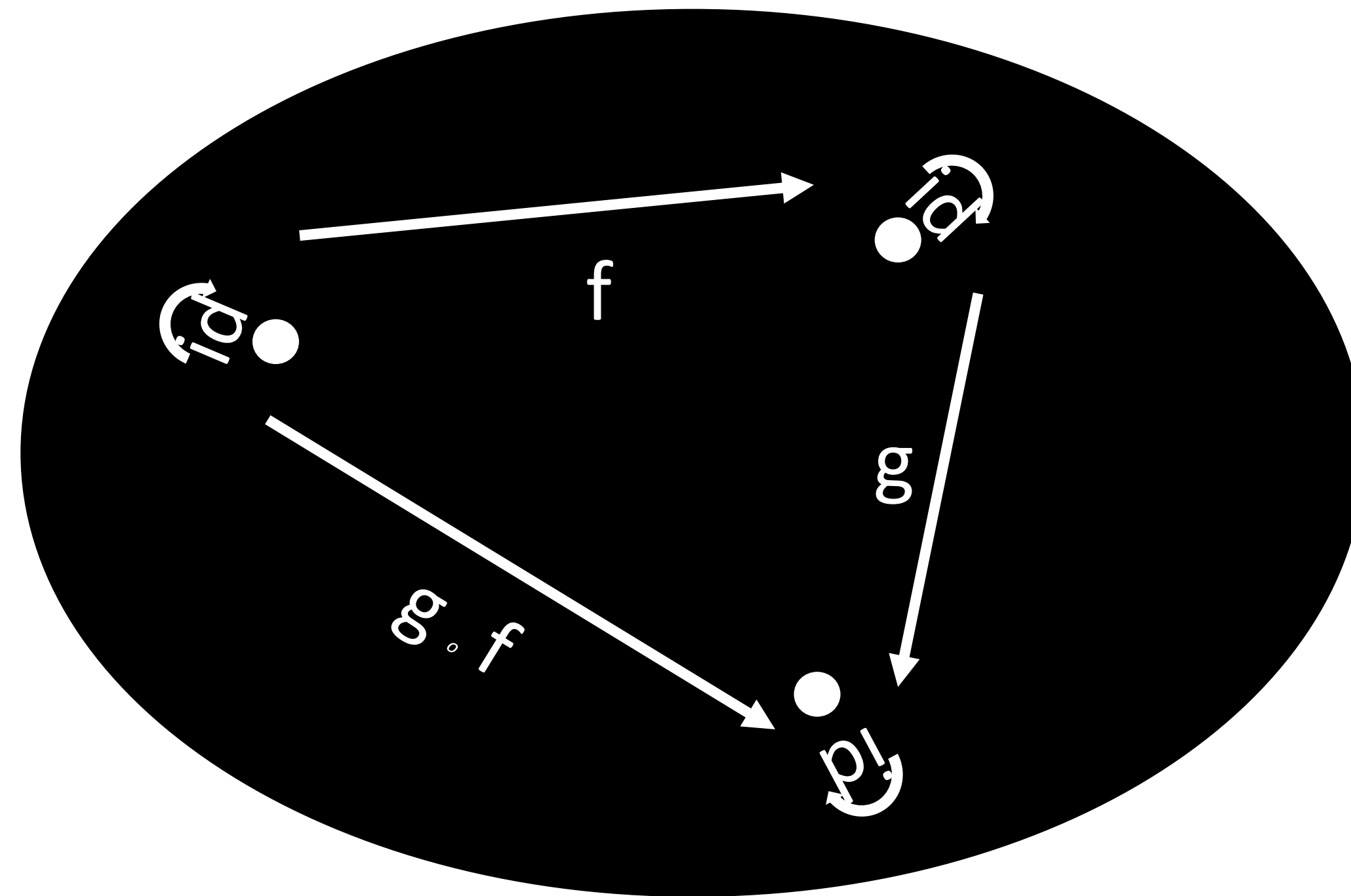
1 usage Romain BERTHON

```
private static bool IsSuspiciousAmount(AccountLine line) => // # 2 Morphisms: Loss of information
    line.Amount.Value > 10_000m;
```

1 usage Romain BERTHON

```
public Amount GetTotalAmount(IReadOnlyList<AccountLine> lines) =>
    lines // IReadOnlyList<AccountLine>
        .Select(line => line.Amount) // # 3 Functor & map
        .Aggregate(Amount.Zero, Amount.Add); // Monoid
```

What's Category Theory ?

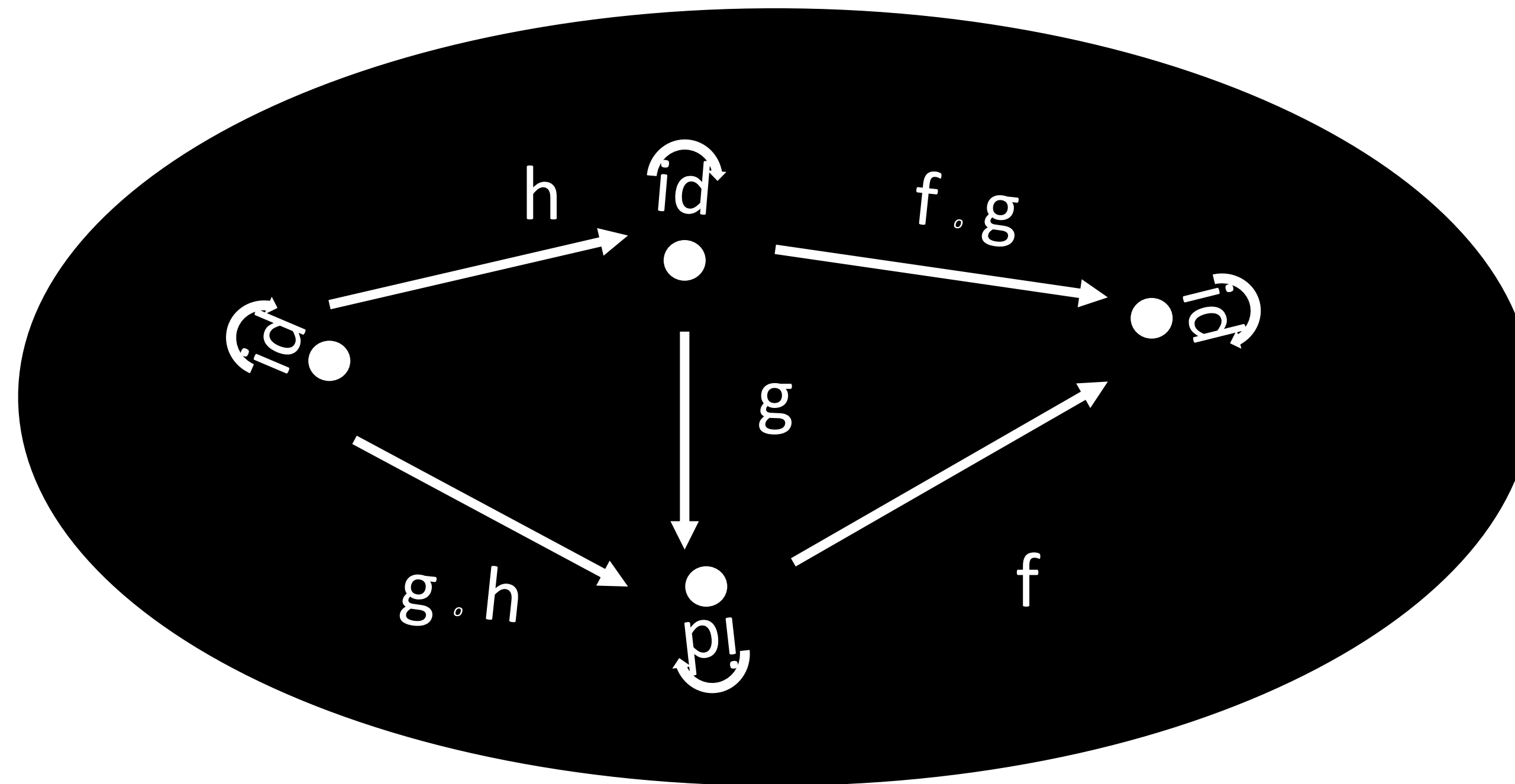


A bunch of **morphisms**, they start and finish with an **object**

Identity: $(a \rightarrow a)$

Composition: $(a \rightarrow b) \rightarrow (b \rightarrow c) = (a \rightarrow c)$

What's Category Theory ?



Identity law: $(id \circ f) = (f \circ id) = f$

Associative law: $(f \circ g) \circ h = f \circ (g \circ h)$

Essence of Category Theory ?

Composition

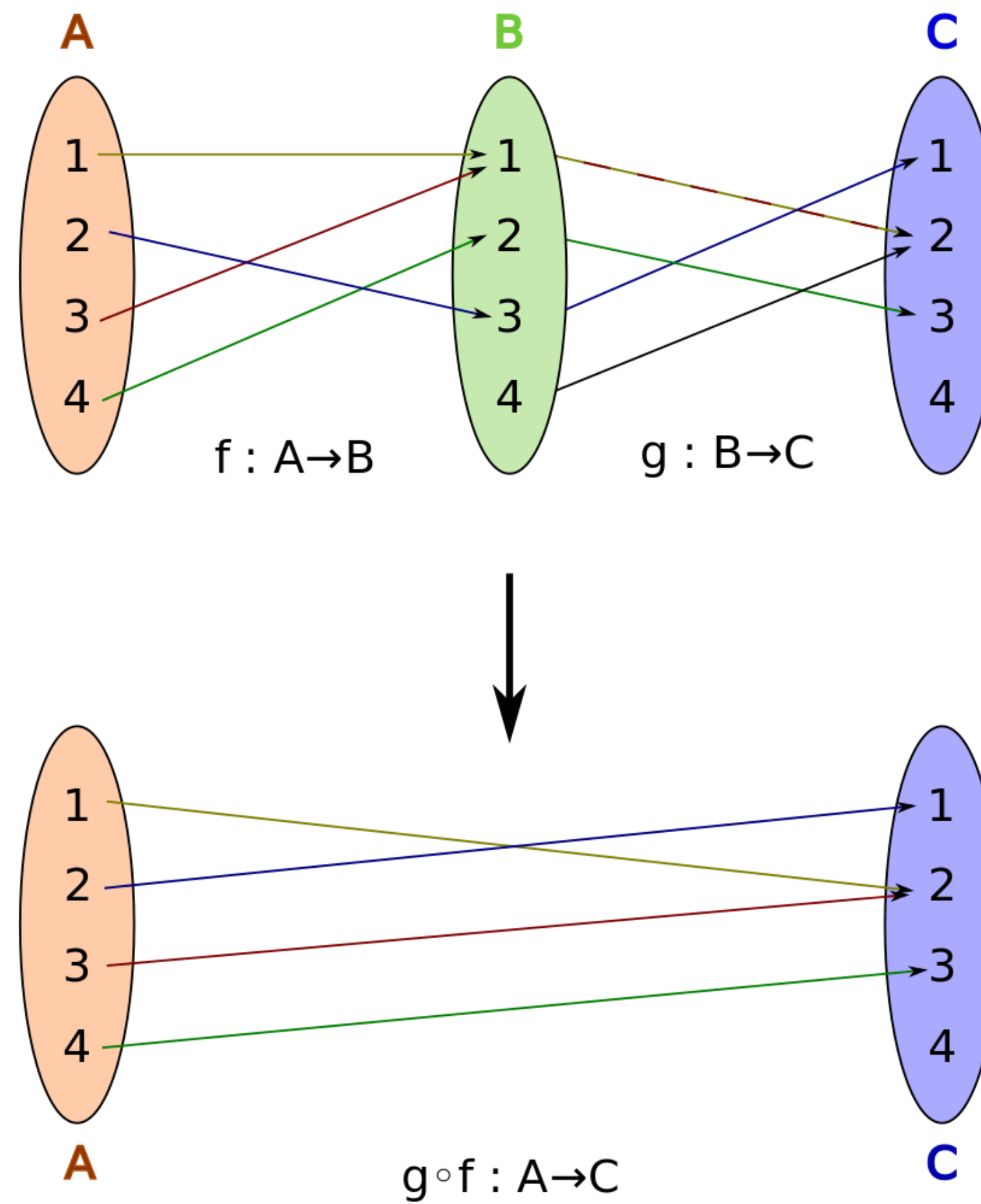
Identity

Abstraction

(Functional)

Programmers know it!

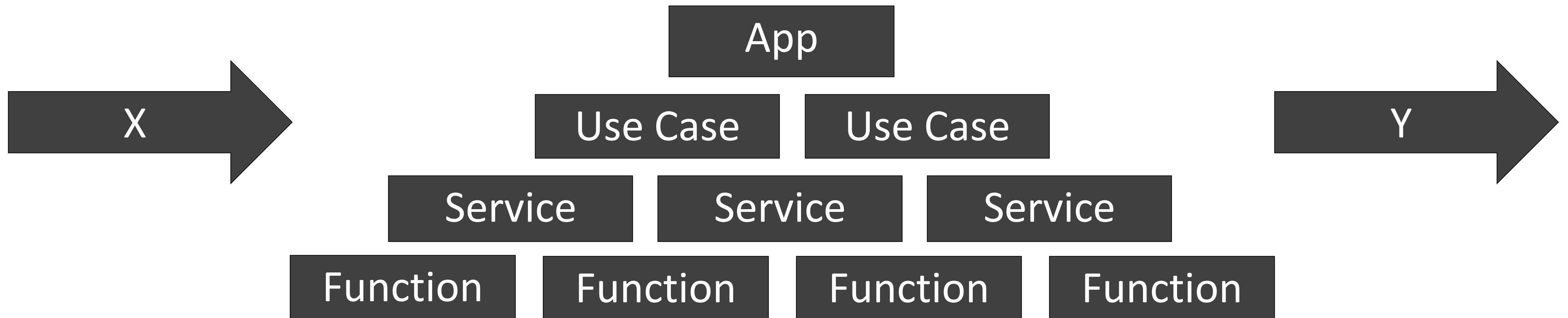
Composition



Composition

Fundamental in functional programming:

- The result of each function is the argument of the next
 - The result of the last one is the result of the whole



Composition

Chunk big problems into smaller ones

Requires no side effect



Example ?

Romain BERTHON

```
public class Composition
```

```
{
```

1 usage Romain BERTHON

```
private int f(decimal value) => (int)value;
```

1 usage Romain BERTHON

```
private bool g(int value) => value % 2 == 0;
```

Romain BERTHON

```
public bool g_after_f(decimal value) => g(f(value));
```

```
}
```

decimal -> int Romain BERTHON

```
let f (x: decimal) = int x
```

int -> bool Romain BERTHON

```
let g (x: int) = x % 2 == 0
```

decimal -> bool Romain BERTHON

```
let g_after_f = f >> g
```

Composition

Composition allows to solve big problems with little composed solutions

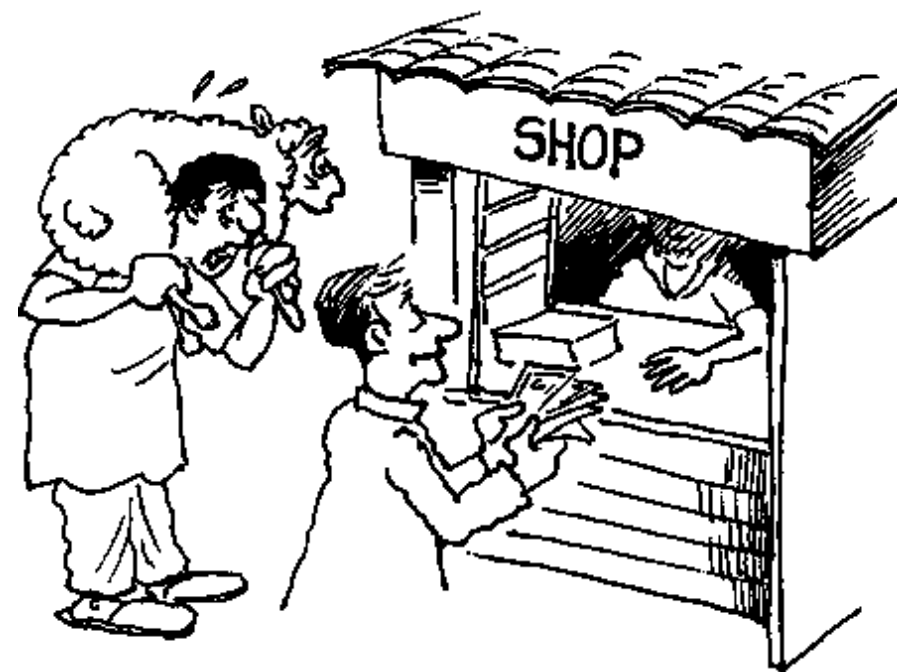


We need composition because reality is too hard to deal with

Abstraction

How do we build software ?

Reality → Software Model → Code



Product

Money

Person

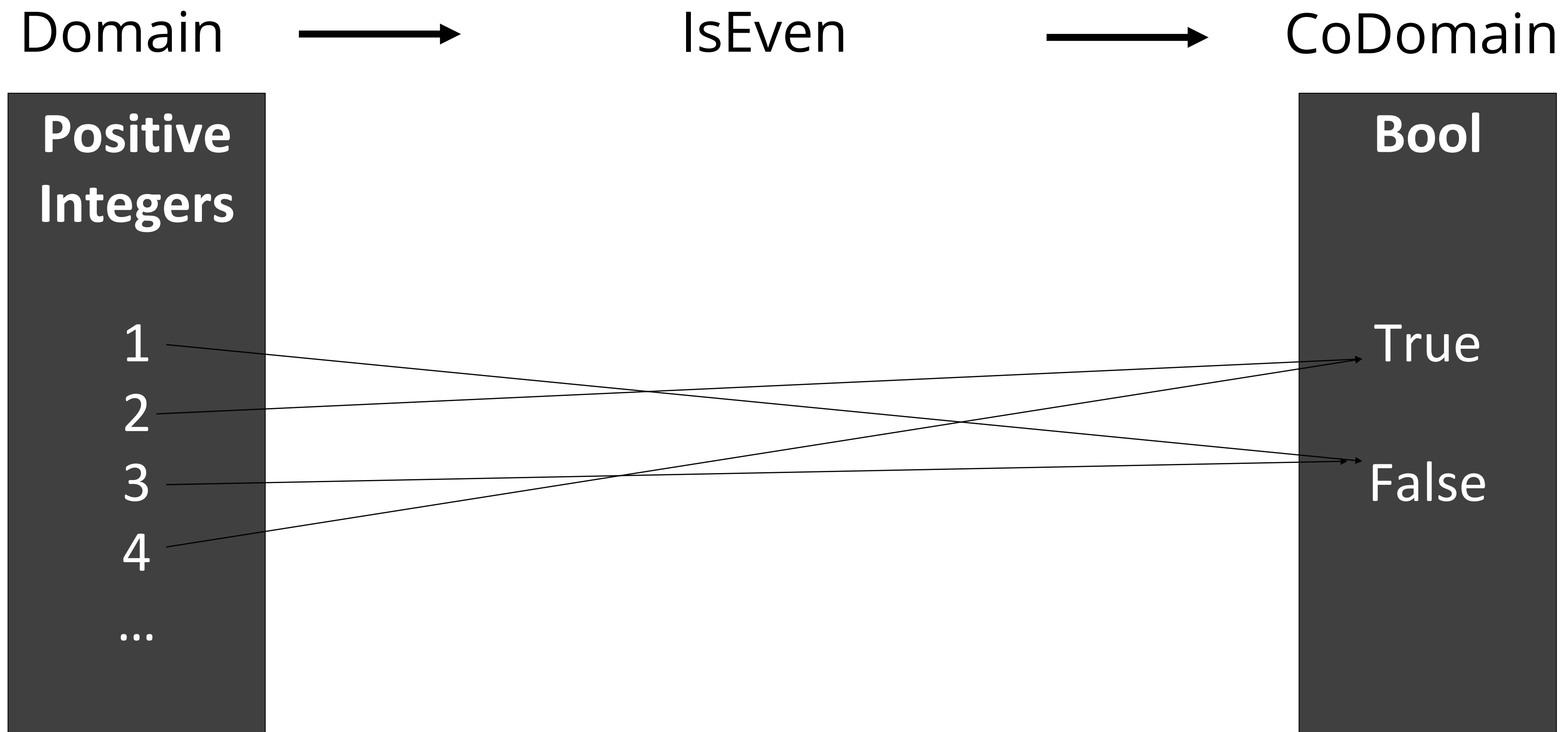
Seller

Buyer

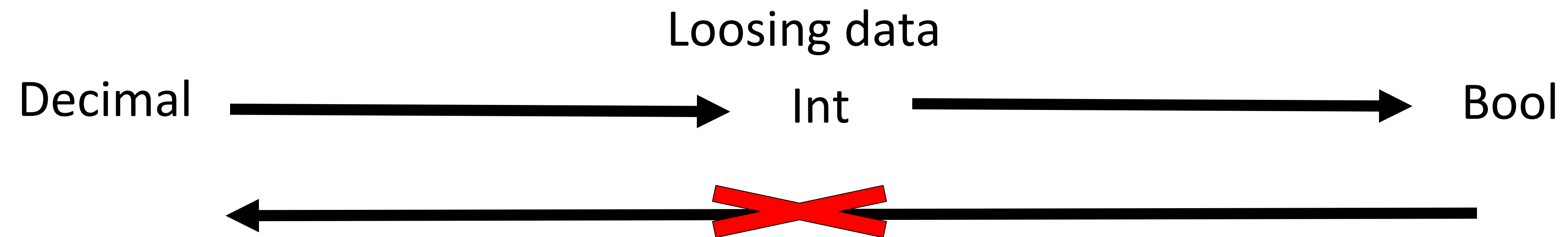
```
class Checkout {  
    private double tax;  
  
    public Checkout(double tax) {  
        this.tax = tax;  
    }  
  
    double total(final List<Product> products) {  
        return products  
            .stream()  
            .mapToDouble(Product::getPrice)  
            .sum() * tax;  
    }  
}
```


Abstraction

How do we abstract stuff ?



Example ?



```
AmountState EvaluateAmountState(AccountLine line)
```

Abstraction

Abstracting is hiding (loosing) useless details in order to apply common rules

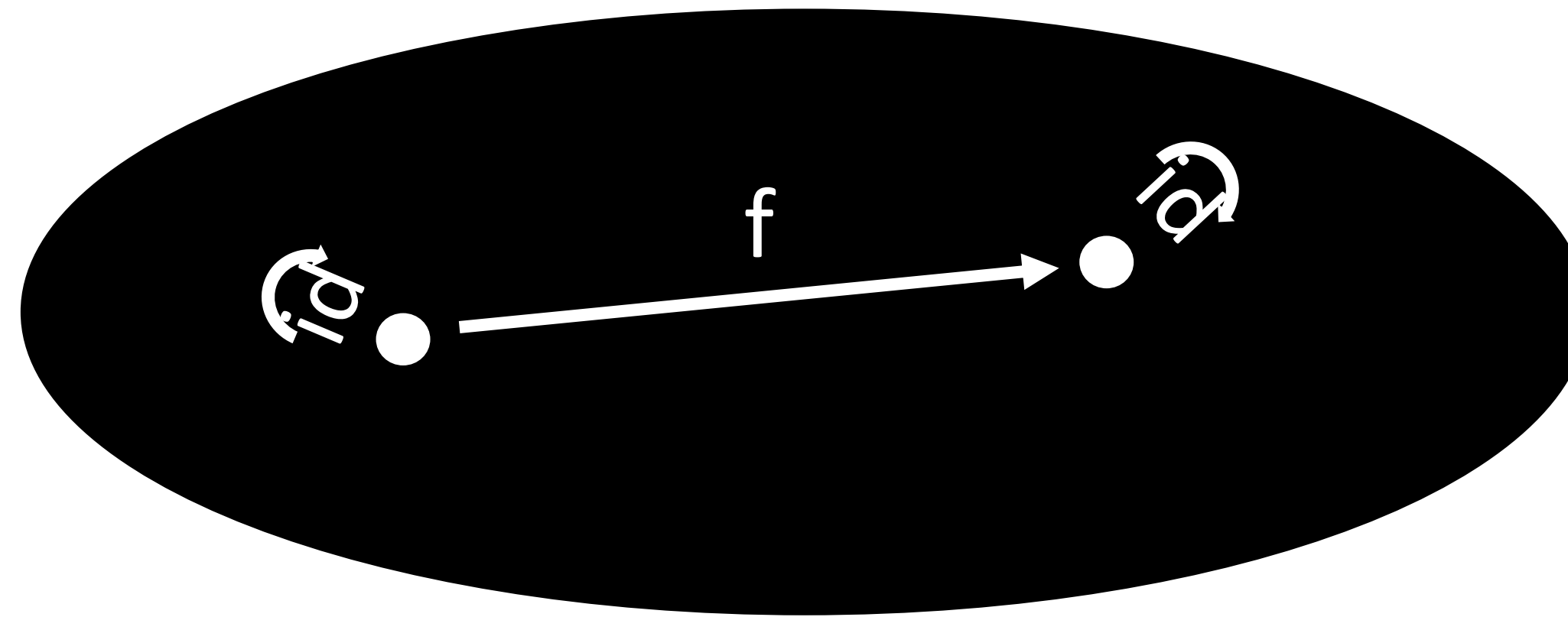


We need abstraction because reality is too hard to deal with

Identity

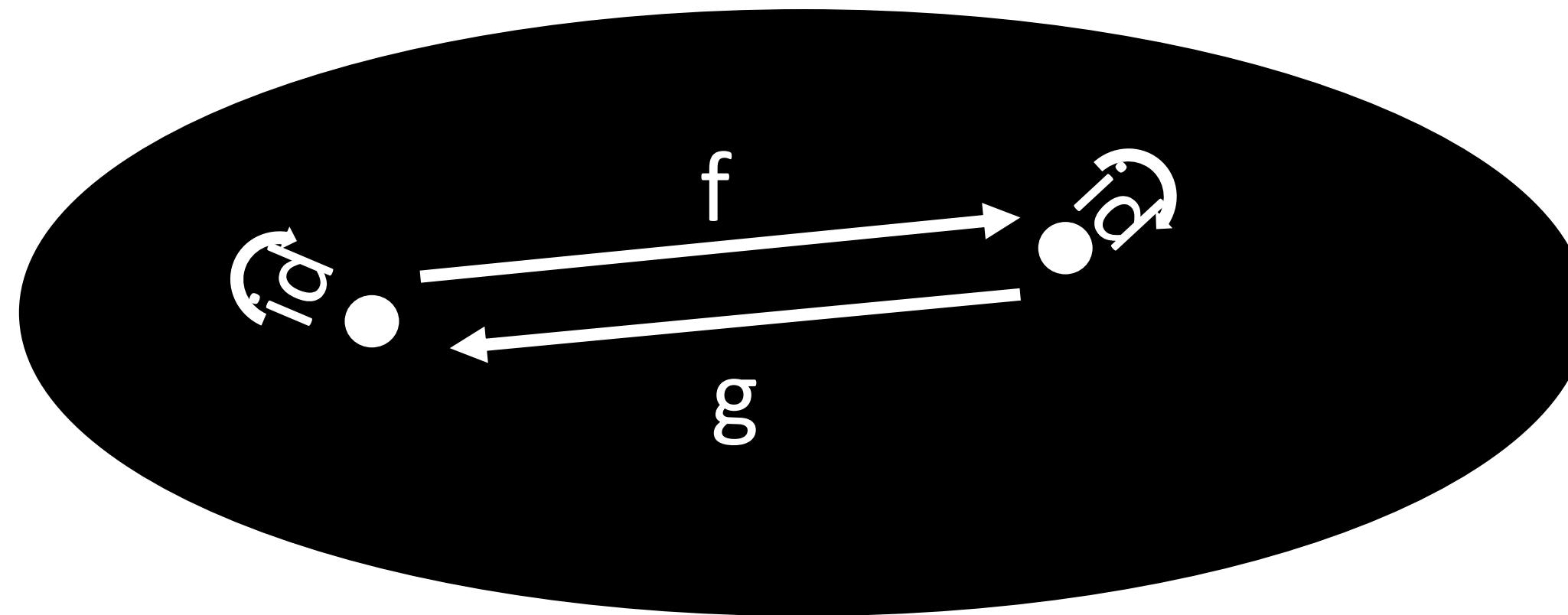


Identity



Identity law: $(\text{id} \circ f) = (f \circ \text{id}) = f$

Identity



Identity law: $(\text{id} \circ f) = (f \circ \text{id}) = f$

$$g \circ f = \text{Id}_{\text{Left}}$$

$$f \circ g = \text{Id}_{\text{Right}}$$

Identity

Does (string,int) = (int,string) ?

Isomorphic

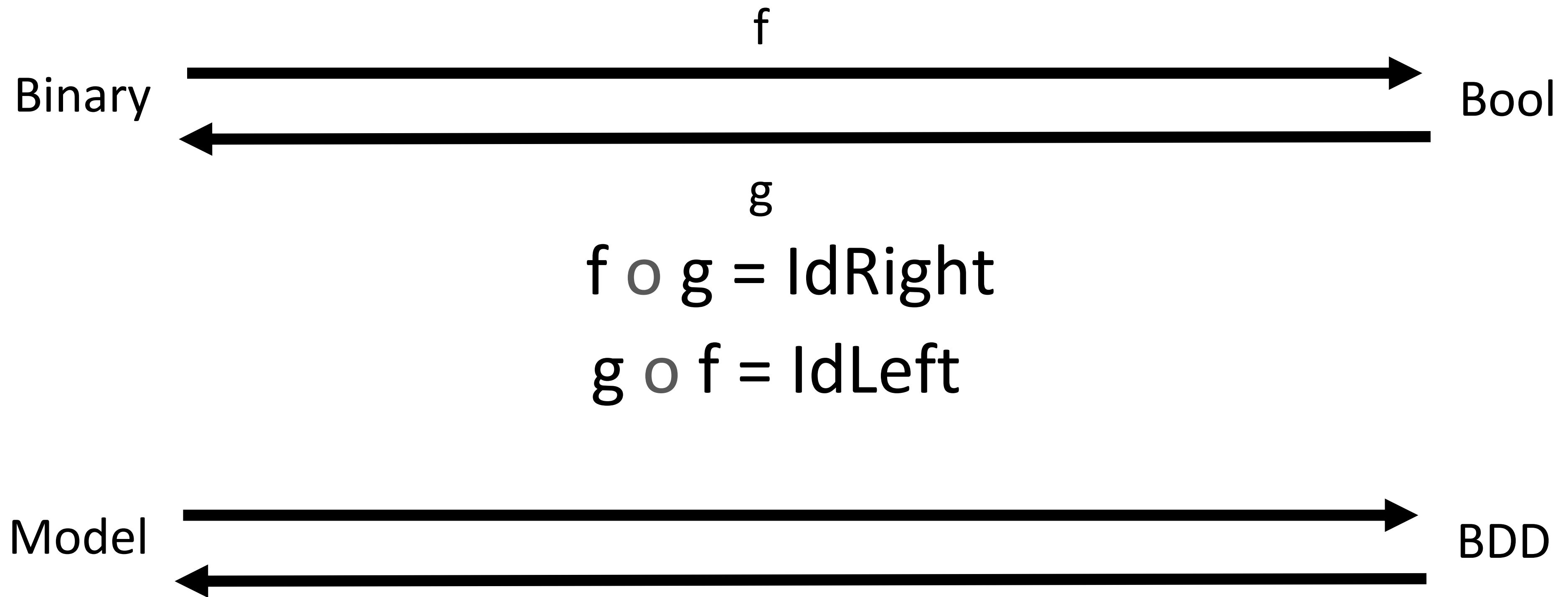
$$\left(\begin{array}{c} \text{f} \\ \circ \\ \text{b0} \end{array} \right) (\text{string}, \text{int}) \begin{array}{c} \text{f} \\ \cong \\ \text{g} \end{array} (\text{int}, \text{string}) \left(\begin{array}{c} \text{f} \\ \circ \\ \text{a0} \end{array} \right)$$

$$g \circ f = \text{Id}_{\text{Left}}$$

$$f \circ g = \text{Id}_{\text{Right}}$$

Example ?

Not loosing data



Identity

Identity is important to define an object in a given context

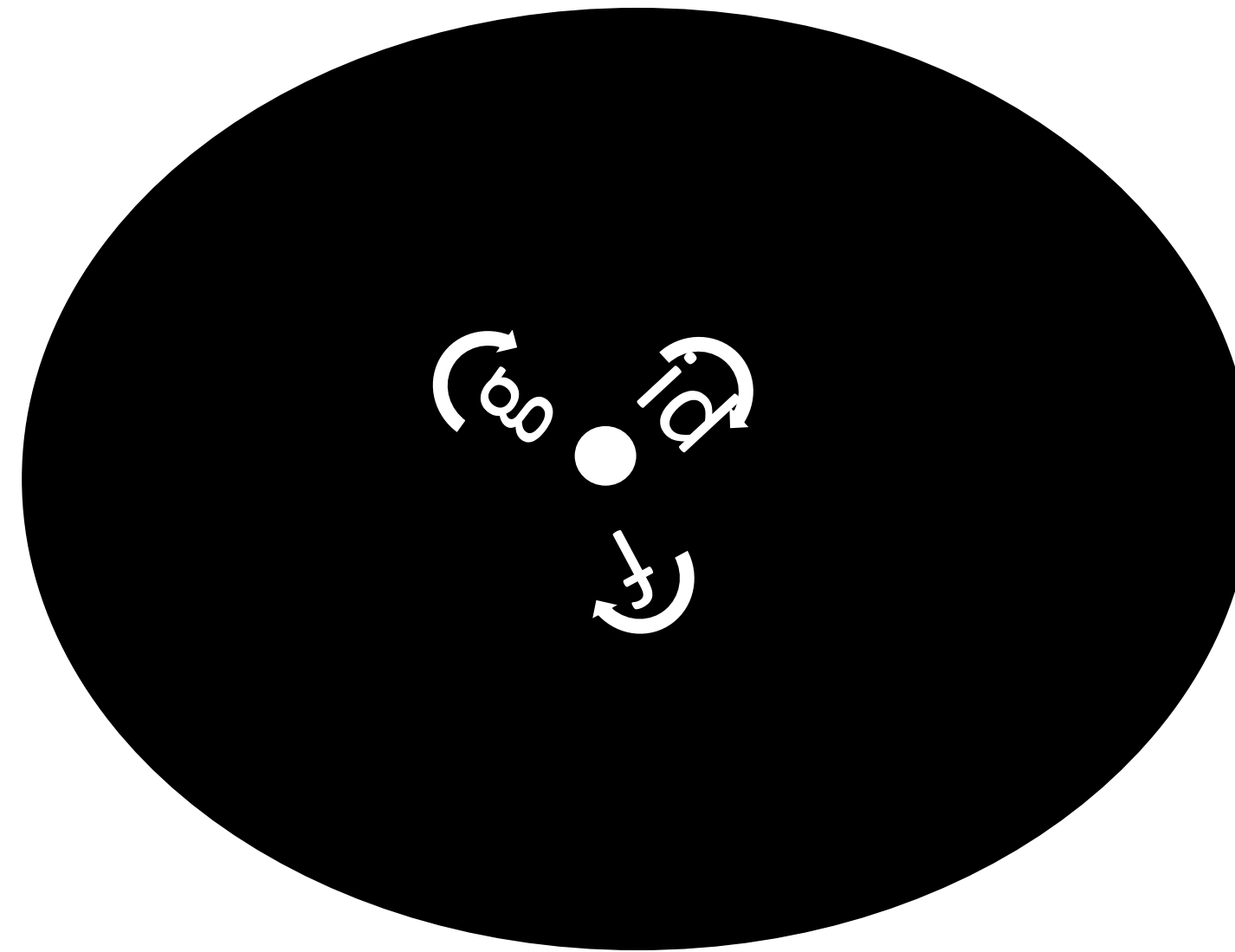


We need identity because reality is too hard to deal with

You already know it!



Monoids

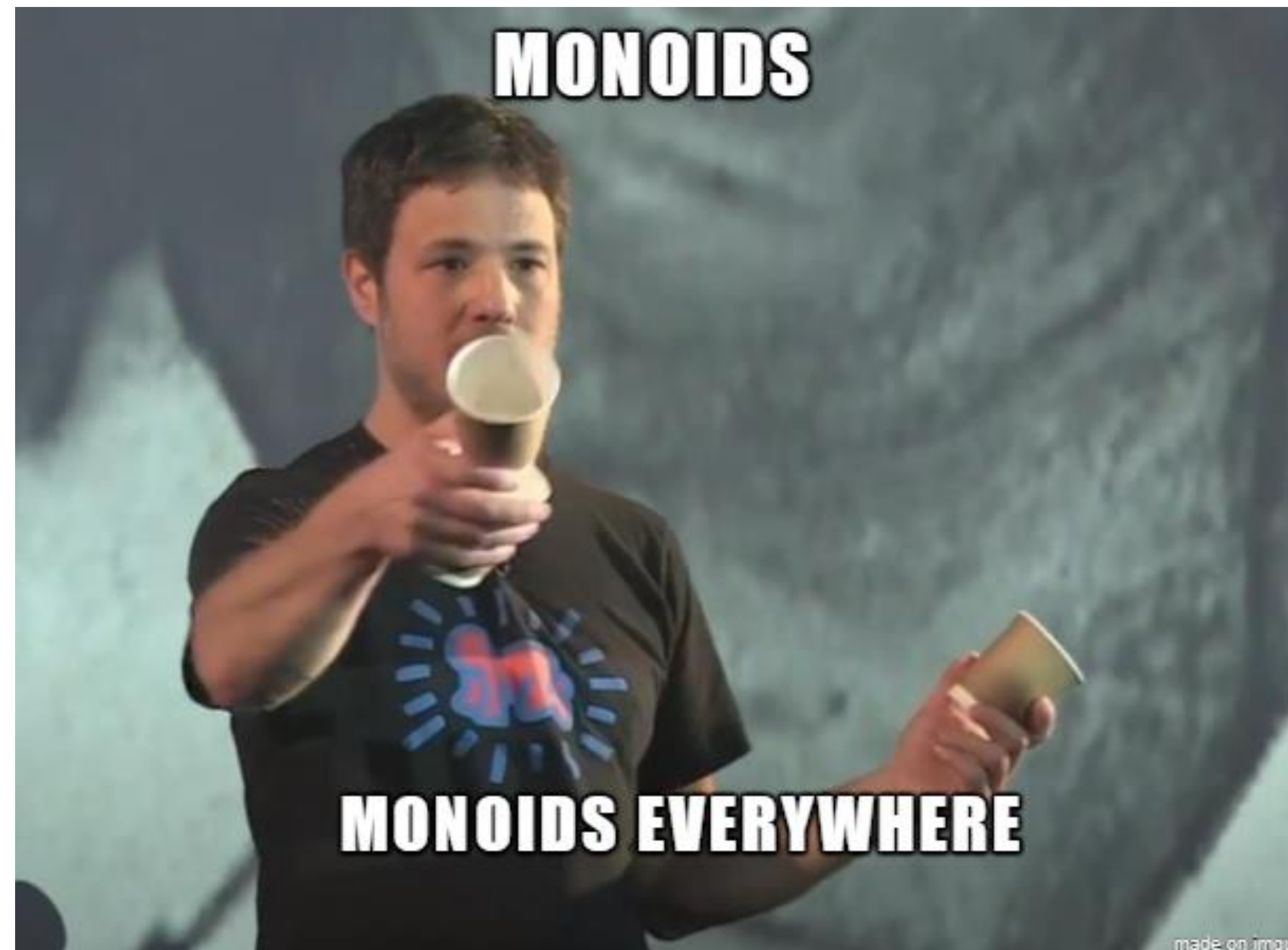


A category with **only one** object

=> **Compose! Associative !**

=> Structure containing only monoids are monoids !

Monoids



<https://youtu.be/J9UwWo2qifg>

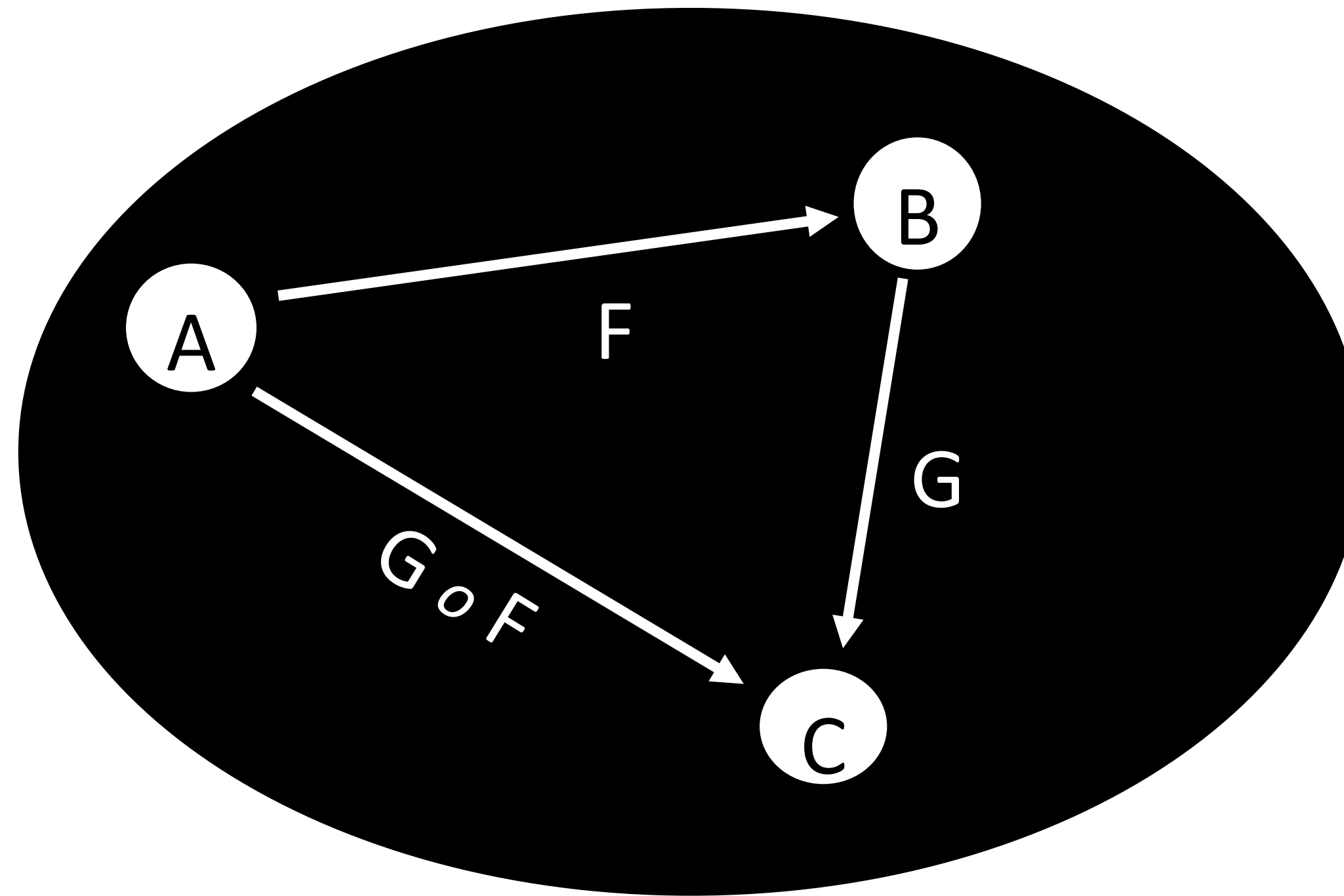
Example ?

```
public record Amount(decimal Value)
{
    public static Amount Add(Amount left, Amount right) => new(left.Value + right.Value);
    public static readonly Amount Zero = new(0m);           // Monoid's Neutral element
}
```

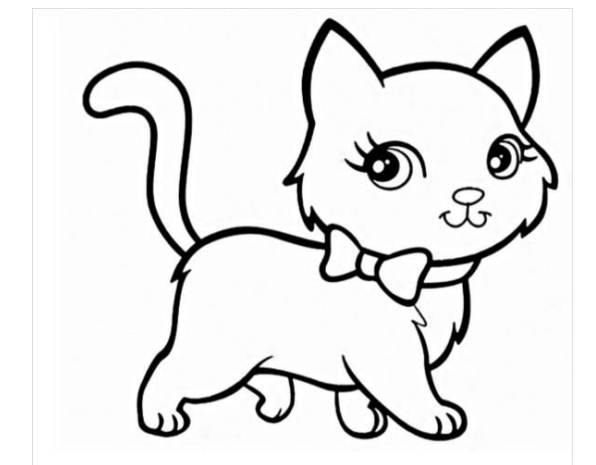
Break



Category you already know



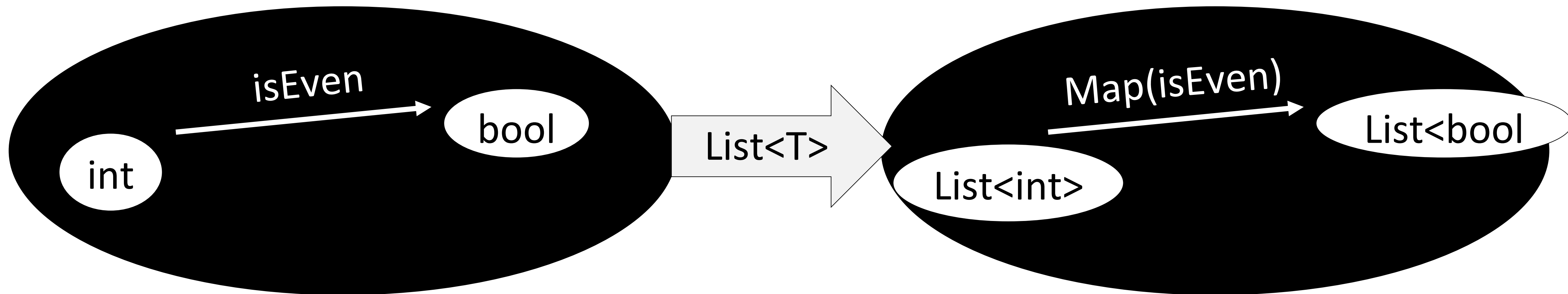
Category of category (cat)



You know category in code

List<T> + map

A functor is a container with a map preserving structure



Example ?

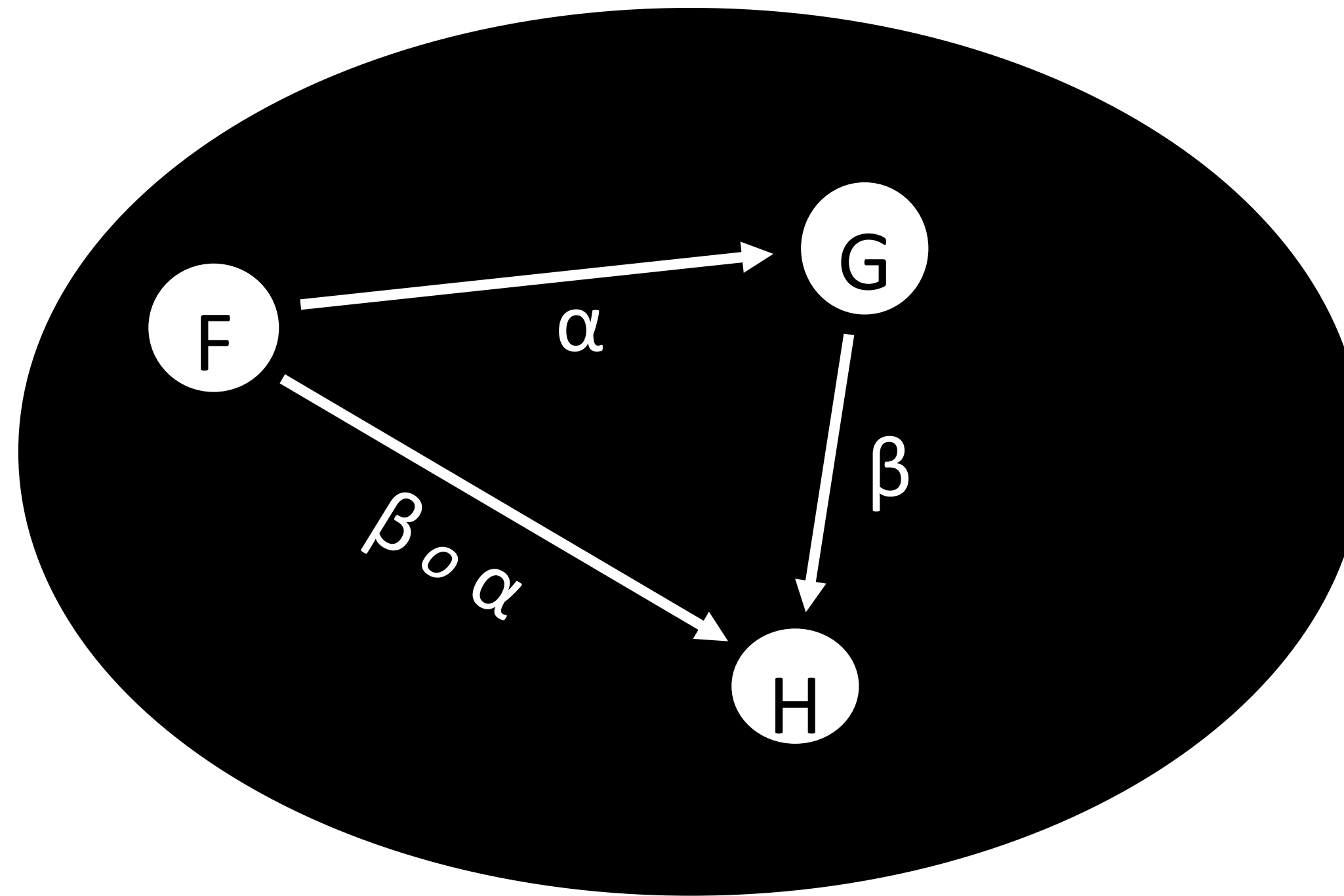
List<T>.map(isEven)



Select in LINQ = map



Category you already know

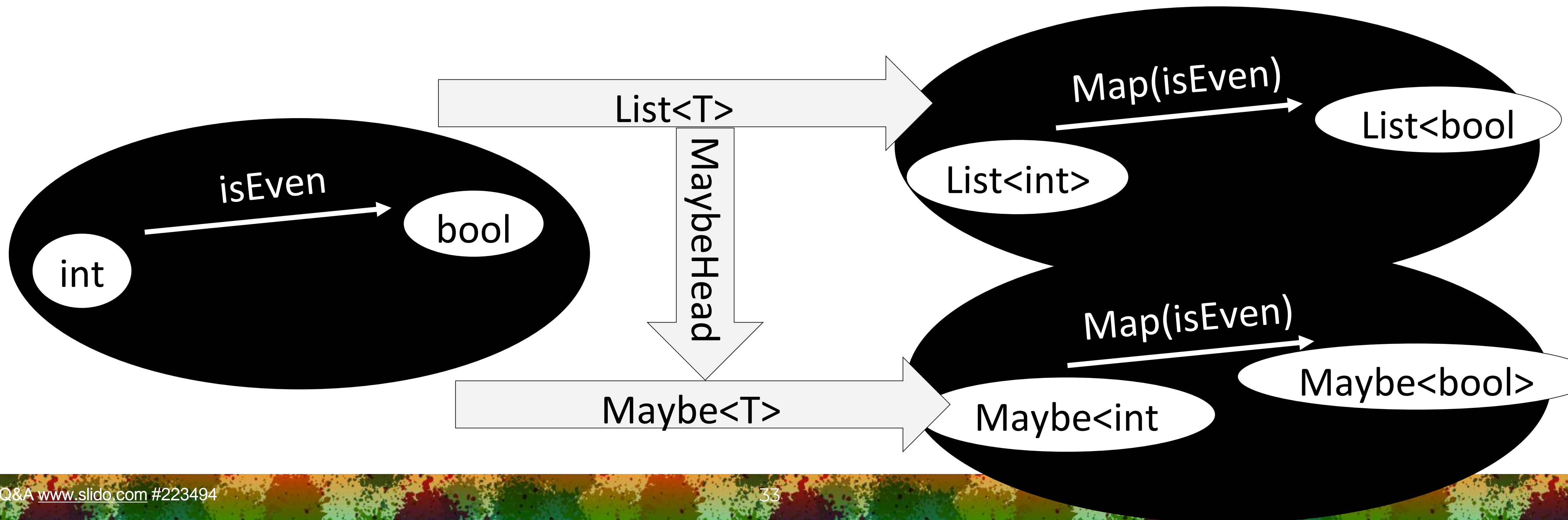


Category of functors

You know category in code

List<T> -> Maybe<T>

A natural transformation is a generic function between two functors



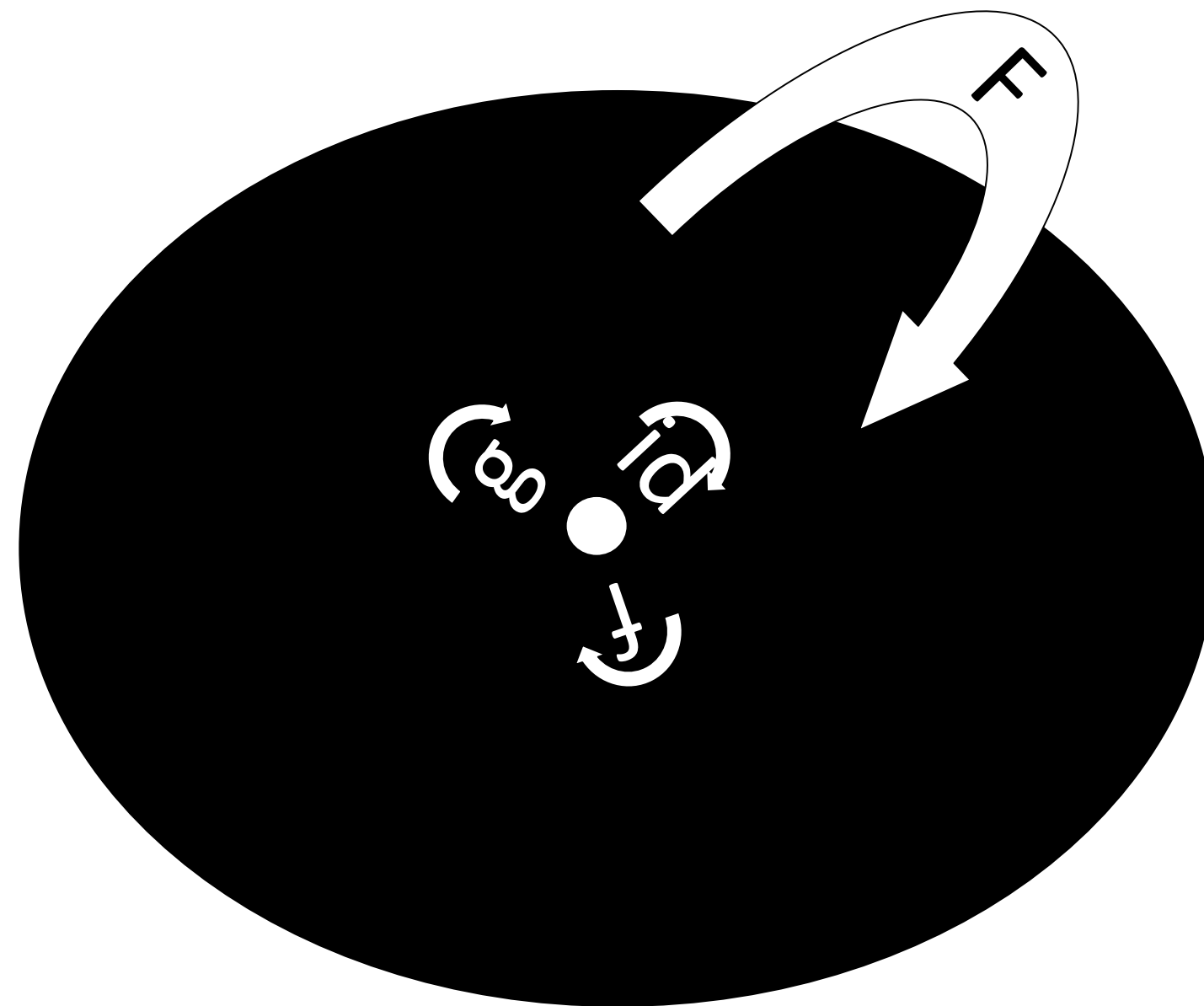
Example ?

List<T>  Maybe<T>
MaybeHead

```
public static Maybe<T> MaybeHead<T>(List<T> list) =>  
    list.Any() ? Maybe<T>.Some(list.First()) : Maybe<T>.None();
```

```
let maybeHead = function  
    | [] -> None  
    | head::tail -> Some head
```

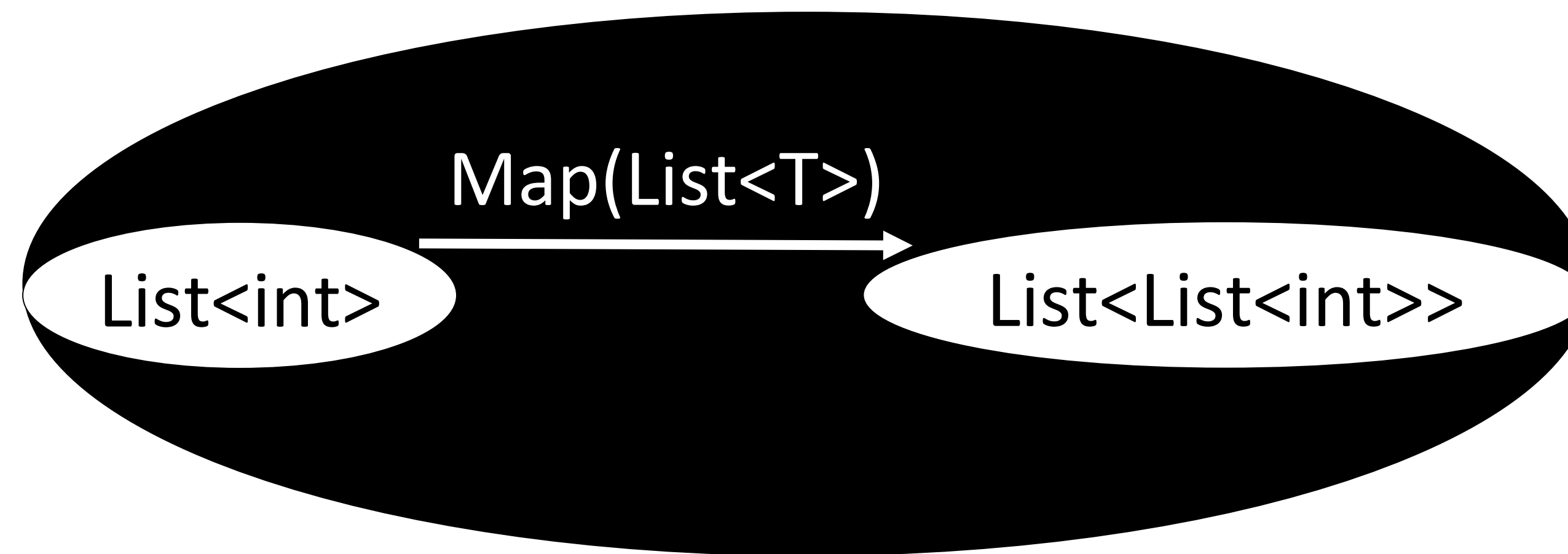

Endofunctor



A functor that map a **category to itself**

Endofunctors are interesting because they do a good job of representing **structures** inside categories that work for **any object**

Endofunctor



An endofunctor can apply to itself

"All the **Functors** we are dealing with in functional **programming** are **Endofunctors**" - <https://blog.softwaremill.com/monoid-in-the-category-of-endofunctors-b85bab43587b>

Example ?

Maybe<T> $\xrightarrow{\text{Map(Maybe<T>)}}$ Maybe<Maybe<T>>

1 usage

```
static Maybe<TOut> Map<TValue, TOut>(Maybe<TValue> maybe, Func<TValue, TOut> morphism) =>
    maybe.Match(
        someMorphism: value => Maybe<TOut>.Some(morphism(value)),
        Maybe<TOut>.None); // Maybe<TOut>
```

```
public static Maybe<Maybe<T>> F<T>(Maybe<T> maybe) => Map(maybe, Maybe<T>.Some);
```

How can I come back ?



Monads



A monad is a monoid in the category of endofunctors

Example ?

```
let safeStringIsPositiveInt = safeStringToInt >> bind safeIsPositive >> map intToString
```

How to compose function with **side effects** ?

=> By **encapsulation** in a monad!

No Silver bullet

Mathematical abstraction are « **easy** to build »

Abstraction from real world will be **harder** to build.



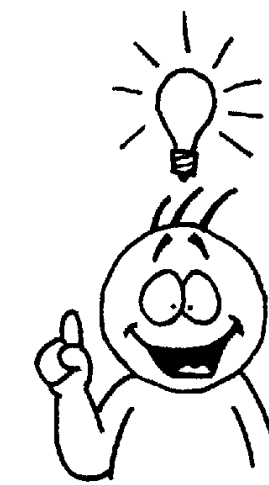
Conclusion

Category theory is **not mandatory** to code...

..But it could help to find **clever solution** to complex programming problems.

[Category theory] does not itself solve hard problems [...] It puts the hard problems in clear relief and makes their solution possible.

—*The Last Mathematician (Hilbert Gottingen)*



Thanks



References

- Category Theory for beginners (Ken Scambler)
- Category Theory for programmers (Bartosz Milewski)
- Philip Wadler 's Blog: <http://homepages.inf.ed.ac.uk/wadler/>
- Robb Seaton's Blog: <http://rs.io/why-category-theory-matters/>
- Category Theory for the working mathematicians (MacLane, Saunders)

History

- 1940s: Einlenberg and MacLane formalize Category Theory
- 1958: Monads discovered by Godement
- 1990: Moggi, Wadler apply monads to functional programming

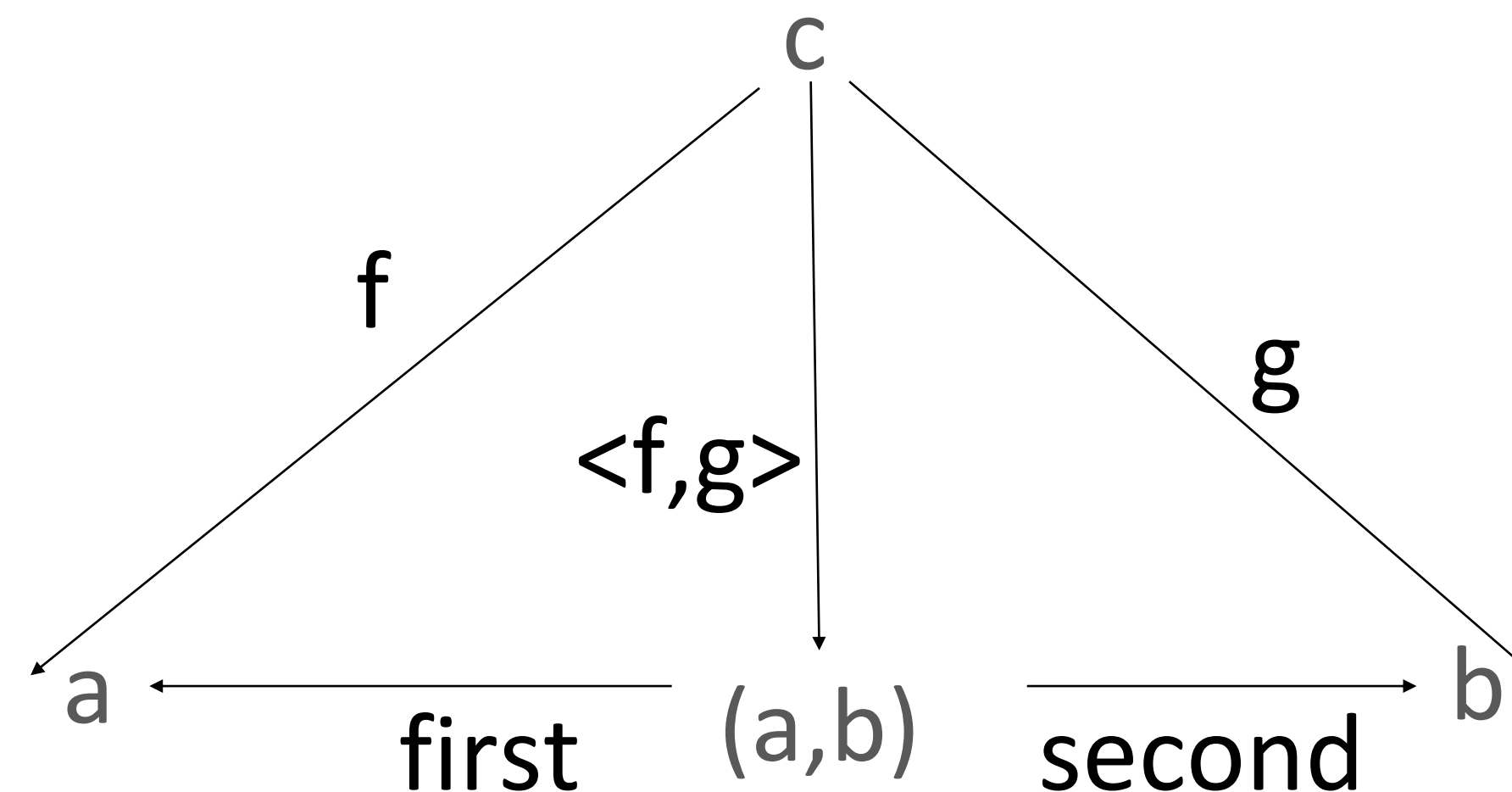


You know category in math

Product

$$A^C * B^C = (A * B)^C$$

$$(C \rightarrow A, C \rightarrow B) = C \rightarrow (A, B)$$

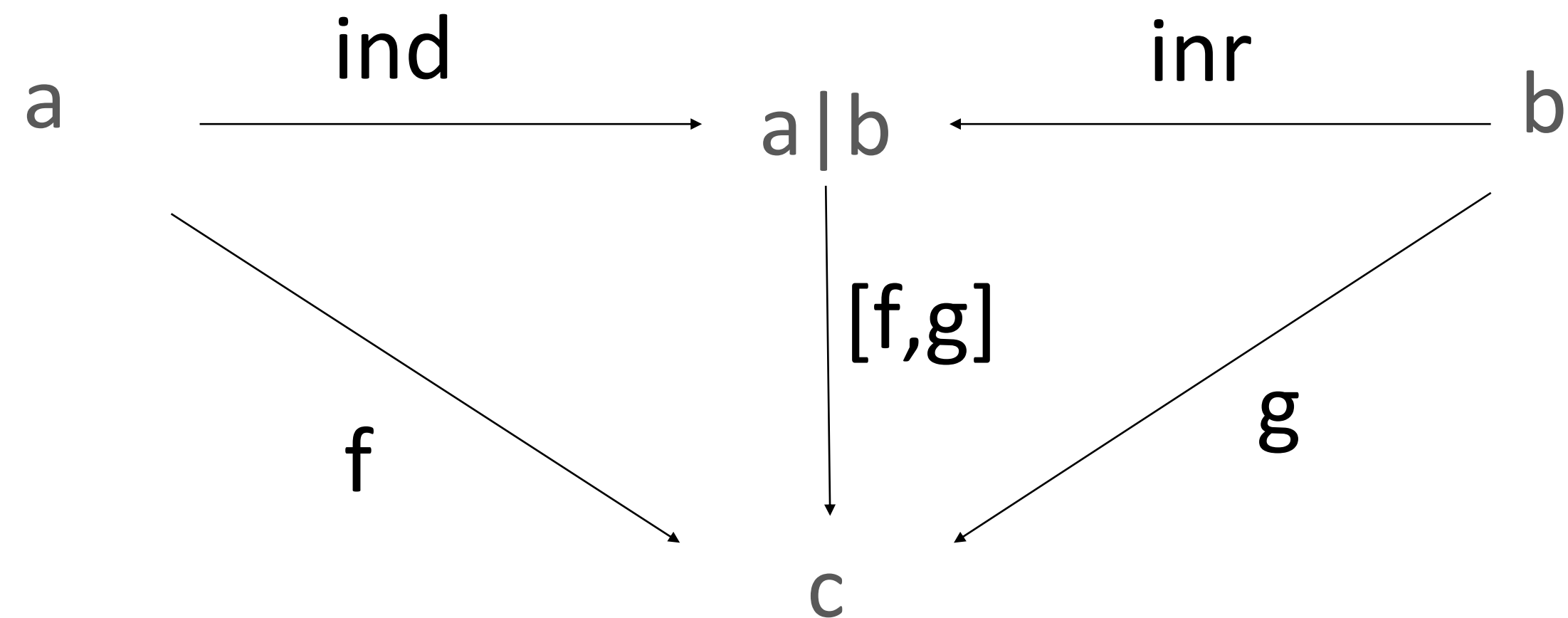


You know category in math

Co Product

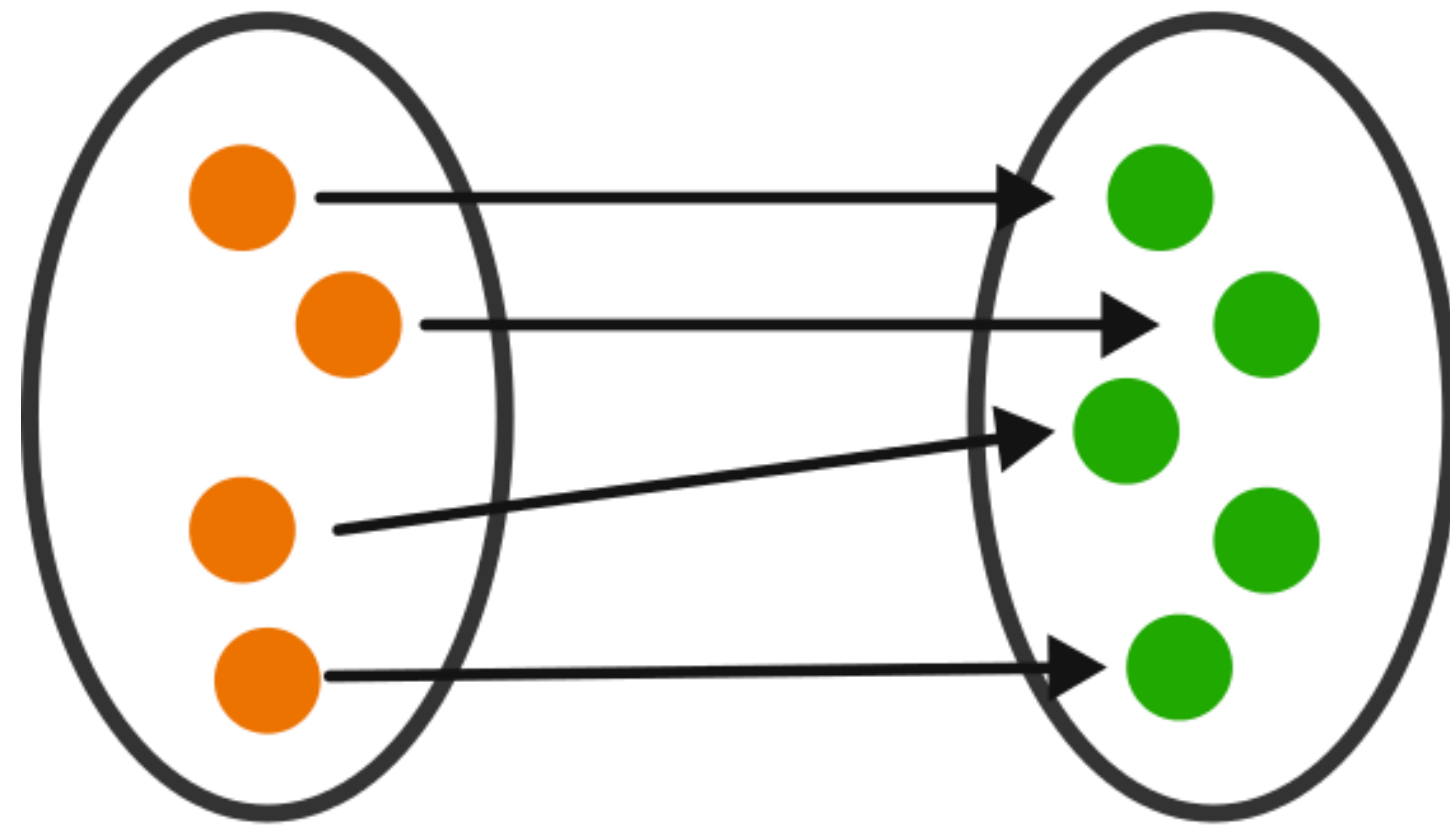
$$(A + B) * C = A * C + B * C$$

$$(A | B , C) = (A, C) | (B, C)$$

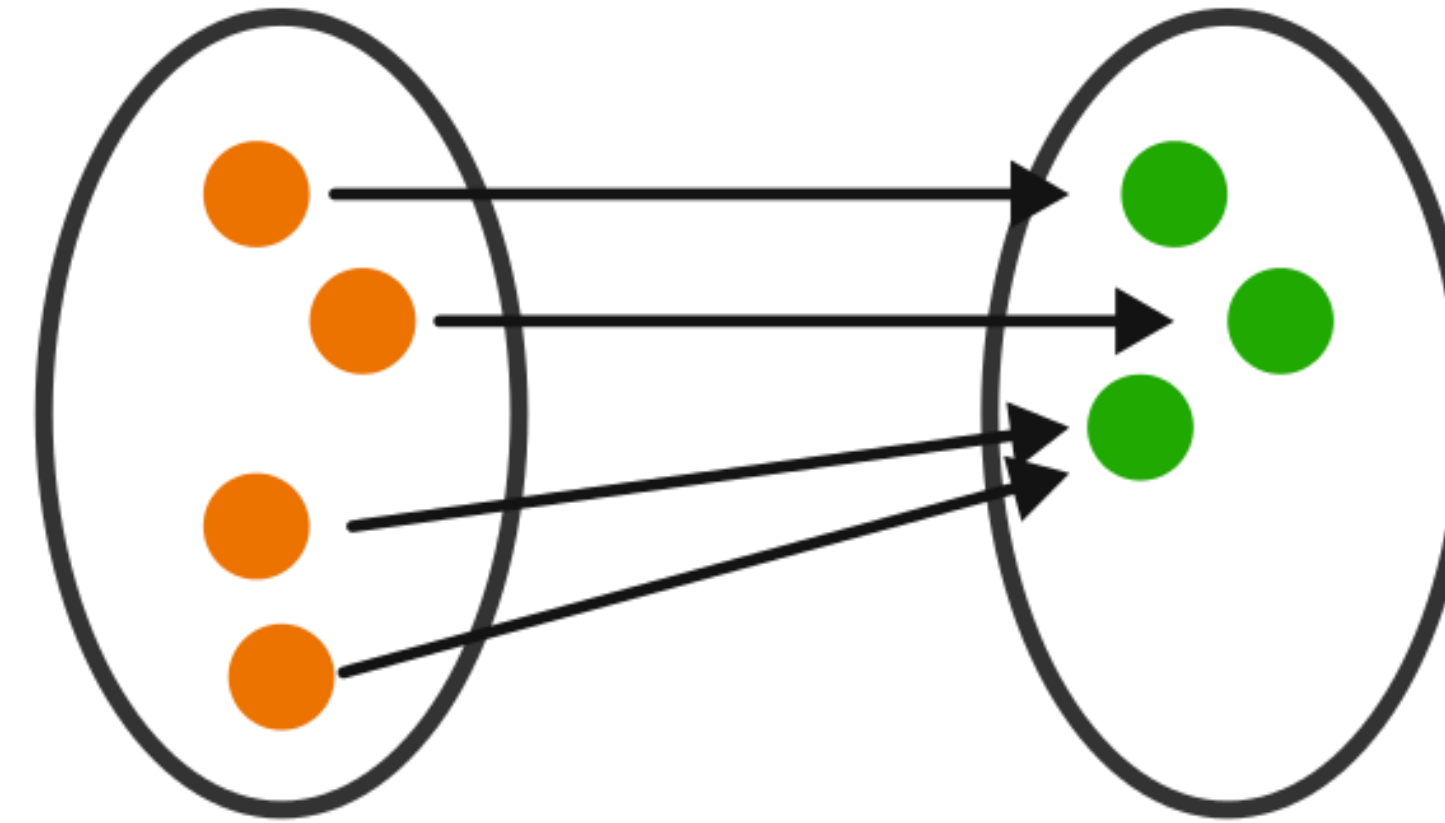


Surjective/Bijective

Injection (One-to-One)



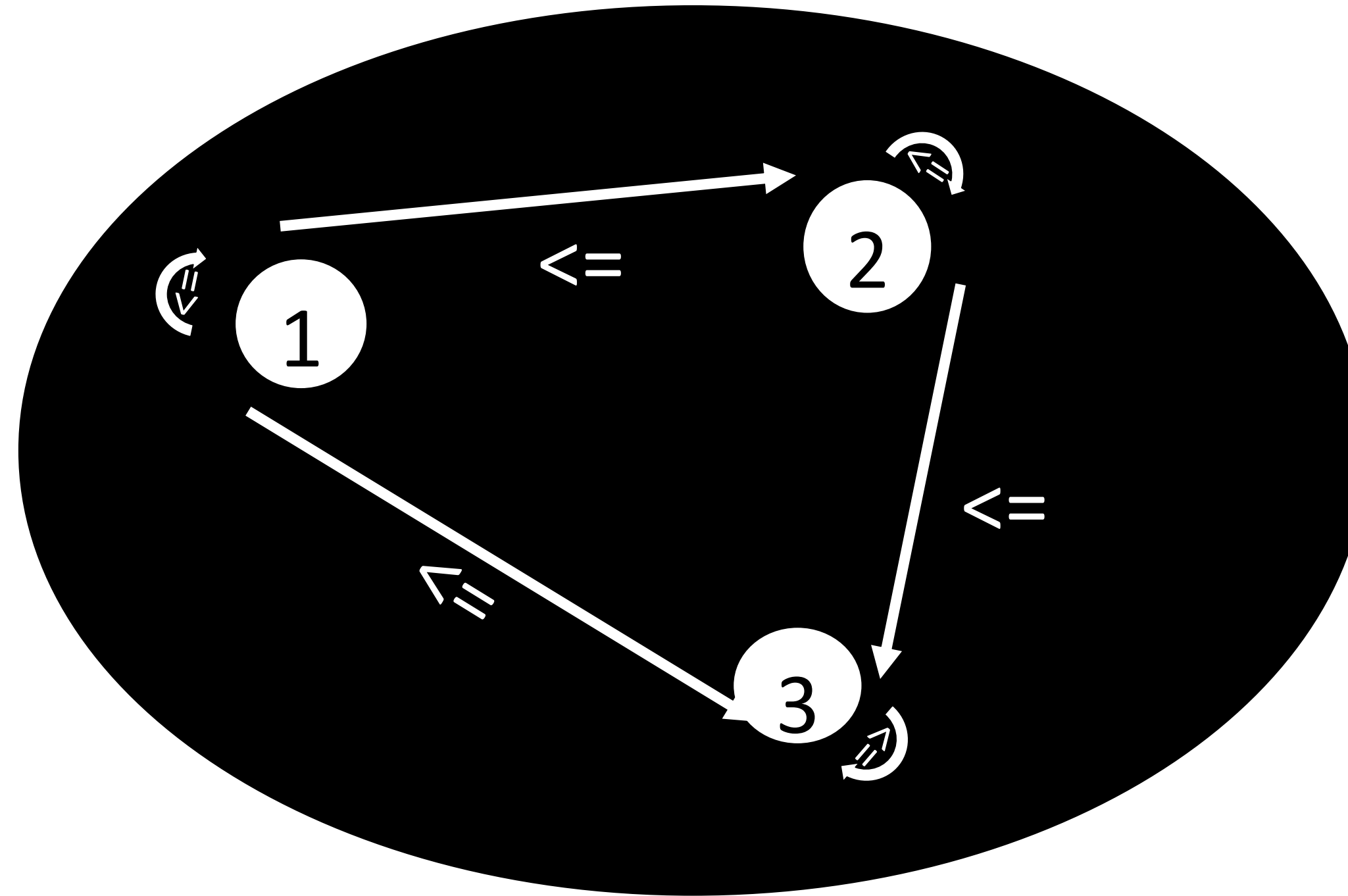
Surjection (Onto)



Bijection (One-to-One and Onto)

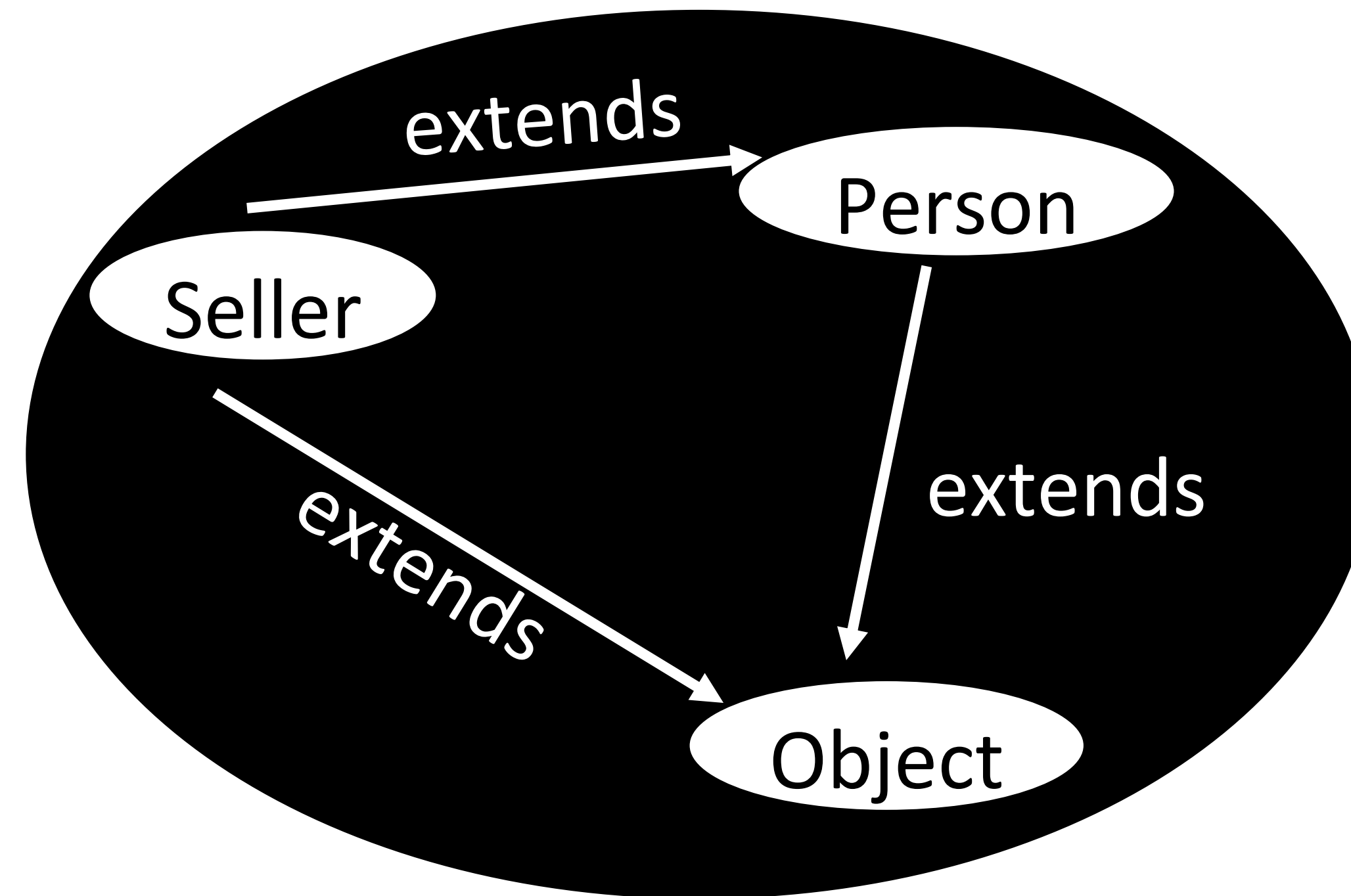


Category you already know



Ordered set

Category you already know



Class hierarchy