

Introduction to Program Synthesis (WS 2024/25)

Chapter 3.3 - Traditional Methodologies (Enumerative and Stochastic search)

Dr. rer. nat. Roman Kalkreuth

Chair for AI Methodology (**AIM**), Department of Computer Science,
RWTH Aachen University, Germany



Center for
Artificial Intelligence

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Traditional Methodologies

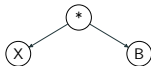
Traditional Methodologies: Stochastic HC

- ▶ Basic search methods for SR → **Stochastic hill climbing** and **enumerative branching**
- ▶ Stochastic HC → Requires a neighbourhood function \mathcal{N}
- ▶ Adapting enumerative search to SR → Depth-first search with **deep cloning**
 - ↪ Consider neighbours sampled by \mathcal{N} as potential next steps

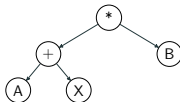
Traditional Methodologies

Traditional Methodologies: Stochastic HC

- **Neighbourhood function** $\mathcal{N} \rightarrow$ replace a leaf node with a randomly generated subtree: $\mathcal{N}(\Psi) \mapsto \Psi'$



$$\Psi = X * B$$



$$\Psi' = (A + X) * B$$

- Stochastic hill climbing with restart is used as a search algorithm:
 - ~ $\mathcal{R} \rightarrow$ **Replacement function**: Selects a neighbour **with better cost value uniformly at random**
 - ~ If no neighbour has a better cost value \rightarrow Return and keep \mathcal{P} (no replacement occurs)

Traditional Methodologies

Traditional Methodologies: Stochastic HC

Algorithm Stochastic Hill Climbing (S-HC)

```
1: while  $t < trials$  do
2:    $t \leftarrow i \leftarrow x \leftarrow 0$ 
3:    $\mathcal{P} \sim U$ 
4:    $\mathcal{C}(\mathcal{P})$ 
5:   while  $i < iterations$  do
6:     while  $x < neighbours$  do
7:        $\mathcal{Q} \leftarrow \mathcal{N}(\mathcal{P}, n)$ 
8:       Prediction( $\mathcal{Q}$ )
9:       Evaluation( $\mathcal{Q}, \mathcal{C}$ )
10:       $\mathcal{P} \leftarrow \mathcal{R}(\mathcal{P}, \mathcal{Q})$ 
11:       $x \leftarrow x + 1$ 
12:    end while
13:     $i \leftarrow i + 1$ 
14:  end while
15:   $t \leftarrow t + 1$ 
16: end while
17: return  $\mathcal{P}$ 
```

▷ Sample initial point

▷ Calculate cost of \mathcal{P}

▷ Sample n neighbours

▷ Make predictions

▷ Evaluate neighbours

▷ Replace with better partial solution or keep \mathcal{P}

Traditional Methodologies

Traditional Methodologies: Steepest HC

- ▶ Steepest HC → Modification of stochastic HC
- ▶ Selects the best (steepest) improvement among the sampled neighbours
- ▶ Suitable for landscapes that are well suited for *greedy* approaches

Traditional Methodologies

Traditional Methodologies: Improvement Strategies

- ▶ Hill climbing is prone to getting stuck in **local optima**
- ▶ Requires advanced strategies to prevent **search stagnation**:
 - ~> **Backtracking** → Rejection of infeasible partial solutions
 - ~> **Backjumping** → Stochastic roll-backs in the space of partial solutions
 - ~> **Population-based approach** → Simultaneous evolving partial solutions
- ~> Enumerative search → Expensive in large symbolic search spaces
 - ~> **Branch and Bound** → Reject solution that are not located in the range of predefined lower/upper bounds
 - ~> Specific **pre-evaluation** of candidate solution
 - ~> Can be integrated into **enumerative recursive search** (ERS)

Traditional Methodologies

Traditional Methodologies: Branch and Bound

- ▶ Branch and bound prunes the search space
 - ↪ Clipping partial solutions within semantic boundaries
- ▶ Rejection of infeasible solutions
- ▶ Assumption that following feasible search trajectories is more effective

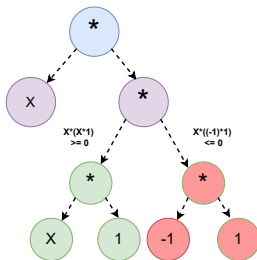


Figure: Branch-and-Bound example: Let $f(x) = x^2 \mapsto \mathbb{R}_{\geq 0}, x \in \mathbb{R}$ be the objective function. Candidate solutions are bounded in $\mathbb{R}_{\geq 0}$. The red branch will be rejected since the evaluation will lead to negative semantic values in \mathbb{R}^- .

Traditional Methodologies

Traditional Methodologies: Branch and Bound

Algorithm ERS with Branch and Bound

Arguments

S_n : Initial or intermediate state

\mathcal{P} : Problem instance

\mathcal{B}_l : Lower bound

\mathcal{B}_u : Upper bound

\mathcal{G} : Goal state

s : Number of steps

Return

S_g : Goal or base case

```
1: procedure ERS( $S_n$ ,  $\mathcal{P}$ ,  $\mathcal{B}_{lower}$ ,  $\mathcal{B}_{upper}$ ,  $n_{neighbour}$ )
2:    $\Omega \leftarrow \mathcal{N}(s)$  ▷ Determine the neighbourhood
3:   for  $\omega$  in  $\Omega$  do ▷ Consider each neighbour
4:      $p_\omega \leftarrow \text{Prediction}(\omega)$  ▷ Make prediction
5:      $c_\omega \leftarrow \mathcal{C}(\omega, \mathcal{P})$  ▷ Calculate cost
6:      $b_{lower}, b_{upper} \leftarrow \text{boundaries}(c_\omega)$  ▷ Determine the bounds
7:     if  $\mathcal{B}_l \leq b_{lower}$  and  $\mathcal{B}_u \geq b_{upper}$  then ▷ Check against the global bounds
8:       if  $S_n == \mathcal{G}$  then ▷ If goal state is reached
9:          $\text{output}(S_n)$  ▷ Output the solution
10:      else ▷ Consider the current step as next one
11:         $S_{n+1} \leftarrow (\omega, c)$ 
12:        ERS( $S_{n+1}$ , ...) ▷ Trigger next recursion step
13:      end if
14:    end if
15:  end for
16: end procedure
```
