

Задачи разрешимости логических формул и приложения

Лекция 12. Анализ программ

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Пример

```
1 void ReadBlocks(int data[], int cookie)
2 {
3     int i = 0;
4     while (true)
5     {
6         int next;
7         next = data[i];
8         if (!(i < next && next < N)) return;
9         i = i + 1;
10        for (; i < next; i = i + 1) {
11            if (data[i] == cookie)
12                i = i + 1;
13            else
14                Process(data[i]);
15        }
16    }
17 }
```

Line	Kind	Instruction or condition
3	Assignment	<code>i = 0;</code>
7	Assignment	<code>next = data[i];</code>
8	Branch	<code>i < next && next < N</code>
9	Assignment	<code>i = i + 1;</code>
10	Branch	<code>i < next</code>
11	Branch	<code>data[i] != cookie</code>
14	Function call	<code>Process(data[i]);</code>
10	Assignment	<code>i = i + 1;</code>
10	Branch	<code>!(i < next)</code>
7	Assignment	<code>next = data[i];</code>
8	Branch	<code>!(i < next && next < N)</code>

Static single assignment

Line	Kind	Instruction or condition
3	Assignment	$i_1 = 0;$
7	Assignment	$next_1 = data_0[i_1];$
8	Branch	$i_1 < next_1 \ \&\& \ next_1 < N_0$
9	Assignment	$i_2 = i_1 + 1;$
10	Branch	$i_2 < next_1$
11	Branch	$data_0[i_2] \neq cookie_0$
14	Function call	$Process(data_0[i_2]);$
10	Assignment	$i_3 = i_2 + 1;$
10	Branch	$!(i_3 < next_1)$
7	Assignment	$next_2 = data_0[i_3];$
8	Branch	$!(i_3 < next_2 \ \&\& \ next_2 < \underline{N_0})$

Path constraint

$$\begin{array}{ll} i_1 = 0 & \wedge \\ next_1 = data_0[i_1] & \wedge \\ (i_1 < next_1 \wedge next_1 < N_0) & \wedge \\ i_2 = i_1 + 1 & \wedge \\ i_2 < next_1 & \wedge \\ data_0[i_2] = cookie_0 & \wedge \\ i_3 = i_2 + 1 & \wedge \\ i_4 = i_3 + 1 & \wedge \\ \neg(i_4 < next_1) & \wedge \\ \neg(0 \leq i_4 \wedge i_4 < N_0) & \end{array}$$

Пример

```
1 void ReadBlocks(int data[], int cookie)
2 {
3     int i = 0;
4     while (true)
5     {
6         int next;
7         next = data[i];
8         if (!(i < next && next < N)) return;
9         i = i + 1;
10        for (; i < next; i = i + 1) {
11            if (data[i] == cookie)
12                i = i + 1;
13            else
14                Process(data[i]);
15        }
16    }
17 }
```

$$i_1 = 0 \wedge \neg(0 \leq i_1 \wedge i_1 < N_0)$$

$$i_1 = 0 \wedge \neg(0 \leq i_1 \wedge i_1 < N_0) \\ \{i_1 := 0, N_0 := 0\}$$

Line	Kind	Instruction or condition
3	Assignment	<code>i = 0;</code>
7	Assignment	<code>next = data[i];</code>
8	Branch	<code>i < next && next < N</code>
9	Assignment	<code>i = i + 1;</code>
10	Branch	<code>i < next</code>
11	Branch	<code>data[i] = cookie</code>
12	Assignment	<code>i = i + 1;</code>
10	Assignment	<code>i = i + 1;</code>
10	Branch	<code>!(i < next)</code>
7	Assertion	<code>0 <= i && i < N</code>

Path constraint

$$\begin{array}{ll} i_1 = 0 & \wedge \\ next_1 = data_0[i_1] & \wedge \\ (i_1 < next_1 \wedge next_1 < N_0) & \wedge \\ i_2 = i_1 + 1 & \wedge \\ i_2 < next_1 & \wedge \\ data_0[i_2] = cookie_0 & \wedge \\ i_3 = i_2 + 1 & \wedge \\ i_4 = i_3 + 1 & \wedge \\ \neg(i_4 < next_1) & \wedge \\ \neg(0 \leq i_4 \wedge i_4 < N_0) & \end{array}$$

$$\{i_1 \mapsto 0, N_0 \mapsto 3, next_1 \mapsto 2, data_0 \mapsto \langle 2, 6, 5 \rangle, \\ i_2 \mapsto 1, cookie_0 \mapsto 6, i_3 \mapsto 2, i_4 \mapsto 3\},$$

Bounded Program Analysis

```
1 if (i < next) {
2   if (data[i] == cookie)
3     i = i + 1;
4   else
5     Process(data[i]);
6
7   i = i + 1;
8
9   if (i < next) {
10    if (data[i] == cookie)
11      i = i + 1;
12    else
13      Process(data[i]);
14
15    i = i + 1;
16  }
17 }
```

```
1  $\gamma_1 = (i_0 < next_0);$ 
2  $\gamma_2 = (data_0[i_1] == cookie_0);$ 
3  $i_1 = i_0 + 1;$ 
4
5
6  $i_2 = \gamma_2 ? i_1 : i_0;$ 
7  $i_3 = i_2 + 1;$ 
8
9  $\gamma_3 = (i_3 < next_0);$ 
10  $\gamma_4 = (data_0[i_3] == cookie_0);$ 
11  $i_4 = i_3 + 1;$ 
12
13
14  $i_5 = \gamma_4 ? i_4 : i_3;$ 
15  $i_6 = i_5 + 1;$ 
16  $i_7 = \gamma_3 ? i_6 : i_3;$ 
17  $i_8 = \gamma_1 ? i_7 : i_0;$ 
```

Unbounded Program Analysis

```
1 int i=0, j=0;
2
3 while(data[i] != '\n')
4 {
5     i++;
6     j=i;
7 }
8
9 assert(i == j);
```

→

```
1 int i=0, j=0;
2
3 if(data[i] != '\n')
4 {
5     i=*;
6     j=*;
7     i++;
8     j=i;
9 }
10
11 assume(data[i] == '\n')
12
13 assert(i == j);
```

Path constraint

$$i_1 = 0 \quad \wedge$$

$$j_1 = 0 \quad \wedge$$

$$\gamma_1 = (data_0[i_1] \neq '\backslash n') \quad \wedge$$

$$i_3 = i_2 + 1 \quad \wedge$$

$$j_3 = i_3 \quad \wedge$$

$$i_4 = \gamma_1 ? i_3 : i_1 \quad \wedge$$

$$j_4 = \gamma_1 ? j_3 : j_1 \quad \wedge$$

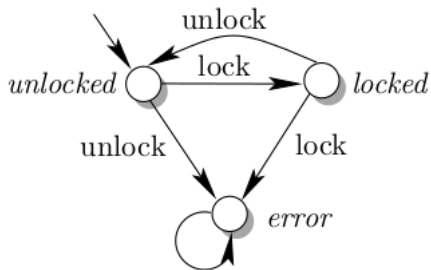
$$data_0[i_4] = '\backslash n' \quad \wedge$$

$$i_4 \neq j_4$$

Неудачный пример аппроксимации

```
1  do {
2      lock();
3      old_count = count;
4      request = GetNextRequest();
5      if (request != NULL) {
6          ReleaseRequest(request);
7          unlock();
8          ProcessRequest(request);
9          count = count + 1;
10     }
11 }
12 while(old_count != count);
13 unlock();
```


Конечный автомат механизма взятия лока



Добавление assert-ов

```
1  state_of_lock = unlocked;
2  do {
3      assert(state_of_lock == unlocked);
4      state_of_lock = locked;
5      old_count = count;
6      request = GetNextRequest();
7      if (request != NULL) {
8          ReleaseRequest(request);
9          assert(state_of_lock == locked);
10         state_of_lock = unlocked;
11         ProcessRequest(request);
12         count = count + 1;
13     }
14 }
15 while (old_count != count);
16 assert(state_of_lock == locked);
17 state_of_lock = unlocked;
```

```
1 state_of_lock = unlocked;
2
3 state_of_lock = *;
4 old_count = *;
5 count = *;
6 request = *;
7
8 assert(state_of_lock == unlocked);
9 state_of_lock = locked;
10 old_count = count;
11 request = GetNextRequest();
12 if (request != NULL) {
13     ReleaseRequest(request);
14     assert(state_of_lock == locked);
15     state_of_lock = unlocked;
16     ProcessRequest(request);
17     count = count + 1;
18 }
19
20 assume(old_count == count);
21
22 assert(state_of_lock == locked);
23 state_of_lock = unlocked;
```

```

1 state_of_lock = unlocked;
2
3 assert(state_of_lock == unlocked); // induction base case
4
5 state_of_lock = *;
6 old_count = *;
7 count = *;
8 request = *;
9
10 assume(state_of_lock == unlocked); // induction hypothesis
11
12 assert(state_of_lock == unlocked); // property
13 state_of_lock = locked;
14 old_count = count;
15 request = GetNextRequest();
16 if (request != NULL) {
17     ReleaseRequest(request);
18     assert(state_of_lock == locked); // property
19     state_of_lock = unlocked;
20     ProcessRequest(request);
21     count = count + 1;
22 }
23
24 // induction step case
25 assert(old_count != count  $\implies$  state_of_lock == unlocked);
26
27 assume(old_count == count);
28
29 assert(state_of_lock == locked); // property
30 state_of_lock = unlocked;

```

