

SAT/SMT solvers 10. Program Analysis

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Lalambda

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Why

- automatic test generation
- detection of dead code
- verification of properties given in the form of assertions

Example

```
void ReadBlocks(int data[], int cookie)
2 {
    int i = 0;
   while (true)
5
      int next;
      next = data[i];
      if (!(i < next && next < N)) return;
8
     i = i + 1;
     for (; i < next; i = i + 1) {
10
          if (data[i] == cookie)
11
             i = i + 1;
12
          else
13
             Process(data[i]);
14
15
16
17 }
```

Example

Line	Kind	Instruction or condition
3	Assignment	i = 0;
3 7 8 9	Assignment	<pre>next = data[i];</pre>
8	Branch	i < next && next < N
9	Assignment	i = i + 1;
10	Branch	i < next
11	Branch	data[i] != cookie
14	Function call	Process(data[i]);
10	Assignment	i = i + 1;
10 10 7	Branch	!(i < next)
	Assignment	<pre>next = data[i];</pre>
8	Branch	!(i < next && next < N)

Static single assignment

Line	Kind	Instruction or condition
3	Assignment	$i_1 = 0;$
7	Assignment	$next_1 = data_0[i_1];$
8	Branch	$i_1 < \text{next}_1 \&\& \text{next}_1 < N_0$
9	Assignment	$i_2 = i_1 + 1;$
10	Branch	$i_2 < next_1$
11	Branch	data ₀ [i ₂] != cookie ₀
14	Function call	Process(data ₀ [i ₂]);
10	Assignment	$i_3 = i_2 + 1;$
10	Branch	$!(i_3 < next_1)$
3 7 8 9 10 11 10 10 10 10 10 10 10 10	Assignment	$next_2 = data_0 [i_3];$
8	Branch	$!(i_3 < next_2 \&\& next_2 < N_0)$

Path constraint

$$\begin{array}{lll} i_1 = 0 & & \wedge \\ next_1 = data_0[i_1] & & \wedge \\ (i_1 < next_1 \ \wedge \ next_1 < N_0) & \wedge \\ i_2 = i_1 + 1 & & \wedge \\ i_2 < next_1 & & \wedge \\ data_0[i_2] = cookie_0 & & \wedge \\ i_3 = i_2 + 1 & & \wedge \\ i_4 = i_3 + 1 & & \wedge \\ \neg (i_4 < next_1) & & \wedge \\ \neg (0 \leq i_4 \ \wedge \ i_4 < N_0) & & \end{array}$$

Example

```
void ReadBlocks(int data[], int cookie)
2 {
    int i = 0;
   while (true)
5
      int next;
      next = data[i];
      if (!(i < next && next < N)) return;
8
     i = i + 1;
     for (; i < next; i = i + 1) {
10
          if (data[i] == cookie)
11
             i = i + 1;
12
          else
13
             Process(data[i]);
14
15
16
17 }
```

Error

$$\textit{i}_1 = 0 \land \neg (0 \leq \textit{i}_1 \land \textit{i}_1 < \textit{N}_0)$$

Error

$$i_1 = 0 \land \neg (0 \le i_1 \land i_1 < N_0)$$

 $\{i_1 := 0, N_0 := 0\}$

Another trace

Line	Kind	Instruction or condition
3 7 8 9	Assignment	i = 0;
7	Assignment	<pre>next = data[i];</pre>
8	Branch	i < next && next < N
9	Assignment	i = i + 1;
10	Branch	i < next
11	Branch	data[i] = cookie
12 10	Assignment	i = i + 1;
10	Assignment	i = i + 1;
10	Branch	!(i < next)
7	Assertion	0 <= i && i < N

Another trace

$$\begin{array}{lll} i_1 = 0 & & \wedge \\ next_1 = data_0[i_1] & & \wedge \\ (i_1 < next_1 \ \wedge \ next_1 < N_0) & \wedge \\ i_2 = i_1 + 1 & & \wedge \\ i_2 < next_1 & & \wedge \\ data_0[i_2] = cookie_0 & & \wedge \\ i_3 = i_2 + 1 & & \wedge \\ i_4 = i_3 + 1 & & \wedge \\ \neg (i_4 < next_1) & & \wedge \\ \neg (0 \leq i_4 \ \wedge \ i_4 < N_0) & & \end{array}$$

Assigment of input

$$\{i_1 \mapsto 0, \ N_0 \mapsto 3, \ next_1 \mapsto 2, \ data_0 \mapsto \langle 2, 6, 5 \rangle,$$

 $i_2 \mapsto 1, \ cookie_0 \mapsto 6, \ i_3 \mapsto 2, \ i_4 \mapsto 3 \},$

Bounded Program Analysis

```
1 if (i < next) {
  if (data[i] == cookie)
   i = i + 1;
3
4 else
   Process (data[i]);
6
   i = i + 1;
8
    if (i < next) {
    if (data[i] == cookie)
10
      i = i + 1;
11
   else
12
13 Process(data[i]);
14
    i = i + 1;
15
16 }
17 }
```

```
_{1} \gamma_{1} = (i_{0} < \text{next}_{0});
 _{2} \gamma_{2} = (data<sub>0</sub>[i<sub>1</sub>] == cookie<sub>0</sub>);
 3 i_1 = i_0 + 1;
 4
 6 i_2 = \gamma_2 ? i_1 : i_0;
 7 i_3 = i_2 + 1:
 9 \gamma_3 = (i_3 < next_0);
\gamma_4 = (data_0[i_3] == cookie_0);
11 i_4 = i_3 + 1;
12
13
14 i_5 = \gamma_4 ? i_4 : i_3;
15 i_6 = i_5 + 1;
16 i_7 = \gamma_3 ? i_6 : i_3;
17 i_8 = \gamma_1 ? i_7 : i_0;
```

Unbounded Program Analysis

```
i = 0, j = 0;
                               2
i = 0, j = 0;
                               3 if (data[i] != '\n')
                               4 {
2
3 while(data[i] != '\n')
                               i = *;
                               j = *;
4 {
                               7 i++;
5 i++;
6 j=i;
                               s j=i;
7 }
                               9 }
8
                               10
9 assert(i == j);
                               11 assume(data[i] == '\n')
                               12
                               13 assert (i == j);
```

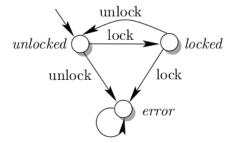
Path constraint

$$\begin{array}{lll} i_1 = 0 & & \wedge \\ j_1 = 0 & & \wedge \\ \gamma_1 = (data_{\theta}[i_1] \neq ' \setminus \texttt{n'}) & \wedge \\ i_3 = i_2 + 1 & & \wedge \\ j_3 = i_3 & & \wedge \\ i_4 = \gamma_1 ? i_3 : i_1 & & \wedge \\ j_4 = \gamma_1 ? j_3 : j_1 & & \wedge \\ data_{\theta}[i_4] = ' \setminus \texttt{n'} & & \wedge \\ i_4 \neq j_4 & & \end{array}$$

Bad approximation

```
do {
1
       lock();
2
       old_count = count;
3
       request = GetNextRequest();
4
       if (request != NULL) {
5
          ReleaseRequest (request);
6
          unlock();
7
          ProcessRequest (request);
8
          count = count + 1;
9
10
11
     while (old_count != count);
12
     unlock();
13
```

Automat of lock



Asserts

```
state_of_lock = unlocked;
1
     do {
2
          assert(state_of_lock == unlocked);
3
           state_of_lock = locked;
4
          old_count = count;
5
          request = GetNextRequest();
6
           if (request != NULL) {
7
              ReleaseRequest (request);
              assert (state of lock == locked);
9
              state_of_lock = unlocked;
10
              ProcessRequest (request);
11
              count = count + 1;
12
13
14
     while (old_count != count);
15
     assert(state_of_lock == locked);
16
     state_of_lock = unlocked;
17
```

Invariants

Invariants

```
\begin{array}{ll} {}_1A;\\ {}_2\text{ while}(C) & \{\\ {}_3& \text{assert}(I);\\ {}_4& B;\\ {}_5\end{array}\}
```

```
old_count = *;
4
     count = *;
5
     request = *;
6
7
     assert(state_of_lock == unlocked);
8
     state_of_lock = locked;
9
     old count = count;
10
     request = GetNextRequest();
11
     if (request != NULL) {
12
       ReleaseRequest (request);
13
       assert(state_of_lock == locked);
14
       state_of_lock = unlocked;
15
       ProcessRequest (request);
16
       count = count + 1;
17
18
19
     assume(old_count == count);
20
21
     assert(state_of_lock == locked);
22
     state_of_lock = unlocked;
23
                                               <● > (車) (車) (車) のQで
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```

state_of_lock = unlocked;

state of lock = *;

1 2

3

```
state_of_lock = unlocked;
2
     assert(state_of_lock == unlocked); // induction base case
3
4
     state of lock = \star:
5
     old count = *;
6
     count = *;
7
     request = *;
8
9
     assume(state of lock == unlocked); // induction hypothesis
10
11
     assert (state of lock == unlocked); // property
12
     state_of_lock = locked;
13
     old count = count;
14
     request = GetNextRequest();
15
     if (request != NULL) {
16
       ReleaseRequest (request);
17
       assert(state of lock == locked); // property
18
       state_of_lock = unlocked;
19
       ProcessRequest (request);
20
       count = count + 1;
21
22
23
     // induction step case
24
25
     assert (old_count != count \improx state_of_lock == unlocked);
26
     assume(old_count == count);
27
28
     assert(state_of_lock == locked); // property
29
     state_of_lock = unlocked;
30
```

