

4.3 Newton interpolation:

Recall Lagrange interpolation:

$$p(u) = \sum_{i=0}^n \alpha_i L_i^n(u)$$

such that $L_i^n(u_j) = \delta_{ij}$

consequence: collocation matrix $\bar{\Phi} = \text{Id}$
thus $\alpha_i = p_i$

The idea of Newton interpolation is to use basis functions $P_0(u), \dots, P_n(u)$ with $P_i(u)$ being of degree i with roots at u_0, \dots, u_{i-1} :

I.e. $P_0(u) = 1$

$$P_i(u) = (u - u_0)(u - u_1) \dots (u - u_{i-1})$$

Consequently, the collocation matrix is lower triangular

$$\bar{\Phi} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ | & * & 0 & \dots & 0 \\ | & | & * & \dots & 0 \\ | & | & | & \dots & 0 \\ 1 & * & * & \dots & * \end{bmatrix}$$

Example 34: (cf. Example 33)

u_i	0	1	2
p_i	2	4	3

• 1 point interpolation:

$$p(u_0) = p_0, \quad \text{degree}(p) = 0$$

$$\text{then } p(u) = \alpha P_0(u) = \alpha \cdot 1 = \alpha \stackrel{!}{=} p_0$$

$$\text{so } \boxed{p(u_0) = p_0 = 2.}$$

In Newton interpolation we introduce space notation. For α we write $p[u_0]$, so $p(u) = p[u_0] P_0$

Here: $p[u_0] = 2$, so $\boxed{p(u) = 2 \cdot 1 = 2}$.

• Now: 2 point interpolation:

The weights are denoted $p[u_0]$ and $p[u_0, u_1]$. So:

$$\text{Find } p(u) = p[u_0] P_0(u) + p[u_0, u_1] P_1(u)$$

such that $\left. \begin{array}{l} p(u_0) = P_0 \\ p(u_1) = P_1 \end{array} \right\} \text{interpolation conditions.}$

The first condition yields

$$\begin{aligned} P_0 &= p[u_0] P_0(u_0) + p[u_0, u_1] \underbrace{P_1(u_0)}_{=0 \text{ by construction}} \\ &\downarrow \\ &= p[u_0] P_0(u_0) \end{aligned}$$

But this is fulfilled due to the 1-point interpol.

The 2nd interpolation condition yields:

$$P_1 = p(u_1) = p[u_0] P_0(u_1) + p[u_0, u_1] P_1(u_1)$$

$$\Leftrightarrow P_1 = p[u_0] \cdot 1 + p[u_0, u_1] (u_1 - u_0)$$

$$\Rightarrow 4 = 2 \cdot 1 + p[u_0, u_1] (1 - 0)$$

$$\Leftrightarrow p[u_0, u_1] = \frac{4-2}{1-0} = \underline{\underline{2}}$$

In summary:

$$p(u) = 2 + 2 P_1(u) = \boxed{2 + 2u}$$

is the 2 point interpolation polynomial.

• 3 point interpolation:

$$\begin{aligned} \text{Find } p(u) &= p[u_0] P_0(u) + p[u_0, u_1] P_1(u) \\ &\quad + p[u_0, u_1, u_2] P_2(u) \end{aligned}$$

Here, the first and 2nd terms have been determined by the 1 point and 2 point interpol.

$$\dots + p[u_2, u_1, u_0] P_2(u)$$

$$\text{thus, } p(u) = 2 + 2u + p[u_0, u_1, u_2](u-u_0)(u-u_1) \\ = 2 + 2u + p[u_0, u_1, u_2](u-0)(u-1)$$

Use the 3rd interpolation condition to solve for

$$p[u_0, u_1, u_2]: p(u_2) = p_2$$

$$\text{So: } p_2 = 2 + 2u_2 + p[u_0, u_1, u_2](u_2 - u_0)(u_2 - u_1)$$

$$\Rightarrow 3 = 2 + 2 \cdot 2 + p[u_0, u_1, u_2](2-0)(2-1)$$

$$\Leftrightarrow \frac{3-6}{2} = p[u_0, u_1, u_2]$$

To summarise: We get the interpolation poly

$$p(u) = 2 + 2u - \frac{3}{2}u(u-1)$$

$$\Leftrightarrow p(u) = 2P_0(u) + 2P_1(u) - \frac{3}{2}P_2(u)$$

The resulting polynomial is the same as for Lagrange interpolation and for the basis $1, x, x^2, \dots$

The weights $p[u_0, \dots, u_i]$ are obtained by back substitut in the lower triangular matrix.

We obtain (proof: no credit exercise)

$$p[u_0, \dots, u_i] = \sum_{k=0}^i \frac{p_k}{\prod_{\substack{j=0 \\ j \neq k}}^i (u_k - u_j)}$$

which can be computed by the following recursion

$$p[u_0] = p_0$$

$$p[u_0, u_1] = \frac{p_0}{u_0 - u_1} + \frac{p_1}{u_1 - u_0}$$

$$\vdots \quad \vdots \quad \vdots \quad p[u, \dots] = p[u_0]$$

$$= \frac{p_1 - p_0}{u_1 - u_0} = \frac{p(u_1) - p(u_0)}{u_1 - u_0}$$

and in general

$$p[u_0, \dots, u_i] = \frac{p[u_1, \dots, u_i] - p[u_0, \dots, u_{i-1}]}{u_i - u_0}$$

The coefficients are called divided differences of order i at nodes u_0, \dots, u_i .

Example 35: (cf. Example 34)

u_i	p_i	$p[u_i]$	$p[u_i, u_{i+1}]$	$p[u_i, u_{i+1}, u_{i+2}]$
0	2	2	$\frac{4-2}{1-0} = 2$	$\frac{-1-2}{2-0} = -\frac{3}{2}$
1	4	4	$\frac{3-4}{2-1} = -1$	
2	3	3		

The resulting coefficients result as the top row of this scheme, so $2, 2, -\frac{3}{2}$.

Note that the final divided differences are invariant under permutations of the nodes.

4.4 Error analysis:

All of the previous approaches lead to the same polynomial. Just the algorithm differs.

Theorem 36: Let $f \in C^{n+1}([a, b])$ and $p(u)$ a polynomial of degree $\leq n$ that interpolates f at $n+1$ distinct nodes $u_0, \dots, u_n \in [a, b]$. Then for any $u \in [a, b]$ there exists a $\xi \in (a, b)$ such that

$$f(u) - p(u) = \frac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n (u - u_i)$$

$$f^{(n)}(x) = \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{i=0}^n (x - x_i)$$

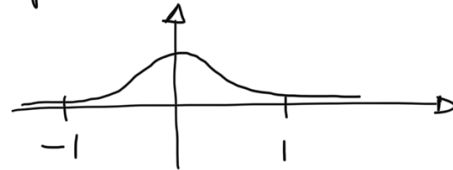
For equidistant nodes $x_i = a + i \cdot \frac{b-a}{n}$ we have

$$|f(x) - p(x)| \leq \frac{1}{4(n+1)} M \left(\frac{b-a}{n} \right)^{n+1}$$

where $M = \max_{[a,b]} |f^{(n+1)}(x)|$

Example 37: Runge's phenomenon

$$f(x) = \frac{1}{1+x^2}$$



then, for equidistant nodes:

$$\lim_{n \rightarrow \infty} \max_{[-1,1]} |f(x) - p(x)| = \infty$$

The interpolation polynomials oscillate and don't provide meaningful interpolations.

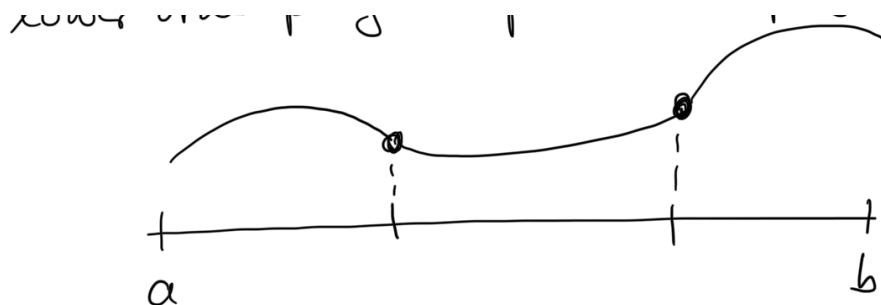
- Higher order polynomial interpolation on equidistant nodes can be problematic
- Use non-equidistant points, e.g.

Chebyshev nodes $\cos\left[\left(\frac{i}{n}\right)\pi\right], i=0, \dots, n$
over $[-1, 1]$.

4.5 Piecewise polynomial interpolation

Another remedy of the problems of high order poly interpolation is piecewise polynomial interpolation.

Subdivide $[a,b]$ into smaller pieces on which lower order poly interpolation is performed.



Problem encountered: Pieces of interpolation fit together continuously but not differentiable.

Solution: Use Hermite interpolation in which additional constraints for continuity of derivatives is used (see homework #4)
