

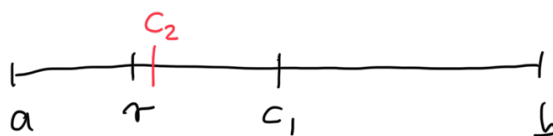
Bisection method:

Producing a sequence of intervals  $[a_i, b_i]$  such that the root  $r$  is inside of these interval.  
 Given is a starting interval  $[a, b] = [a_0, b_0]$ .

Theorem 22: The bisection method, when applied to an interval  $[a, b]$  and  $f \in C^0([a, b])$  with  $f(a)f(b) < 0$  after  $n$  steps has computed an approximation  $c_n$  of the root  $r$  with an error

$$|r - c_n| < \frac{b-a}{2^n}.$$

Proof:



we have

$$|r - c_1| < |a - c_1| = \frac{b-a}{2}$$

and then

$$|r - c_2| < |a - c_2| = \frac{b-a}{4}$$

so iteratively, we get

$$|r - c_n| < \frac{b-a}{2^n} \quad \square$$

Example 23: Let  $[a, b] = [0, 1]$ . How many iterations are needed to decrease the error below  $2^{-20}$ ?

Find  $n \in \mathbb{N}$  such that

$$\frac{b-a}{2^n} < 2^{-20}$$

$$\Leftrightarrow \frac{b-a}{2^{-20}} < 2^n$$

$$\Rightarrow \log_2\left(\frac{b-a}{2^{-20}}\right) < n$$

↑  
because  $\log_2$   
strictly increasing

Here:  $\log_2\left(\frac{1}{2^{-20}}\right) < n$

$$\Leftrightarrow \log_2(1) - \log_2(2^{-20}) < n$$

$$\Leftrightarrow \boxed{n > 20}$$

This is slow convergence: For every iteration step we get one binary digit as accuracy increase. Recall, for the approximation of  $\cos(x)$  by its Taylor poly, we got 2 decimal digits per ten of the poly.

Def 24: Suppose the sequence  $x_n$  converges to  $\tau$  as  $n \rightarrow \infty$ . Then, the sequence

(a) converges linearly to  $\tau$ , if  $\exists M \in (0, 1)$  such that

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - \tau|}{|x_n - \tau|} = M$$

For bisection method,  
 $M = 1/2$

(b) converges super-linearly to  $\tau$  if

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - \tau|}{|x_n - \tau|} = 0$$

(c) converges sub-linearly to  $\tau$  if

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - \tau|}{|x_n - \tau|} = 1$$

(d) converges with order  $q > 1$  if  $\exists M > 0$  s.th

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - r|}{|x_n - r|^q} < M$$

if  $q=2$  we call the convergence quadratic  
 $q=3$  ————— " ————— cubic

### 3.2 Newton's method:

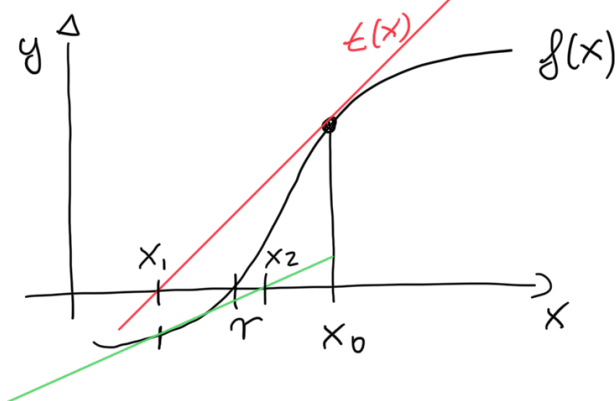
Let  $f \in C'([a, b])$ , then at every  $x_0 \in (a, b)$  there exists a tangent to the graph of  $f$ :

tangent:  $t(x) = f(x_0) + f'(x_0)(x - x_0)$

the tangent at  $x_0$  is the first order Taylor poly

obviously:  $t(x_0) = f(x_0)$

$$t'(x) = f'(x_0)$$



use  $t(x)$  as an approximation to  $f$  whose root can be easily found:

$$0 = t(x_1) = f(x_0) + f'(x_0)(x_1 - x_0)$$

$$\Leftrightarrow \boxed{x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}}$$

Iterating this yields the Newton sequence (Newton-Raphson iteration)

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}, \quad x_0 = \text{initial guess}$$

Theorem 25: (cf. Calculus I/II notes)

When the Newton sequence converges, it converges to a root of  $f$ .

Theorem 26: (cf. Calculus I/II)

Let  $f \in C'([a, b])$  fulfill:

- (1)  $f(a)f(b) < 0$
- (2)  $f$  has no critical pt in  $(a, b)$
- (3)  $f''$  exists, is continuous and either  $f'' > 0$  OR  $f'' < 0$  in whole  $(a, b)$

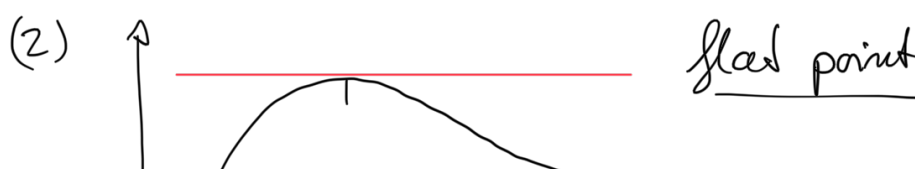
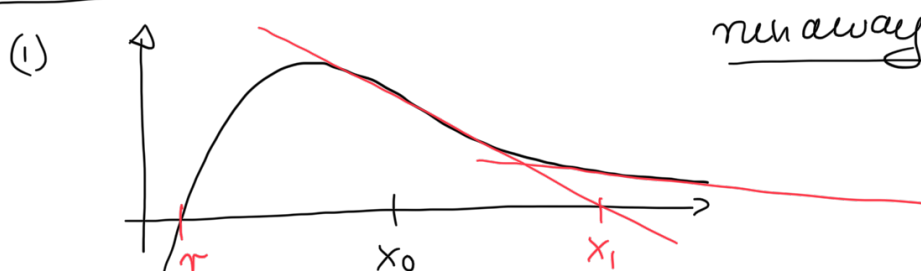
Then  $f(x) = 0$  has exactly one solution  $\tau$ . The Newton sequence always converges to  $\tau$  as  $n \rightarrow \infty$ , when the initial guess is chosen accord to:

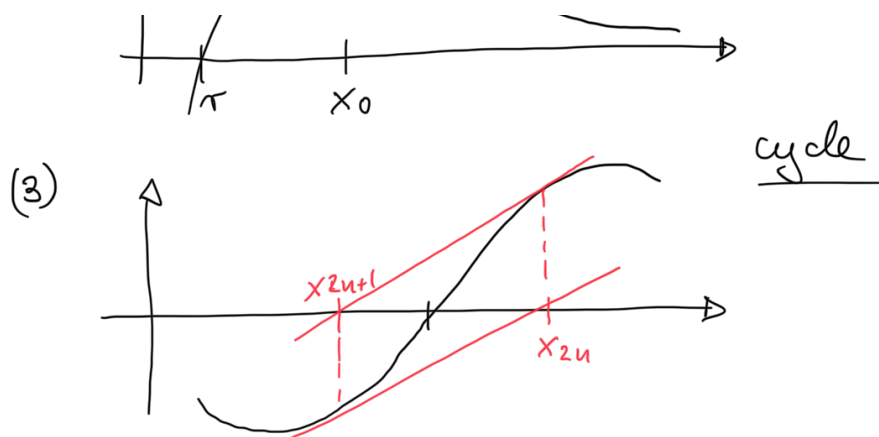
- if  $f(a) < 0, f'' < 0$  OR  $f(a) > 0, f'' > 0$  then  $x_0 \in [a, \tau]$   
e.g.  $x_0 = a$ .
- if  $f(a) < 0, f'' > 0$  OR  $f(a) > 0, f'' < 0$  then  $x_0 \in [\tau, b]$ , e.g.  $x_0 = b$ .

In any case we have the estimate

$$|x_n - \tau| < \frac{f(x_n)}{\min_{[a, b]} |f'(x)|}$$

Problems / violations of assumptions:





### Theorem 27: Convergence of Newton's method

Let  $f: [a, b] \rightarrow \mathbb{R}$  fulfill

(1)  $f(a)f(b) < 0$

(2)  $f$  has no critical pt in  $(a, b)$

(3)  $f''$  exists and is continuous

and  $x_0$  needs to be "close enough" to the true root  $\pi$ . Then, the Newton sequence converges quadratically.

Proof: To show:

$$|x_{n+1} - \pi| \leq M |x_n - \pi|^2 \text{ as } n \rightarrow \infty$$

By Taylor series approximation:

$$f(x) = f(x_n) + f'(x_n)(x - x_n) + \frac{1}{2} f''(\xi_n)(x - x_n)^2$$

Since  $\pi$  is a root we have

$$0 = f(\pi) = f(x_n) + f'(x_n)(\pi - x_n) + \frac{f''(\xi_n)}{2}(\pi - x_n)^2$$

rearranging terms:

$$\frac{f(x_n)}{f'(x_n)} + (\pi - x_n) = - \frac{f''(\xi_n)}{2f'(x_n)} (\pi - x_n)^2$$

$\swarrow$   
 $-x_{n+1}$

$$\Leftrightarrow (r - x_{n+1}) = - \frac{f''(\xi_n)}{2f'(x_n)} (r - x_n)^2$$

$$\Rightarrow |r - x_{n+1}| = \left( \frac{f''(\xi_n)}{2f'(x_n)} \right) |r - x_n|^2$$

$$\text{let } M := \sup_{x_n, \xi_n} \frac{f''(\xi_n)}{2f'(x_n)} \quad \text{then}$$

$$|r - x_{n+1}| \leq M |r - x_n|^2$$

which is what we had to show.

### 3.3 Secant method:

When the conditions for Newton's method are not fulfilled but something more sophisticated than bisection shall be done, we approximate derivatives by secants / difference quotient.

$$f'(x) \approx \frac{f(x+h) - f(x)}{h}$$